Optimization using a SAT solver

Bjørnar Luteberget

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SAT

- Boolean Satisfiability problem: SAT : F → {true, false} deciding if there is an assignment of to the variables of a Boolean formula such that the formula is satisfied (true).
- Consider the formula $\phi = (\mathbf{a} \vee \mathbf{b}) \wedge (\neg \mathbf{a} \vee \neg \mathbf{c})$
 - The assignment $b = \top$ and $c = \bot$ satisfies the formula.
 - $SAT(\phi) = true$
- Can be considered a form of mathematical programming, but with a different history and focus.

Satisfiability and validity

Asking whether

$$\phi(\vec{\mathbf{x}})$$

is satisfiable is the same as asking whether

$$\exists \vec{\mathbf{x}}.\phi(\vec{\mathbf{x}})$$

is true.

Satisfiability and validity

In automated theorem proving, it is often more interesting to check validity:

$$\forall \mathbf{x}.\phi(\mathbf{x})$$

Transform to a SAT problem:

$$\forall x. \phi(x) \Leftrightarrow \neg \exists x \neg \phi(x)$$
 and check $SAT(\neg \phi(x))$

• Special case, optimization (minimization):

$$\forall x (\phi(x) \to obj(x) > c) \Leftrightarrow \neg \exists x (\phi(x) \land obj(x) < c)$$



SAT solvers

- SAT solver: given a formula, find an assignment to the variables, or report that none exist.
- More recently: given a formula, determine its satisfiability and produce either a satisfying assignment or a proof that none exist.
- For a theorem $\forall x.\phi(x)$:
 - SAT $(\neg \phi(x)) = \top$ corresponds to a counter-example
 - SAT $(\neg \phi(\mathbf{x})) = \bot$ corresponds to a proof
- In constrast, in OR, infeasibility is often "just a bug".
 - 4% of instances of MIPLIB2017 are infeasible
 - 53% of instances of the SAT Competition 2022 were unsatisfiable

Conjunctive normal form

Most SAT solvers take only formulas given in conjunctive normal form

$$(x \vee \neg y \vee \ldots) \wedge (z \vee \neg w \vee \ldots) \wedge \ldots$$

- Variables and negated variables are called literals.
- Disjunctions of literals are called clauses.

$$\mathbf{x} \lor \mathbf{y} \lor \neg \mathbf{z} \qquad (\Leftrightarrow \mathbf{x} + \mathbf{y} + (1 - \mathbf{z}) \ge 1)$$

- CNF is a conjunction of clauses.
- By adding some variables, we can transform any propositional formula to CNF (with worst-case linear increase in size).



CNF encoding

- Good formulations (encodings) are still very important.
 - · Cardinality constraints are almost a sub-field in itself

$$x + y + \neg z + \ldots \leq k$$

Pseudo-Boolean constraints (almost binary ILP a.k.a 0, 1-programs)

$$a_1x_1+a_2(\neg x_2)+\ldots\leq k$$

- We are not worrying about this today
- $x + y \le 1$ can be encoded $\neg x \lor \neg y$
- $x + y + z \le 1 \Rightarrow (\neg x \lor \neg y) \land (\neg x \lor \neg z) \land (\neg y \lor \neg z)$

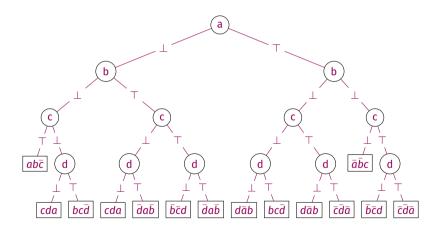


An example computation

Consider the formula

$$\phi(a,b,c,d) = (a \lor b \lor \neg c) \land \\ (b \lor c \lor \neg d) \land \\ (c \lor d \lor a) \land \\ (d \lor \neg a \lor b) \land \\ (\neg a \lor \neg b \lor c) \land \\ (\neg b \lor \neg c \lor d) \land \\ (\neg c \lor \neg d \lor \neg a) \land \\ (\neg d \lor a \lor \neg b)$$

Tree search (case splitting)

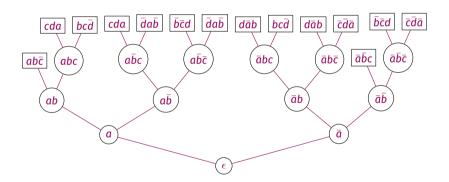


How to prove something

- Implicitly by algoritm: the algorithm is correct for all formulas. It returned UNSAT for this formula, so the formula is unsatisfiable.
 - What if the solver is incorrect?
 - Maybe an explicit proof is interesting in itself?
- ... or **explicitly** by defining a proof structure.
- Need a set of proof rules.
- For propositional logic, one is enough:

$$\frac{\mathbf{x}\vee\phi\quad\neg\mathbf{x}\vee\psi}{\phi\vee\psi} \text{ Resolution}$$

Proof tree: resolution steps



Conflict-driven clause learning

- Conflict-driven clause learning (CDCL) is an algorithm that attacks the problem from both sides: search and proof.
- Ingredients:
 - Heuristic assignment
 - Unit propagation
 - $\bullet \ \ \text{Conflict analysis} \to \text{clause learning} \\$

Trail (level=1):

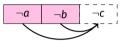
 $\neg a$

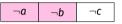
Guessing $\neg a$

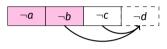
Trail (level=2):



Guessing $\neg b$

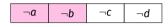








Trail (level=2):



Conflict clause: $a \lor c \lor d$

Trail (level=2):



Conflict clause: $a \lor c \lor d$

Trail (level=2):



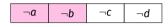
Conflict clause: $a \lor b \lor c$

Trail (level=2):

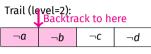


Conflict clause: $a \lor b \lor c$

Trail (level=2):



Conflict clause: $a \lor b$



Learnt new clause: $a \lor b$

 $\begin{array}{c} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \hline a \lor b \end{array}$



 $\begin{array}{l} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \neg d \lor a \lor \neg b \\ \hline a \lor b \end{array}$

Trail (level=1):

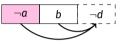
 $\neg a$

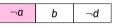


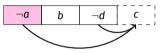
 $\begin{array}{l} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \neg d \lor a \lor \neg b \\ \hline a \lor b \end{array}$



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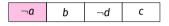


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 $\begin{array}{c} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \neg d \lor a \lor \neg b \\ \hline a \lor b \end{array}$

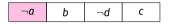
Trail (level=1):



Conflict clause: $\neg b \lor d \lor \neg c$

 $\begin{array}{c} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \neg d \lor a \lor \neg b \\ \hline a \lor b \end{array}$

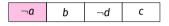
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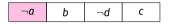
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Trail (level=1):



Conflict clause: $a \lor \neg b \lor d$

Trail (level=1):



Conflict clause: $a \lor \neg b \lor d$

 $\begin{array}{c} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \hline a \lor b \end{array}$

Trail (level=1):



Conflict clause: $a \lor \neg b$

 $\begin{array}{c} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \neg d \lor a \lor \neg b \\ \hline a \lor b \end{array}$

Trail (level=1):



Conflict clause: $a \lor \neg b$

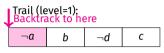
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Trail (level=1):



Conflict clause: a

 $\begin{array}{c} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \hline a \lor b \end{array}$



Learnt new clause: a

 $\begin{array}{l} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \hline a \lor b \end{array}$

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Trail (level=1):



 $\begin{array}{l} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \neg d \lor a \lor \neg b \\ a \\ \end{array}$

Trail (level=0):

 $\begin{array}{l} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \hline a \lor b \end{array}$

а

Trail (level=0):

a ;

 $a \lor b \lor \neg c$ $b \lor c \lor \neg d$ $c \lor d \lor a$ $d \lor \neg a \lor b$ $\neg a \lor \neg b \lor c$ $\neg b \lor \neg c \lor d$ $\neg c \lor \neg d \lor \neg a$ $\neg d \lor a \lor \neg b$ $a \lor b$

Trail (level=0):

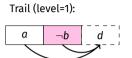
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 $\begin{array}{l} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \neg d \lor a \lor \neg b \\ a \\ \end{array}$

Trail (level=1):

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Guessing $\neg b$



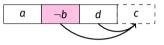
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Trail (level=1):

a -b d

 $\begin{array}{c} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \neg d \lor a \lor \neg b \\ a \lor b \end{array}$

Trail (level=1):



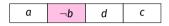
 $\begin{array}{l} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \hline a \lor b \end{array}$

а

Trail (level=1):

 $\begin{array}{l} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \neg d \lor a \lor \neg b \\ a \lor b \end{array}$

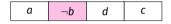
Trail (level=1):



Conflict clause: $\neg a \lor \neg d \lor \neg c$

 $\begin{array}{l} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \neg d \lor a \lor \neg b \\ a \lor b \end{array}$

Trail (level=1):



Conflict clause: $\neg a \lor \neg d \lor \neg c$

 $\begin{array}{c} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \hline a \lor b \end{array}$

Trail (level=1):



Conflict clause: $\neg a \lor b \lor \neg d$

 $a \lor b \lor \neg c$ $b \lor c \lor \neg d$ $c \lor d \lor a$ $d \lor \neg a \lor b$ $\neg a \lor \neg b \lor c$ $\neg b \lor \neg c \lor d$ $\neg c \lor \neg d \lor \neg a$ $\neg d \lor a \lor \neg b$ $a \lor b$

Trail (level=1):



Conflict clause: $\neg a \lor b \lor \neg d$

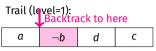
 $\begin{array}{l} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \neg d \lor a \lor \neg b \\ a \\ \end{array}$

Trail (level=1):



Conflict clause: $\neg a \lor b$

```
a \lor b \lor \neg c
b \lor c \lor \neg d
c \lor d \lor a
d \lor \neg a \lor b
\neg a \lor \neg b \lor c
\neg b \lor \neg c \lor d
\neg c \lor \neg d \lor \neg a
\neg d \lor a \lor \neg b
a
```



Learnt new clause: $\neg a \lor b$

 $\begin{array}{l} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \neg d \lor a \lor \neg b \\ a \\ \end{array}$

 $\neg a \lor b$

Trail (level=1):

 $\begin{array}{l} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \hline a \lor b \\ a \end{array}$

 $\neg a \lor b$

Trail (level=0):

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 $\begin{array}{l} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \hline a \lor b \\ a \end{array}$

 $\neg a \lor b$

Trail (level=0):



 $\begin{array}{c} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \neg d \lor a \lor \neg b \\ a \\ \end{array}$

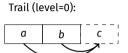
 $\neg a \lor b$

Trail (level=0):

a b

 $\begin{array}{c} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \neg d \lor a \lor \neg b \\ a \lor b \end{array}$

 $\neg a \lor b$



 $\begin{array}{l} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \neg d \lor a \lor \neg b \\ a \\ \end{array}$

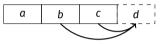
 $\neg a \lor b$

Trail (level=0):

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 $\begin{array}{c} a \lor b \lor \neg c \\ b \lor c \lor \neg d \\ c \lor d \lor a \\ d \lor \neg a \lor b \\ \neg a \lor \neg b \lor c \\ \neg b \lor \neg c \lor d \\ \neg c \lor \neg d \lor \neg a \\ \hline a \lor b \\ a \\ \neg a \lor b \end{array}$

Trail (level=0):



a $\neg a \lor b$

Trail (level=0):

a b c d

 $a \vee b \vee \neg c$ $b \lor c \lor \neg d$ $c \lor d \lor a$ $d \vee \neg a \vee b$ $\neg a \lor \neg b \lor c$ $\neg b \lor \neg c \lor d$ $\neg c \lor \neg d \lor \neg a$ $\neg d \lor a \lor \neg b$

а $\neg a \lor b$

 $a \vee b$

Trail (level=0):



Conflict clause: $\neg c \lor \neg d \lor \neg a$

Conflict on decision level zero, UNSAT!

CDCL proofs

The learnt clauses, in order, are a proof in the following sense:

- Take the first learnt clause, and assert its negation:
 - $x \lor y \Rightarrow \neg x \land \neg y$
 - A contradiction is derived using unit propagation only
- Consider the clause now a part of the formula (it is implied)
- After the last learnt clause has been added to the formula, a contradiction can be found using unit propagation.
- The formula, including the *original* clauses and some *implied* clauses, can be shown inconsistent by unit propagation ⇒ the original formula is *UNSAT*.

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World's longest maths proof: Solution to a 30-year-old problem would take 10 BILLION years to read - all for a prize of just \$100

- In 1980s American mathematician offered a prize to solve a maths problem
- Problem involves assigning sides of a right-angled triangle two colours
- Now 30 years later a team have solved it, with the help of a supercomputer
- But the solution is 200 terabytes of data, which is too large to read

By ABIGAIL BEALL FOR MAILONLINE PUBLISHED: 16:13 GMT, 11 July 2016 | UPDATED: 18:44 GMT, 11 July 2016





















In the 1980s an American mathematician named Robert Graham offered a prize of \$100 (677) to apyone able to solve a brain teasor that he could not solve himself



Search

Optimization

Finally: optimization

Optimization

- A few different problem definitions
 - MaxSAT: satisfy the maximum number of clauses of a SAT problem
 - (weighted) (partial) MaxSAT
 - Pseudo-boolean Optimization (0,1 ILP)
 - Weighted Boolean optimization
- The differences are not too interesting

min
$$\sum w_i x_i$$

s.t. $\phi(x_1, x_2, ...)$

Main MaxSAT approaches

- Branch and bound
- Stochastic local search
- SAT-based algorithms
 - Iterative SAT solving
 - Core-guided algorithms

Iterative SAT solving

- Find some way to translate the constraint $(obj(x) \le k) = (\sum w_i x_i \le k)$ to propositional logic
- Solve a sequence of SAT problems to find a k value s.t.
 - $SAT(\phi(x) \wedge obj(x) \leq k) = false$, and
 - $SAT(\phi(x) \land obj(x) \le k+1) = true$.
- Drawback: the objective can become a large formula

Core-guided algorithms

- (Assume $w_i \in \{0, 1\}$, and let S be the set of i's for which $w_i = 1$.)
- The special case k = 0 ($\Sigma_{i \in S} x_i \le 0$) doesn't require any large encodings:

$$\phi(\mathbf{x}) \wedge \bigwedge_{a \in A_0} \mathbf{a}$$

$$A_0 = \{ \neg \mathbf{x}_i \mid i \in S \}$$

$$A_0 = \{ \mathbf{x}_i < 0 \mid i \in S \}$$

• If SAT(...) = true - good, the optimum was zero.



Core-guided algorithms

• If UNSAT and the proof contains at least one clause

$$\neg C \subseteq A_0$$

- Example $A_0 = \{\neg x, \neg y, \neg z\}$ and $C = x \lor y$.
- We know that $obj(x) \ge 1$
- Can produce a new formula representing all $obj(x) \le 1$ solutions

$$\phi(\mathbf{X}) \wedge \bigwedge_{\mathbf{a} \in \mathbf{A}_1} \mathbf{a}$$
$$\mathbf{A}_1 = (\mathbf{A}_0 \setminus \neg \mathbf{C}) \cup$$

$$A_1 = (A_0 \setminus \neg C) \cup \{\Sigma x_i \le k + 1 \mid \Sigma x_i \le k \in \neg C\} \cup \{\Sigma_{x \in C} x \le 1\}$$

Assumptions

- We can make sure that $C \subseteq A$ by setting $x \in A$ as the first decisions.
- Call this assumptions
- Stop when a conflict at decision level $\leq |A|$
- The last learnt clause C will be a subset of A.

Example:

• Schedule three activities A, B, C with a resource limit of 2.

Activity	Duration	Resources	Earliest	Latest
A	1	1	0	3
В	1	2	0	3
C	2	1	0	2

• Optimize sum of starting times

 $\min \Sigma s_i$

CNF encoding of start time

- Unary encoding of start time of activities.
 - Boolean variables represent lower bounds:
 - $t_A^x := (s_A \ge x)$ $\neg t_A^x := (s_A < x)$
 - $t_A^1, t_A^2, t_A^3, t_B^1, t_B^2, t_B^3, t_C^1, t_C^2$
 - Constraints:
 - $t_A^3 \Rightarrow t_A^2$, $t_A^3 \Rightarrow t_A^2$, $t_B^3 \Rightarrow t_B^2$, $t_B^3 \Rightarrow t_B^2$, $t_C^2 \Rightarrow t_C^1$
 - Examples:
 - $[t_A^1, t_A^2, t_A^3] = [0, 0, 0] \Rightarrow s_A = 0$
 - $[t_{\Delta}^{1}, t_{\Delta}^{2}, t_{\Delta}^{3}] = [1, 0, 0] \Rightarrow s_{A} = 1$
 - $[t_A^1, t_A^2, t_A^3] = [1, 1, 0] \Rightarrow s_A = 2$
 - $[t_A^1, t_A^2, t_A^3] = [1, 1, 1] \Rightarrow s_A = 3$

CNF encoding of resource constraints

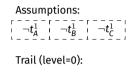
- Resource limit is 2
 - ⇒ activity B cannot run at the same time as A or C.
- Convert this to clauses over t's:
 - B start at time 0:
 - $t_B^1 \vee t_A^1$
 - $t_B^1 \vee t_C^1$
 - B start at time 1:
 - $t_{\mathsf{R}}^2 \vee \neg t_{\mathsf{R}}^1 \vee \neg t_{\mathsf{A}}^1 \vee t_{\mathsf{A}}^2$
 - $t_B^2 \vee \neg t_B^1 \vee t_C^2$
 - B start at time 2:
 - $t_B^3 \vee \neg t_B^2 \vee \neg t_A^2 \vee t_A^3$
 - $t_B^3 \vee \neg t_B^2 \vee \neg t_C^1$
 - B start at time 3:
 - $\neg t_{P}^{3} \vee \neg t_{A}^{3}$
 - $\neg t_B^3 \lor \neg t_C^2$



Assumptions

• UB=0 solutions:

$$\begin{aligned} \mathbf{A}_0 &= \{ (\mathbf{s}_{\mathsf{A}} \leq 0), (\mathbf{s}_{\mathsf{B}} \leq 0), (\mathbf{s}_{\mathsf{C}} \leq 0) \} \\ &= \left\{ \neg \mathbf{t}_{\mathsf{A}}^1, \neg \mathbf{t}_{\mathsf{B}}^1, \neg \mathbf{t}_{\mathsf{C}}^1 \right\} \end{aligned}$$









Trail (level=1):

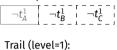


Using assumption $\neg t_A^1$

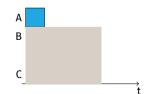


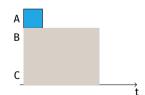


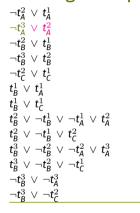
Assumptions:



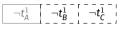


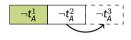


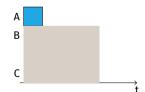




Assumptions:

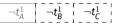




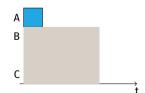


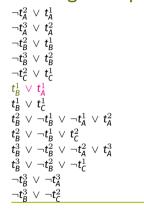


Assumptions:

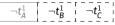


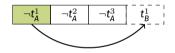
$$\neg t_A^1 \quad \neg t_A^2 \quad \neg t_A^3$$

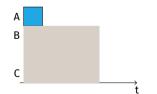


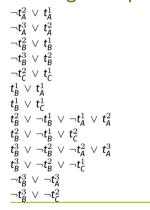






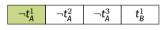


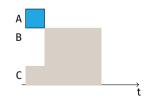


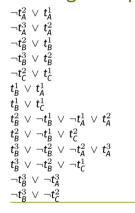


Assumptions:

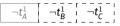




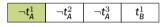




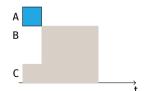
Assumptions:



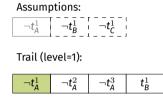
Trail (level=1):



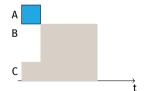
Conflict clause: $\neg t_B^1 \lor t_B^1$



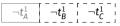




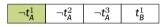
Conflict clause: $\neg t_B^1 \lor t_B^1$







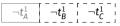
Trail (level=1):



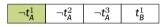
Conflict clause: $t_A^1 \vee t_B^1$







Trail (level=1):



Conflict clause: $t_A^1 \vee t_B^1$



UB=1

$$\begin{split} A_0 &= \left\{ \neg t_A^1, \neg t_B^1, \neg t_C^1 \right\}, \qquad C = t_A^1 \lor t_B^1 \\ A_1 &= (A_0 \setminus \neg C) \cup \\ \left\{ \Sigma x_i \leq k + 1 \mid \Sigma x_i \leq k \in \neg C \right\} \cup \\ \left\{ \Sigma_{\mathbf{X} \in C} \mathbf{X} \leq 1 \right\} \\ &= \left\{ \neg t_C^1 \right\} \cup \left\{ \mathbf{s}_A \leq 1, \mathbf{s}_B \leq 1 \right\} \cup \left\{ t_A^1 + t_B^1 \leq 1 \right\} \\ &= \left\{ \neg t_C^1, \neg t_A^2, \neg t_B^2, \neg v_{A1}^{B1} \right\} \end{split}$$

```
Assumptions: \begin{bmatrix} -\tau_C^1 & -\tau_A^2 & -\tau_B^2 & -\tau_B^2 \\ -t_C^1 & -t_A^2 & -\tau_B^2 & -\tau_A^2 \end{bmatrix}
Trail (level=0):
```



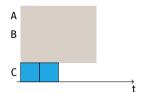
Assumptions:



Trail (level=1):



Using assumption $\neg t_{\mathcal{C}}^1$

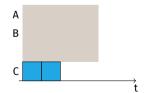


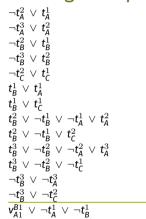


Assumptions:







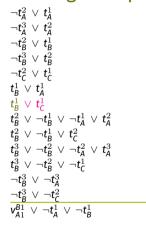


Assumptions:



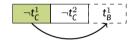
$$\neg t_{\mathcal{C}}^1 \qquad \neg t_{\mathcal{C}}^2$$

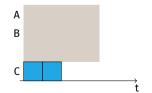




Assumptions:



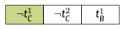


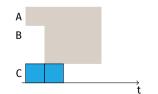


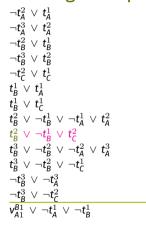


Assumptions:



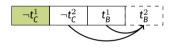


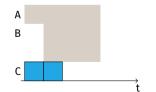




Assumptions:



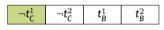






Assumptions:





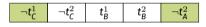




Assumptions:

$\neg t_{\mathcal{C}}^1$	$\neg t_A^2$	$\neg t_{B}^2$ $\neg v_{A1}^{B1}$

Trail (level=2):



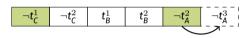
Using assumption $\neg t_A^2$





Assumptions:









Assumptions:









Assumptions:

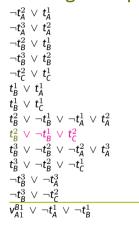
$\neg t_C^1$	$\neg t_A^2$	•	-	_ _;	$t_{B}^{\overline{2}}$	_	1	-	¬V	Б А1	-	1
			_	_	_	_		_	_	_	_	•

Trail (level=2):



Conflict clause: $\neg t_B^2 \lor t_B^2$

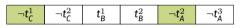




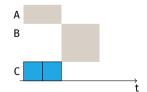
Assumptions:

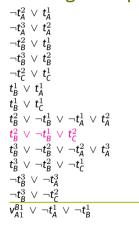


Trail (level=2):



Conflict clause: $\neg t_B^2 \lor t_B^2$





Assumptions:



Trail (level=2):



Conflict clause: $t_{\mathcal{C}}^2 \vee \neg t_{\mathcal{B}}^1 \vee t_{\mathcal{B}}^2$



Assumptions:



Trail (level=2):



Conflict clause: $t_{\it C}^2 \lor \lnot t_{\it B}^1 \lor t_{\it B}^2$





Assumptions:

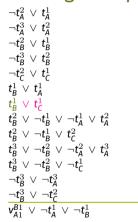


Trail (level=2):



Conflict clause: $t_{\it C}^1 \lor \neg t_{\it B}^1 \lor t_{\it B}^2$

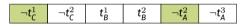




Assumptions:

$\neg t_C^1$	$\neg t_A^2$	•	-	_ _;	$t_{B}^{\overline{2}}$	_	1	-	¬V	Б А1	-	1
			_	_	_	_		_	_	_	_	•

Trail (level=2):



Conflict clause: $t_{\mathcal{C}}^1 \vee \neg t_{\mathcal{B}}^1 \vee t_{\mathcal{B}}^2$



Assumptions:

$\neg t_{\mathcal{C}}^1$	$\neg t_A^2$	$\neg t_B^2$ $\neg v_A1^B1$

Trail (level=2):



Conflict clause: $t_{\it C}^1 \lor t_{\it B}^2$



Assumptions:

$\neg t_{\mathcal{C}}^1$	$\neg t_A^2$	$\neg t_B^2$ $\neg v_A1^B1$

Trail (level=2):



Conflict clause: $t_{\it C}^1 \lor t_{\it B}^2$



UB=2

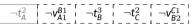
$$\begin{aligned} & A_1 = \left\{ \neg t_C^1, \neg t_A^2, \neg t_B^2, \neg v_{A1}^{B1} \right\}, \qquad C = t_C^1 \lor t_B^2 \\ & A_2 = \left\{ \neg t_A^2, \neg v_{A1}^{B1} \right\} \cup \left\{ s_C \le 1, s_B \le 2 \right\} \cup \left\{ t_C^1 + t_B^2 \le 1 \right\} \\ & = \left\{ \neg t_A^2, \neg v_{A1}^{B1}, \neg t_C^2, \neg t_B^3, \neg v_{C1}^{B2} \right\} \end{aligned}$$

```
\begin{array}{l} \neg t_{A}^{2} \ \lor \ t_{A}^{1} \\ \neg t_{A}^{3} \ \lor \ t_{A}^{2} \\ \neg t_{B}^{2} \ \lor \ t_{B}^{1} \\ \neg t_{B}^{3} \ \lor \ t_{C}^{2} \\ \neg t_{C}^{2} \ \lor \ t_{C}^{1} \end{array}
v_{B2}^{C1} \vee \neg t_B^2 \vee \neg t_C^1
```




```
\begin{array}{l} \neg t_{A}^{2} \ \lor \ t_{A}^{1} \\ \neg t_{A}^{3} \ \lor \ t_{A}^{2} \\ \neg t_{B}^{2} \ \lor \ t_{B}^{1} \\ \neg t_{B}^{3} \ \lor \ t_{C}^{2} \\ \neg t_{C}^{2} \ \lor \ t_{C}^{1} \end{array}
 v_{B2}^{C1} \vee \neg t_{B}^{2} \vee \neg t_{C}^{1}
```

Assumptions:

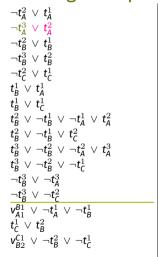


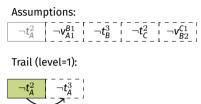
Trail (level=1):



Using assumption $\neg t_A^2$



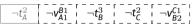






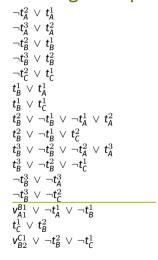
```
\begin{array}{l} \neg t_{A}^{2} \ \lor \ t_{A}^{1} \\ \neg t_{A}^{3} \ \lor \ t_{A}^{2} \\ \neg t_{B}^{2} \ \lor \ t_{B}^{1} \\ \neg t_{B}^{3} \ \lor \ t_{C}^{2} \\ \neg t_{C}^{2} \ \lor \ t_{C}^{1} \end{array}
v_{B2}^{C1} \vee \neg t_B^2 \vee \neg t_C^1
```

Assumptions:

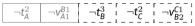


$$\neg t_A^2 \qquad \neg t_A^3$$

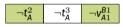




Assumptions:



Trail (level=2):



Using assumption $\neg v_{A1}^{B1}$





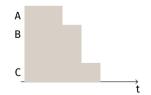
Assumptions:



Trail (level=3):



Using assumption $\neg t_B^3$





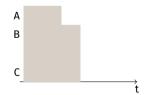
Assumptions:

$\neg t_A^2 \neg v_{A1}^{B1} \neg t_B^3 \neg t_C^2$	$\neg v_{B2}^{C1}$
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Trail (level=4):



Using assumption $\neg t_{\it C}^2$





Assumptions:

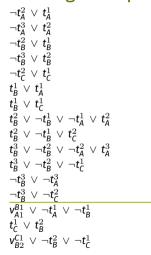
$\neg t_A^2$	$\neg v_{A1}^{B1}$	$\neg t_B^3$	$\neg t_{\mathcal{C}}^2$	$\neg v_{B2}^{C1}$
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Trail (level=5):



Using assumption $\neg v_{B2}^{C1}$

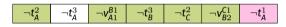




Assumptions:

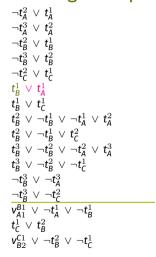
$\neg t_A^2$	$\neg v_{A1}^{B1}$	$\neg t_B^3$	$\neg t_{\mathcal{C}}^2$	$\neg v_{B2}^{C1}$
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Trail (level=6):



Guessing $\neg t_A^1$

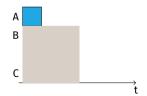


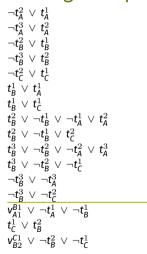


Assumptions:

$\neg t_A^2$	$\neg V_{A1}^{B1}$	$\neg t_B^3$	$\neg t_{\mathcal{C}}^2$	$\neg V_{B2}^{C1}$
--------------	--------------------	--------------	--------------------------	--------------------



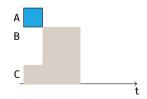


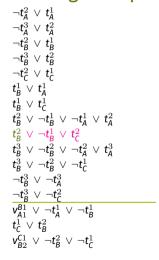


Assumptions:

$\neg t_A^2 \qquad \neg v_{A1}^{B1}$	$\neg t_B^3$	$\neg t_{\mathcal{C}}^2$	$\neg v_{B2}^{C1}$
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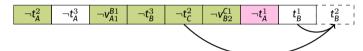
	$\neg t_{A}^2$	$\neg t_A^3$	$\neg v_{A1}^{B1}$	$\neg t_B^3$	$\neg t_{\mathcal{C}}^2$	$\neg v_{B2}^{C1}$	$\neg t^1_A$	t_B^1]
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Assumptions:

$\neg t_A^2 \qquad \neg v_{A1}^{B1}$	$\neg t_B^3$	$\neg t_{\mathcal{C}}^2$	$\neg v_{B2}^{C1}$
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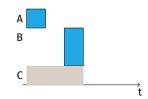


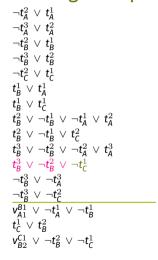




Assumptions:

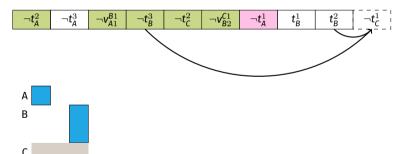
	$ eg t_{A}^2$	$\neg t_{A}^{3}$	$ eg \textit{v}_{A1}^{B1}$	$\neg t_B^3$	$\neg t_{\mathcal{C}}^2$	$\neg v_{B2}^{C1}$	$ eg t_A^1$	t_B^1	t_B^2	
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Assumptions:

$\neg t_A^2$	$\neg v_{A1}^{B1}$	$\neg t_B^3$	$\neg t_{\mathcal{C}}^2$	$\neg V_{B2}^{C1}$
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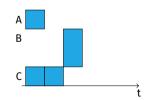




Assumptions:

$\neg t_A^2$	$\neg V_{A1}^{B1}$	$\neg t_B^3$	$\neg t_{\mathcal{C}}^2$	$\neg V_{B2}^{C1}$
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$\neg t_{A}^2$	$\neg t_A^3$	$ eg \textit{v}_{A1}^{B1}$	$\neg t_B^3$	$\neg t_{\mathcal{C}}^2$	$\neg v_{B2}^{C1}$	$\neg t^1_A$	t_B^1	t_B^2	$\neg t_{\mathcal{C}}^1$
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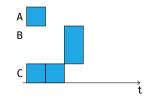


Assumptions:

Trail (level=6):

	$\neg t_A^2$	$\neg t_{A}^3$	$ eg \textit{v}_{A1}^{B1}$	$\neg t_B^3$	$\neg t_{\mathcal{C}}^2$	$\neg v_{B2}^{C1}$	$\neg t^1_A$	t_B^1	t_B^2	$ eg t_{\mathcal{C}}^1$	
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SAT!



Conclusions

- We formulated the problem using infeasibility for useful information.
- No explicitly encoded/formulated objective Σs_a (only "lazily").
- The *cores* are valid constraints/cuts, could be used in other algorithms?
- Problem formulation can be built incrementally: adding new variables and constraints preserves:
 - Cores and lower bounds
 - Learnt clauses
 - Adaptive branching/value heuristics
- It's not the whole explicit UNSAT proof that's interesting (we used only the resulting cores), but getting new perspectives on how to optimize something.