

Topological Models

Step 1:

The logic of
open sets

Logika

 \top \perp $\varphi \wedge \psi$ $\varphi \vee \psi$ $\varphi \Rightarrow \psi$ $\forall a:A . \varphi(a)$ $\exists a:A . \varphi(a)$

Topologija

 \mathcal{X} \emptyset $[\varphi] \cap [\psi]$ $[\varphi] \cup [\psi]$ $\text{Int}([\neg\psi] \cup [\varphi]^c)$ $\text{Int}(\bigcap_{a \in A} [\varphi(a)])$ $\bigcup_{a \in A} [\varphi(a)]$

Theorem: \mathcal{X} validates $L\bar{E}M$ iff
 \mathcal{X} is discrete.

Proof.

If \mathcal{X} is discrete then $\mathcal{O}\mathcal{X}$ is Boolean.

If $\mathcal{X} \models L\bar{E}M$ then $\forall U \in \mathcal{X}. U \vee \neg U$

$$\neg U = U \Rightarrow \perp = \text{Int}(\emptyset \vee U^c) = \text{Int}(U^c).$$

Then if $\mathcal{X} = U \vee \text{Int}(U^c)$ we must have

$$U^c = \text{Int}(U^c),$$

so U is closed.

□

Step 2:

Heyting valued
sets

Let \mathcal{F} be a sheaf.

Define $F := \sum_{u \in \mathcal{P}_x} \mathcal{F}(u)$ and $\|a\| = \text{pr}_1 a$.

Let $F' \subseteq F$. Then F' generates \mathcal{F} when

(a) \mathcal{F} is the least subsheaf containing F'

(b) $\forall f \in \mathcal{F}(u), \exists f_i \in F', u_i \in \mathcal{P}_u, f_i|_{u_i} = f|_{u_i}$

$$f = f_1|_{u_1} \cup \dots \cup f_n|_{u_n}$$

Definition: An \mathcal{L} -set is a set A with

$$[- = -]: A \times A \longrightarrow \mathcal{L} \quad \text{s.t.}$$

$$1. [a = b] \leq [b = a]$$

$$2. [a = b] \wedge [b = c] \leq [a = c]$$

Definition. An \mathcal{L} -morphism is a map

$$[- = f(-)]: B \times A \longrightarrow \mathcal{L} \quad \text{s.t.}$$

$$1. \quad [b = f(a)] \leq \|b\| \wedge \|a\|$$

$$2. \quad [b' = b] \wedge [b = f(a)] \wedge [a = a'] \leq [b' = f(a')]$$

$$3. \quad [b = f(a)] \wedge [b' = f(a)] \leq [b = b']$$

$$4. \quad \|a\| \leq \bigvee_{b \in B} [b = f(a)]$$

Step 3:
Internal language

Logika

 \top \perp $\varphi \wedge \psi$ $\varphi \vee \psi$ $\varphi \Rightarrow \psi$ $\forall a:A . \varphi(a)$ $\exists a:A . \varphi(a)$ $a = b$

Topologija

 \mathcal{X} \emptyset $[\varphi] \cap [\psi]$ $[\varphi] \cup [\psi]$ $\text{Int}([\psi] \cup [\varphi]^c)$ $\text{Int}(\bigcap_{a \in A} [\|a\| \Rightarrow \varphi(a)])$ $\bigcup_{a \in A} [\|a\| \wedge \varphi(a)]$ $[a = b]$

Definition: An \mathcal{L} -set is a set A with

$$[- = -]: A \times A \longrightarrow \mathcal{L} \quad \text{s.t.}$$

$$1. \vdash a = b \Rightarrow b = a$$

$$2. \vdash a = b \wedge b = c \Rightarrow a = c$$

Definition. An \mathcal{L} -morphism is a map

$$[- = f(-)]: B \times A \longrightarrow \mathcal{L} \quad \text{s.t.}$$

$$1. \vdash b = f(a) \Rightarrow \|b\| \wedge \|a\|$$

$$2. \vdash b' = b \wedge b = f(a) \wedge a = a' \Rightarrow b' = f(a')$$

$$3. \vdash b = f(a) \wedge b' = f(a) \Rightarrow b = b'$$

$$4. \vdash \forall a:A. \exists b:B. b = f(a)$$

Facts:

- $[- \equiv_A -]$ is id_A .

- $c = g \circ f(a)$ is $\exists b : B. c = g(b) \wedge b = f(a)$.

- $R(f(a))$ is $\exists b : B. b = f(a) \wedge R(b)$.

$\hookrightarrow f(a) = g(a)$ is $\exists b : B. b = f(a) \wedge b = g(a)$

$$\text{Ex. } \mathbb{1} := \{*\}, \mathbb{1} * \mathbb{1} = \top$$

$$\text{Ex. } \underline{A} = A, \llbracket a = a' \rrbracket = \bigvee \{ \top \mid a = a' \}.$$

$$\text{Ex. } A|_u := A, \llbracket a = a' \rrbracket = \llbracket a = a' \rrbracket \wedge u.$$

$$\text{Ex. } \Omega = \mathcal{L}, \llbracket p = q \rrbracket = p \Leftrightarrow q.$$

$$\text{Ex. } B^A = A \rightarrow B, \llbracket f = g \rrbracket = \llbracket \forall a:A. f(a) = g(a) \rrbracket$$

$$\text{Ex. } A/\sim = A, \llbracket a = a' \rrbracket = a \sim a'.$$

Lemma (funext). $f = g$ iff $\vdash f =_{B^A} g$.

Proof. Assume $\vdash f = g$.

Let $a:A, b:B$ s.t. $b = f(a)$.

By ass, $\exists b':B. b' = f(a) \wedge b' = g(a)$

Then $b = f(a) = b'$, so $b = b' = g(a)$. \square

Theorem. $f: A \rightarrow B$ is mono iff

$$+ \forall x, y: A. f(x) = f(y) \Rightarrow x = y.$$

Proof. (\Leftarrow) Let $g, h: C \rightarrow A$, $f \circ g = f \circ h$.

Then $f(g(c)) = f(h(c))$, so $g = h$.

(\Rightarrow) Let $x, y: B$ st. $f(x) = f(y)$.

Define $\hat{x}, \hat{y}: \mathbb{1} \rightarrow B$, $\hat{x}(*) := x$, $\hat{y}(*) := y$

Then $f \circ \hat{x} = f \circ \hat{y}$, so $\hat{x} = \hat{y}$.

□

Theorem $f: A \rightarrow B$ is epi iff

$$\vdash \forall b: B \exists a: A. b = f(a).$$

Proof. (\Leftarrow) let $g, h: B \rightarrow C$, $g \circ f = h \circ f$.

Let $b: B$ Then $\exists a: A. b = f(a)$

$$\text{Then } g(b) = g(f(a)) = h(f(a)) = h(b).$$

(\Rightarrow) let $b: B$ and $i_1, i_2: B \rightarrow \text{cohet}$.

$$\text{Then } i_1 \circ f = i_2 \circ f \rightsquigarrow i_1 = i_2 \rightsquigarrow \exists a: A. b = f(a). \quad \square$$

Corollary. Isomorphisms are bijections.











