CS 211: High Performance Computing Project 1

Performance Optimization via Register and Cache Reuse

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1. Register Reuse

Part #1.

For n=1000, the time my computer spends to finish dgemm0 is $\underline{5.5672(s)}$; the time my computer spends to finish dgemm1 is $\underline{4.2780(s)}$. The time wasted on accessing operands that are not in registers is $\underline{1.2892(s)}$.

The time spend in the triple loop for each algorithm (dgemm0, dgemm1) on TARDIS with n = 64, 128, 256, 512, 1024, 2048 is: (in seconds)

| algorithm | n=64 | n=128 | n=256 | n=512 | n=1024 | n=2048 |
|-----------|------|-------|-------|-------|--------|--------|
| dgemm0 | 0.00 | 0.04 | 0.37 | 3.95 | 34.19 | 548.45 |
| dgemm1 | 0.00 | 0.02 | 0.23 | 2.73 | 24.10 | 351.63 |

The performance of each algorithms on TARDIS with $n=64,\,128,\,256,\,512,\,1024,\,2048$ is: (in Gflops)

| algorithm | n=64 | n=128 | n=256 | n=512 | n=1024 | n=2048 |
|-----------|------|-------|-------|-------|--------|--------|
| dgemm0 | NA | 0.10 | 0.09 | 0.07 | 0.06 | 0.03 |
| dgemm1 | NA | 0.21 | 0.15 | 0.10 | 0.09 | 0.05 |

Part #2.

The time spend in the algorithm *dgemm2* on TARDIS with n = 64, 128, 256, 512, 1024, 2048 is:

The performance of the algorithm *dgemm2* on TARDIS with n = 64, 128, 256, 512, 1024, 2048 is:

| algorithm | n=64 | n=128 | n=256 | n=512 | n=1024 | n=2048 |
|-----------|------|-------|-------|-------|--------|--------|
| dgemm2 | 0.00 | 0.02 | 0.18 | 1.99 | 19.05 | 205.10 |

Part #3.

The performance comparisons of dgemm0, dgemm1, dgemm2, dgemm3:

| algorithm | n=64 | n=128 | n=256 | n=512 | n=1024 | n=2048 |
|-----------|------|-------|-------|-------|--------|--------|
| dgemm0 | NA | 0.10 | 0.09 | 0.07 | 0.06 | 0.03 |
| dgemm1 | NA | 0.21 | 0.15 | 0.10 | 0.09 | 0.05 |
| dgemm2 | NA | 1.68 | 1.49 | 1.45 | 0.90 | 0.67 |
| dgemm3 | NA | 3.36 | 3.36 | 2.08 | 1.72 | 1.13 |

2. Cache Reuse

Part 1.

For 10*10 matrices with ijk, ikj, jik, jki, kji, kji algorithm, cache misses of each elements are the same:

$$A_{[1,1]}$$
: 1, $A_{[1,2]}$: 0, $A_{[1,3]}$: 0, ..., $A_{[1,10]}$: 0
 $A_{[2,1]}$: 1, $A_{[2,2]}$: 0, $A_{[2,3]}$: 0, ..., $A_{[2,10]}$: 0

$$A_{[2,1]}: 1, A_{[2,2]}: 0, A_{[2,3]}: 0, ..., A_{[2,10]}: 0 A_{[3,1]}: 1, A_{[3,2]}: 0, A_{[3,3]}: 0, ..., A_{[3,10]}: 0$$

$$A_{[3,1]}$$
: 1, $A_{[3,2]}$: 0, $A_{[3,3]}$: 0, ..., $A_{[3,10]}$: 0

 $A_{[10,1]}$: 1, $A_{[10,2]}$: 0, $A_{[10,3]}$: 0, ..., $A_{[10,10]}$: 0

$$\begin{split} &B_{[1,1]} \colon 1, B_{[1,2]} \colon 0, B_{[1,3]} \colon 0, \dots, B_{[1,10]} \colon 0 \\ &B_{[2,1]} \colon 1, B_{[2,2]} \colon 0, B_{[3,3]} \colon 0, \dots, B_{[3,10]} \colon 0 \\ & \dots \\ &B_{[10,1]} \colon 1, B_{[10,2]} \colon 0, B_{[10,3]} \colon 0, \dots, B_{[10,10]} \colon 0 \\ &C_{[1,1]} \colon 1, C_{[1,2]} \colon 0, C_{[1,3]} \colon 0, \dots, C_{[1,10]} \colon 0 \\ &C_{[2,1]} \colon 1, C_{[2,2]} \colon 0, C_{[2,3]} \colon 0, \dots, C_{[2,10]} \colon 0 \\ & \dots \\ &C_{[10,1]} \colon 1, C_{[10,2]} \colon 0, C_{[10,3]} \colon 0, \dots, C_{[10,10]} \colon 0 \end{split}$$

The percentage of read cache misses: 1.4%

For 10000*10000 matrices with **ijk&jik** algorithm, cache misses of each elements:

$$\begin{array}{l} A_{[1,1]} \colon 10000, \, A_{[1,2]} \colon 0, \, \dots, A_{[1,11]} \colon 10000, \, A_{[1,12]} \colon 0, \, \dots, \, A_{[1,10000]} \colon 0 \\ A_{[2,1]} \colon 10000, \, A_{[2,2]} \colon 0, \, \dots, A_{[2,11]} \colon 10000, \, A_{[2,12]} \colon 0, \, \dots, \, A_{[2,10000]} \colon 0 \\ \dots \\ A_{[10000,1]} \colon 10000, \, A_{[10000,2]} \colon 0, \, \dots, A_{[10000,11]} \colon 10000, \, A_{[10000,12]} \colon 0, \, \dots, \, A_{[10000,10000]} \colon 0 \\ B_{[1,1]} \colon 10000, \, B_{[1,2]} \colon 10000, \, B_{[1,3]} \colon 10000, \, \dots, \, B_{[1,10000]} \colon 10000 \\ B_{[2,1]} \colon 10000, \, B_{[2,2]} \colon 10000, \, B_{[3,3]} \colon 10000, \, \dots, \, B_{[3,10000]} \colon 10000 \\ \dots \\ B_{[10000,1]} \colon 10000, \, B_{[10000,2]} \colon 10000, \, B_{[10000,3]} \colon 10000, \, \dots, \, B_{[10000,10000]} \colon 10000 \\ C_{[1,1]} \colon 1, \, C_{[1,2]} \colon 1, \, C_{[1,3]} \colon 1, \, \dots, \, C_{[1,10000]} \colon 1 \\ \dots \\ C_{[10000,1]} \colon 1, \, C_{[10000,2]} \colon 1, \, C_{[10000,3]} \colon 1, \, \dots, \, C_{[10000,10000]} \colon 1 \\ \end{array}$$

The percentage of read cache misses: 1.4%

For 10000*10000 matrices with **ikj&kij** algorithm, cache misses of each elements:

$$\begin{array}{l} A_{[1,1]} \colon 1, A_{[1,2]} \colon 0, \dots, A_{[1,11]} \colon 1, A_{[1,12]} \colon 0, \dots, A_{[1,10000]} \colon 0 \\ A_{[2,1]} \colon 1, A_{[2,2]} \colon 0, \dots, A_{[2,11]} \colon 1, A_{[2,12]} \colon 0, \dots, A_{[2,10000]} \colon 0 \\ \dots \\ A_{[10000,1]} \colon 1, A_{[10000,2]} \colon 0, \dots, A_{[10000,11]} \colon 1, A_{[10000,12]} \colon 0, \dots, A_{[10000,10000]} \colon 0 \\ B_{[1,1]} \colon 10000, B_{[1,2]} \colon 0, \dots, B_{[1,11]} \colon 10000, B_{[1,12]} \colon 0, \dots, B_{[1,10000]} \colon 0 \\ B_{[2,1]} \colon 10000, B_{[2,2]} \colon 0, \dots, B_{[2,11]} \colon 10000, B_{[2,12]} \colon 0, \dots, B_{[2,10000]} \colon 0 \\ \dots \\ B_{[10000,1]} \colon 10000, B_{[10000,2]} \colon 0, \dots, B_{[10000,11]} \colon 10000, B_{[10000,12]} \colon 0, \dots, B_{[10000,10000]} \colon 0 \end{array}$$

$$\begin{split} &C_{[1,1]}\text{: }1,C_{[1,2]}\text{: }0,...,C_{[1,11]}\text{: }1,C_{[1,12]}\text{: }0,...,C_{[1,10000]}\text{: }0\\ &C_{[2,1]}\text{: }1,C_{[2,2]}\text{: }0,...,C_{[2,11]}\text{: }1,C_{[2,12]}\text{: }0,...,C_{[2,10000]}\text{: }0\\ &...\\ &C_{[10000,1]}\text{: }1,C_{[10000,2]}\text{: }0,...,C_{[10000,11]}\text{: }1,C_{[10000,12]}\text{: }0,...,C_{[10000,10000]}\text{: }0 \end{split}$$

The percentage of read cache misses: 1.4%

For 10000*10000 matrices with **jki&kji** algorithm, cache misses of each elements:

$$\begin{array}{l} A_{[1,1]} \colon 10000, \, A_{[1,2]} \colon 10000, \, A_{[1,3]} \colon 10000, \dots, \, A_{[1,10000]} \colon 10000 \\ A_{[2,1]} \colon 10000, \, A_{[2,2]} \colon 10000, \, A_{[3,3]} \colon 10000, \dots, \, A_{[3,10000]} \colon 10000 \\ \dots \\ A_{[10000,1]} \colon 10000, \, A_{[10000,2]} \colon 10000, \, A_{[10000,3]} \colon 10000, \dots, \, A_{[10000,10000]} \colon 10000 \\ B_{[1,1]} \colon 1, \, B_{[1,2]} \colon 1, \, B_{[1,3]} \colon 1, \dots, \, B_{[1,10000]} \colon 1 \\ B_{[2,1]} \colon 1, \, B_{[2,2]} \colon 1, \, B_{[2,3]} \colon 1, \dots, \, B_{[2,10000]} \colon 1 \\ \dots \\ B_{[10000,1]} \colon 1, \, B_{[10000,2]} \colon 1, \, B_{[10000,3]} \colon 1, \dots, \, B_{[10000,10000]} \colon 1 \\ C_{[1,1]} \colon 10000, \, C_{[1,2]} \colon 10000, \, C_{[1,3]} \colon 10000, \dots, \, C_{[1,10000]} \colon 10000 \\ C_{[2,1]} \colon 10000, \, C_{[2,2]} \colon 10000, \, C_{[2,3]} \colon 10000, \dots, \, C_{[2,10000]} \colon 10000 \\ \dots \\ C_{[10000,1]} \colon 10000, \, C_{[10000,2]} \colon 10000, \, C_{[10000,3]} \colon 10000, \dots, \, C_{[10000,10000]} \colon 10000 \\ \end{array}$$

The percentage of read cache misses: 1.4%

<u>Part 2.</u>

For 10000*10000 matrices with **ijk&jik** blocked version algorithm, cache misses of each elements:

$$\begin{array}{l} A_{[1,1]} \colon 1, A_{[1,2]} \colon 0, A_{[1,3]} \colon 0, \ldots, \ A_{[1,11]} \colon 1, A_{[1,12]} \colon 0, A_{[1,13]} \colon 0, \ldots, A_{[1,10000]} \colon 0 \\ A_{[2,1]} \colon 1, A_{[2,2]} \colon 0, A_{[2,3]} \colon 0, \ldots, \ A_{[2,11]} \colon 1, A_{[2,12]} \colon 0, A_{[2,13]} \colon 0, \ldots, A_{[2,10000]} \colon 0 \\ \ldots \\ A_{[10000,1]} \colon 1, A_{[10000,2]} \colon 0, A_{[10000,3]} \colon 0, \ldots, A_{[10000,11]} \colon 1, A_{[10000,12]} \colon 0, A_{[10000,13]} \colon 0, \ldots, A_{[10000,10000]} \colon 0 \\ B_{[1,1]} \colon 1000, B_{[1,2]} \colon 0, B_{[1,3]} \colon 0, \ldots, B_{[1,11]} \colon 1000, B_{[1,12]} \colon 0, B_{[1,13]} \colon 0, \ldots, B_{[1,10000]} \colon 0 \\ B_{[2,1]} \colon 1000, B_{[2,2]} \colon 0, B_{[2,3]} \colon 0, \ldots, B_{[2,11]} \colon 1000, B_{[2,12]} \colon 0, B_{[2,13]} \colon 0, \ldots, B_{[2,10000]} \colon 0 \end{array}$$

$$\begin{array}{l} C_{[1,1]} \colon 1, C_{[1,2]} \colon 0, C_{[1,3]} \colon 0, \ldots, \ C_{[1,11]} \colon 1, C_{[1,12]} \colon 0, C_{[1,13]} \colon 0, \ldots, C_{[1,10000]} \colon 0 \\ C_{[2,1]} \colon 1, C_{[2,2]} \colon 0, C_{[2,3]} \colon 0, \ldots, \ C_{[2,11]} \colon 1, C_{[2,12]} \colon 0, C_{[2,13]} \colon 0, \ldots, C_{[2,10000]} \colon 0 \end{array}$$

. . .

$$C_{\tiny{[10000,1]}}\!\!: 1, C_{\tiny{[10000,2]}}\!\!: 0, C_{\tiny{[10000,3]}}\!\!: 0, \ldots, C_{\tiny{[10000,11]}}\!\!: 1, C_{\tiny{[10000,12]}}\!\!: 0, C_{\tiny{[10000,13]}}\!\!: 0, \ldots, C_{\tiny{[10000,10000]}}\!\!: 0$$

The percentage of read cache misses: 0.95%

For 10000*10000 matrices with **ikj&kij** blocked version algorithm, cache misses of each elements:

$$\begin{array}{l} A_{[1,1]} \colon 1, A_{[1,2]} \colon 0, A_{[1,3]} \colon 0, \ldots, A_{[1,11]} \colon 1, A_{[1,12]} \colon 0, A_{[1,13]} \colon 0, \ldots, A_{[1,10000]} \colon 0 \\ A_{[2,1]} \colon 1, A_{[2,2]} \colon 0, A_{[2,3]} \colon 0, \ldots, A_{[2,11]} \colon 1, A_{[2,12]} \colon 0, A_{[2,13]} \colon 0, \ldots, A_{[2,10000]} \colon 0 \end{array}$$

. . .

$$A_{[10000,1]}\!\!: 1, A_{[10000,2]}\!\!: 0, A_{[10000,3]}\!\!: 0, \ldots, A_{[10000,11]}\!\!: 1, A_{[10000,12]}\!\!: 0, A_{[10000,13]}\!\!: 0, \ldots, A_{[10000,10000]}\!\!: 0$$

$$\begin{array}{l} B_{[1,1]}\text{: }1000, B_{[1,2]}\text{: }0, B_{[1,3]}\text{: }0, \ldots, B_{[1,11]}\text{: }1000, B_{[1,12]}\text{: }0, B_{[1,13]}\text{: }0, \ldots, B_{[1,10000]}\text{: }0 \\ B_{[2,1]}\text{: }1000, B_{[2,2]}\text{: }0, B_{[2,3]}\text{: }0, \ldots, B_{[2,11]}\text{: }1000, B_{[2,12]}\text{: }0, B_{[2,13]}\text{: }0, \ldots, B_{[2,10000]}\text{: }0 \end{array}$$

...

$$\begin{array}{l} C_{[1,1]} \colon 1000, C_{[1,2]} \colon 0, C_{[1,3]} \colon 0, \ldots, C_{[1,11]} \colon 1000, C_{[1,12]} \colon 0, C_{[1,13]} \colon 0, \ldots, C_{[1,10000]} \colon 0 \\ C_{[2,1]} \colon 1000, C_{[2,2]} \colon 0, C_{[2,3]} \colon 0, \ldots, C_{[2,11]} \colon 1000, C_{[2,12]} \colon 0, C_{[2,13]} \colon 0, \ldots, C_{[2,10000]} \colon 0 \end{array}$$

. . .

The percentage of read cache misses: 0.95%

For 10000*10000 matrices with **kji&jki** blocked version algorithm, cache misses of each elements:

$$\begin{array}{l} A_{[1,1]} \colon 1000, A_{[1,2]} \colon 0, A_{[1,3]} \colon 0, \ldots, A_{[1,11]} \colon 1000, A_{[1,12]} \colon 0, A_{[1,13]} \colon 0, \ldots, A_{[1,10000]} \colon 0 \\ A_{[2,1]} \colon 1000, A_{[2,2]} \colon 0, A_{[2,3]} \colon 0, \ldots, A_{[2,11]} \colon 1000, A_{[2,12]} \colon 0, A_{[2,13]} \colon 0, \ldots, A_{[2,10000]} \colon 0 \end{array}$$

...

$$B_{[1,1]} \colon 1000, B_{[1,2]} \colon 0, B_{[1,3]} \colon 0, \dots, B_{[1,11]} \colon 1000, B_{[1,12]} \colon 0, B_{[1,13]} \colon 0, \dots, B_{[1,10000]} \colon 0$$

$$\begin{split} &B_{[2,1]}\text{: }1000,B_{[2,2]}\text{: }0,B_{[2,3]}\text{: }0,...,B_{[2,11]}\text{: }1000,B_{[2,12]}\text{: }0,B_{[2,13]}\text{: }0,...,B_{[2,10000]}\text{: }0\\ &\dots\\ &B_{[10000,1]}\text{: }1000,\ B_{[10000,2]}\text{: }0,\ B_{[10000,3]}\text{: }0,\ \dots,\ B_{[10000,11]}\text{: }1000,\ B_{[10000,12]}\text{: }0,\ B_{[10000,13]}\text{: }0\\ &C_{[1,1]}\text{: }1,C_{[1,2]}\text{: }0,C_{[1,3]}\text{: }0,...,C_{[1,11]}\text{: }1,C_{[1,12]}\text{: }0,C_{[1,13]}\text{: }0,...,C_{[1,10000]}\text{: }0\\ &C_{[2,1]}\text{: }1,C_{[2,2]}\text{: }0,C_{[2,3]}\text{: }0,...,C_{[2,11]}\text{: }1,C_{[2,12]}\text{: }0,C_{[2,13]}\text{: }0,...,C_{[2,10000]}\text{: }0\\ &\dots\\ &C_{[10000,1]}\text{: }1,C_{[10000,2]}\text{: }0,C_{[10000,3]}\text{: }0,...,C_{[10000,11]}\text{: }1,C_{[10000,12]}\text{: }0,C_{[10000,13]}\text{: }0,...,C_{[10000,10000]}\text{: }0 \end{split}$$

The percentage of read cache misses: 0.95%

Part 3.
n=2048

| | Simple Method | Block 8 | Block 16 | Block 32 | Block 64 |
|-----|------------------|---------|----------|----------|----------|
| ijk | 480.54 | 106.57 | 98.50 | 90.92 | 87.68 |
| jik | 596.16 | 107.93 | 104.78 | 91.83 | 88.38 |
| ikj | 837.26 | 116.31 | 105.85 | 112.72 | 102.61 |
| kij | 829.00 | 88.73 | 97.94 | 140.71 | 129.77 |
| jki | 567.85 | 155.22 | 131.80 | 123.86 | 122.00 |
| kji | 612.17 | 148.37 | 136.96 | 129.84 | 128.67 |

optimal block size is 64.

<u>Part 4.</u>

| | O0 | O1 | O2 | О3 |
|-------------------------|--------|-------|-------|--------------------|
| Register&Cache Reuse | 309.41 | 80.80 | 76.37 | <mark>76.19</mark> |