Distributionally Robust Optimization and its Tractable Approximations with ROME

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Distributionally Robust Optimization

- · An approach toward optimization under uncertainty
 - · Partially characterized distribution
 - Historically called minimax stochastic programming (Žáčková, 1966)
 - Motivation: in practice we will never know the actual uncertainty distribution
- Contrast with classical Robust Optimization (RO)
 - Uncertainties in RO characterized by uncertainty set (support)
 - Ben-Tal and Nemirovski (1998), Bertsimas and Sim (2004)

Research Objectives

- Construct an approximate solution framework for a class of DRO problems.
 Features:
 - Multi-stage decisions
 - Distribution-free, but partially characterized by support, moments, and directional deviations (Chen, Sim, and Sun, 2007)
 - More flexible decision rules compared to AARC (Ben-Tal, Goryashko, Guslitzer, and Nemirovski, 2004), aka LDRs.
- To build a software modeling tool to solve DRO problems within this framework (ROME)

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Notation

- $\tilde{z} \in \Re^N$: uncertainties with distribution $\mathbb{P} \in \mathbb{F}$.
- $x \in \Re^n$: "here-and-now" decision variables.
- $\mathbf{y}^k(\tilde{\mathbf{z}}) \in \Re^{m_k}$: "wait-and-see" decision rules, $\forall k \in [K]$.
- Each decision rule need not depend on the full uncertainty vector. We introduce K information index sets, to capture the dependency structure, denoted by $I_k \subseteq [N]$, $\forall k \in [K]$.

$$\boldsymbol{y}^k \in \mathcal{Y}(m_k, N, I_k) \quad \forall k \in [K]$$

$$\mathcal{Y}(m,N,I) \triangleq \left\{ oldsymbol{f}: \Re^N
ightarrow \Re^m : oldsymbol{f} \left(oldsymbol{z} + \sum_{i
otin I} \lambda_i oldsymbol{e}^i
ight) = oldsymbol{f}(oldsymbol{z}), orall oldsymbol{\lambda} \in \Re^N
ight\}$$

- $\bullet \text{ E.g. } \frac{I_1 = \{2,3\}}{\boldsymbol{y}^1 \in \mathcal{Y}(m,N,I_1)} \implies \boldsymbol{y}^1(\tilde{\boldsymbol{z}}) = \boldsymbol{y}^1(\tilde{\boldsymbol{z}}_2,\tilde{\boldsymbol{z}}_3)$
- For multi-stage decisions with progressive information revelation, $I_1 \subset I_2 \subset \ldots \subset I_K$

Model of Uncertainty

- ullet Actual uncertainty distribution ${\mathbb P}$ lies in a family of distributions ${\mathbb F}$
- Partially characterized by distributional properties:
 - Support \mathcal{W} : Tractable conic representable, full-dim.
 - Mean Support $\hat{\mathcal{W}}$: Tractable conic representable.
 - Covariance matrix Σ : Positive definite
 - Upper bounds on Directional Deviations for stochastically independent components $(\boldsymbol{H}_{\sigma}, \boldsymbol{\sigma_f}, \boldsymbol{\sigma_b})$.

Formulation of General Model

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Linear Approximation of General Model

- General model is intractable.
 - How to choose recourse decisions?
- Following Ben-Tal et. al. (2004), we restrict ourselves to LDRs (or AARCs), where the recourse decisions are affinely dependent on the model uncertainties.
- Define the space of LDRs

$$\mathcal{L}(m,N,I) \triangleq \left\{ \boldsymbol{f}: \Re^N \to \Re^m: \exists \left(\boldsymbol{y}^0,\boldsymbol{Y}\right) \in \Re^m \times \Re^{m \times N}: \begin{array}{c} \boldsymbol{f}(\boldsymbol{z}) = \boldsymbol{y}^0 + \boldsymbol{Y}\boldsymbol{z} \\ \boldsymbol{Y}\boldsymbol{e}^i = \boldsymbol{0}, \forall i \notin I \end{array} \right\}$$

• Use this to approximate $\mathcal{Y}(m, N, I)$.

LDR Approximation of General Model

$$\begin{split} & \mathbf{Z_{LDR}} = \\ & \min_{\boldsymbol{x}, \{\boldsymbol{y}^k(\cdot)\}_{k=1}^K} \quad \boldsymbol{c^{0'}\boldsymbol{x}} + \sup_{\mathbb{P} \in \mathbb{F}} \mathrm{E}_{\mathbb{P}} \left(\sum_{k=1}^K \boldsymbol{d^{0,k'}\boldsymbol{y}^k(\tilde{\boldsymbol{z}})} \right) \\ & \text{s.t.} \qquad \boldsymbol{c^{l'}\boldsymbol{x}} + \sup_{\mathbb{P} \in \mathbb{F}} \mathrm{E}_{\mathbb{P}} \left(\sum_{k=1}^K \boldsymbol{d^{l,k'}\boldsymbol{y}^k(\tilde{\boldsymbol{z}})} \right) \leq b_l \quad \forall l \in [M] \\ & \boldsymbol{T(\tilde{\boldsymbol{z}})\boldsymbol{x}} + \sum_{k=1}^K \boldsymbol{U^k}\boldsymbol{y}^k(\tilde{\boldsymbol{z}}) = \boldsymbol{v}(\tilde{\boldsymbol{z}}) \\ & \underline{\boldsymbol{y}}^k \leq \boldsymbol{y}^k(\tilde{\boldsymbol{z}}) \leq \overline{\boldsymbol{y}}^k \qquad \forall k \in [K] \\ & \underline{\boldsymbol{x}} \geq \boldsymbol{0} \\ & \boldsymbol{y}^k \in \boldsymbol{\mathcal{L}}(\boldsymbol{m_k}, N, I_k) \qquad \forall k \in [K] \end{split}$$

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LDR Approx. of General Model (Explicit)

$$Z_{LDR} = \min_{\boldsymbol{x}, \left\{\boldsymbol{y}^{0,k}, \boldsymbol{Y}^{k}\right\}_{k=1}^{K}} \quad \boldsymbol{c}^{0'}\boldsymbol{x} + \sum_{k=1}^{K} \boldsymbol{d}^{0,k'}\boldsymbol{y}^{0,k} + \sup_{\hat{\boldsymbol{z}} \in \hat{\mathcal{W}}} \left(\sum_{k=1}^{K} \boldsymbol{d}^{0,k'}\boldsymbol{Y}^{k} \hat{\boldsymbol{z}}\right)$$
s.t.
$$\boldsymbol{c}^{l'}\boldsymbol{x} + \sum_{k=1}^{K} \boldsymbol{d}^{l,k'}\boldsymbol{y}^{0,k} + \sup_{\hat{\boldsymbol{z}} \in \hat{\mathcal{W}}} \left(\sum_{k=1}^{K} \boldsymbol{d}^{l,k'}\boldsymbol{Y}^{k} \hat{\boldsymbol{z}}\right) \leq b_{l} \quad \forall l \in [M]$$

$$\boldsymbol{T}^{0}\boldsymbol{x} + \sum_{k=1}^{K} \boldsymbol{U}^{k}\boldsymbol{y}^{0,k} = \boldsymbol{v}^{0}$$

$$\boldsymbol{T}^{j}\boldsymbol{x} + \sum_{k=1}^{K} \boldsymbol{U}^{k}\boldsymbol{Y}^{k}\boldsymbol{e}^{j} = \boldsymbol{v}^{j} \qquad \forall j \in [N]$$

$$\boldsymbol{y}^{k} \leq \boldsymbol{y}^{0,k} + \boldsymbol{Y}^{k}\boldsymbol{z} \leq \boldsymbol{\overline{y}}^{k} \qquad \forall \boldsymbol{z} \in \mathcal{W} \quad \forall k \in [K]$$

$$\boldsymbol{Y}^{k}\boldsymbol{e}^{j} = \boldsymbol{0} \qquad \forall j \notin I_{k}, \forall k \in [K]$$

$$\boldsymbol{x} \geq \boldsymbol{0}$$

Is LDR Approximation too Conservative?

- LDR approximation of General Model is tractable, but is it too conservative?
- We will proceed by using the basic LDR approximation as a starting point, and aim to find other more complex decision rules which are better.

$$Z_{GEN}^* \leq ??? \leq ??? \leq Z_{LDR}$$

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Segregating the Uncertainties

- Idea: to re-map the uncertainties into a higher-dimensional space, and apply an LDR on the new uncertainties, for more flexibility.
- E.g. for a scalar uncertainty, we can segregate into positive and negative parts (as was done by Chen, Sim, Sun, and Zhang, 2008).

$$\tilde{z} = \underbrace{\tilde{z}^+}_{\tilde{\zeta}_1} - \underbrace{\tilde{z}^-}_{\tilde{\zeta}_2}$$

• Can segregate more generally into distinct segments of the real line.

SLDR Approximation of General Model

$$\begin{split} & Z_{SLDR} = \\ & \min_{\boldsymbol{x}, \left\{\boldsymbol{r}^{0,k}, \boldsymbol{R}^{k}\right\}_{k=1}^{K}} \quad \boldsymbol{c}^{0'}\boldsymbol{x} + \sum_{k=1}^{K} \boldsymbol{d}^{0,k'}\boldsymbol{r}^{0,k} + \sup_{\hat{\boldsymbol{\zeta}} \in \hat{\boldsymbol{V}}} \left(\sum_{k=1}^{K} \boldsymbol{d}^{0,k'}\boldsymbol{R}^{k} \hat{\boldsymbol{\zeta}}\right) \\ & \text{s.t.} \qquad \boldsymbol{c}^{l'}\boldsymbol{x} + \sum_{k=1}^{K} \boldsymbol{d}^{l,k'}\boldsymbol{r}^{0,k} + \sup_{\hat{\boldsymbol{\zeta}} \in \hat{\boldsymbol{V}}} \left(\sum_{k=1}^{K} \boldsymbol{d}^{l,k'}\boldsymbol{R}^{k} \hat{\boldsymbol{\zeta}}\right) \leq b_{l} \quad \forall l \in [M] \\ & \boldsymbol{\mathcal{T}}^{0}\boldsymbol{x} + \sum_{k=1}^{K} \boldsymbol{U}^{k}\boldsymbol{r}^{0,k} = \boldsymbol{\nu}^{0} \\ & \boldsymbol{\mathcal{T}}^{j}\boldsymbol{x} + \sum_{k=1}^{K} \boldsymbol{U}^{k}\boldsymbol{R}^{k}\boldsymbol{e}^{j} = \boldsymbol{\nu}^{j} \qquad \forall j \in [N_{E}] \\ & \boldsymbol{y}^{k} \leq \boldsymbol{r}^{0,k} + \boldsymbol{R}^{k}\boldsymbol{\zeta} \leq \boldsymbol{\overline{y}}^{k} \qquad \forall \boldsymbol{\zeta} \in \boldsymbol{\mathcal{V}} \quad \forall k \in [K] \\ & \boldsymbol{R}^{k}\boldsymbol{e}^{j} = \boldsymbol{0} \qquad \qquad \forall j \notin \boldsymbol{\Phi}_{k}, \forall k \in [K] \\ & \boldsymbol{x} \geq \boldsymbol{0} \end{split}$$

Properties of SLDR

Theorem

For the general problem, $Z_{SLDR} \leq Z_{LDR}$

- Increases the size of the problem (and computational effort needed), but retains the linear structure.
- Can we do better?

$$Z_{GEN}^* \le ???? \le Z_{SLDR} \le Z_{LDR}$$

• Will base further improvement on the SLDR.

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Review of Deflected LDR (DLDR)

- Originally introduced by Chen, Sim, Sun, and Zhang (2008) to circumvent restrictiveness imposed by the LDR.
 - The structure of the set of constraints might allow some piecewise-linear decision rules to be used instead.
 - Apply bounds on expected positive part of an LDR to exploit piecewise-linearity.
- Bi-deflected LDR seeks to improve and extend the DLDR.
 - Improve: Reduce the objective even further.
 - Extend: Explore modeling scenarios (non-anticipativity, expectations in constraints) which the DLDR didn't handle.

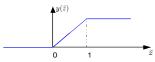
Motivating Example

- Scalar uncertainty \tilde{z} with unknown distribution $\mathbb P$ in a family $\mathbb F$.
- $\bullet \ \mathbb{F}$ defined by infinite support, zero mean, and unit variance. Consider the following problem:

$$\begin{aligned} \min_{y(\cdot)} & \sup_{\mathbb{P} \in \mathbb{F}} \operatorname{E}_{\mathbb{P}} \left(|y(\tilde{z}) - \tilde{z}| \right) \\ & 0 \leq y(\tilde{z}) \leq 1 \\ & y \in \mathcal{Y} \left(1, 1, \{1\} \right) \end{aligned}$$

• Solution is pretty straightforward: piecewise-linear recourse function:

$$y_{sol}(\tilde{z}) = \left\{ \begin{array}{ll} \tilde{z} & \text{if } 0 \leq \tilde{z} \leq 1 \\ 0 & \text{if } \tilde{z} \leq 0 \\ 1 & \text{if } \tilde{z} \geq 1 \end{array} \right.$$



Example Reformulated into DRO Framework

- Apply the identities $x = x^+ x^-, |x| = x^+ + x^-.$
- Model reformulates into DRO framework:

$$\begin{aligned} & \underset{y(\cdot)}{\min} & & \sup_{\mathbb{P} \in \mathbb{F}} \mathbb{E}_{\mathbb{P}} \left(|y(\tilde{z}) - \tilde{z}| \right) \\ & & 0 \leq y(\tilde{z}) \leq 1 \\ & & y \in \mathcal{Y} \left(1, 1, \{1\} \right) \\ & & & \vdots \\ & & \sup_{y(\cdot), u(\cdot), v(\cdot)} & \underset{\mathbb{P} \in \mathbb{F}}{\sup} \mathbb{E}_{\mathbb{P}} \left(u(\tilde{z}) + v(\tilde{z}) \right) \\ & \text{s.t.} & & u(\tilde{z}) - v(\tilde{z}) = y(\tilde{z}) - \tilde{z} \\ & & 0 \leq y(\tilde{z}) \leq 1 \\ & & u(\tilde{z}), v(\tilde{z}) \geq 0 \\ & & y, u, v \in \mathcal{Y} \left(1, 1, \{1\} \right) \end{aligned}$$

Using Different Decision Rules

- Using LDR, problem is infeasible, objective $= +\infty$.
- Using DLDR (of Chen et. al. 2008), we get the piecewise-linear decision rule:

$$\begin{array}{rcl} \hat{u}_{D}(z) & = & \left(u^{0} + uz\right)^{+} + \left(v^{0} + vz\right)^{-} \\ \hat{v}_{D}(z) & = & \left(v^{0} + vz\right)^{+} + \left(u^{0} + uz\right)^{-} \\ y(z) & = & y^{0} + yz \end{array}$$

After applying bounds, becomes an SOCP

$$Z_{DLDR} = \min_{u^{0}, u, v^{0}, v} \left\| \begin{pmatrix} u^{0} \\ u \end{pmatrix} \right\|_{2} + \left\| \begin{pmatrix} v^{0} \\ v \end{pmatrix} \right\|_{2}$$
s.t. $u - v = -1$
 $0 < u^{0} - v^{0} < 1$

• Objective = 1.

Consider this hypothetical decision rule, which satisfies the model constraints:

$$\hat{u}(z) = (u^{0} + uz)^{+} + (v^{0} + vz)^{-} + (y^{0} + yz)^{-}
\hat{v}(z) = (v^{0} + vz)^{+} + (u^{0} + uz)^{-} + (y^{0} - 1 + yz)^{+}
\hat{y}(z) = (y^{0} + yz)^{+} - (y^{0} - 1 + yz)^{+}$$

Applying robust bounds, problem transforms into the SOCP:

$$\min_{\substack{y^0,y,u^0,u,v^0,v\\\text{s.t.}}} \quad \left\| \left(\begin{array}{c} u^0 \\ u \end{array} \right) \right\|_2 + \left\| \left(\begin{array}{c} v^0 \\ v \end{array} \right) \right\|_2 + \frac{1}{2} \left\| \left(\begin{array}{c} y^0 \\ y \end{array} \right) \right\|_2 + \frac{1}{2} \left\| \left(\begin{array}{c} y^0-1 \\ y \end{array} \right) \right\|_2 - \frac{1}{2}$$

- Objective = 0.707 (better!)
- Aim to find algorithm to automatically predict this decision rule

Two-stage Bi-deflected LDR Problem Setup

Consider a simpler two-stage model first:

$$\begin{aligned} \min_{\boldsymbol{x}, \boldsymbol{r}(\cdot)} \quad & \boldsymbol{c}' \boldsymbol{x} + \sup_{\mathbb{P} \in \mathbb{F}} \mathbb{E}_{\mathbb{P}} \left(\boldsymbol{d}' \boldsymbol{r}(\tilde{\boldsymbol{\zeta}}) \right) \\ \text{s.t.} \quad & \boldsymbol{\mathcal{T}}(\tilde{\boldsymbol{\zeta}}) \boldsymbol{x} + \boldsymbol{U} \boldsymbol{r}(\tilde{\boldsymbol{\zeta}}) = \boldsymbol{\nu}(\tilde{\boldsymbol{\zeta}}) \\ & \underline{\boldsymbol{y}} \leq \boldsymbol{r}(\tilde{\boldsymbol{\zeta}}) \leq \overline{\boldsymbol{y}} \\ & \boldsymbol{r} \in \mathcal{L}(m, N_E, [N_E]) \end{aligned}$$

• Define the index sets of non-infinite bounds:

$$\frac{\underline{J}}{\overline{J}} = \left\{ i \in [m] : \underline{y}_i > -\infty \right\}$$

$$\overline{J} = \left\{ i \in [m] : \overline{y}_i < +\infty \right\}$$

Two-stage BDLDR Sub-problems

• We solve a series of sub-problems:

$$\begin{array}{lll} \min\limits_{\boldsymbol{p}} & \boldsymbol{d'p} & \qquad & \frac{\forall i \in \underline{J}}{J} \\ \min\limits_{\boldsymbol{p}} & \boldsymbol{d'p} & \qquad & \min\limits_{\boldsymbol{q}} & \boldsymbol{d'q} \\ \text{s.t.} & \boldsymbol{Up} = \mathbf{0} & \text{s.t.} & \boldsymbol{Uq} = \mathbf{0} \\ & p_i = 1 & \qquad & q_i = -1 \\ & p_j \geq 0 & \forall j \in \underline{J} \\ & p_j \leq 0 & \forall j \in \underline{J} \setminus \{i\} & \qquad & q_j \leq 0 & \forall j \in \overline{J} \setminus \{i\} \end{array}$$

Two-stage BDLDR

Consider an LDR which satisfies the relaxed constraints:

$$\begin{split} \mathcal{T}(\tilde{\zeta})x + Ur(\tilde{\zeta}) &= \nu(\tilde{\zeta}) \\ r_j(\tilde{\zeta}) &\geq \underline{y}_j & \forall j \in \underline{J} \setminus \underline{J}^\circ \\ r_j(\tilde{\zeta}) &\leq \overline{y}_j & \forall j \in \overline{J} \setminus \overline{J}^\circ \end{split}$$

The two-stage BDLDR is defined as:

$$\hat{\boldsymbol{r}}(\tilde{\boldsymbol{\zeta}}) \triangleq \underbrace{\boldsymbol{r}(\tilde{\boldsymbol{\zeta}})}_{\text{Linear Part}} + \sum_{i \in \underline{J}^{\circ}} \left(r_i(\tilde{\boldsymbol{\zeta}}) - \underline{y}_i\right)^{-} \underbrace{\bar{\boldsymbol{p}}^i}_{j} + \sum_{i \in \overline{J}^{\circ}} \left(r_i(\tilde{\boldsymbol{\zeta}}) - \overline{y}_i\right)^{+} \underbrace{\bar{\boldsymbol{q}}^i}_{j}$$
Optimal solution of the i^{th} sub-problem (Deflected Part)

Two-stage BDLDR Properties

Theorem

The Bi-Deflected Linear Decision Rule, $\hat{r}(\tilde{\zeta})$, satisfies the following properties

- 1. $U\hat{r}(\tilde{\zeta}) = Ur(\tilde{\zeta})$
- 2. $\mathbf{y} \leq \hat{\mathbf{r}}(\tilde{\boldsymbol{\zeta}}) \leq \overline{\mathbf{y}}$
 - For Reference:

$$\hat{\boldsymbol{r}}(\tilde{\boldsymbol{\zeta}}) \triangleq \underbrace{\boldsymbol{r}(\tilde{\boldsymbol{\zeta}})}_{\mathsf{Linear\ Part}} + \sum_{i \in \underline{J}^{\circ}} \left(r_{i}(\tilde{\boldsymbol{\zeta}}) - \underline{y}_{i}\right)^{-} \underbrace{\bar{\boldsymbol{p}}^{i}}_{} + \sum_{i \in \overline{J}^{\circ}} \left(r_{i}(\tilde{\boldsymbol{\zeta}}) - \overline{y}_{i}\right)^{+} \underbrace{\bar{\boldsymbol{q}}^{i}}_{}$$
 Optimal solution of the i^{th} sub-problem (Deflected Part)

BDLDR Approximation of Two-stage Problem

• Using the BDLDR, the problem becomes:

$$\begin{split} Z_{BDLDR} &= \\ & \min_{\boldsymbol{x}, \boldsymbol{r}(\cdot)} \quad \boldsymbol{c}' \boldsymbol{x} + \sup_{\mathbb{P} \in \mathbb{F}} \mathrm{E}_{\mathbb{P}} \left(\boldsymbol{d}' \boldsymbol{r}(\tilde{\boldsymbol{\zeta}}) \right) \\ & + \sum_{i \in J_{R}^{\circ}} \sup_{\mathbb{P} \in \mathbb{F}} \mathrm{E}_{\mathbb{P}} \left(\left(r_{i}(\tilde{\boldsymbol{\zeta}}) - \underline{y}_{i} \right)^{-} \right) \boldsymbol{d}' \bar{\boldsymbol{p}}^{i} \\ & + \sum_{i \in \overline{J}_{R}^{\circ}} \sup_{\mathbb{P} \in \mathbb{F}} \mathrm{E}_{\mathbb{P}} \left(\left(r_{i}(\tilde{\boldsymbol{\zeta}}) - \overline{y}_{i} \right)^{+} \right) \boldsymbol{d}' \bar{\boldsymbol{q}}^{i} \\ & \text{s.t.} \\ & \boldsymbol{\mathcal{T}}(\tilde{\boldsymbol{\zeta}}) \boldsymbol{x} + \boldsymbol{U} \boldsymbol{r}(\tilde{\boldsymbol{\zeta}}) = \boldsymbol{\nu}(\tilde{\boldsymbol{\zeta}}) \\ & r_{j}(\tilde{\boldsymbol{\zeta}}) \geq \underline{y}_{j} & \forall j \in \underline{J} \setminus \underline{J}^{\circ} \\ & r_{j}(\tilde{\boldsymbol{\zeta}}) \leq \overline{y}_{j} & \forall j \in \overline{J} \setminus \overline{J}^{\circ} \\ & \boldsymbol{r} \in \mathcal{L}(m, N_{E}, [N_{E}]) \end{split}$$

Performance of Two-stage BDLDR

Theorem

For the two-stage problem, $Z_{BDLDR} \leq Z_{DLDR} \leq Z_{LDR}$.

- Proof is a bit tedious, but the basic idea is that the BDLDR includes the DLDR as a special case, and hence is more flexible and therefore yields a lower objective.
- Next, to generalize the BDLDR for use together with non-anticipativity requirements and expectation constraints.

BDLDR for General Model

Define the following sets $\forall k \in [K]$:

$$N^{+}(k) = \{j \in [K] : \Phi_{k} \subseteq \Phi_{j}\}$$

$$N^{-}(k) = \{j \in [K] : \Phi_{j} \subseteq \Phi_{k}\}$$

Similarly, we have two sub-problems:

ilarly, we have two sub-problems:
$$\begin{array}{ll} \forall k \in [K], \forall i \in \underline{J}_k \\ \overline{\sum_{j \in N^+(k)} \boldsymbol{U}^j \boldsymbol{p}^{i,k,j}} &= & \mathbf{0} \\ p_l^{i,k,j} &\geq & 0 & \forall l \in \underline{J}_j & \forall j \in N^+(k) \\ p_l^{i,k,j} &\leq & 0 & \forall l \in \overline{J}_j & \forall j \in N^+(k) \setminus \{k\} \\ p_l^{i,k,k} &\leq & 0 & \forall l \in \overline{J}_k \setminus \{i\} \\ p_l^{i,k,k} &\leq & 0 & \forall l \in \overline{J}_k \setminus \{i\} \\ p_l^{i,k,k} &= & 1 \\ \end{array}$$

$$\frac{\forall k \in [K], \forall i \in \overline{J}_k}{\sum_{j \in N^+(k)} \boldsymbol{U}^j \boldsymbol{q}^{i,k,j}} = \boldsymbol{0} \\ q_l^{i,k,j} &\leq & 0 & \forall l \in \overline{J}_j & \forall j \in N^+(k) \\ q_l^{i,k,j} &\geq & 0 & \forall l \in \underline{J}_j & \forall j \in N^+(k) \setminus \{k\} \\ q_l^{i,k,k} &\geq & 0 & \forall l \in \underline{J}_k \setminus \{i\} \\ q_l^{i,k,k} &= & -1 \end{array}$$

Constructing the BDLDR

Consider an LDR which satisfies the relaxed constraints

$$egin{aligned} \mathcal{T}(ilde{\zeta}) x + \sum_{k=1}^K U^k r^k(ilde{\zeta}) &= oldsymbol{
u}(ilde{\zeta}) \ r_j^k(ilde{\zeta}) &\geq y_j^k & orall k \in [K] \,, orall j \in \underline{J}_k \setminus \underline{J}_k^\circ \ r_j^k(ilde{\zeta}) &\leq \overline{y}_j^k & orall k \in [K] \,, orall j \in \overline{J}_k \setminus \overline{J}_k^\circ \ r^k \in \mathcal{L}(m_k, N_E, \Phi_k) & orall k \in [K] \end{aligned}$$

The non-anticipative BDLDR is defined as:

$$\hat{\boldsymbol{r}}^k(\tilde{\boldsymbol{\zeta}}) \triangleq \underbrace{\boldsymbol{r}^k(\tilde{\boldsymbol{\zeta}})}_{\text{Linear Part}} + \sum_{j \in N^-(k)} \left(\sum_{i \in \underline{J}_j^o} \left(r_i^j(\tilde{\boldsymbol{\zeta}}) - \underline{y}_i^j \right)^- \underline{\boldsymbol{p}}^{i,j,k} + \sum_{i \in \overline{J}_j^o} \left(r_i^j(\tilde{\boldsymbol{\zeta}}) - \overline{y}_i^j \right)^+ \underline{\boldsymbol{q}}^{i,j,k} \right)$$
Optimal solution of i^{th} sub-problem (Deflected Part)

BDLDR Properties

Theorem

Each non-anticipative BDLDR, $\hat{m{r}}^k(ilde{m{\zeta}})$ satisfies the following properties:

- 1. $\sum_{k=1}^K \boldsymbol{U}^k \boldsymbol{\hat{r}}^k(\boldsymbol{\tilde{\zeta}}) = \sum_{k=1}^K \boldsymbol{U}^k \boldsymbol{r}^k(\boldsymbol{\tilde{\zeta}})$
- 2. $\underline{\boldsymbol{y}}^k \leq \hat{\boldsymbol{r}}^k(\tilde{\boldsymbol{\zeta}}) \leq \overline{\boldsymbol{y}}^k, \quad \forall k \in [K]$
 - For reference:

$$\hat{\boldsymbol{r}}^k(\tilde{\boldsymbol{\zeta}}) \triangleq \underbrace{\boldsymbol{r}^k(\tilde{\boldsymbol{\zeta}})}_{\text{Linear Part}} + \sum_{j \in N^-(k)} \left(\sum_{i \in \underline{J}_j^\circ} \left(r_i^j(\tilde{\boldsymbol{\zeta}}) - \underline{y}_i^j \right)^- \underbrace{\boldsymbol{p}^{i,j,k}}_{i \in \overline{J}_j^\circ} + \sum_{i \in \overline{J}_j^\circ} \left(r_i^j(\tilde{\boldsymbol{\zeta}}) - \overline{y}_i^j \right)^+ \underbrace{\boldsymbol{q}^{i,j,k}}_{i \in \overline{J}_j^\circ} \right)^- \underbrace{\boldsymbol{p}^{i,j,k}}_{i \in \overline{J}_j^\circ} \left(r_i^j(\tilde{\boldsymbol{\zeta}}) - \overline{y}_i^j \right)^+ \underbrace{\boldsymbol{q}^{i,j,k}}_{i \in \overline{J}_j^\circ} \right)^- \underbrace{\boldsymbol{p}^{i,j,k}}_{i \in \overline{J}_j^\circ} \left(r_i^j(\tilde{\boldsymbol{\zeta}}) - \overline{y}_i^j \right)^+ \underbrace{\boldsymbol{q}^{i,j,k}}_{i \in \overline{J}_j^\circ} \left(r_i^j(\tilde{\boldsymbol{\zeta}}) - \overline{y}_i^j \right)^+ \underbrace{\boldsymbol{q$$

Optimal solution of i^{th} sub-problem (Deflected Part)

BDLDR Performance

Theorem

For the general problem, we have

$${Z_{GEN}}^* \le {Z_{BDLDR}} \le {Z_{SLDR}} \le {Z_{LDR}}$$

- The k^{th} BDLDR sums over $j \in N^-(k)$,
 - These are the indices which are contained by the current (k^{th}) information index set.
 - BDLDR does not violate non-anticipative of the LDR which it is based upon.

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Designing Software for RO Problems

- RO programs, especially with more complex decision rules (e.g. BDLDR), can be messy, typically involving the following steps:
 - Constructing robust counterparts,
 - Finding deflected components (non-anticipative),
 - Constructing robust bounds.
- Software Design Goals:
 - Mathematically intuitive modeling of RO problems
 - Rapid prototyping of RO ideas
 - Ease transition from RO theory to practice
 - Allow numerical studies

Robust Optimization Made Easy with ROME

- ROME: an algebraic modeling language in the MATLAB environment for modeling Robust Optimization (RO) problems
- ROME is primarily designed for structured RO problems, within the DRO framework presented.
 - Other MATLAB-based modeling languages (e.g. CVX, YALMIP) solve more general types of problems.
 - ROME focuses on modeling phenonena more specific to RO, e.g. uncertainty description, decision rules, non-anticipativity.
- ROME is a modeling language, does not do actual solving.
 - Calls underlying solver engines to do perform solve step
 - Present version supports CPLEX, MOSEK, and SDPT3 solver engines

Example: Robust Inventory Model

- Model a distribution-free, multi-period, inventory control problem with service constraints.
- Demand is exogenous, with unknown, but partially characterized distribution in a family F, defined by:
 - Covariance Matrix: Temporal demand correlation can be modeled by a non-diagonal covariance matrix.
 - Mean: Assume fixed mean (for simplicity).
 - Support: Maximum demand in each period.
- Backorders allowed, but in some applications, a penalty cost might not be a good model for stockouts.
 - In our model, we avoid stockouts with a constraint on the fill-rate.
 - Fill-rate constraint acts as a service guarantee to the consumers.

Model Parameters

Parameters

 $\begin{array}{lll} \text{Num Periods} & : & T \in \mathbb{N} \\ \text{Order Cost} & : & \boldsymbol{c} \in \Re^T \end{array}$

Holding Cost : $oldsymbol{h} \in \Re^T$

Uncertainties

Demand : $\tilde{\pmb{z}} \in \Re^T$

Decisions

Order Quantity : $x(\tilde{z}): x_t(\tilde{z}) \in \mathcal{L}(1,T,[t-1])$

Inventory Level : $y(\tilde{z}): y_t(\tilde{z}) \in \mathcal{L}(1,T,[t])$

Robust Fill Rate Constraint

$$\mathsf{Fill} \ \mathsf{Rate} = \frac{\mathsf{Expected} \ \mathsf{Sales}}{\mathsf{Expected} \ \mathsf{Demand}} \geq \beta$$

• Using our notation, the robust (worst-case) version:

$$\inf_{\mathbb{P}\in\mathbb{F}} \operatorname{E}_{\mathbb{P}} \left(\min \left\{ \tilde{z}_t, y_{t-1}(\tilde{z}) + x_t(\tilde{z}) \right\} \right) \ge \beta_t \mu_t$$

Apply inventory balance equation and re-arrange:

$$\sup_{\mathbb{P}\in\mathbb{F}} \mathcal{E}_{\mathbb{P}} \left(y_t(\tilde{z})^- \right) \le (1 - \beta_t) \mu_t$$

Model Formulation

• Family of uncertainties:

$$\mathbb{F} = \left\{ \mathbb{P} : \mathrm{E}_{\mathbb{P}} \left(\tilde{\boldsymbol{z}} \right) = \boldsymbol{\mu}, \mathrm{E}_{\mathbb{P}} \left(\tilde{\boldsymbol{z}} \tilde{\boldsymbol{z}}' \right) = \boldsymbol{\Sigma} + \boldsymbol{\mu} \boldsymbol{\mu}', \mathbb{P} \left(0 \leq \tilde{\boldsymbol{z}} \leq \boldsymbol{z}^{MAX} \right) = 1 \right\}$$

Robust Inventory Model with Fill Rate constraints:

$$\begin{aligned} \min_{\boldsymbol{x}(\cdot),\boldsymbol{y}(\cdot)} & \sup_{\mathbb{P} \in \mathbb{F}} \mathbb{E}_{\mathbb{P}} \left(\boldsymbol{c}' \boldsymbol{x}(\tilde{\boldsymbol{z}}) + \boldsymbol{h}' \left(\boldsymbol{y}(\tilde{\boldsymbol{z}}) \right)^{+} \right) \\ \text{s.t.} & \sup_{\mathbb{P} \in \mathbb{F}} \mathbb{E}_{\mathbb{P}} \left(\boldsymbol{y}(\tilde{\boldsymbol{z}})^{-} \right) \leq (\boldsymbol{I} - \operatorname{diag} \left(\boldsymbol{\beta} \right)) \boldsymbol{\mu} \\ & \boldsymbol{D} \boldsymbol{y} \left(\tilde{\boldsymbol{z}} \right) - \boldsymbol{x} \left(\tilde{\boldsymbol{z}} \right) = -\tilde{\boldsymbol{z}} \\ & \boldsymbol{0} \leq \boldsymbol{x} \left(\tilde{\boldsymbol{z}} \right) \leq \boldsymbol{x}^{MAX} \\ & \boldsymbol{x}_{t} \in \mathcal{L}(1,T,[t-1]) \quad \forall t \in [T] \end{aligned} \qquad \boldsymbol{D} = \begin{bmatrix} 1 & 0 & 0 & \dots & 0 \\ -1 & 1 & 0 & \dots & 0 \\ 0 & -1 & 1 & \dots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \dots & 1 \end{bmatrix}$$

Model Transformed into DRO Framework

After linearizing piecewise-linear terms,

$$\begin{aligned} & \min_{\boldsymbol{x}(\cdot), \boldsymbol{y}(\cdot), \boldsymbol{r}(\cdot), \boldsymbol{s}(\cdot)} & \sup_{\mathbb{P} \in \mathbb{F}} \mathbb{E}_{\mathbb{P}} \left(\boldsymbol{c}' \boldsymbol{x}(\tilde{\boldsymbol{z}}) + \boldsymbol{h}' \boldsymbol{r}(\tilde{\boldsymbol{z}}) \right) \\ & \text{s.t.} & \sup_{\mathbb{P} \in \mathbb{F}} \mathbb{E}_{\mathbb{P}} \left(\boldsymbol{s}(\tilde{\boldsymbol{z}}) \right) \leq \left(\boldsymbol{I} - \mathsf{diag} \left(\boldsymbol{\beta} \right) \right) \boldsymbol{\mu} \\ & \boldsymbol{r}(\tilde{\boldsymbol{z}}) \geq \boldsymbol{y}\left(\tilde{\boldsymbol{z}} \right) \\ & \boldsymbol{s}\left(\tilde{\boldsymbol{z}} \right) \geq - \boldsymbol{y}\left(\tilde{\boldsymbol{z}} \right) \\ & \boldsymbol{s}\left(\tilde{\boldsymbol{z}} \right) \geq - \boldsymbol{y}\left(\tilde{\boldsymbol{z}} \right) \\ & \boldsymbol{r}\left(\tilde{\boldsymbol{z}} \right), \boldsymbol{s}\left(\tilde{\boldsymbol{z}} \right) \geq \boldsymbol{0} \\ & \boldsymbol{D} \boldsymbol{y}\left(\tilde{\boldsymbol{z}} \right) - \boldsymbol{x}\left(\tilde{\boldsymbol{z}} \right) = - \tilde{\boldsymbol{z}} \\ & \boldsymbol{0} \leq \boldsymbol{x}\left(\tilde{\boldsymbol{z}} \right) \leq \boldsymbol{x}^{MAX} \\ & \boldsymbol{x}_{t} \in \mathcal{L}(1, T, [t-1]) & \forall t \in [T] \\ & \boldsymbol{r}_{t}, \boldsymbol{s}_{t}, \boldsymbol{y}_{t} \in \mathcal{L}(1, T, [t]) & \forall t \in [T] \end{aligned}$$

ROME Code for Inventory Problem

```
% model parameters
T = 10:
                       % planning horizon
c = 1 *ones(T. 1): % order cost rate
hcost = 2*ones(T. 1): % holding cost rate
beta = 0.50*ones(T, 1); % minimum fillrate in each period
xMax = 100*ones(T, 1); % maximum order quantity in each period
alpha = 0.5;
               % temporal autocorrelation factor
L = alpha * tril(ones(T), -1) + eye(T); % autocorrelation matrix
% numerical uncertainty parameters
zMax = 105*ones(T, 1); % maximum demand in each period
zMean = 30*ones(T, 1); % mean demand in each period
zCovar = 20*(L * L'); % temporal demand covariance
% differencing matrix
D = eye(T) - diag(ones(T-1, 1), -1);
% dependency structure
pX = logical([tril(ones(T)), zeros(T, 1)]);
```

ROME Code for Inventory Problem (cont.)

```
% Step 3: BDLDR Method
h = rome_begin('Robust Inventory (BDLDR)');
tic;
% declare uncertainties
newvar z(T) uncertain nonneg;
% define uncertainty parameters
% define LDR variables
newvar x(T, z, 'Pattern', pX) linearrule; % order quantity
newvar v(T, z) linearrule;
                                                   % inventory level
                                                  r\left(	ilde{oldsymbol{z}}
ight), s\left(	ilde{oldsymbol{z}}
ight) \geq \mathbf{0}
% define auxilliary variables
newvar r(T, z) s(T, z) linearrule nonneg; x_t \in \mathcal{L}(1,T,[t-1]) \quad \forall t \in [T] r_t, s_t, y_t \in \mathcal{L}(1,T,[t]) \quad \forall t \in [T]
```

ROME Code for Inventory Problem (cont.)

```
% auxilliary constraints
rome_constraint(r >= y); % since r >= y^+
rome_constraint(s >= -y); % since s >= y^-
% fillrate constraint
rome_constraint(mean(s) <= diag(ones(T, 1) - beta) * zMean);
                                                                    min
                                                                                      \sup \mathcal{E}_{\mathbb{P}}\left(\boldsymbol{c}'\boldsymbol{x}(\tilde{\boldsymbol{z}}) + \boldsymbol{h}'\boldsymbol{r}(\tilde{\boldsymbol{z}})\right)
                                                            x(\cdot), y(\cdot), r(\cdot), s(\cdot)
% inventory balance constraint
                                                                                      \sup E_{\mathbb{P}}(s(\tilde{z})) \leq (I - \mathsf{diag}(\beta))\mu
                                                                    s.t.
rome_constraint(D*v == x - z);
                                                                                      Dy(\tilde{z}) - x(\tilde{z}) = -\tilde{z}
% order quantity constraints
                                                                                      r(\tilde{z}) > y(\tilde{z})
rome_box(x, 0, xMax);
                                                                                      s\left(\tilde{z}\right) \geq -y\left(\tilde{z}\right)
                                                                                      \mathbf{0} < \boldsymbol{x}(\boldsymbol{\tilde{z}}) < \boldsymbol{x}^{MAX}
% objective
rome_minimize(c'*mean(x) + hcost'*mean(r));
% solve and display optimal objective
h.solve_deflected:
disp(sprintf('BDLDR Obj = %0.2f, time = %0.2f secs', h.ObjVal, toc));
x_sol_bdldr = h.eval(x)
rome_end:
```

ROME Output (MOSEK Solver)

```
EDU>> inventory fillrate example
Status: OPTIMAL
LDR Ob\dot{\gamma} = 1747.50, time = 0.25 secs
x sol ldr =
    97.500
    0.000 + 0.952 * z1
 -0.000 + 0.007*z1 + 0.945*z2
    0.000 + 0.005*z1 + 0.010*z2 + 0.937*z3
-0.000 + 0.004*z1 + 0.006*z2 + 0.013*z3 + 0.929*z4
    0.000 + 0.003*z1 + 0.005*z2 + 0.008*z3 + 0.015*z4 + 0.921*z5
-0.000 + 0.004*z1 + 0.004*z2 + 0.006*z3 + 0.009*z4 + 0.018*z5 + 0.911*z6
-0.000 + 0.003*z1 + 0.005*z2 + 0.006*z3 + 0.008*z4 + 0.012*z5 + 0.023*z6 + 0.897*z7
    0.000 + 0.004 \times 21 + 0.005 \times 22 + 0.006 \times 23 + 0.009 \times 24 + 0.011 \times 25 + 0.016 \times 26 + 0.031 \times 27 + 0.870 \times 28
-0.000 + 0.007*z1 + 0.007*z2 + 0.009*z3 + 0.011*z4 + 0.015*z5 + 0.019*z6 + 0.031*z7 + 0.058*z8 + 0.795*z9
Status: NEAR OPTIMAL
BDLDR Obi = 300.48, time = 0.58 secs
x sol bdldr =
      15.333
-0.000 + 1.003*z1 + 1.00(-0.000 + 1.003*z1)^- - 1.00(100.000 - 1.003*z1)^-
-0.000 + 0.000*z1 + 1.004*z2 + 1.00(-0.000 + 0.000*z1 + 1.004*z2)^- - 1.00(100.000 - 0.000*z1 - 1.004*z2)^-
-0.000 + 0.000*z1 - 0.000*z2 + 1.004*z3 + 1.00(-0.000 + 0.000*z1 - 0.000*z2 + 1.004*z3)^- - 1.00(100.000 - 0.000*z2)
-0.000 - 0.000*z1 - 0.000*z2 + 0.000*z3 + 1.004*z4 + 1.00(-0.000 - 0.000*z1 - 0.000*z2 + 0.000*z3 + 1.004*z4)^{-1}
   0.000 - 0.000*z1 - 0.000*z2 - 0.000*z3 + 0.000*z4 + 1.004*z5 + 1.00(0.000 - 0.000*z1 - 0.000*z2 - 0.000*z3 + 0.000*z3 +
-0.000 + 0.000*z1 - 0.000*z2 + 0.000*z3 + 0.000*z4 - 0.000*z5 + 1.004*z6 + 1.00(-0.000 + 0.000*z1 - 0.000*z2 + 0.000*z2
    0.000 - 0.000 \times z1 - 0.000 \times z2 + 0.000 \times z3 + 0.000 \times z4 - 0.000 \times z5 - 0.000 \times z6 + 1.004 \times z7 + 1.00(0.000 - 0.000 \times z1 - 0.000 \times 
-0.000 + 0.000*z1 - 0.000*z2 + 0.000*z3 + 0.000*z4 + 0.000*z5 - 0.000*z6 - 0.000*z7 + 1.004*z8 + 1.00(-0.000 + 0.000*z7)
-0.000 + 0.000*z1 + 0.000*z2 - 0.000*z3 + 0.000*z4 + 0.000*z5 + 0.000*z6 + 0.000*z7 + 0.000*z8 + 1.005*z9 +
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 Involuction
 Framework
 Segregated LDR
 Bi-Deflected LDR
 ROME
 Conclusion

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Outline

Introduction

Framework

Segregated LDF

Ri-Deflected I DR

ROME

Conclusion

Summary

- Different decision rules for Distributionally Robust Optimization:
 - LDRs: Tractable approximation, but poor performance
 - SLDRs: Improves performance, and retains linear structure
 - BDLDRs: Improves performance beyond SLDR, but makes problem messy
- ROME: Robust Optimization Made Easy
 - MATLAB-based modeling language for Distributionally Robust Optimization problems
 - Aims to allow robust optimization problems to be modeled and solved in a mathematically intuitive way
 - Freely-available for academic use from www.RobustOpt.com (together with User's Guide)