



- We extend our Cuppa3 language to Cuppa4 with the addition of a type system with four types:
 - int
 - float
 - string
 - void
- We also assume that int is a subtype of float and float is a subtype of string, that is, a compiler/interpreter is allowed to insert widening conversions and should flag errors for narrowing conversions.





We want to be able to write programs such as these:

```
int inc(int x) return x+1;
int y = inc(3);
put "the result is" + y;
```

```
float pow(float b,int p) {
  if (p == 0)
    return 1.0;
  else
    return b*pow(b,p-1);
float v:
get v;
int p;
get p;
float result = pow(v,p);
put v + " to the power of " + p +" is "+result;
```

Type system implementation: Syntax



```
program : stmt list
stmt list : stmt stmt list
          empty
stmt : VOID TYPE ID '(' opt formal args ')' stmt
       data_type ID '(' opt_formal_args ')' stmt
       data type ID opt init opt semi
       ID '=' exp opt semi
       GET ID opt semi
       PUT exp opt semi
       ID '(' opt actual args ')' opt semi
       RETURN opt exp opt semi
      WHILE '(' exp ')' stmt
      IF '(' exp ')' stmt opt else
       '{' stmt list '}'
data type : INTEGER TYPE
            FLOAT TYPE
            STRING TYPE
```

```
opt formal args : formal args
                  empty
formal args : data type ID ',' formal args
            data type ID
opt init: '=' exp
           empty
opt actual args : actual args
                  empty
actual_args : exp ',' actual_args
             exp
opt exp : exp
          empty
opt else : ELSE stmt
         empty
opt semi : ';'
          empty
```

Type system implementation: Syntax



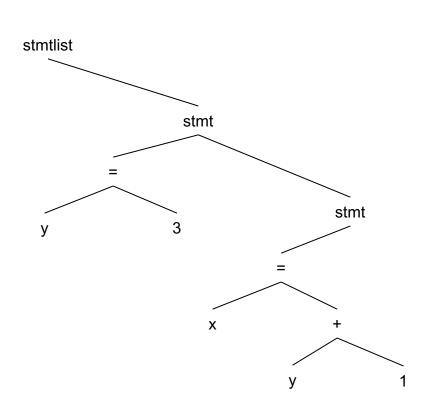
```
exp : exp PLUS exp
| exp MINUS exp
| exp TIMES exp
| exp DIVIDE exp
| exp EQ exp
| exp LE exp
| INTEGER
| FLOAT
| STRING
| ID
| ID '(' opt_actual_args ')'
| '(' exp ')'
| MINUS exp %prec UMINUS
| NOT exp
```



- At the semantic level we annotate all ASTs with type information
- We use type propagation to check that expressions/statements are properly typed.
 - Type propagation is the systematic tagging of an AST from leafs up with type information.

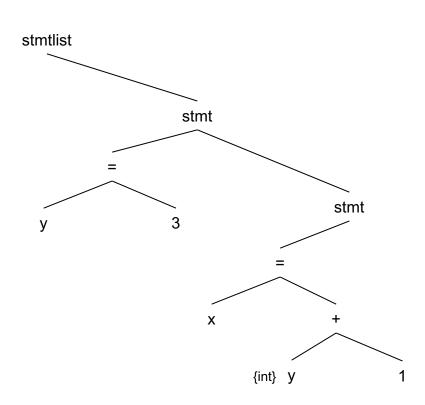


```
int y;
int x;
y = 3;
x = y + 1;
```



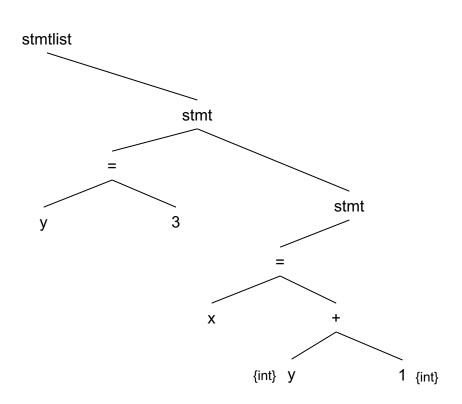


```
int y;
int x;
y = 3;
x = y + 1;
```



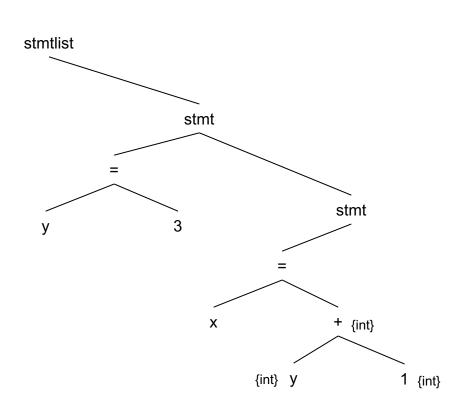


```
int y;
int x;
y = 3;
x = y + 1;
```



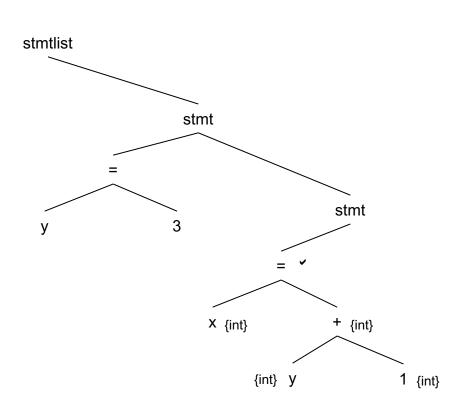


```
int y;
int x;
y = 3;
x = y + 1;
```



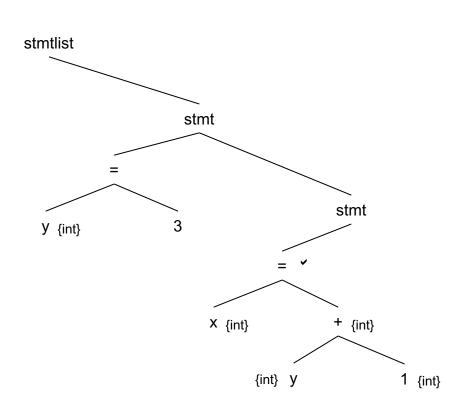


```
int y;
int x;
y = 3;
x = y + 1;
```



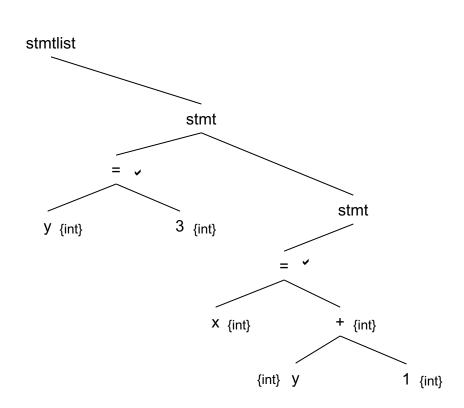


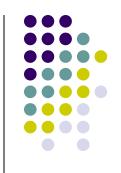
```
int y;
int x;
y = 3;
x = y + 1;
```



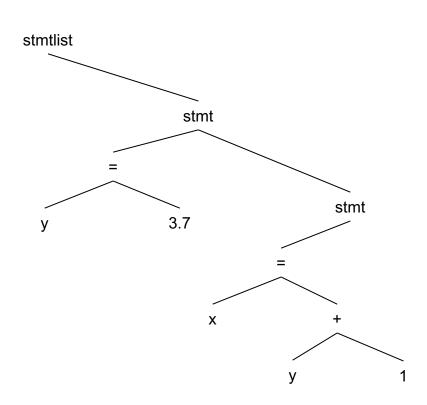


```
int y;
int x;
y = 3;
x = y + 1;
```



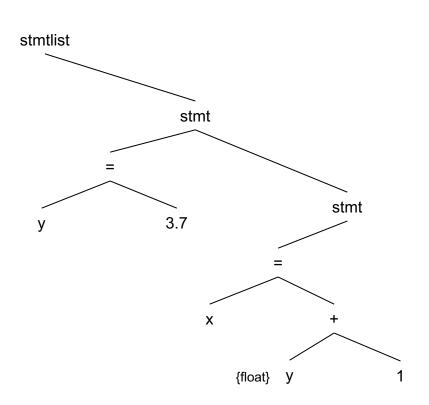


```
float y;
int x;
y = 3.7;
x = y + 1;
```



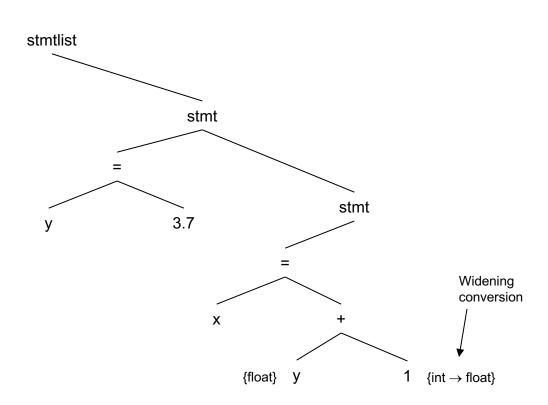


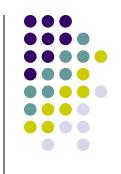
```
float y;
int x;
y = 3.7;
x = y + 1;
```



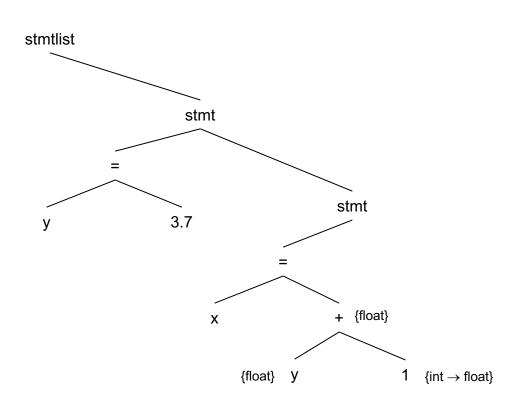


```
float y;
int x;
y = 3.7;
x = y + 1;
```



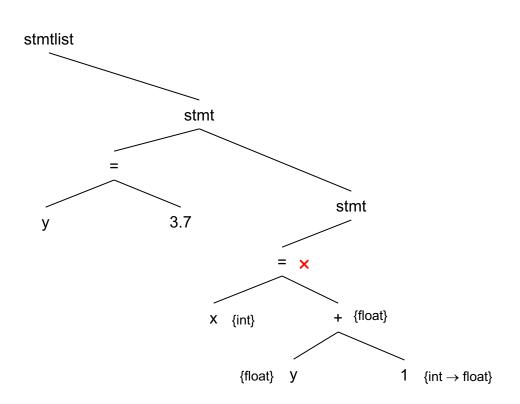


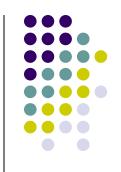
```
float y;
int x;
y = 3.7;
x = y + 1;
```



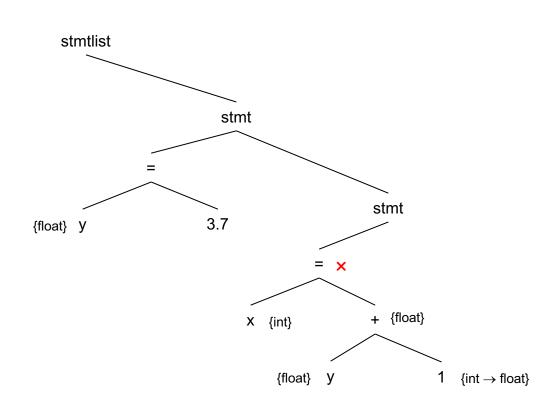


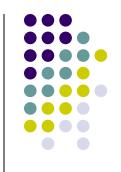
```
float y;
int x;
y = 3.7;
x = y + 1;
```



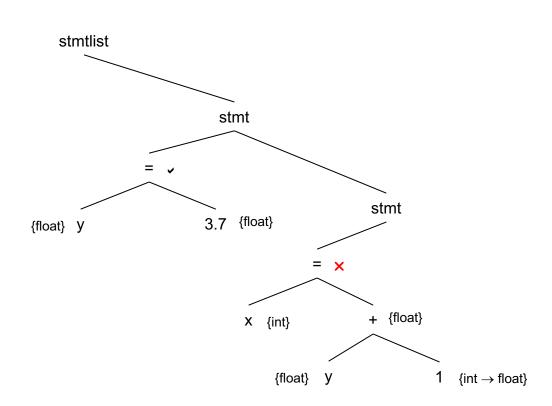


```
float y;
int x;
y = 3.7;
x = y + 1;
```



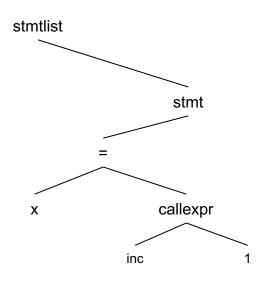


```
float y;
int x;
y = 3.7;
x = y + 1;
```





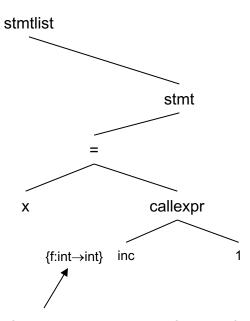
```
int inc(int i) return i+1;
int x;
x = inc(1);
```





Here is an example with a function call:

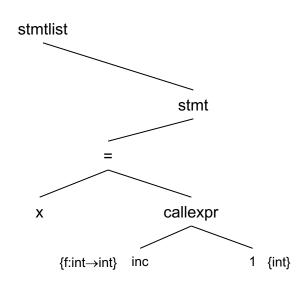
```
int inc(int i) return i+1;
int x;
x = inc(1);
```



We have to track function symbols, both for their formal parameter types and return types.

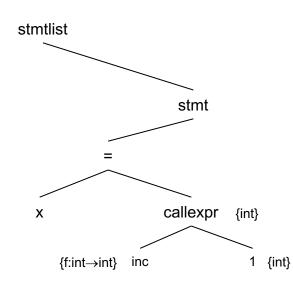


```
int inc(int i) return i+1;
int x;
x = inc(1);
```



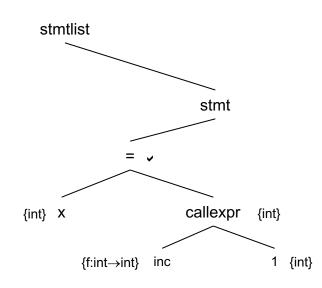


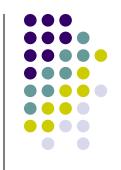
```
int inc(int i) return i+1;
int x;
x = inc(1);
```





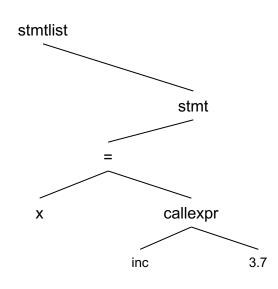
```
int inc(int i) return i+1;
int x;
x = inc(1);
```





 Here is an example with a function call and a type error:

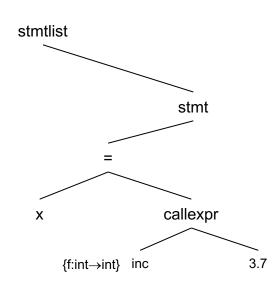
```
int inc(int i) return i+1;
int x;
x = inc(3.7);
```

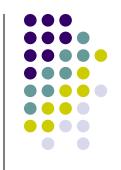




 Here is an example with a function call and a type error:

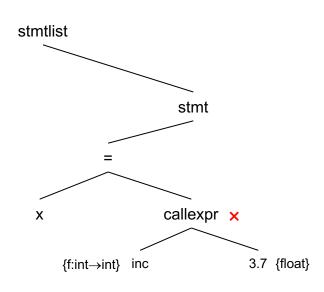
```
int inc(int i) return i+1;
int x;
x = inc(3.7);
```





 Here is an example with a function call and a type error:

```
int inc(int i) return i+1;
int x;
x = inc(3.7);
```







- Even though Cuppa4 is statically typed we will implement a dynamic type checker
 - Type propagation is done as part of the interpretation!
- Central to our implementation is the <u>type</u> <u>promotion table</u> that implements our type hierarchy.
- We use the type promotion table to implement our type propagation and type checking

Type system implementation: Type Promotion Table



```
type promotion tables for builtin primitive types. these tables
# implement the type hierarchy
              integer < float < string
              void
promote table = {
  'string': {'string': 'string', 'float': 'string', 'integer': 'string', 'void': 'void'}, 'float' : {'string': 'string', 'float': 'float', 'integer': 'float', 'void': 'void'},
  'integer': {'string': 'string', 'float': 'float', 'integer': 'integer', 'void': 'void'},
  'void' : ['string': 'void', 'float': 'void', 'integer': 'void', 'void': 'void'],
conversion table = {
 'string': {'string': str, 'float': str, 'integer': str, 'void': None},
'float': {'string': str, 'float': float, 'integer': float, 'void': None},
  'integer': {'string': str, 'float': float, 'integer': int, 'void': None},
  'void' : { 'string': None, 'float': None, 'integer': None, 'void': None},
safe assign table = {
 'string': {'string': True, 'float': True, 'integer': True, 'void': False}, 'float': {'string': False, 'float': True, 'integer': True, 'void': False},
 'integer': {'string': False, 'float': False, 'integer': True, 'void': False},
  'void' : { 'string': False, 'float': False, 'integer': False, 'void': False},
def promote(type1, type2):
    return promote table.get(type1).get(type2)
def conversion fun(ltype, rtype):
    return conversion table.get(ltype).get(rtype)
def safe assign(ltype, rtype):
         return safe assign table.get(ltype).get(rtype)
```

Type system implementation: Symbol Table



```
# symbol functions
# types of symbol values:
     ('scalar-val', type, val)
      ('function-val', return data type, arglist, body, context)
def declare scalar(self, sym, data type, init val):
    declare a symbol in the current scope.
    # first we need to check whether the symbol was already declared
   # at this scope
    if sym in self.scoped symtab[CURR SCOPE]:
        raise ValueError("symbol {} already declared".format(sym))
    # enter the symbol with its value in the current scope
    value = ('scalar-val', data type, init val)
    self.scoped symtab[CURR SCOPE].update({sym : value})
def declare fun(self, sym, return data type, arglist, body, context):
    declare a symbol in the current scope.
    # first we need to check whether the symbol was already declared
   # at this scope
    if sym in self.scoped symtab[CURR SCOPE]:
        raise ValueError("symbol {} already declared".format(sym))
    # enter the symbol with its value in the current scope
    value = ('function-val', return data type, ('lambda', arglist, body, context))
    self.scoped symtab[CURR SCOPE].update({sym : value})
```

Type system implementation: Walk



```
def plus_exp(node):
    (PLUS,c1,c2) = node
    assert_match(PLUS, '+')

    (t1, v1) = walk(c1)
    (t2, v2) = walk(c2)

    type = promote(t1, t2)

if type in ['integer', 'float']:
        return (type, v1 + v2)
    elif type == 'string':
        return ('string', str(v1) + str(v2))
    else:
        raise ValueError('unsupported type {} in + operator'.format(type))
```





```
In [2]: from cuppa4 interp import interp
          str concat = \
In [13]:
          float x = 3.0 + 1;
          put "x = " + x;
          1.1.1
          interp(str concat)
          x = 4.0
        add = \
In [4]:
         1.1.1
         float add (float a, float b)
            return a+b;
        float x = add(3.0,2);
        put x;
         1.1.1
        interp(add)
```

```
In [9]: factrec = \
        // recursive implementation of factorial
        int fact(int x)
             if (x <= 1)
                return 1;
             else
                return fact(x-1) * x;
        put fact(3);
        interp(factrec)
        6
```