## COMENIUS UNIVERSITY IN BRATISLAVA FACULTY OF MATHEMATICS, PHYSICS AND INFORMATICS

# NEW APPROACHES TO NOWHERE-ZERO FLOW PROBLEMS

DIPLOMA THESIS

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Bc. Lukáš Gáborik

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DIPLOMA THESIS

Study Programme: Computer Science Field of Study: Computer Science

Department: Department of Computer Science

Supervisor: Mgr. Jozef Rajník, PhD.

Bratislava, 2026 Bc. Lukáš Gáborik





#### Univerzita Komenského v Bratislave Fakulta matematiky, fyziky a informatiky

#### ZADANIE ZÁVEREČNEJ PRÁCE

Meno a priezvisko študenta:	Bc. Lukáš Gáborik
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**Študijný program:** informatika (Jednoodborové štúdium, magisterský II. st.,

denná forma)

Študijný odbor:informatikaTyp záverečnej práce:diplomováJazyk záverečnej práce:anglickýSekundárny jazyk:slovenský

**Názov:** New approaches to nowhere-zero flow problems

Nové prístupy k problémom o nikde-nulových tokoch

Anotácia: Táto práca nadväzuje na výsledky bakalárskej práce rovnakého autora. Novo

zavedený pojem viacrozmerných Manhattanských a Čebyševovských tokov ponecháva priestor pre ďalší výskum, ako napríklad hľadanie dolných odhadov pre tokové čísla. Jedným z hlavných výsledkov spomínanej bakalárskej práce je predstavenie hypotézy tvrdiacej, že každý bezmostový graf pripúšťa (1,2)-cirkulačnú dekompozíciu, t. j. 2-cirkuláciu a 4-cirkuláciu takú, že zakaždým keď je 2-cirkulácia nulová na nejakej hrane, tak 4-cirkulácia nemôže nadobúdať 0, +1 alebo -1. To ponúka bohatý priestor na výskum vrátane rôznych zovšeobecnení, v ktorých kladieme ďalšie požiadavky na tokové hodnoty.

Ciel':

- 1. Dokázať netriviálne dolné odhady pre dvojrozmerné Čebyševovské tokové číslo grafu.
- 2. Preskúmať možné spôsoby, ako dokázať hypotézu, že každý bezmostový graf pripúšťa (1, 2)-cirkulačnú dekompozíciu, a iné súvisiace hypotézy. Dokázať túto hypotézu pre niektoré nekonečné triedy snarkov, prípadne pre niektoré snarky, ktoré majú ďaleko od toho, aby boli zafarbiteľné (napr. s nepárnosťou
- 2, indexom perfektného párenia 4, ...).
- 3. Preskúmať ďalšie zovšeobecnenia tokov potenciálne užitočné v kontexte vyššie uvedenej hypotézy.

**Vedúci:** Mgr. Jozef Rajník, PhD.

**Katedra:** FMFI.KI - Katedra informatiky **Vedúci katedry:** prof. RNDr. Martin Škoviera, PhD.

**Dátum zadania:** 21.11.2024

**Dátum schválenia:** 05.12.2024 prof. RNDr. Rastislav Kráľovič, PhD.

garant študijného programu

študent	vedúci práce





#### Comenius University Bratislava Faculty of Mathematics, Physics and Informatics

#### THESIS ASSIGNMENT

Name and Surname: Bc. Lukáš Gáborik

**Study programme:** Computer Science (Single degree study, master II. deg., full

time form)

Field of Study: Computer Science Type of Thesis: Diploma Thesis

**Language of Thesis:** English **Secondary language:** Slovak

**Title:** New approaches to nowhere-zero flow problems

**Annotation:** This work builds on the results of the bachelor thesis of the same author. The

newly introduced notion of multidimensional Manhattan and Chebyshev flows still leaves some possibilities for further research like finding lower bounds on flow numbers. One of the main results of the mentioned bachelor thesis is the introduction of the conjecture asserting that each bridgeless graph admits a (1,2)-circulation decomposition, that is a 2-circulation and a 4-circulation such that whenever the 2-circulation is zero on any edge, the 4-circulations can not attain 0, +1 or -1. This offers wide possibilities of exploration including various

generalisations where further requirements are posed on the flow values.

Aim: 1. Prove nontrivial lower bounds on 2-dimensional Chebyshev flow number of

a graph.

2. Explore possible ways of proving the conjecture that each bridgeless graph admits a (1, 2)-circulation decomposition, and other related conjectures. Prove this conjecture for some infinite families of snarks, eventually for some snarks that are far from being colourable (e.g. with oddness 2, perfect matching index

4, ...)

3. Research other generalisation of flows potentially useful in the context of the

abovementioned conjecture.

**Supervisor:** Mgr. Jozef Rajník, PhD.

**Department:** FMFI.KI - Department of Computer Science

**Head of** prof. RNDr. Martin Škoviera, PhD.

department:

**Assigned:** 21.11.2024

**Approved:** 05.12.2024 prof. RNDr. Rastislav Kráľovič, PhD.

Guarantor of Study Programme

Student	Supervisor

 ${\bf Acknowledgments:}$ 

#### Abstrakt

Kľúčové slová:

#### Abstract

Keywords:



### Contents



### Introduction

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### Chapter 1

### Flows

In this chapter, we aim to introduce all concepts and current knowledge about flows and related graph theory terms. Mainly, we tend to explain Chebyshev nowhere-zero flows, which have arisen in the bachelor thesis of the author [TODO].

### Chapter 2

### Bounds on Chebyshev flow number

#### 2.1 Upper bound on generalised Blanuša snarks

#### 2.2 General lower bound

**Proposition 2.1.** Let  $\Gamma$  denote a snark of order n. Then  $\Phi_2^{\infty}(\Gamma) \geq 2 + 1/\xi$ , where  $\xi = \lfloor \frac{n-2}{4} \rfloor$ .

*Proof.* Assume by contradiction that there exists a 2-dimensional flow  $\varphi$  of  $\Gamma$  such that  $\varphi = (\varphi_1(e), \varphi_2(e))$  for each edge  $e \in E(\Gamma)$ , with  $\|\varphi(e)\|_{\infty} \geq 1$  and  $\varphi_i(e) \in (-1 - 1/\xi, 1 + 1/\xi)$  for i = 1, 2.

We say an edge  $e \in E(\Gamma)$  to be *good* with respect to  $\varphi_i$  if  $|\varphi_i(e)| \in [1, 1 + 1/\xi)$ , bad otherwise. Observe that an edge e can be good with respect to both  $\varphi_1$  and  $\varphi_2$ , but it cannot be bad with respect to both  $\varphi_1$  and  $\varphi_2$ , for otherwise  $||\varphi(e)||_{\infty} < 1$ .

Denote by  $B_i$  the subgraph of  $\Gamma$  induced by the bad edges with respect to  $\varphi_i$  and by  $G_i$  the one induced by the good edges with respect to  $\varphi_i$ , i = 1, 2. By previous observation at least one between  $B_1$  and  $B_2$  has at most  $\lfloor |E(\Gamma)|/2 \rfloor = \lfloor \frac{3n}{4} \rfloor$  edges, say  $B_1$ .

 $B_i$  is a spanning subgraph of  $\Gamma$  and  $\Delta(B_i) \leq 2$ , for i = 1, 2.

Proof. Observe that  $\Delta(G_i) \leq 2$ , because the sum of three real numbers all with absolute value in the interval  $[1, 1 + 1/\xi)$  cannot give 0 as a result, making the Kirkoff's law impossible to be satisfied by  $\varphi$  around a vertex of  $\Gamma$ . Hence  $B_i$  is spanning, for otherwise  $\Delta(G_i) = 3$  and  $\Delta(B_i) \leq 2$ , for otherwise  $\Delta(G_{3-i}) = 3$ .

If  $C \subseteq E(G_i)$  is an odd edge-cut of  $\Gamma$  for i = 1, 2, then  $|C| \ge 2\xi + 3$ .

*Proof.* Consider an odd edge-cut C of  $\Gamma$  containing 2k+1 good edges, separating components  $\Gamma_1, \Gamma_2$ . WLOG assume  $\varphi_i(e) \geq 0$  for every edge e of C and more edges are directed towards  $\Gamma_2$ . Then, there are at least k+1 edges towards  $\Gamma_2$ , resulting

in total inflow at least  $(k+1) \cdot 1$ . Analogously, the total outflow is strictly less than  $k \cdot (1+1/\xi)$ . Together with the Kirkoff's law, this leads to  $k+1 < k \cdot (1+1/\xi)$ , which is equivalent to  $k > \xi$ .

A path of length 2k cannot be a connected component of  $B_i$ , for  $k = 1, 2, ..., \xi - 1$  and i = 1, 2.

*Proof.* For the sake of contradiction assume there is a path of a length 2k,  $k < \xi$  in  $B_i$ . Then the edges adjacent to this path are in  $G_i$  and they form an odd edge-cut of  $\Gamma$ , containing at most  $2k + 3 < 2\xi + 3$  edges, which is in contradiction with the Claim 2.

A cycle of length 2k+1 cannot be a connected component of  $B_i$ , for  $k=1,2,\ldots,\xi$  and i=1,2.

*Proof.* For the sake of contradiction assume there is a cycle of a length 2k+1,  $k < \xi + 1$  in  $B_i$ . Then the edges adjacent to this cycle are in  $G_i$  and they form an odd edge-cut of  $\Gamma$ , containing at most  $2k+1 < 2\xi + 3$  edges, which is in contradiction with the Claim 2.

 $E(B_i)$  cannot contain a perfect matching of  $\Gamma$ , for i = 1, 2.

Proof. For the sake of contradiction assume that  $E(B_i)$  contains a perfect matching M of  $\Gamma$ . Then also  $E(G_{3-i})$  contains a perfect matching M of  $\Gamma$ . Next,  $F = E(\Gamma) \setminus M$  is a 2-factor of  $\Gamma$ . Note that F must contain an odd cycle of length  $2k + 1 \leq \frac{n}{2}$ . This is equivalent with  $k \leq \xi$ . Then the edges adjacent to this cycle are in  $G_{3-i}$  and they form an odd edge-cut of  $\Gamma$ , containing at most  $2k + 1 \leq 2\xi + 1$  edges, which is in contradiction with the Claim 2.

Note that by the Claim 3, each even path in  $B_1$  contains at least  $2\xi$  edges. Similarly by the Claim 4, each odd cycle in  $B_1$  contains at least  $2\xi + 3$  edges. Let *odd components* denote odd cycles and even paths. Since  $4\xi = 4\lfloor \frac{n-2}{4} \rfloor > \lfloor \frac{3n}{4} \rfloor \geq E(B_1)$  obviously holds for any  $n \geq 10$ ,  $B_1$  may contain at most one odd component. On the other hand,  $B_1$  contains even number of odd componentss. As a result,  $B_1$  contains only odd paths and even cycles, but then also a perfect matching of  $\Gamma$ , which is in contradiction with the Claim 5.

### Conclusion

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