

## → Prime check and Prime Factorization 53

- Brute Force Method to check prime.
- Time Complexity  $O(N)$

```

BOOL PRIMECHECK (INT n)
{

```

```

    BOOL ISPRIME = TRUE;

```

→ If  $n$  is 1 return false as 1 is a non-prime number

```

    IF (n == 1)
    {

```

```

        RETURN FALSE;
    }

```

```

    FOR (INT i = 2; i < n; i++)
    {

```

```

        IF (n % i == 0)
        {

```

```

            RETURN FALSE;
        }
    }

```

```

    RETURN TRUE;
}

```

→ But we can observe a pattern here, just like in finding divisors, we could use  $\sqrt{n}$  till  $\text{sqrt}(n)$ , here we can do the same, because if a number is divisible by a number coming before  $\text{sqrt}(n)$ , it will automatically be divisible by numbers coming afterwards.

→ eg: For  $n = 36$

→  $2 \rightarrow 36 / 2 = 18$

→  $3 \rightarrow 36 / 3 = 12$

→  $4 \rightarrow 36 / 4 = 9$

(So if a number is divisible by 2, it will automatically be divisible by 18)



→ 6 →  $36/6 = 6$

NOTE: The

- Square Root Method to check Prime
- Time Complexity  $O(\sqrt{n})$  or  $O(\sqrt{n})$

```

BOOL SQRTPRIMECHK (INT n)
{

```

```

    BOOL ISPRIME = TRUE;

```

→ If  $n$  is 1 return false as 1 is a non prime number.

```

    IF (n == 1)

```

```

    {

```

```

        RETURN FALSE;

```

```

    }

```

```

    FOR (INT i = 2; i * i < n; i++)

```

```

    {

```

```

        IF (n % i == 0)

```

```

        {

```

```

            RETURN FALSE;

```

```

        }

```

```

    }

```

```

    RETURN TRUE;

```

```

}

```

→ Prime Factors:

Approach: First we have to find the smallest prime number of a given  $n$ , and then we keep on dividing that

$n$  with that prime numbers, until it is no more divisible, now we have to find this new  $n$ 's prime numbers, and follow the same process until  $n$  becomes 1.

→ Brute Force Approach

→ Time Complexity  $O(n)$

```

INT n;
CIN n;
VECTOR<INT> PRIME_FACT-B;
FOR (INT i=2; i ≤ n; i++)
{
    WHILE (n % i == 0)
    {
        PRIME_FACT-B.PB(i);
        n /= i;
    }
}

```

→ As in prime factorization, we always consider the smallest possible prime number which is divisor of  $n$ , so, we use  $n$  till  $\text{sqrt}(n)$ , because if a number is divisible by a number coming before  $\text{sqrt}(n)$ , it will automatically be divisible by numbers coming afterwards. This will work fine until we have  $n$  composite, but as we get  $n$  a prime no., we have to apply an if check, bc if  $n = 3$ , then  $\text{sqrt}(3)$  condition will



not work.

- > Square Root Method for prime factorization
- > Time complexity  $O(\sqrt{n})$  or  $O(\sqrt{n})$

~~for~~

int n;

cin >> n;

vector<int> prime\_facts;

for (int i=2; i\*i <= n; i++)

{

while (n % i == 0)

{

prime\_facts.pb(i);

n /= i;

}

}

if (n > 1)

{

prime\_facts.pb(n);

}