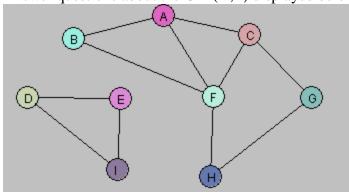
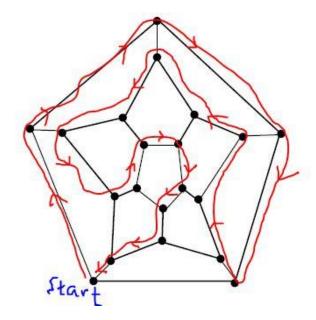
## **Lab 11 Solutions**

1. Answer questions about the G = (V,E) displayed below.



- A. Is the graph G connected? If not, what are the connected components for G? Solution: G is not connected. It has two connected components...
- B. Draw a spanning tree/forest for G. Solution: T = {DE, EI, FB, FA, FC, FH, GH}
- C. Is G a Hamiltonian graph? Solution: No, it has no Hamiltonian Cycle.
- D. Is there a Vertex Cover of size less than or equal to 5 for G? If so, what is the Vertex Cover?
  Solution: Yes. C={D, E, F, A, G}
- 2. Hamiltonian Graphs. The following graph has a Hamiltonian cycle. Find it.



3. Express in pseudo-code an algorithm which accepts as input a graph G and which outputs a vertex cover for G of smallest possible size. You may make use of the PowerSet algorithm without showing any pseudo-code details indicating how it works. Also, you may assume that your algorithm can make use of these operations freely:

computeEndpoints (e) //returns the two endpoints of the edge e belongs To(x, U) // returns true if vertex x belongs to set U; false otherwise Follow the rules for the pseudo-code language as completely as possible.

## **Solution:**

```
Algorithm: SmallestVertexCover
   Input: A graph G whose set of vertices is denoted V and set
     of edges is denoted E
  Output: Smallest size of a vertex cover U for G
  pow ← PowerSet(V)
  minCover ← V
  minVal ← |V|
   for each U in pow do
     isCover ← true
      //verify U is a vertex cover
     for each e in \mathbb{E} do
         (u,v) \leftarrow computeEndpoints(e)
        if( !(belongsTo(u,U) and !belongsTo(v,U))
           isCover ← false
      if(isCover and U.size() < minCover.size()) then</pre>
           minCover ← U
           minVal ← |U|
   return minVal
```