

Topological Models

Step 1:

The logic of
open sets

Logika

\top

\perp

$\varphi \wedge \neg$

$\varphi \vee \neg$

$\varphi \Rightarrow \neg$

$\forall a:A . \varphi(a)$

$\exists a:A . \varphi(a)$

Topologija

\mathcal{X}

\emptyset

$[\varphi] \cap [\neg]$

$[\varphi] \cup [\neg]$

$\text{Int}([\neg] \cup [\varphi]^c)$

$\text{Int}(\bigcap_{a \in A} [\varphi(a)])$

$\bigcup_{a \in A} [\varphi(a)]$

Theorem: \mathcal{X} validates LEM iff
 \mathcal{X} is discrete.

Proof. LEM says

"every truth value is complemented."
open set closed

Points are closed. By assumption
they are open.

□

Step 2:

Heyting valued
sets

Let \mathcal{F} be a sheaf.

Define $F := \sum_{U \in \mathcal{O}_X} \mathcal{F}(U)$ and $\|a\| = \text{pr}_1 a$

Let $F' \subseteq F$. Then F' generates \mathcal{F} when

(a) \mathcal{F} is the least subsheaf containing F'

(b) $\forall f \in \mathcal{F}(U). \exists f_i \in F', U_i \overset{\text{op}}{\subseteq} U, f_i|_{U_i} = f|_{U_i}$

$$f = f_1|_{U_1} \cup \dots \cup f_n|_{U_n}$$

Definition: An L -set is a set A with

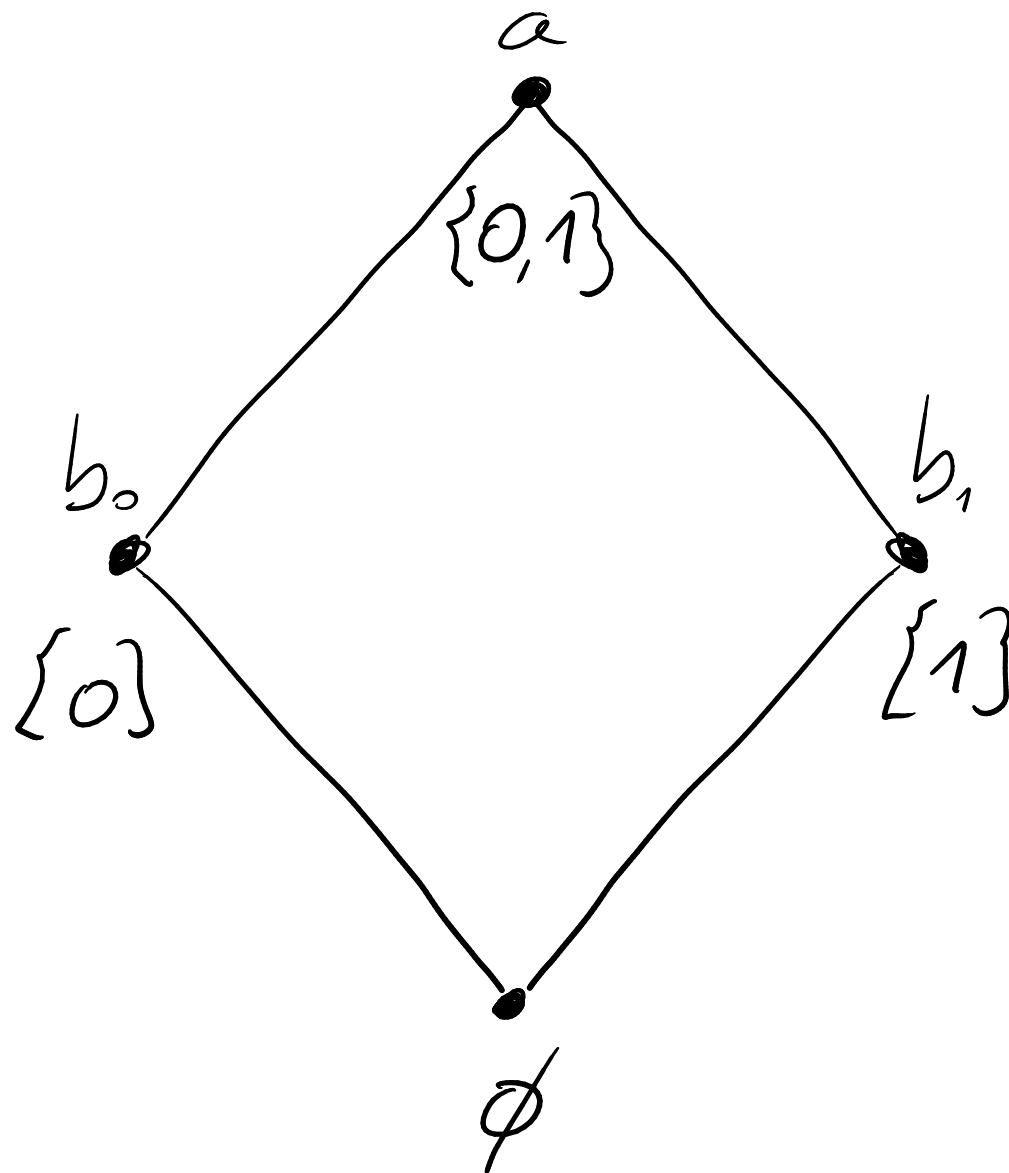
$$[- = -]: A \times A \longrightarrow L \text{ s.t.}$$

$$1. [a = b] \leq [b = a]$$

$$2. [a = b] \wedge [b = c] \leq [a = c]$$

Define $\|a\| := [a = a]$.

$$A = \{a\}, B = \{b_0, b_1\}$$



Definition. An \mathcal{L} -morphism is a map

$$[- = f(-)]: B \times A \longrightarrow \mathcal{L} \quad \text{s.t.}$$

$$1. \quad [b = f(a)] \leq \|b\| \wedge \|a\|$$

$$2. \quad [b' = b] \wedge [b = f(a)] \wedge [a = a'] \leq [b' = f(a')]$$

$$3. \quad [b = f(a)] \wedge [b' = f(a)] \leq [b = b']$$

$$4. \quad \|a\| \leq \bigvee_{b \in B} [b = f(a)]$$

Step 3:
Internal language

Logika

\top

\perp

$\varphi \wedge \neg$

$\varphi \vee \neg$

$\varphi \Rightarrow \neg$

$\forall a:A . \varphi(a)$

$\exists a:A . \varphi(a)$

$a = b$

Topologija

\mathcal{X}

\emptyset

$[\varphi] \cap [\neg]$

$[\varphi] \cup [\neg]$

$\text{Int}([\neg] \cup [\varphi]^c)$

$\text{Int}(\bigcap_{a \in A} [\|a\| \Rightarrow \varphi(a)])$

$\bigcup_{a \in A} [\|a\| \wedge \varphi(a)]$

$[a = b]$

Logika

$$\forall a:A . \varphi(a)$$

$$\exists a:A . \varphi(a)$$

$$a = b$$

$$R(a_1, \dots, a_n)$$

$$b = f(a)$$

Topologija

$$\text{Int}(\bigcap_{a \in A} [\|a\| \Rightarrow \varphi(a)])$$

$$\bigcup_{a \in A} [\|a\| \wedge \varphi(a)]$$

$$[a = b]$$

$$R(a_1, \dots, a_n)$$

$$[b = f(a)]$$

Definition: An \mathcal{L} -set is a set A with

$$[- = -]: A \times A \longrightarrow \mathcal{L} \quad \text{s.t.}$$

$$1. \vdash a = b \Rightarrow b = a$$

$$2. \vdash a = b \wedge b = c \Rightarrow a = c$$

Definition. An \mathcal{L} -morphism is a map

$$[- = f(-)]: B \times A \longrightarrow \mathcal{L} \quad \text{s.t.}$$

$$1. \vdash b = f(a) \Rightarrow \|b\| \wedge \|a\|$$

$$2. \vdash b' = b \wedge b = f(a) \wedge a = a' \Rightarrow b' = f(a')$$

$$3. \vdash b = f(a) \wedge b' = f(a) \Rightarrow b = b'$$

$$4. \vdash \forall a:A. \exists b:B. b = f(a)$$

Step 4:

Properties of L -sets
and morphisms

Facts:

- $[- =_A -]$ is id_A .

- $c = g \circ f(a)$ is $\exists b : B. c = g(b) \wedge b = f(a)$.

- $R(f(a))$ is $\exists b : B. b = f(a) \wedge R(b)$.

$\hookrightarrow f(a) = g(a)$ is $\exists b : B. b = f(a) \wedge b = g(a)$

$$\text{Ex. } \mathbb{1} := \{*\}, \quad \|*\| := \top, \quad \mathbb{0} := \emptyset$$

$$\text{Ex. } \underline{A} := A, \quad \llbracket a = a' \rrbracket := \bigvee \{ \top \mid a = a' \}.$$

$$\text{Ex. } \Omega := \mathcal{L}, \quad \llbracket p = q \rrbracket := p \Leftrightarrow q.$$

$$\text{Ex. } A \times B := \{ (a, b) \in A \times B \mid \|a\| = \|b\| \}$$

$$\text{Ex. } B^A := A \rightarrow B, \quad \llbracket f = g \rrbracket := \llbracket \forall a:A. f(a) = g(a) \rrbracket$$

$$\text{Ex. } A_{\sim} := A, \quad \llbracket a = a' \rrbracket := a \sim a'.$$

Lemma (funext). $f = g$ iff $\vdash f =_{B^A} g$.

Proof. Assume $\vdash f = g$.

Let $a:A, b:B$ s.t. $b = f(a)$.

By ass., $\exists b':B. b' = f(a) \wedge b' = g(a)$

Then $b = f(a) = b'$, so $b = b' = g(a)$. □

Theorem. $f: A \rightarrow B$ is mono iff
 $\vdash \forall x, y: A. f(x) = f(y) \Rightarrow x = y.$

Proof. (\Leftarrow) Let $g, h: C \rightarrow A, f \circ g = f \circ h.$

Then $f(g(c)) = f(h(c)),$ so $g = h.$

(\Rightarrow) Let $x, y: B$ st. $f(x) = f(y).$

Define $\hat{x}, \hat{y}: \mathbb{1} \rightarrow B, \hat{x}(*) := x, \hat{y}(*) := y$

Then $f \circ \hat{x} = f \circ \hat{y},$ so $\hat{x} = \hat{y}.$

□

Theorem $f: A \rightarrow B$ is epi iff

$$\vdash \forall b: B \exists a: A. b = f(a).$$

Proof. (\Leftarrow) let $g, h: B \rightarrow C$, $g \circ f = h \circ f$.

Let $b: B$ Then $\exists a: A$ $b = f(a)$

$$\text{Then } g(b) = g(f(a)) = h(f(a)) = h(b).$$

(\Rightarrow) let $b: B$ and $i_1, i_2: B \rightarrow \text{cohet}$

$$\text{Then } i_1 \circ f = i_2 \circ f \rightsquigarrow i_1 = i_2 \rightsquigarrow \exists a: A. b = f(a). \square$$

✧ Karakterizacija ✧ podobjektov

1. monomorphisms $S \rightarrow A$

2. sections of $P(A) = \Omega^A$

3. predicates on A

$$f: A \rightarrow \Omega:$$

$$1. \vdash U = f(a) \Rightarrow \|U\| \wedge \|a\|$$

$$2. \vdash U' = U \wedge U = f(a) \wedge a = a' \Rightarrow U' = f(a')$$

$$3. \vdash U = f(a) \wedge U' = f(a) \Rightarrow U = U'$$

$$4. \vdash \forall a:A, \exists U:\Omega. U = f(a)$$

Define $\hat{f}(a) := \llbracket T = f(a) \rrbracket$

$$1. \hat{f}(a) \Rightarrow \|a\|$$

$$2. \hat{f}(a) \wedge a = a' \Rightarrow \hat{f}(a')$$

Image and preimage constructions

Let $f: A \rightarrow B$, $S \subseteq A$, $T \subseteq B$.

Define $f_*(S) := \{y \in B \mid \exists a \in S. y = f(a)\}$

$f^*(T) := \{x \in A \mid \exists b \in T. b = f(x)\}$

If $S = \{a\}$ and $T = \{b\}$

- $f_*(S)(y) = \llbracket y = f(a) \rrbracket$

- $f^*(T)(x) = \llbracket b = f(x) \rrbracket$

(Sub)singletons

Consider $f: \mathbb{1} \rightarrow A$.

$$1. \vdash a \in f(*) \Rightarrow \|a\| \wedge \|*\|$$

$$2. \vdash a' = a \wedge a \in f(*) \Rightarrow a' \in f(*)$$

$$3. \vdash a \in f(*) \wedge a' \in f(*) \Rightarrow a = a'$$

$$4. \vdash \|*\| \Rightarrow \exists a: A. a \in f(*)$$

Definition: $S \subseteq A$ is a (sub)singleton when

$$\vdash x \in S \wedge y \in S \Rightarrow x = y.$$

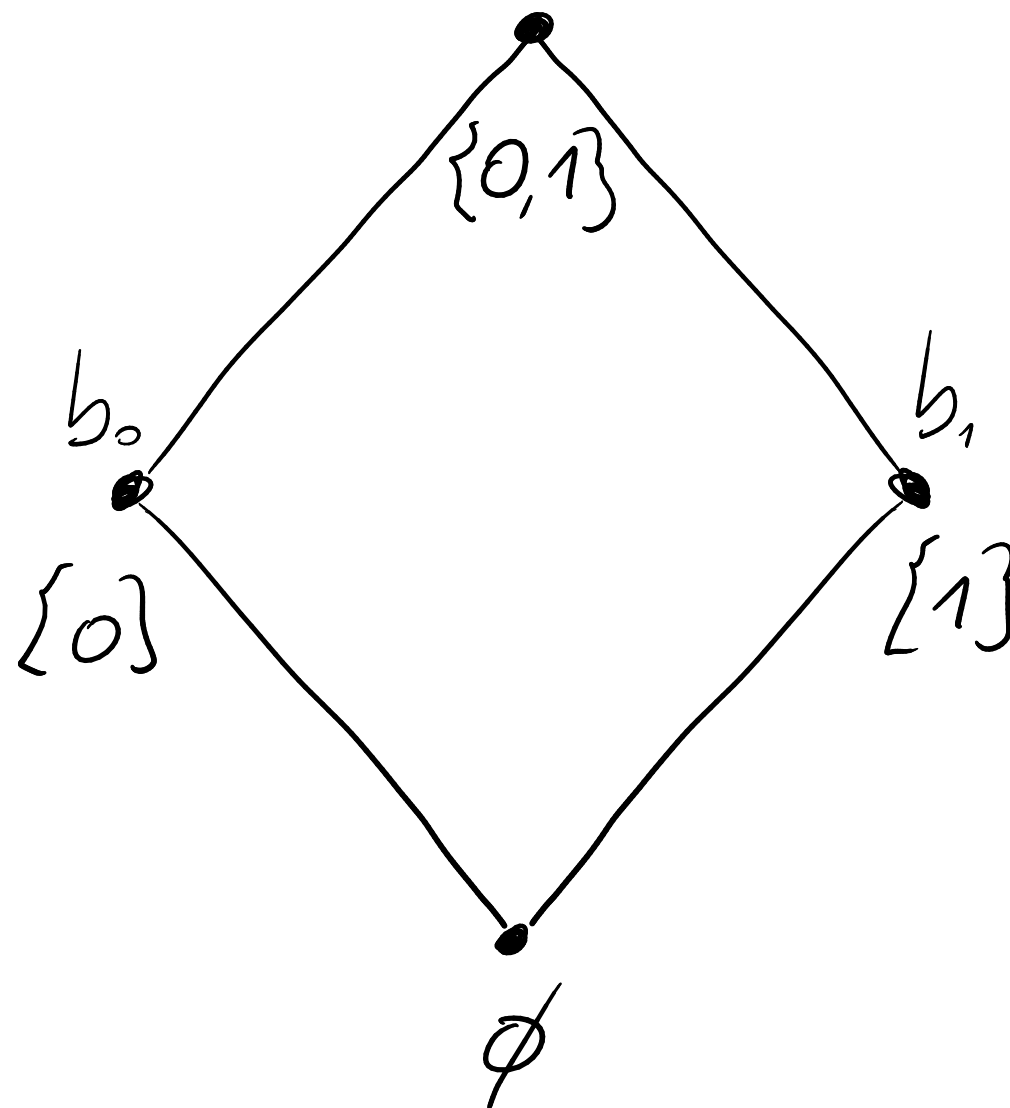
$$A = \{a\}, B = \{b_0, b_1\}$$

$$- A = \mathbb{1}_T \leadsto f(*) = \{b_0 + b_1\}$$

$$- A = \mathbb{1}_{\{0\}} \leadsto f(*) = \{b_0\}$$

$$- A = \mathbb{1}_{\{1\}} \leadsto f(*) = \{b_1\}$$

$$- A = \mathbb{1}_\perp \leadsto f(*) = \{\emptyset\}$$



Definition. $\sigma(A) = \{\sigma \subseteq A \mid \sigma \text{ singleton}\}$ is an \mathcal{L} -set
with $\sigma = \tau \iff \exists a \in A. a \in \sigma \wedge a \in \tau$.

Lemma. 1. $\vdash \sigma = \{a\} \iff a \in \sigma$,

2. $\vdash \{a\} = \{b\} \iff a = b$,

3. $\vdash a = b \iff \exists \sigma \in \sigma(A). a \in \sigma \wedge b \in \sigma$.

Definition. A is complete when the map
 $\{-\}: A \longrightarrow \sigma(A)$ is an isomorphism.

Restrictions: Let $a \in A, V \subseteq \|a\|$.

Define $x \in \sigma \Leftrightarrow x \in \{a\} \wedge V$

By completeness $\sigma = \{a'\}$.

$$a = a' \Leftrightarrow \{a\} = \{a'\} \Leftrightarrow \exists x:A. x \in \{a\} \wedge x \in \{a'\}$$

$$\Leftrightarrow \exists x:A. x \in \{a\} \wedge V \Leftrightarrow \|a\| \wedge V \Leftrightarrow V.$$

Gluing: Let $\{a_i\}_{i \in I}$ be a compatible family

Define $x \in \sigma \Leftrightarrow \exists i \in I. x \in \{a_i\}$

If $x \in \{a_i\}$ and $y \in \{a_j\}$ then $x = a_i = a_j = y$.

So $\exists! a: A. \sigma = \{a\}$.

$$\begin{aligned} x \in a \restriction_{U_i} &\Leftrightarrow x \in a \cap U_i \\ &\Leftrightarrow \exists j \in I. x \in \{a_j\} \cap U_i \\ &\Leftrightarrow x \in \{a_j\} \cap U_i \\ &\Leftrightarrow x \in \{a_j\} \end{aligned}$$

Theorem. If B is complete then

every $f: A \rightarrow B$ comes from an $\hat{f}: A \rightarrow B$.

Proof. For $a: A$ $f(a)$ is a singleton

$$\text{so } \hat{f}(a) := \{f(a)\}^{-1}$$

□

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$$\text{Ex. } \underline{A} := A, \quad \llbracket a = a' \rrbracket := \bigvee \{ \top \mid a = a' \}.$$

$$\text{Ex. } \Omega := \mathcal{L}, \quad \llbracket p = q \rrbracket := p \Leftrightarrow q.$$

$$\text{Ex. } A \times B := \{ (a, b) \in A \times B \mid \|a\| = \|b\| \}$$

$$\text{Ex. } B^A := A \rightarrow B, \quad \llbracket f = g \rrbracket := \llbracket \forall a:A. f(a) = g(a) \rrbracket$$

$$\text{Ex. } A_{\sim} := A, \quad \llbracket a = a' \rrbracket := a \sim a'.$$

Theorem. $\sigma(A)$ is complete.

Proof. Let $\Sigma: \sigma(A)$.

Define $a \in \tau \Leftrightarrow \{a\} \in \Sigma$.

Take $a: A$ s.t. $a \in \tau$.

Then $\tau = \{a\}$ and $\{a\} \in \Sigma$ so $\tau \in \Sigma$, or $\Sigma = \{\tau\}$. \square

Theorem. $A \cong \sigma(A)$.

Proof. Define $\eta(\sigma, a) := a \in \sigma$.

3. $a \in \sigma \wedge a \in \tau \Rightarrow \exists a : A. a \in \sigma \wedge a \in \tau$,

4. $a \in \{a\} \Rightarrow \exists \sigma : \sigma(A). a \in \sigma$,

inj. $a \in \sigma \wedge b \in \sigma \Rightarrow a = b$,

surj. $\forall \sigma : \sigma(A). \sigma = \sigma$, so $\exists a : A. a \in \sigma$.

□

Natural numbers

$\sigma(\underline{N})$:

$$\sigma: \underline{N} \rightarrow \mathcal{L}$$

$$1. \vdash n \in \sigma \wedge m \in \sigma \Rightarrow n = m,$$

$$2. \vdash n = m \wedge m \in \sigma \Rightarrow n \in \sigma.$$

For $x \in [\exists n: \underline{N}. n \in \sigma]$, $x \in \bigcup_{n \in \underline{N}} \sigma(n)$

[illegible]

$$\text{Thm. } f = \hat{f} \quad \hat{f}(u, a) = \llbracket u =_A \llbracket \top = f(a) \rrbracket \rrbracket$$

$$\text{Pf. Let } a:A. \quad f(a) = \hat{f}(a) \quad \text{is } f(a) = \llbracket \top = f(a) \rrbracket \\ \text{or } f(a) = (\top \Rightarrow f(a))$$

$$\text{Thm. } \hat{f} = \hat{g} \Rightarrow f = g.$$

$$\text{Pf. Let } a:A. \text{ Show } f(a) = g(a)$$

$$\text{By ass. } \llbracket \top = f(a) \rrbracket = \hat{f}(a) = \hat{g}(a) = \llbracket \top = g(a) \rrbracket.$$

$$\text{Then } f(a) = g(a).$$

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