

CMSC 15200 Final 3A

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TOTAL POINTS

69 / 100

QUESTION 1

1 Data Representation 14 / 16

- 0 pts Correct
- ✓ - 1 pts DR a) 3
 - 1 pts Unsigned Min
 - 1 pts Unsigned Max
 - 1 pts Signed Min
 - 1 pts Signed Max
 - 1 pts Bitwise (1)
 - 1 pts Bitwise (2)
 - 1 pts Bitwise(3)
 - 1 pts Bitwise (4)
 - 1 pts Bitwise (5)
 - 1 pts Bitwise (6)
- ✓ - 1 pts Bitwise (7)
 - 1 pts DR a) 1
 - 2 pts DR a) 1
 - 7 pts Bitwise

QUESTION 2

2 Computational Complexity 3 / 10

- 0 pts Correct
- ✓ - 1 pts a) identify what variables influence execution time
 - 1 pts a) circle the instructions that will be executed most
 - 1 pts a) define an equation for execution time
- ✓ - 2 pts a) asymptotic complexity (big-O)
 - 1 pts b) identify what variables influence execution time
- ✓ - 1 pts b) circle the instructions that will be executed most
- ✓ - 1 pts b) define an equation for execution time
- ✓ - 2 pts b) asymptotic complexity (big-O)

QUESTION 3

3 Heaps 20 / 20

- ✓ - 0 pts Correct
 - 5 pts a
 - 5 pts b
 - 5 pts c
 - 5 pts d

QUESTION 4

4 Heap Array-Based Implementation 18 / 20

- 0 pts Correct
- 2 pts Bubble should start at 1, not 0
- 6 pts Not compare with parent
- ✓ - 2 pts Not swapping the last element
 - 2 pts Incorrect return
 - 2 pts Incorrect while loop
 - 2 pts Not decreasing num_filled
 - 20 pts Irrelevant to question
 - 2 pts no or incorrect boundary check
 - 4 pts the if statement should be in the while loop
 - 16 pts No bubbling or incorrect bubbling
 - 2 pts root is data[1]
 - 8 pts Incorrect bubbling for left and right separately
 - 2 pts Unnecessary swap during bubbling (swap only once on the same parent)
 - 8 pts Not swapping elements during bubbling
 - 6 pts Incorrect computation of children

QUESTION 5

5 Static variables implementation 12 / 25

- 0 pts Correct
- 1 pts undeclared variables
- 2 pts no static char array
- ✓ - 2 pts wrong buffer size
 - 3 pts operator handling
 - 3 pts set operation

- **3 pts** append operation
- ✓ - **5 pts** remove operation
- ✓ - **6 pts** insert operation
- **2 pts** print result
- **3 pts** buffer not allocated properly
- **1 pts** overly complicated code
- **1 pts** out of bounds access
- **2 pts** buffer declaration / initialization
- **2 pts** memory leak
- **5 pts** use of undefined functions
- **3 pts** inexistent string library function names

QUESTION 6

6 Hash Tables 2 / 9

- **0 pts** Correct
- ✓ - **2 pts** question (a)
- ✓ - **2 pts** question (b) change 1
- **2 pts** question (b) change 2
- ✓ - **1 pts** question (c) change 1
- ✓ - **1 pts** question (c) change 2
- ✓ - **1 pts** question (d)

Your name: Alex Compit

CS 152
Exam 3a
Winter 2019

Data Representation	/ 16
Algorithmic Complexity	/ 10
Heaps	/ 20
Array-based Implementation	/ 20
Static variables	/ 25
Hash Tables	/ 9

Total /100

1. No one may leave for any reason and come back to work on his/her test. If you need to go to the restroom, go now.
2. You may have 1 sheet of notes, hand-written, double-sided
3. Your backpack must be zipped up.
4. You may not use any electronic devices of any kind. Make sure your cell phone is turned off so it does not ring during the test.
5. You may not wear hats, hoods, sunglasses, or do anything that would obscure your eyes.
6. You may not sit near your partner if you are partnered in programming or lab (within 3 seats in any direction)
8. We will not answer questions during the exam.
9. Do not hold your test up. It must stay on your desk.
10. If you seem to be looking around, you will be moved to the front to allow you to look around without having other students' test in your field of vision.

$1111\ 1111\ 1101\ 1101$
 $2 \cdot 17 = 34$
 $17 = 16 + 1$
 $2 \cdot 17 = 34$
 $2 \cdot 17 = 34$
 $16 = 16$
 $32 = 32$
 $0001\ 0001\ 0001\ 0001$
 $0000\ 0000\ 0000\ 0000$

1. Data Representation (16 pts)

a) (5 pts) Each row represents the same number expressed in different formats. Fill in the blank spots.

Decimal	Unsigned binary (16 bits)	Signed 2's comp (16 bits)	Hex (4 digs)	Octal (6 digs)
34	0000 0000 0010 0010	X	0x0022	0.000042
-34	X	1111 1111 1101 1101	X	X
103	X	X	0x0037	X

b) (4 pts) If your computer represents integers in 11 bits, what are the ranges for integers? Express in decimal.

	MIN	MAX
unsigned	0	2 $2^{11} - 1$
signed 2's complement	-2^{10}	$2^{10} - 1$

Bitwise arithmetic – remember that bits are numbered from right to left, and numbering starts with 0.

c) (7 pts) For these questions, I have written a line of code in C or described what I want accomplished. Write one or a few lines of *efficient* C code that will accomplish the task. If you are supposed to set or extract a bit in a variable, do only that bit, not any in the rest of the number.

Code / description w/out bitwise ops	Code with bitwise ops
$x = y / 64;$ $64 = 2^6$	$x = y / 64;$ $x = y >> 6$ (shift y right 6 spots)
$x = y \% 16;$ $16 = 2^4$	$x = y \& 15$ (CLEAR ALL BITS TO LEFT OF FINAL FOUR bits)
Set bits 2 and 4 to 1 in variable a;	$a = a 20$ (because 20 represented as 10100 - 2 & 4 set to 1)
$x = y * 256;$ 2^8	$x = y << 8$
$x = y \% 32;$ $32 = 2^5$	$x = y \& 31$
Set bit 3 to 0 in variable a;	$a = a \& \sim (1 < 3)$
extract bit 7 in variable a and store into variable y	$a = (a >> 7) \& 1$ If $(a == 1) \{ y = y (1 < 7) \}$ // if a is 1, set bit 7 in y If $(a == 0) \{ y = y \& \sim (1 < 7) \}$ // if a is 0, clear bit 7

$2 \& 4 = 0$
 $4 \& 2 = 0$
 $1 \& 0 = 0$
 $16 + 4 = 20$
 $64 \& 2 = 0$
 $12 \& 8 = 8$
 $1 \& 2 = 0$

$0\ 32\ 10\ 10\ 10\ 10$
 $1111\ 1111\ 1101\ 1101$
 $1101\ 1101$
 $2/4 = 2$
 $1\ 32\ 64\ 7$
 103

$4\ 3\ 2\ 1\ 0$
 1111
 8421
 15
 $16 = 31$
 $0\ 0\ 0\ 0\ 4\ 2$
 $0000\ 0000\ 0010\ 0010$
 $0\ 0\ 2\ 2$
 $0x0022$
 $7 = 4 + 2 + 1 = 0111$
 $3 = 2 + 1 = 0011$
 $0011\ 0011$
 1432

2. Computational Complexity: (10 pts - 5 each)

Below are simplifications of actual algorithms used frequently in computer science:

For each problem:

- Identify what variable(s) is the problem's size (i.e. whose value(s) influence(s) the execution time)
- Circle the instruction that will be executed the most times.
- Create an equation to express how many times that instruction will be executed.
- Express the computational complexity in big-O notation

a)

```
for (c = 0; c < m; c++) {
  for (d = 1; d <= q; d *= 2) {
    sum += d * c;
  }
}
```

→ this will be executed the most times

→ the outer loop is $O(m)$

→ the inner loop is $O(\log_2 q)$

first loop executed m times

second loop executes $\log_2(q)$

Problem size:

c, m, d, q
 (c & d being incremented, m & q being bounds)

Equation:

$$m \cdot \log_2(q)$$

Complexity:

~~$O(m \log q)$~~
 $O(m \log n)$

Problem size:

n ~~d & q~~
~~is not~~

Equation (show your work)

the recursive portion
 executed n times

b)

```
void func(int *array, int n) {
  if (n <= 1)
    return;
  func(array, n-1); →  $O(n)$ 
  int tmp = array[n-1];
  int i = n-1;
  while ((i > 0) && (tmp < array[i-1])) {
    array[i] = array[i-1];
    i--;
  }
  array[i] = tmp;
}
```

→ $O(n)$

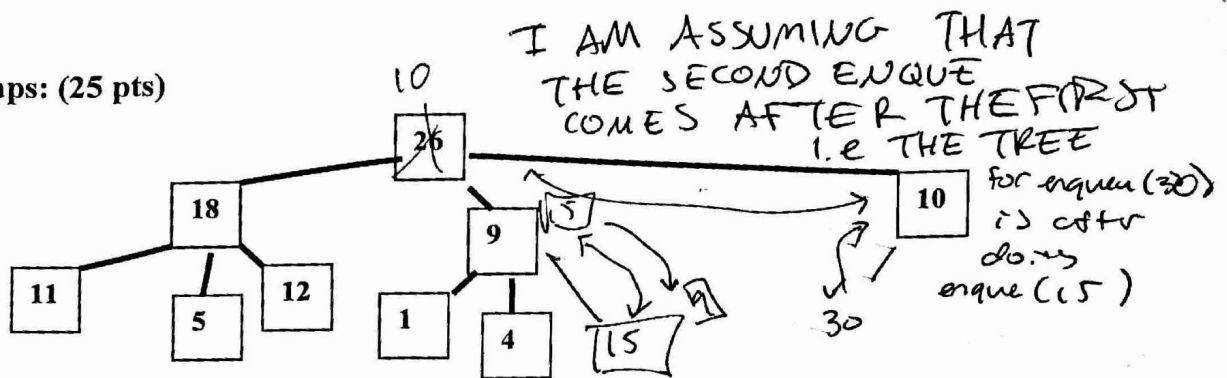
the while loop
 how many gets executed

Complexity:

$$O(n)$$

the second portion
 will never execute
 because it returns
 before it ever gets
 to the while loop,
 i.e. n keeps decreasing
 and once it gets to
 1, the function stops
 running

3. Heaps: (25 pts)

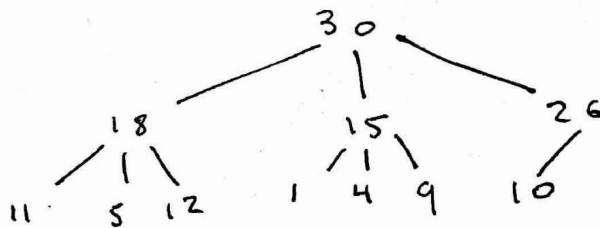


(5 pts each) Imagine that the heap is stored as a ternary tree in which every node has up to three children. Enqueue two numbers below. Draw the tree after each enqueue.

a) Enqueue(15)



b) Enqueue(30)

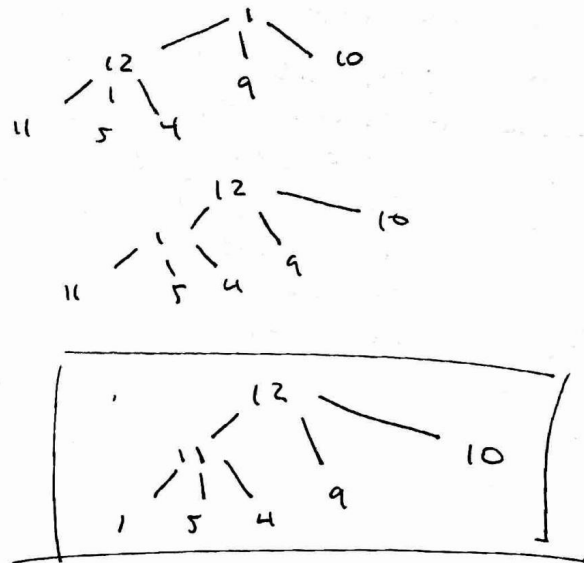
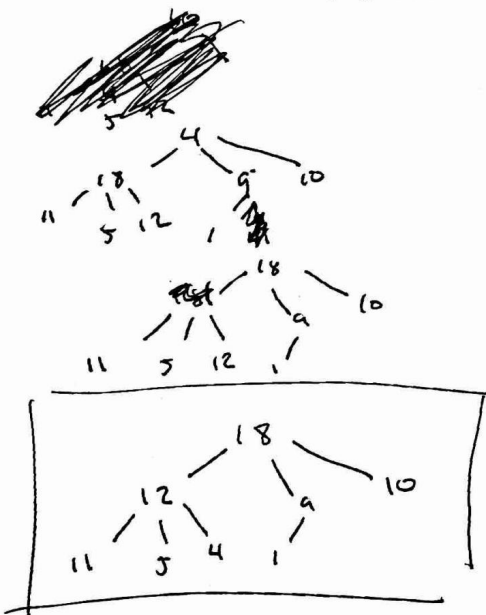


Now dequeue twice, but do so from the original tree above. Draw the tree after each dequeue

c) Dequeue()

~~AGAIN DEQUEUE~~

d) Dequeue()



6. Hash Tables (9 pts)

Provide the following explanations for how hash tables maintain $O(1)$ performance.

a. (2 pts) What causes the hash table to make a major change?

If by make a change this means correcting for issues in the hash table, we say that collisions in hashing cause something "extra" to be done to the hash table / function & running out of slots in the table cause something extra to be done.

b. (4 pts) What two changes occur to the hash table?

Change 1: collision resolution (e.g. linear probing, quadratic probing, double hashing).

Change 2: Resizing (increase or decrease the size of the hash table)

c. (4 pts) What root problem is each change solving? In other words, why do each of change 1 and change 2?

Change 1: ~~we need to~~ Since we have lots of (large) keys that are hashed to indices, it is possible that more than one key maps to the same place. Collision resolution allows us to find a new way to store the key if the slot it is mapped to is already taken.

Change 2: We need to provide $O(1)$ access, but this can only happen if the # of operations on the table is proportional to the size of the table. Otherwise, insertions, deletions, and searches border on taking closer to $O(n)$ time.

d. (1 pt) Why do these changes cost a lot of time?

These changes require us to perform operations with linear run time. On average, (ie based on the frequency with which they are carried out) however, we obtain an amortized run time of $O(1)$.

4. Heap Array-Based Implementation (20 pts)

Imagine you are implementing a function for an array-based heap. Each element in the array stores a pointer to the data. If there is no node for that particular location, then the pointer is set to NULL. Implement the dequeue function for a heap.

```
typedef unsigned int uint;
```

```
typedef struct {
    void **data;           // an array of void pointers, each pointing to a data item
    uint num_allocated;    // length of the array
    uint num_filled;       // number of items currently in the heap
    int (*compare)(void *, void *); // comparison function for items in the array
} heap_t;
```

```
} heap;
```

To access the data at the root: `h->data[1];`

```
void* dequeue(heap *h)
```

```

// If array is not allocated, then
// print "empty array"
if (ch == '\0')
{
    printf("empty array\n");
}

```

if (h → num_allocated == 1)

```

3 XXXXXXXXXX h → num - filled - - ;
    h → num - filled - - ;
    m = h → data [i];
    h → data [i] = NULL;
    return m; }

```

11 extract min

```
min = h → data[i];
```

```
11 set new data
```

$h \rightarrow data[1] = h \rightarrow data[h \rightarrow num_allocated - 1]$

$\text{h} \rightarrow \text{num} - \frac{\text{size}}{2}$

11 Reform bubble

bubble (h)

return pointer to deque element
return min;

(should be manifested)

* DOESN'T SAY IF THIS IS A MIN OR MAX HEAP - used min & max
 → void bubble (heap * h)

void bubble (h_{top} * h)

```
{
    uint i=1
    uint tmp=1
    uint h, v, w, l
    void tmp2;
```

while (~~h~~ h \rightarrow num-filled)

```

// start while loop
{
    l = 2 * i + 1;
    r = 2 * i + 2;
    if (l < h) {
        swap(arr[l], arr[r]);
        l++;
        r++;
    }
}

```

$$\{ \text{true} = \text{true} \}$$

$\{ \text{tmp} = l; j \}$
 if ($r < h \rightarrow s: \text{be}$ & $\text{comp}(h \rightarrow \text{data}[r], h \rightarrow \text{data}[i] = -1)$)

{ temp > j; } set a minimum
 (know temp = i)
 if (temp is not minimum, swap)
 temp = ~~temp~~ data[i];

```
h->data[i] = h->data[tmp];
```

```
h->data[tmp] = tmp2;
```

this is the bubble part

 $i = \text{temp}$

3

~~void delete (h)~~
void delqueue (heap)

NOTES

{ // first set to null, decrease size of array

~~h->data[i] = NULL~~ ~~h->num_allocated--~~ ~~tmp = h->data[i]~~
while (i > 0)

if (~~h->num_allocated~~ <= 0)

return; // error

if (h->num_allocated == 1)

{ h->num_allocated = 0;

return h->data[i];

void ~~delete~~ r = h->data[i];

~~h->data[i] = h->data[h->size-1];~~

h->size--;

HEAP BUBBLE

{ return h->data[i];

void bubble (heap *h)

uint i = 1; void tmp2

uint l = 2*i+1

uint r = 2*i+2

uint tmp = l

if (l < h->size && ~~h->data[l] < h->data[i]~~ ~~cmp()~~ >= -1)

{ tmp = l;

if ~~tmp~~ r < tmp for r
tmp = r

if (tmp != i)

{ ~~tmp2~~ tmp2 = h->data[i];

h->data[tmp] = h->data[tmp2];

h->data[tmp2] = tmp2;

bubble (tmp); }

5. Static variables implementation (25 pts)

Static variables were an integral part of iterators. Now we will use them for a different purpose. You are writing an editor that stores one giant string – it can support 1024 characters. The editor can perform several actions, and each action modifies the string. There is no return value – the function prints the new string after each operation. Not all operations use all inputs. All operations use op. Set uses str, Insert uses str and start_index, Remove uses start_index and num_chars, and append uses str. You MAY use string library functions.

typedef unsigned int uint;

enum Operation { SET, INSERT, REMOVE, APPEND };

int calculate(enum Operation op, int rvalue);

int main()

```
{
    edit(SET, "How are you?");           // prints "How are you?"
    edit(APPEND, " Diana", 0, 0);        // prints "How are you? Diana"
    edit(REMOVE, NULL, 11, 1);           // prints "How are you Diana"
    edit(INSERT, ",", 11, 0);            // prints "How are you, Diana"
}
```

void edit(enum Operation op, char *str, uint start_index, uint num_chars)

```
{
    // initialize static variable
    Static static char out = NULL;
    // use switch statement to check match
    operation to code
    switch (op) {
        case SET:
            out = strcpy strcpy(out, str);
            break;
        case APPEND:
            out = strcat(out, str);
            break;
        case REMOVE:
            out = strdel(out, start_index, num_chars);
            out = strdel(out, start_index, num_chars);
            break;
        case INSERT:
            out = strinsrt(out, str, start_index);
            break;
    }
    // finally, print out the string
    printf("%s\n", out);
}
```

(not correct
library name,
but should
work + at
index is
second parameter)

~~Soft~~ Li
only keep

