

COMENIUS UNIVERSITY IN BRATISLAVA  
FACULTY OF MATHEMATICS, PHYSICS AND INFORMATICS

NEW APPROACHES TO NOWHERE-ZERO FLOW  
PROBLEMS  
DIPLOMA THESIS

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BC. LUKÁŠ GÁBORIK



COMENIUS UNIVERSITY IN BRATISLAVA  
FACULTY OF MATHEMATICS, PHYSICS AND INFORMATICS

NEW APPROACHES TO NOWHERE-ZERO FLOW  
PROBLEMS

DIPLOMA THESIS

Study Programme: Computer Science  
Field of Study: Computer Science  
Department: Department of Computer Science  
Supervisor: Mgr. Jozef Rajník, PhD.

Bratislava, 2026  
Bc. Lukáš Gáborik





Univerzita Komenského v Bratislave  
Fakulta matematiky, fyziky a informatiky

## ZADANIE ZÁVEREČNEJ PRÁCE

**Meno a priezvisko študenta:** Bc. Lukáš Gáborik  
**Študijný program:** informatika (Jednoodborové štúdium, magisterský II. st., denná forma)  
**Študijný odbor:** informatika  
**Typ záverečnej práce:** diplomová  
**Jazyk záverečnej práce:** anglický  
**Sekundárny jazyk:** slovenský

**Názov:** New approaches to nowhere-zero flow problems  
*Nové prístupy k problémom o nikde-nulových tokoch*

**Anotácia:** Táto práca nadväzuje na výsledky bakalárskej práce rovnakého autora. Novo zavedený pojem viacrozmerých Manhattanských a Čebyševovských tokov ponecháva priestor pre ďalší výskum, ako napríklad hľadanie dolných odhadov pre tokové čísla. Jedným z hlavných výsledkov spomínanej bakalárskej práce je predstavenie hypotézy tvrdiacej, že každý bezmostový graf pripúšťa (1,2)-cirkulačnú dekompozíciu, t. j. 2-cirkuláciu a 4-cirkuláciu takú, že zakaždým keď je 2-cirkulácia nulová na nejakej hrane, tak 4-cirkulácia nemôže nadobúdať 0, +1 alebo -1. To ponúka bohatý priestor na výskum vrátane rôznych zovšeobecnení, v ktorých kladieme ďalšie požiadavky na tokové hodnoty.

**Cieľ:**

1. Dokázať netriviálne dolné odhady pre dvojrozmerné Čebyševovské tokové číslo grafu.
2. Preskúmať možné spôsoby, ako dokázať hypotézu, že každý bezmostový graf pripúšťa (1, 2)-cirkulačnú dekompozíciu, a iné súvisiace hypotézy. Dokázať túto hypotézu pre niektoré nekonečné triedy snarkov, prípadne pre niektoré snarky, ktoré majú ďaleko od toho, aby boli zafarbitelné (napr. s nepárnosťou 2, indexom perfektného párenia 4, ...).
3. Preskúmať ďalšie zovšeobecnenia tokov potenciálne užitočné v kontexte vyššie uvedenej hypotézy.

**Vedúci:** Mgr. Jozef Rajník, PhD.  
**Katedra:** FMFI.KI - Katedra informatiky  
**Vedúci katedry:** prof. RNDr. Martin Škoviera, PhD.  
**Dátum zadania:** 21.11.2024

**Dátum schválenia:** 05.12.2024

prof. RNDr. Rastislav Kráľovič, PhD.  
garant študijného programu

.....  
študent

.....  
vedúci práce



## THESIS ASSIGNMENT

**Name and Surname:** Bc. Lukáš Gáborik  
**Study programme:** Computer Science (Single degree study, master II. deg., full time form)  
**Field of Study:** Computer Science  
**Type of Thesis:** Diploma Thesis  
**Language of Thesis:** English  
**Secondary language:** Slovak

**Title:** New approaches to nowhere-zero flow problems

**Annotation:** This work builds on the results of the bachelor thesis of the same author. The newly introduced notion of multidimensional Manhattan and Chebyshev flows still leaves some possibilities for further research like finding lower bounds on flow numbers. One of the main results of the mentioned bachelor thesis is the introduction of the conjecture asserting that each bridgeless graph admits a  $(1,2)$ -circulation decomposition, that is a 2-circulation and a 4-circulation such that whenever the 2-circulation is zero on any edge, the 4-circulations can not attain 0, +1 or -1. This offers wide possibilities of exploration including various generalisations where further requirements are posed on the flow values.

**Aim:**

1. Prove nontrivial lower bounds on 2-dimensional Chebyshev flow number of a graph.
2. Explore possible ways of proving the conjecture that each bridgeless graph admits a  $(1, 2)$ -circulation decomposition, and other related conjectures. Prove this conjecture for some infinite families of snarks, eventually for some snarks that are far from being colourable (e.g. with oddness 2, perfect matching index 4, ...)
3. Research other generalisation of flows potentially useful in the context of the abovementioned conjecture.

**Supervisor:** Mgr. Jozef Rajník, PhD.  
**Department:** FMFI.KI - Department of Computer Science  
**Head of department:** prof. RNDr. Martin Škoviera, PhD.

**Assigned:** 21.11.2024

**Approved:** 05.12.2024

prof. RNDr. Rastislav Kráľovič, PhD.  
Guarantor of Study Programme

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Student

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Supervisor

**Acknowledgments:**

# Abstrakt

Klíčové slová:



# Abstract

Keywords:



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# Introduction



# Chapter 1

## Flows

In this chapter, we aim to introduce all concepts and current knowledge about flows and related graph theory terms. Mainly, we explain Chebyshev nowhere-zero flows, which have arisen in the Student Science Conference paper of the author [3]. Moreover, we present a concept of two smaller flows as a sufficient condition for a graph to have a larger nowhere-zero flow and some Chebyshev flow, also stated in that paper. Then we connect this idea to the theory of group connectivity of a graph established by Jaeger et al. [4] and  $k$ -closures and  $k$ -bases introduced by Seymour [6] in the paper, where he proved the 6-flow theorem.

### 1.1 Integral nowhere-zero flows

In this thesis, we consider only unoriented unweighted simple graphs  $\Gamma = (V, E)$ . By a *dart* we represent an edge with an additional orientation. This means, for any edge  $e = uv \in E$  we have two darts  $\vec{uv}$ ,  $\vec{vu}$  belonging to that edge. By a *dart-set* of  $S \subseteq V$  we mean the set of darts with initial vertex in the set  $S$ .

**Definition.** Let  $\Gamma = (V, E)$  be a graph and  $k \geq 2$  be an integer. A  $k$ -flow on  $\Gamma$  is an assignment  $\varphi: D(V) \rightarrow \{-(k-1), -(k-2), \dots, k-2, k-1\}$  of flow values to the darts satisfying  $\varphi(\vec{uv}) = -\varphi(\vec{vu})$  for any dart  $\vec{uv} \in D(V)$  and

$$\sum_{\vec{d} \in D(v)} \varphi(\vec{d}) = 0 \tag{1.1}$$

for each vertex  $v$  of  $G$ . The equality (1.1) is often referred to as the *flow conservation constraint*. Moreover, a *nowhere-zero  $k$ -flow* (or NZ  $k$ -flow) is a  $k$ -flow such that no flow value is zero, i.e.  $\varphi(\vec{d}) \neq 0$  for any  $\vec{d} \in D(V)$ .

We also establish a notation  $\varphi(e)$  for  $e = uv \in E$ , where  $\varphi(e) = \max\{\varphi(\vec{uv}), \varphi(\vec{vu})\}$ .

The flow conservation constraint (1.1) is generalisable to any subset of vertices. Hence, having a graph  $\Gamma$  with a bridge  $e$ , the interpretation of the flow conservation constraint for one of the components of  $\Gamma - e$  results in the following observation.

*Claim.* If  $e \in E$  is a bridge in  $\Gamma$ , then  $\varphi(e) = 0$ . Moreover,  $\Gamma$  allows a NZ  $k$ -flow for no natural  $k$ .

Therefore, NZ flows are examined only on bridgeless graphs.

We can clearly see that graph with an NZ  $k$ -flow also admits an NZ  $(k + 1)$ -flow. That means, graphs with an NZ 2-flow are somehow the “simplest” ones and the existence of NZ flows relates with some complexity of a graph. So we define a flow number, which is a rate of this complexity.

**Definition.** Let  $\Gamma$  be a bridgeless graph. A *flow number* of  $\Gamma$  is

$$\Phi(\Gamma) := \min\{k \mid \exists \text{ NZ } k\text{-flow on } \Gamma\}.$$

However, it turns out that this rate of complexity achieves only few values.

**Theorem 1.1** (6-flow theorem). [6, p. 133] *There exists an NZ 6-flow on any bridgeless graph.*

Seymour constructed this NZ 6-flow from two simpler flows from the fact, that if  $a \in \{0, \pm 1\}$ ,  $b \in \{0, \pm 1, \pm 2\}$  and not both  $a, b$  are zero, then  $3a + b \in \{\pm 1, \pm 2, \pm 3, \pm 4, \pm 5\}$ .

**Lemma 1.2.** [6, p. 132] *For any bridgeless graph, there exist a 2-flow  $\varphi_2$  and a 3-flow  $\varphi_3$  such that for each edge  $e$ , at least one value from  $\varphi_2(e), \varphi_3(e)$  is non-zero.*

On the other hand, there is no known graph with  $\Phi(\Gamma) = 6$ , all known graphs have  $\Phi(\Gamma) \leq 5$ . Therefore, the conjecture about the NZ 5-flow by Tutte still holds unresolved.

**Conjecture 1.3** (5-flow conjecture). [8, p. 83] *There exists an NZ 5-flow on any bridgeless graph.*

As for many problems in the graph theory, it is sufficient to explore the problem only on the smaller subset of graphs. There exists a folklore reduction of nowhere-zero flows to 3-regular (or cubic) graphs.

**Proposition 1.4.** *There exists an NZ  $k$ -flow on any bridgeless graph if and only if there exists an NZ  $k$ -flow on any cubic bridgeless graph.*

Hence, by considering only cubic graphs, we still solve the universal problem. Moreover, in terms of finding a counterexample to Tutte’s 5-flow conjecture, we can also avoid 3-edge-colourable ones and limit research to the cubic graphs, that are not 3-edge-colourable, as a corollary of the next proposition.

**Proposition 1.5.** [2, pp. 160, 161] *A cubic graph has an NZ 4-flow if and only if it is 3-edge-colourable.*



The class of type II cubic graphs is widely used, so its members have special name.

**Definition.** A *snark* is a bridgeless cubic graph, which has no cycles of lengths 3 and 4, is cyclically 4-edge connected (at least 4 edges must be removed from the graph to get two components containing a cycle) and is not 3-edge-colourable.

## 1.2 Chebyshev flows and sufficient flow-pairs

There were some attempts to generalise nowhere-zero flows to real numbers and more dimensions. By generalising integer nowhere-zero flows to real numbers, researchers found out, that restricting only zero flow value is not sufficient, since any flow consisting of non-zero flow values can be made “smaller” and then, the notion of flow number disappears. Therefore, we need to determine a larger interval of real values forbidden as flow values, and hence, one-dimensional real nowhere-zero flows are allowed to use flow values from the interval  $[1, r - 1]$ . Then the first attempts of generalisation real flows to more dimensions [TODO] tried to pick only vectors with the Euclidean norm from the mentioned interval  $[1, r - 1]$ . It worked, but it did not have such nice properties, as we would expect. However, changing a norm turned out to be a good idea [3], so we will consider only Chebyshev multidimensional nowhere-zero flows.

**Definition.** Let  $\Gamma = (V, E)$  be a bridgeless graph,  $d$  be a positive integer and  $r \geq 2$  be a real constant. A  *$d$ -dimensional Chebyshev nowhere-zero  $r$ -flow* (or  $dD$  ChNZ  $r$ -flow) satisfies the conservation constraint (1.1) and the condition  $1 \leq \|\varphi(e)\|_\infty \leq r - 1$  for each edge  $e \in E$ .

The flow number can be defined in a similar way to integral nowhere-zero flows. The only problem is, that the set of real numbers does not need to contain a minimal element, so formally correct definition would use infimum instead of minimum. However, the infimum of the flow number candidates is always feasible [TODO], so we can define the flow number straightforward with minimum.

**Definition.** Let  $\Gamma$  be a bridgeless graph and  $d$  be a positive integer. A  *$d$ -dimensional Chebyshev flow number* of  $\Gamma$  is

$$\Phi_d^\infty(\Gamma) := \min\{r \mid \exists dD \text{ ChNZ } r\text{-flow on } \Gamma\}.$$

Now we recall the Seymour’s flows from Lemma ?? – a 2-flow and a 3-flow with a property, that no edge has both flow values equal to zero. By simply using the value of the 2-flow as the first coordinate and the value of the 3-flow as the second coordinate, we get a Chebyshev nowhere-zero flow. Moreover, this flow gives us a nice bound on the Chebyshev flow number.

**Proposition 1.6.** [3, p. 343] *Each bridgeless graph has a 2D ChNZ 3-flow and hence,  $\Phi_2^\infty(\Gamma) \leq 3$ .*

Again, no graph with  $\Phi_2^\infty(\Gamma) = 3$  is known. Furthermore, there is no known graph with  $\Phi_2^\infty(\Gamma) > 5/2$ . Therefore, the next conjecture is interesting to examine.

**Conjecture 1.7.** [3, p. 344] *Each bridgeless graph has a 2D ChNZ  $5/2$ -flow and hence,  $\Phi_2^\infty(\Gamma) \leq 5/2$ .*

And similarly as for integral flows, the problem can be reduced to cubic graphs and then to snarks, since 3-edge-colourable cubic graphs have small flow number.

**Theorem 1.8.** [3, p. 343] *A cubic graph has a 2D ChNZ 2-flow if and only if it is 3-edge-colourable.*

As an NZ 6-flow and a 2D ChNZ 3-flow can be constructed from two simpler flows, something similar can be made for smaller flow numbers. However, the existence of two simpler flows is only a sufficient condition to the existence of the original flow, the condition is not necessary. Hence we will use a term  *$p/q$ -sufficient flow-pair* instead of a former  *$(p, q)$ -circulation decomposition*.

**Definition.** Let  $p \leq q$  be the positive integers. A  *$p/q$ -sufficient flow-pair* of a graph  $\Gamma$  is a 2-flow  $\varphi_2$  and a  $(p + q + 1)$ -flow  $\varphi_{p+q+1}$  such that whenever  $\varphi_2(e)$  is zero, the value  $\varphi_{p+q+1}(e)$  is at least  $q$ .

One possible ambiguity of this definition is, that since  $p/q$  is written as a fraction, the  $2/6$ -sufficient flow-pair can be mixed up with the  $1/3$ -sufficient flow-pair. However, this is not a bug, this is a feature, since we can WLOG assume  $p, q$  relatively prime.

**Lemma 1.9.** *A bridgeless graph  $\Gamma$  with a  $dp/dq$ -sufficient flow-pair also has a  $p/q$ -sufficient flow-pair and vice versa.*

We mentioned, that the  $p/q$ -sufficient flow-pair is a sufficient condition to the existence of some nowhere-zero flows. Those flows are somehow related with the value  $p/q$  and that is the reason to write it as a fraction in the definition of the sufficient flow-pair,

**Proposition 1.10.** [3, p. 344] *Consider a bridgeless graph  $\Gamma$  with a  $p/q$ -sufficient flow-pair. Then, there are a 2D ChNZ  $(2 + p/q)$ -flow and an NZ  $(4 + 2p/q)$ -flow on  $\Gamma$ .*

Note that if there is a graph, which is a counterexample to Tutte's 5-flow conjecture, it should not have the  $1/2$ -sufficient flow-pair. Hence it would be interesting to classify graphs with  $1/2$ -sufficient flow-pair. By a computer, it was checked that this flow-pair exists for snarks up to roughly 30 vertices, so it looks like the  $1/2$ -sufficient flow-pair exists always.

**Conjecture 1.11.** *Each bridgeless graph has a  $1/2$ -sufficient flow-pair.*

### 1.3 Closures and bases

**Definition.** [6, p. 132] Let  $\Gamma = (V, E)$  be a graph and  $S \subseteq E$  its edge subset. A  $k$ -closure  $\langle S \rangle_k$  of  $S$  is a minimal  $T$  such that  $S \subseteq T \subseteq E$  and for any circuit  $\mathcal{C} \not\subseteq T$  of  $G$ , the size of  $\mathcal{C} \cap T^c$  is strictly greater than  $k$ .

**Definition.** [1, p. 7] Let  $\Gamma = (V, E)$  be a graph. Then  $S \subseteq E$  is a  $k$ -base of  $\Gamma$  if  $\langle S \rangle_k = E$ .

*Claim.* Any spanning tree of a graph is its inclusion-minimal 1-base.

**Lemma 1.12.** [6, p. 133] *For any 3-connected graph  $\Gamma$  there exists its even-factor which is also its 2-base.*

**Lemma 1.13.** [6, p. 134] *For any 3-connected cubic graph  $\Gamma = (V, E)$  there exists a vertex partition  $E = E_1 \cup E_2$  such that  $E_1, E_2$  are its 1- and 2-bases, respectively.*

### 1.4 Group connectivity



## Chapter 2

### Multipoles generation



# Chapter 3

## Bounds on Chebyshev flow number

### 3.1 Upper bound on generalised Blanuša snarks

### 3.2 General lower bound

In one-dimensional flows, there has been proved a lower bound on the flow number of a snark

$$\Phi_1(\Gamma) \geq 4 + 1 \left\lfloor \frac{n-4}{8} \right\rfloor$$

depending on the number  $n$  of its vertices [5, p. 14]. Moreover, there is also a meaningful lower bound on the two-dimensional Chebyshev flow number of the Petersen graph  $P$  [7, p. 99]. It states that  $\Phi_2^\infty(P) \geq 5/2$ . We provide generalisation of this bound for any snark. Moreover, the bounding number is roughly a half of the bounding number in one dimension.

**Proposition 3.1.** *Let  $\Gamma$  denote a snark of order  $n$ . Then  $\Phi_2^\infty(\Gamma) \geq 2 + 1/\xi$ , where  $\xi = \lfloor \frac{n-2}{4} \rfloor$ .*

*Proof.* Assume by contradiction that there exists a 2-dimensional flow  $\varphi$  of  $\Gamma$  such that  $\varphi = (\varphi_1(e), \varphi_2(e))$  for each edge  $e \in E(\Gamma)$ , with  $\|\varphi(e)\|_\infty \geq 1$  and  $\varphi_i(e) \in (-1 - 1/\xi, 1 + 1/\xi)$  for  $i = 1, 2$ .

We say an edge  $e \in E(\Gamma)$  to be *good* with respect to  $\varphi_i$  if  $|\varphi_i(e)| \in [1, 1 + 1/\xi)$ , *bad* otherwise. Observe that an edge  $e$  can be good with respect to both  $\varphi_1$  and  $\varphi_2$ , but it cannot be bad with respect to both  $\varphi_1$  and  $\varphi_2$ , for otherwise  $\|\varphi(e)\|_\infty < 1$ .

Denote by  $B_i$  the subgraph of  $\Gamma$  induced by the bad edges with respect to  $\varphi_i$  and by  $G_i$  the one induced by the good edges with respect to  $\varphi_i$ ,  $i = 1, 2$ . By previous observation at least one between  $B_1$  and  $B_2$  has at most  $\lfloor |E(\Gamma)|/2 \rfloor = \lfloor \frac{3n}{4} \rfloor$  edges, say  $B_1$ .

*Claim.*  $B_i$  is a spanning subgraph of  $\Gamma$  and  $\Delta(B_i) \leq 2$ , for  $i = 1, 2$ .

*Proof.* Observe that  $\Delta(G_i) \leq 2$ , because the sum of three real numbers all with absolute value in the interval  $[1, 1 + 1/\xi)$  cannot give 0 as a result, making the Kirkoff's law impossible to be satisfied by  $\varphi$  around a vertex of  $\Gamma$ . Hence  $B_i$  is spanning, for otherwise  $\Delta(G_i) = 3$  and  $\Delta(B_i) \leq 2$ , for otherwise  $\Delta(G_{3-i}) = 3$ .  $\square$

*Claim.* If  $C \subseteq E(G_i)$  is an odd edge-cut of  $\Gamma$  for  $i = 1, 2$ , then  $|C| \geq 2\xi + 3$ .

*Proof.* Consider an odd edge-cut  $C$  of  $\Gamma$  containing  $2k + 1$  good edges, separating components  $\Gamma_1, \Gamma_2$ . WLOG assume  $\varphi_i(e) \geq 0$  for every edge  $e$  of  $C$  and more edges are directed towards  $\Gamma_2$ . Then, there are at least  $k + 1$  edges towards  $\Gamma_2$ , resulting in total inflow at least  $(k + 1) \cdot 1$ . Analogously, the total outflow is strictly less than  $k \cdot (1 + 1/\xi)$ . Together with the Kirkoff's law, this leads to  $k + 1 < k \cdot (1 + 1/\xi)$ , which is equivalent to  $k > \xi$ .  $\square$

*Claim.* A path of length  $2k$  cannot be a connected component of  $B_i$ , for  $k = 1, 2, \dots, \xi - 1$  and  $i = 1, 2$ .

*Proof.* For the sake of contradiction assume there is a path of a length  $2k$ ,  $k < \xi$  in  $B_i$ . Then the edges adjacent to this path are in  $G_i$  and they form an odd edge-cut of  $\Gamma$ , containing at most  $2k + 3 < 2\xi + 3$  edges, which is in contradiction with the Claim 2.  $\square$

*Claim.* A cycle of length  $2k + 1$  cannot be a connected component of  $B_i$ , for  $k = 1, 2, \dots, \xi$  and  $i = 1, 2$ .

*Proof.* For the sake of contradiction assume there is a cycle of a length  $2k + 1$ ,  $k < \xi + 1$  in  $B_i$ . Then the edges adjacent to this cycle are in  $G_i$  and they form an odd edge-cut of  $\Gamma$ , containing at most  $2k + 1 < 2\xi + 3$  edges, which is in contradiction with the Claim 2.  $\square$

*Claim.*  $E(B_i)$  cannot contain a perfect matching of  $\Gamma$ , for  $i = 1, 2$ .

*Proof.* For the sake of contradiction assume that  $E(B_i)$  contains a perfect matching  $M$  of  $\Gamma$ . Then also  $E(G_{3-i})$  contains a perfect matching  $M$  of  $\Gamma$ . Next,  $F = E(\Gamma) \setminus M$  is a 2-factor of  $\Gamma$ . Note that  $F$  must contain an odd cycle of length  $2k + 1 \leq \frac{n}{2}$ . This is equivalent with  $k \leq \xi$ . Then the edges adjacent to this cycle are in  $G_{3-i}$  and they form an odd edge-cut of  $\Gamma$ , containing at most  $2k + 1 \leq 2\xi + 1$  edges, which is in contradiction with the Claim 2.  $\square$

Note that by the Claim 3, each even path in  $B_1$  contains at least  $2\xi$  edges. Similarly by the Claim 4, each odd cycle in  $B_1$  contains at least  $2\xi + 3$  edges. Let *odd components* denote odd cycles and even paths. Since  $4\xi = 4\lfloor \frac{n-2}{4} \rfloor > \lfloor \frac{3n}{4} \rfloor \geq E(B_1)$  obviously holds for any  $n \geq 10$ ,  $B_1$  may contain at most one odd component. On the other hand,  $B_1$



contains even number of odd componentss. As a result,  $B_1$  contains only odd paths and even cycles, but then also a perfect matching of  $\Gamma$ , which is in contradiction with the Claim 5.  $\square$



# Conclusion



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