

1/18

## Review

\* Markov decision process

$$\text{MDP} = \langle \mathcal{S}, \mathcal{A}, P(s'|s, a), R(s) \rangle$$

\* Policy = deterministic mapping  $\pi: \mathcal{S} \rightarrow \mathcal{A}$ 

\* Value functions

$$V^\pi(s) = E^\pi \left[ \sum_{t=0}^{\infty} \gamma^t R(s_t) \mid S_0 = s \right]$$

$$Q^\pi(s, a) = E^\pi \left[ \sum_{t=0}^{\infty} \gamma^t R(s_t) \mid S_0 = s, a_0 = a \right]$$

\* Policy evaluation

$$\text{Solve linear equations: } \sum_{s'} [I(s, s') - \gamma P(s'|s, \pi(s))] V^\pi(s') = R(s)$$

\* Policy improvement

$$\text{Greedy policy } \pi'(s) = \underset{a}{\operatorname{argmax}} Q^\pi(s, a)$$

$$\text{Thm: } V^{\pi'}(s) \geq V^\pi(s) \text{ for all states } s$$

\* Policy iteration

$$\pi_0 \xrightarrow{\text{evaluate}} \frac{V^{\pi_0}(s)}{Q^{\pi_0}(s)} \xrightarrow{\text{improve}} \pi_1 \longrightarrow \dots$$

$$\text{Thm: if } \pi'(s) = \underset{a}{\operatorname{argmax}} Q^\pi(s, a) \text{ and } V^{\pi'}(s) = V^\pi(s) \text{ for all } s, \\ \text{then } V^\pi(s) = V^*(s).$$

Proof: 1) From last time:

$$V^\pi(s) = R(s) + \gamma \max_a \sum_{s'} P(s'|s, a) V^\pi(s') : \text{Bellman optimality equation}$$

2) Iterate RHS:

$$V^\pi(s) = R(s) + \gamma \max_a \sum_{s'} P(s'|s, a) [R(s') + \gamma \max_{a'} \sum_{s''} P(s''|s', a') V^\pi(s'')] : \textcircled{B}$$

Imagine iterating above  $t$  times.

Let's show that this iterated expression implies optimality.

Let  $\hat{\pi}(s)$  be any other policy:

$$V^{\hat{\pi}}(s) = R(s) + \gamma \sum_{s'} P(s'|s, \hat{\pi}(s)) V^{\hat{\pi}}(s') : \text{Bellman equation}$$

$$\leq R(s) + \gamma \max_a \sum_{s'} P(s'|s, a) V^{\hat{\pi}}(s') : \text{greedy}$$

$$= R(s) + \gamma \max_a \sum_{s'} P(s'|s, a) [R(s') + \gamma \sum_{s''} P(s''|s', \hat{\pi}(s')) V^{\hat{\pi}}(s'')] : \text{iterate}$$

$$V^{\hat{\pi}}(s) \leq R(s) + \gamma \max_a \sum_{s'} P(s'|s, a) [R(s') + \gamma \max_{a'} \sum_{s''} P(s''|s', a') V^{\hat{\pi}}(s'')] : \text{be greedy}$$

$$= \dots \leq \dots = \dots : \textcircled{A}$$

Consider upper bound on  $V^{\hat{\pi}}(s)$  by iterating  $t$  times. (being greedy then applying Bellman equation.)Compare to equality after  $t$  iterations for  $V^\pi(s)$ .As  $t \rightarrow \infty$ , RHS on upper bound on  $V^{\hat{\pi}}(s)$  converges to RHS of equality for  $V^\pi(s)$ .

Thus, as  $t \rightarrow \infty$ :  $V^{\hat{\pi}}(s) \leq \lim_{t \rightarrow \infty} [\textcircled{A}] = \lim_{t \rightarrow \infty} [\textcircled{B}] = V^{\pi}(s)$

Thus, for all policies  $\hat{\pi}(s)$  and states  $s$ :  $V^{\pi}(s) \geq V^{\hat{\pi}}(s)$

$$V^{\pi}(s) = \max_{\pi} V^{\hat{\pi}}(s) \rightarrow V^{\pi}(s) = V^*(s).$$

Finally, once you have  $V^*(s)$ ,

$$\begin{aligned} \pi^*(s) &= \operatorname{argmax}_a Q^*(s, a) \\ &= \operatorname{argmax}_a \sum_{s'} P(s'|s, a) V^*(s') \end{aligned}$$

\* Pros & Cons of policy iteration.

(+) converges quickly in finite # steps

(-) requires policy evaluation  $O(n^3)$  at each iteration.

### Value Iteration

\* How to compute  $V^*(s)$  directly?

\* Bellman optimality equation.

$$\begin{aligned} V^*(s) &= \max_a Q^*(s, a) \\ &= \max_a [R(s) + \gamma \sum_{s'} P(s'|s, a) V^*(s')] \\ &= R(s) + \gamma \max_a \sum_{s'} P(s'|s, a) V^*(s') \quad (*) \end{aligned}$$

Before we showed that if  $V^{\pi}(s)$  satisfied (\*), then  $\pi$  is optimal.

Here, we show that  $V^*(s)$  must satisfy (\*).

Q: How to solve this set of  $N$  nonlinear equations for  $N$  unknowns?

\* Algorithm

1) initialize:  $V_0(s) = 0$  for all  $s$ .

2) iterate:  $V_{k+1}(s) = R(s) + \gamma \max_a \left[ \sum_{s'} P(s'|s, a) V_k(s') \right]$  for all  $s=1, 2, \dots, n$

Note: this algorithm works on value functions, not policies.

But, incremental policies can be computed as:

$$\pi_{k+1}(s) = \text{greedy}[V_k(s)] = \operatorname{argmax}_a \left[ \sum_{s'} P(s'|s, a) V_k(s') \right]$$

3) after convergence:

$$\pi^*(s) = \operatorname{argmax}_a [Q^*(s, a)] = \operatorname{argmax}_a \left[ \sum_{s'} P(s'|s, a) V^*(s') \right]$$

Does algorithm converge?  $V^*(s)$  is obviously a fixed point.

But are there others? Does it always reach  $V^*(s)$ ?

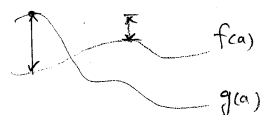
\* Lemma:  $|\max_a [f(a)] - \max_a [g(a)]| \leq \max_a |f(a) - g(a)|$

Proof of lemma:

$$\text{For all } a: f(a) - \max_{a'} g(a') \leq f(a) - g(a)$$

$$\text{Max over } a: \max_a f(a) - \max_{a'} g(a') \leq \max_a [f(a) - g(a)] \leq \max_a |f(a) - g(a)|$$

$$\text{By symmetry: } \max_a g(a) - \max_{a'} f(a') \leq \max_a |f(a) - g(a)|.$$



Combining last two bounds proves lemma.

\* Theorem: Value iteration converges.

$$\lim_{k \rightarrow \infty} [V_k(s)] \rightarrow V^*(s) \text{ for all states } s.$$

\* Proof: let  $\Delta_k = \max_s |V_k(s) - V^*(s)|$  error at  $k^{\text{th}}$  iteration.

$$\begin{aligned} \Delta_{k+1} &= \max_s |V_{k+1}(s) - V^*(s)| \\ &= \max_s \left| [R(s) + \gamma \max_a \sum_{s'} P(s'|s, a) V_k(s')] \right. \\ &\quad \left. - [R(s) + \gamma \max_a \sum_{s'} P(s'|s, a) V^*(s')] \right| \\ &= \gamma \max_s \left| \underbrace{\max_a \sum_{s'} P(s'|s, a) V_k(s')}_{f(a)} - \underbrace{\max_a \sum_{s'} P(s'|s, a) V^*(s')}_{g(a)} \right| \end{aligned}$$

Apply lemma

$$\begin{aligned} \Delta_{k+1} &\leq \gamma \max_s \max_a \left| \sum_{s'} P(s'|s, a) V_k(s') - \sum_{s'} P(s'|s, a) V^*(s') \right| \\ &= \gamma \max_s \max_a \left| \sum_{s'} P(s'|s, a) [V_k(s') - V^*(s')] \right| \\ &\leq \gamma \max_s \max_a \left| \sum_{s'} P(s'|s, a) \max_{s''} |V_k(s'') - V^*(s'')| \right| \\ &= \gamma \max_s \max_a \left| \sum_{s'} P(s'|s, a) \Delta_k \right| \\ &= \gamma \Delta_k \max_s \max_a (1) \end{aligned}$$

Hence:  $\Delta_{k+1} \leq \gamma \Delta_k$  with  $\gamma < 1$ . so-called "contraction" mapping.

By iteration:  $\Delta_k \leq \gamma^k \Delta_0 \rightarrow 0$  as  $k \rightarrow \infty$ .

$$\begin{aligned} \text{Assume rewards are bounded: } \Delta_0 &= \max_s |V_0(s) - V^*(s)| \\ &= \max_s |V^*(s)| \\ &\leq \max_s |R(s)| (1 + \gamma + \gamma^2 + \dots) \\ &= \max_s |R(s)| \cdot \left( \frac{1}{1-\gamma} \right) \end{aligned}$$

More iterations required as  $\gamma \rightarrow 1$ .