Luke Palmer MATH 5000 2006-12-04

(A) If  $\mathcal{A} = (A, +, \cdot)$  is a nonstandard model of Th  $(\omega, +, \cdot)$ , then the twin prime conjecture is true in  $(\omega, +, \cdot)$  if and only if there is a nonstandard  $a \in A$  such that  $\mathcal{A} \models$  "a and a + 2 are prime".

*Proof.* We have seen in an earlier homework how to define < and prime in  $(\omega, +, \cdot)$ . Let twin(x) be shorthand for  $prime(x) \wedge prime(x + 2)$ . Then we can state the twin prime conjecture as follows:  $\forall x \exists y \ (x < y \wedge twin(y))$ .

- ( $\Leftarrow$ ) Suppose there is a nonstandard  $a \in A$  with  $\mathcal{A} \models twin(a)$ . If the twin prime conjecture were false in  $(\omega, +, \cdot)$ , then there would be some  $b \in \omega$  where  $(\omega, +, \cdot) \models \neg \exists x \, (b < x \land twin(x))$ . But, of course, b can be written in the form  $1 + 1 + \cdots + 1$  a finite number of times, so that statement could be written for (and would be true in)  $\mathcal{A}$ , too, contradicting the existence of a.
- (⇒) Suppose that  $(\omega, +, \cdot) \models \forall x \exists y (x < y \land twin(y))$ . This sentence is also true in  $\mathcal{A}$ , so pick some nonstandard  $a \in A$ . Therefore, there must exist  $b \in A$  with a < b (so b is nonstandard) where  $\mathcal{A} \models twin(b)$ .  $\square$
- **(B)** Given an infinite set of primes P. There is a countable model  $\mathcal{B} = (B, +, \cdot)$  of Th  $(\omega, +, \cdot)$  and a  $b \in B$  such that for every prime  $p \in \omega$ ,  $\mathcal{B} \models \bar{p}|b$  if and only if  $p \in P$ .

*Proof.* Let a|b mean  $\exists n \ n \cdot a = b$ . Let  $\mathcal{L} = (+,\cdot,c)$ . Let  $\Sigma = \text{Th}(\omega,+,\cdot) \cup \{(\bar{p}|c)|p \in P\} \cup \{\neg(\bar{p}|c)|p \in \omega - P\}$ .

Given some finite  $F \subseteq \Sigma$ , F will have some sentences from Th  $(\omega, +, \cdot)$  and some sentences saying certain primes do and do not divide c. Build a model of F by letting the universe be  $\omega$  and c be the product of the primes which  $\Sigma$  says should divide c. So by the compactness theorem,  $\Sigma$  has a model  $\mathcal{B}$ , and  $c^{\mathcal{B}}$  has our desired properties by construction.  $\square$ 

**(Di)**  $\{\varphi_0, \varphi_1, \dots, \varphi_{n-1}\} \vdash \psi$  if and only if  $\vdash (\varphi_0 \to (\varphi_1 \to \dots (\varphi_{n-1} \to \psi)))$ .

*Proof.* By induction on n. The atomic case n=0 is trivial.  $\{\varphi_0, \varphi_1, \ldots, \varphi_n\} \vdash \psi$  iff (by IH)  $\{\varphi_0\} \vdash (\varphi_1 \to \cdots (\varphi_n \to \psi))$  iff (by the deduction theorem)  $\vdash (\varphi_0 \to \cdots (\varphi_n \to \psi))$ .

**(Dii)** If  $\Sigma \cup \{\varphi\} \vdash \psi$  and  $\Sigma \cup \{\neg \varphi\} \vdash \psi$ , then  $\Sigma \vdash \psi$ .

*Proof.* By (Di) we have that  $\Sigma \vdash (\varphi \to \psi)$  and that  $\Sigma \vdash (\neg \varphi \to \psi)$ , so  $\psi$  follows as a logical consequence (and can thus be proved by the completeness theorem).