Fine-Grained, Secure and Efficient Data Provenance on Blockchain Systems

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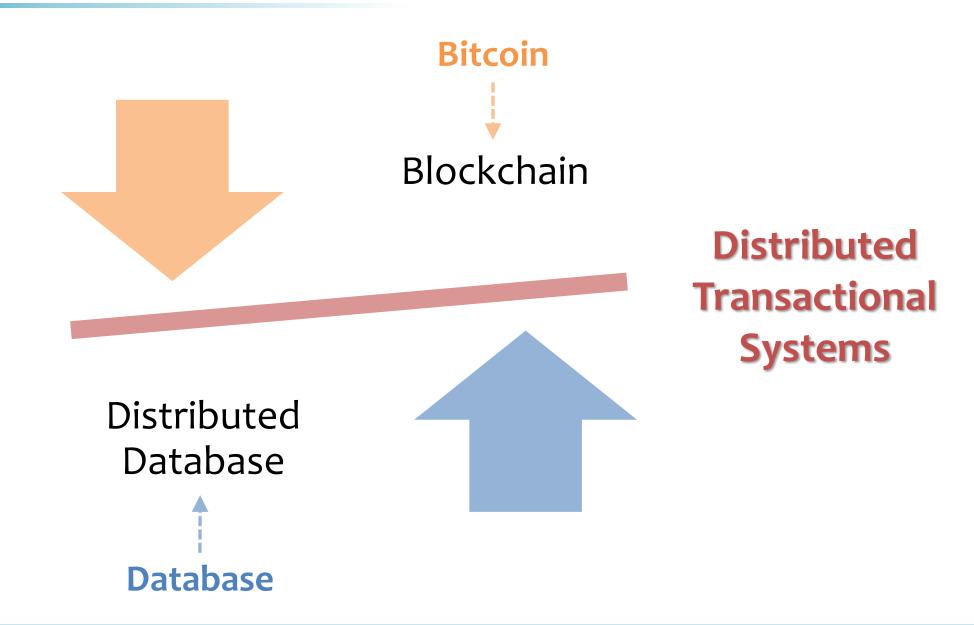






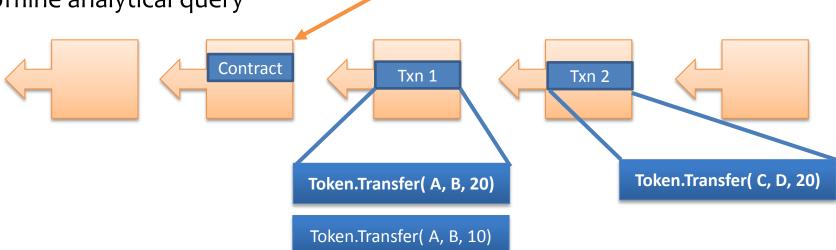


Blockchain Is a Class of Database



Blockchain Basics

- P2P network
 - Asynchronous transaction
- Byzantine environment
 - Mutual distrusting setup
- Distributed ledger
 - Smart contract
- Inherent provenance-preserving
 - ONLY for offline analytical query



contract Token {

} } }

method Transfer(sender, recipient, amount) {

gState[sender] = bal1 - amount; gState[recipient] = bal2 + amount;

bal1 = gState[sender];

if (amount < bal1) {

bal2 = gState[recipient];

Motivation

Expose provenance information to smart contracts both

Efficiently

Securely



Enabler for provenancedependent smart contracts

Enrich the transaction semantics

Provenance-dependent Contracts

Previous transfer precondition:

- Enough balance from the sender
- CURRENT STATE ONLY

New transfer precondition:

- Historical balance > threshold
- Recipient not transacted with certain blacklisted addresses recently
- HISTORICAL STATE & PROVENANCE INFO

```
contract Token {
   method Transfer(sender, recipient, amount) {
     bal1 = gState[sender];
     bal2 = gState[recipient];
     if (amount < bal1) {
        gState[sender] = bal1 - amount;
        gState[recipient] = bal2 + amount;
} }
}</pre>
```

Workarounds

Workaround 1:

- Dump every thing into current state
- Effort-needed, expensive, error-prune

Workaround 2:

- Offline analytics + Online transactions
- Break of serializability
- Transaction-ordering attacks

Workaround 3: Account1_v1: 10

 Minimum system instrumentation

Account1 v2: 20

Account1_v3: 15

Account2_v2: 12

- NOT protocol level (e.g., Hyperledger Fabric v1.0+)
- Data tampering

Holistic Approach:

- Protocol-level enhancement
 - → Secure
- Performance-aware
 - → Efficient

Challenges

NO standardized operations

- With clearly-defined transformation semantics
- E.g
 - Map and reduce in Hadoop
 - Select, join and aggregation in SQL

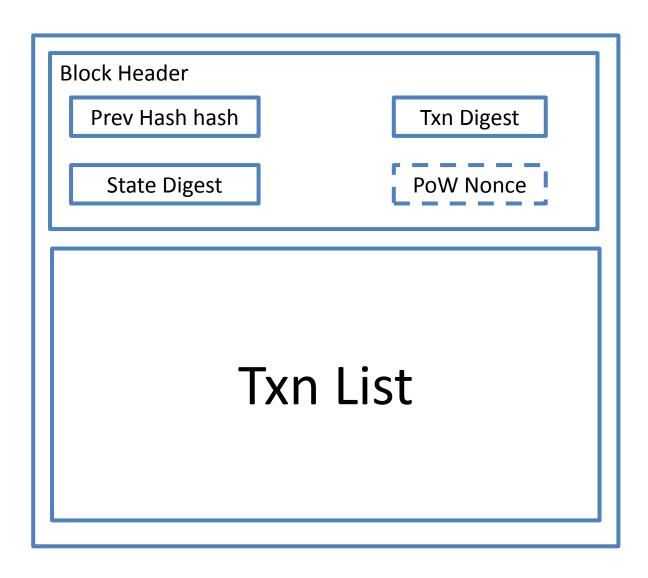
Byzantine environment

- Tamper evidence
- Integrity proof

Ever-growing ledger

- Gas mechanism
- Verifier's dilemma

Block Structure



Enhancement Basis (Merkle Tree Variants)

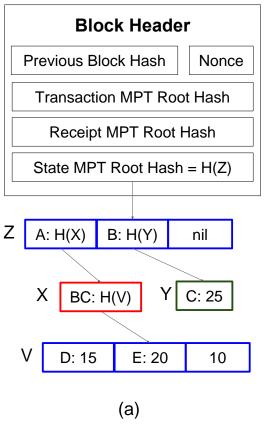
Limitation

Latest State only

Tamper evidence

- Succinct digest (root hash)
- Integrity proof (access path)

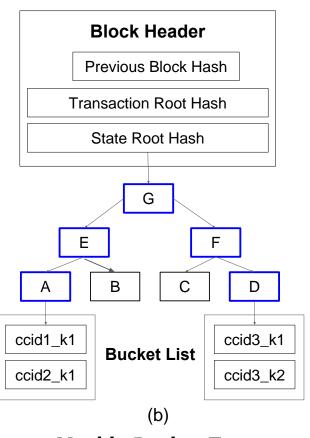
Account Address and Assoicated Balance in Global State:



Merkle Patricia Trie

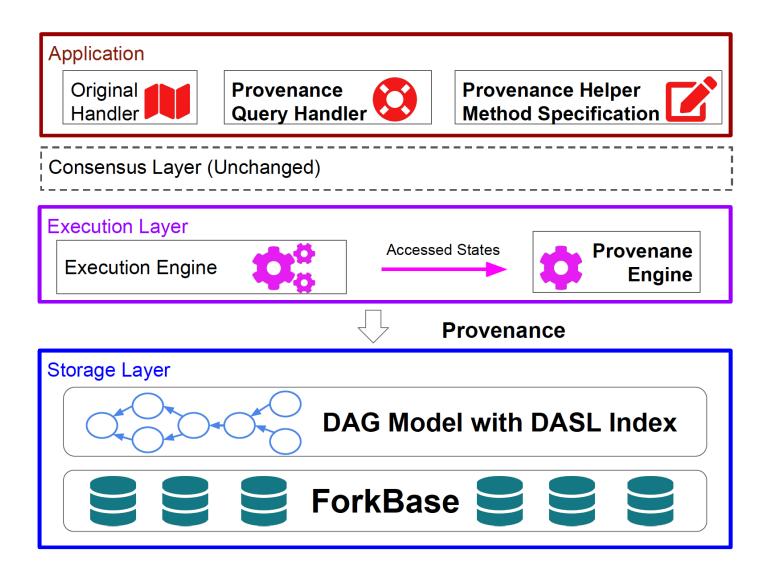
<Updated Chaincode ID>_<Key>:

ccid1_k1 ccid2_k1 ccid3_k1 ccid3_k2



Merkle Bucket Tree

LineageChain Overview



Application Layer

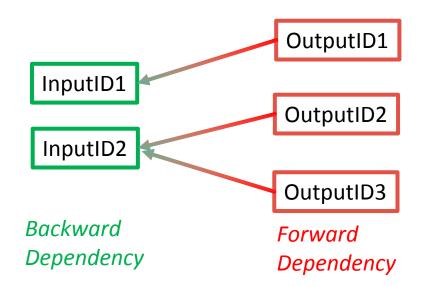




Provenance specification

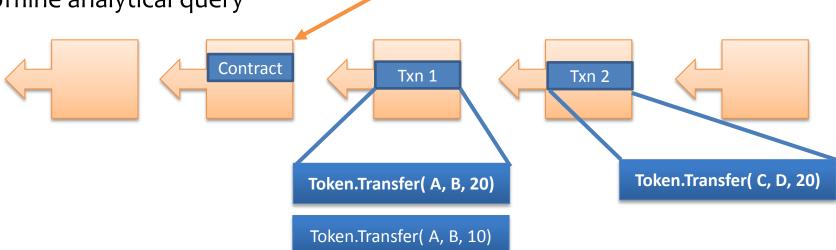
User-defined input-output dependency

Provenance query handler



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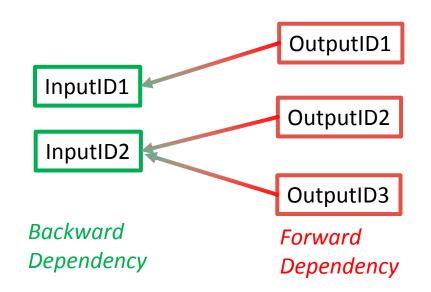
Application Layer



Provenance specification

User-defined input-output dependency

Provenance query handler



Application Layer

```
contract Token {
    method Transfer(...){...} // as above
    method prov_helper(name, reads, writes) {
      if name == "Transfer" {
        for (id, value) in writes {
            if (reads[id] < value) {</pre>
                recipient = id;
            } else {sender = id; }
        // dependency list with a
        // single element.
        dep = [sender]:
        return {recipient:dep};
method Refund(addr) {
  blk := last block in the ledger
  first_blk := first block in this month
  sum = count = 0;
  while (first_blk < blk) 4
    val, startBlk, txnID = Hist(addr, blk)
    blk = startBlk - 1;
    sum += val;
    count += 1;
  avg = sum / count;
  refund_amount := refund amount based on avg
  gState[addr] += refund_amount;
```









Recipient -> Sender

```
method Blacklist(addr) {
  blk := last block in the ledger
  blacklisted = false;
  iterate 5 times {
    val, startBlk, txnID = Hist(addr, blk);
    for (depAddr, depBlk)
        in (Backward(addr, startBlk))
        or Forward(addr, startBlk)) {
        if depAddr in gState["blacklist"] {
            gState["blacklist"].append(addr);
            return;
        }
     }
    blk = startBlk - 1;
}
```

Execution Layer

Receive

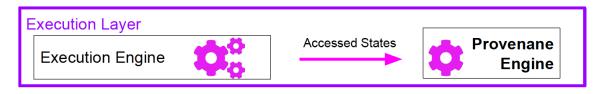
- Contract invocation context
- Provenance specification

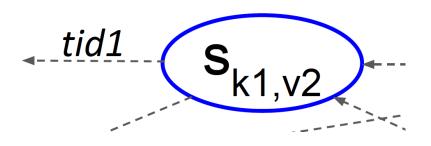
Compute

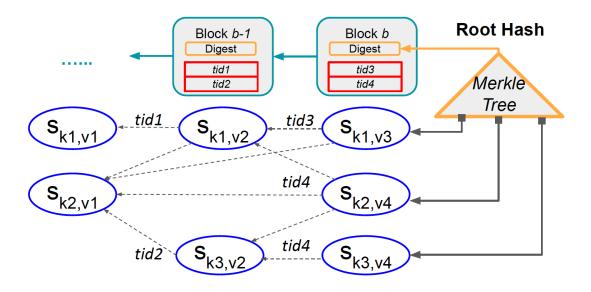
- Transaction results
- Concrete dependency

Prepare Merkle DAG

- Introduce one layer of direction
- Hash reference to encode provenance backward dependency







Execution Layer

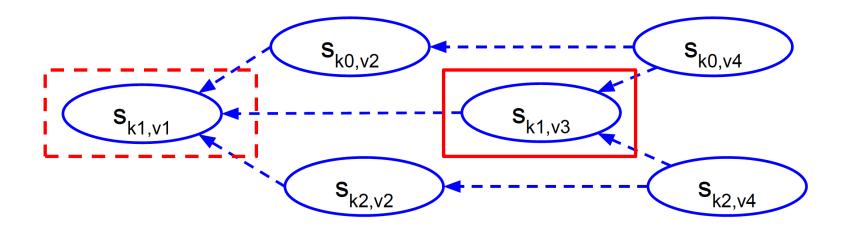


Forward tracking

- Problem: Undecided forward dependency during state update

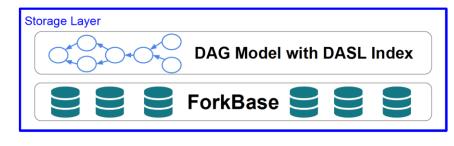
Solution

Lazily store forward dependency on the successor state entry



Storage Layer

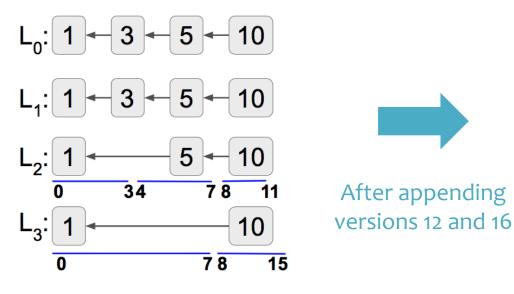
• Problem

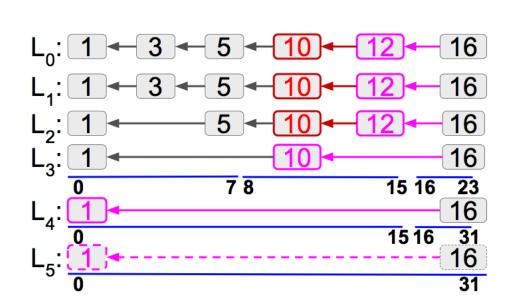


- Efficient version-based (historical) query for a state ID

Solution:

- Deterministic Append-only Skip List
- Hash-based reference





version numbers of states with identifier
$$k$$
, in which $v_i < v_j$ for all $i < j$. A DASL index for k consists of N linked lists $L_0, L_1, ..., L_{N-1}$. Let v_j^{i-1} and v_j^i be the versions in the

 $(j-1)^{th}$ and j^{th} node of list L_i . Let b be the base number, a system-wide parameter. The content of L_i is constructed as follows: 1) $v_0 \in L_i$

Definition 6. Let $V_k = \langle v_0, v_1, ... \rangle$ be the sequence of

2) Given
$$v_{j-1}^i$$
, v_j^i is the smallest version in V_k such that:

$$\left| \frac{v_{j-1}^i}{b^i} \right| < \left| \frac{v_j^i}{b^i} \right| \tag{5}$$

Evaluation

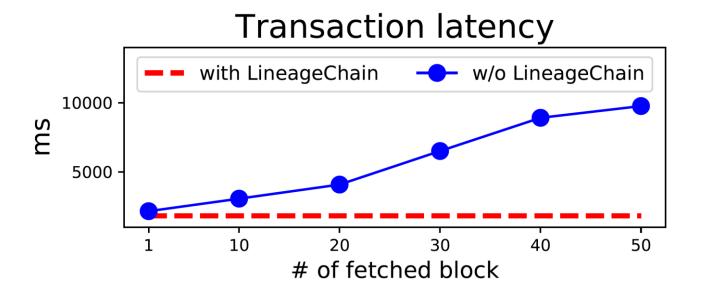
- MICRO benchmarking (vs. flat storage)
 - Preference to recent version query (with DASL)
 - More efficient BFS enabled by backtrack (with ForkBase)

- MACRO benchmarking (applied to Hyperledger Fabric vo.6 and v1.3)
 - Negligible runtime overhead
 - Tiny proportion of latency
 - Negligible storage overhead
 - >70% of space for blocks
 - 25% for historical states
 - 2~4% for DASL indexes and hash pointers

Performance of Provenance Query

vs. Workaround 2

Compute data provenance offline and conditionally trigger online transaction



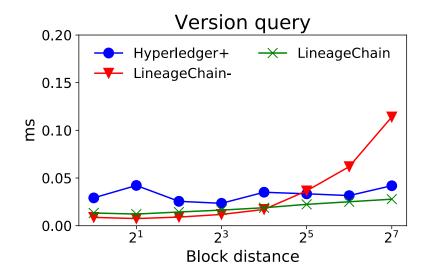
Micro Performance of Provenance Query

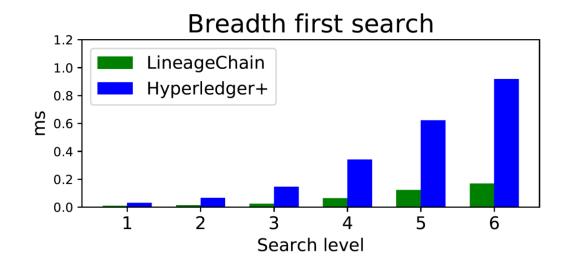
vs. Workaround 1

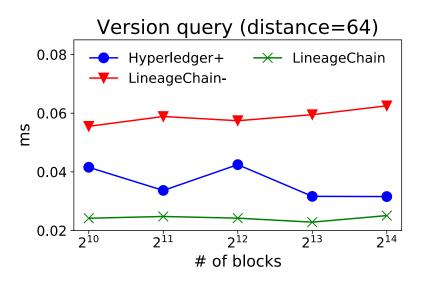
Dump everything into the current state

vs. Workaround 3

Use Hyperledger Fabric's built-in HistoryDB

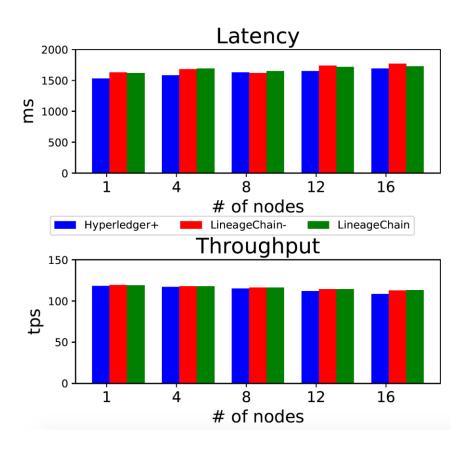




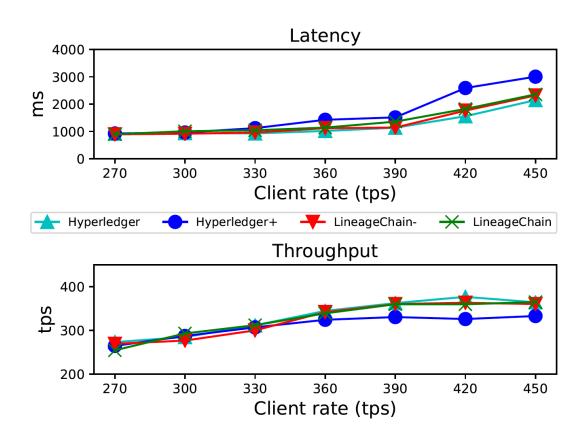


Runtime Overhead

Transaction processing

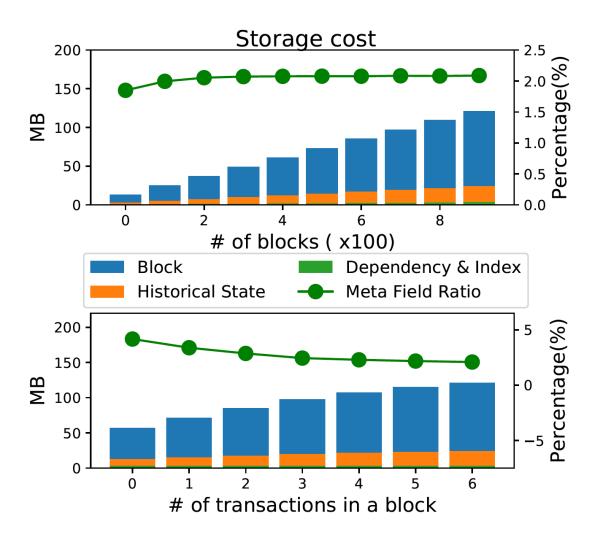


Hyperledger Fabric vo.6



Hyperledger Fabric v1.3

Storage Overhead



Conclusion

LineageChain

- Enabler for provenance-dependent blockchain applications
- Protocol-level enhancement w.r.t. efficiency and security
- Negligible performance and storage overhead

Key designs

- User-defined dependency specification
- Merkle DAG with dependency tracking
- DASL index to accelerate data provenance query
- Adoption in Hyperledger Fabric (vo.6 & v1.3)

Thank You!