

**Exercise 1** (Proving with tactics).

Prove the following formulas using basic tactics for logical connectives. Do not use automation (`auto`, `tauto` or `firstorder`).

1.  $(P \rightarrow Q) \rightarrow \neg Q \rightarrow \neg P.$
2.  $P \rightarrow (P \rightarrow Q) \rightarrow Q.$
3.  $P \rightarrow \neg\neg P.$
4.  $P \vee \neg P \rightarrow \neg\neg P \rightarrow P.$
5.  $\neg\neg(P \rightarrow Q) \rightarrow \neg\neg P \rightarrow \neg\neg Q.$
6.  $(\neg\neg P \rightarrow \neg\neg Q) \rightarrow \neg\neg(P \rightarrow Q).$
7.  $\neg(P \wedge \neg P).$
8.  $(P \vee Q \rightarrow R) \rightarrow (P \rightarrow R) \wedge (Q \rightarrow R).$
9.  $(\exists x.P \wedge Rx) \leftrightarrow P \wedge \exists x.Rx.$
10.  $P \vee (\forall x.Rx) \rightarrow \forall x.P \vee Rx.$
11.  $(\exists x.Rx) \rightarrow \neg(\forall x.\neg Rx).$
12.  $\neg(P \vee Q) \leftrightarrow \neg P \wedge \neg Q.$

**Exercise 2** (Induction on natural numbers).

Prove by induction the following facts about addition. Do not use `lia` or `omega`.

*Hint.* Use the `auto` tactic. It knows some basic facts about natural numbers.

1. Commutativity: `forall x y, x + y = y + x.`
2. Associativity: `forall x y z, (x + y) + z = x + (y + z).`
3. Distributivity of multiplication: `forall x y z, x * (y + z) = x * y + x * z.`

*Hint.* Use the previous two points and the tactic `replace t with t'` which replaces `t` with `t'` in the goal.

**Exercise 3** (Induction on lists).

Consider the following function which reverses a list.

```
Fixpoint reverse {A} (l : list A) :=
  match l with
  | [] => []
  | h :: t => reverse t ++ [h]
  end.
```

This is *not* an efficient implementation of list reversal (quadratic complexity), but it is easier to reason with than the linear implementation from the first lecture.

Prove by induction the following facts about `reverse`.

1. `forall` `l1 l2 : list A, reverse (l1 ++ l2) = reverse l2 ++ reverse l1.`

*Hint.* Use `SearchRewrite` to find helper lemmas in Coq's standard library.

2. `forall` `l : list A, reverse (reverse l) = l.`

3. `forall` `l : list A, reverse l = List.rev l.`

Is it possible to prove `reverse = List.rev`?