

CS580 Homework 9
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December 9, 2022

Question 1.

a We can first observe the following:

- (a) set $A = A \cap B + A \setminus B$.
- (b) set $B = A \cap B + B \setminus A$.
- (c) $A \cup B = A \cap B + A \setminus B + B \setminus A$.

If we define the function $f'(S, T)$ to be the number of edges from S to T and S and T are disjoint, we can see that

- (a) $f(A) = f'(A, V \setminus A)$
- (b) $f(B) = f'(B, V \setminus B)$
- (c) $f(A \cap B) = f'(A \cap B, V \setminus (A \cap B))$
- (d) $f(A \cup B) = f'(A \cup B, V \setminus (A \cup B))$

We break down each of the above to prove the function f is submodular.

$$f'(A, V \setminus A) = f'(A \cap B, V \setminus (A \cup B)) + f'(A \cap B, B \setminus A) + f'(A \setminus B, B \setminus A) + f'(A \setminus B, V \setminus (A \cup B))$$

$$f'(B, V \setminus B) = f'(A \cap B, V \setminus (A \cup B)) + f'(A \cap B, A \setminus B) + f'(B \setminus A, A \setminus B) + f'(B \setminus A, V \setminus (A \cup B))$$

$$f'(A \cup B, V \setminus (A \cup B)) = f'(A \cap B, V \setminus (A \cup B)) + f'(A \setminus B, V \setminus (A \cup B)) + f'(B \setminus A, V \setminus (A \cup B))$$

$$f'(A \cap B, V \setminus (A \cap B)) = f'(A \cap B, V \setminus (A \cup B)) + f'(A \cap B, A \setminus B) + f'(A \cap B, B \setminus A)$$

From the above, we can see that $f(A) + f(B) - f(A \cup B) - f(A \cap B) = f'(A, V \setminus A) + f'(B, V \setminus B) - f'(A \cap B, V \setminus (A \cap B)) - f'(A \cup B, V \setminus (A \cup B)) = f'(A \setminus B, B \setminus A) + f'(B \setminus A, A \setminus B) \geq 0$.

Hence, $f(A) + f(B) \geq f(A \cup B) + f(A \cap B)$ and f is submodular

b We prove both directions

(a) if G is Eulerian, f is symmetric:

if G is Eulerian, we know every vertex has the same number of outgoing edges as incoming edges. We prove f is symmetric using induction on the size of S

Base case: when $|S| = 1$, meaning there is only 1 vertex in the set, $f(S)$ is the number of outgoing edges and $f(V \setminus S)$ is the number of edges going from $V \setminus S$ to S , which is also the number of incoming edges of S . Since G is Eulerian, $f(S) = f(V \setminus S)$

Induction Hypothesis: when $|S| = k \leq |V| - 1$, $f(S) = f(V \setminus S)$

Inductive Step: when $|S| = k + 1$, we added a new vertex v to the previous set S' . We define the function $f'(S, T)$ to be the number of edges from S to T and S and T are disjoint. Hence $f(S) = f(S') + f'(v, V \setminus S) - f'(S', v)$ and $f(V \setminus S) = f(V \setminus S') + f'(V \setminus S, v) - f'(v, S')$. Since $f'(v, V \setminus S) - f'(S', v) = f'(V \setminus S, v) - f'(v, S')$, we can conclude $f(S) = f(V \setminus S)$. Hence, f is symmetric

(b) if f is symmetric, G is Eulerian If f is symmetric, we know $f(S) = f(V \setminus S)$. If $|S| = 1$, S contains only one vertex v and $f(S)$ is the number of outgoing edges of v and $f(V \setminus S)$ is the number of incoming edges to v . Hence, the number of outgoing edges is the same as the number of incoming edges for a single vertex. This is true for all vertices, hence G is Eulerian

c From question 2, we know if G is Eulerian, we have $f(A) = f(V \setminus A)$ and $f(B) = f(V \setminus B)$. Since $f(V \setminus A) \geq f(B \setminus A)$ and $f(V \setminus B) \geq f(A \setminus B)$, we know $f(A) + f(B) \geq f(B \setminus A) + f(A \setminus B)$

Question 2.

Suppose (S, T) is a valid minimum (s, t) -cut of G , where $s \in S$ and $t \in T$. Then u either belongs to S or T . Suppose $u \in S$. In this case, (S, T) is also a minimum cut for (u, t) . Therefore $\lambda(s, t) = \lambda(u, t)$. Suppose $\lambda(s, u) > \lambda(u, t)$, then the goal becomes to prove $\lambda(s, t) \geq \lambda(u, t)$, which is true because $\lambda(s, t) = \lambda(u, t)$. Suppose $\lambda(s, u) \leq \lambda(u, t)$, then the goal becomes to prove $\lambda(s, t) \geq \lambda(s, u)$. Since $\lambda(s, t) = \lambda(u, t)$, the goal is equivalent to proving $\lambda(u, t) \geq \lambda(s, u)$, which is the hypothesis we assumed for this case. Therefore, $\lambda(s, t) \geq \min \lambda(s, u), \lambda(u, t)$. The same proof logic can be used to prove the $u \in T$ case.