

Imperative Programming – Foundations

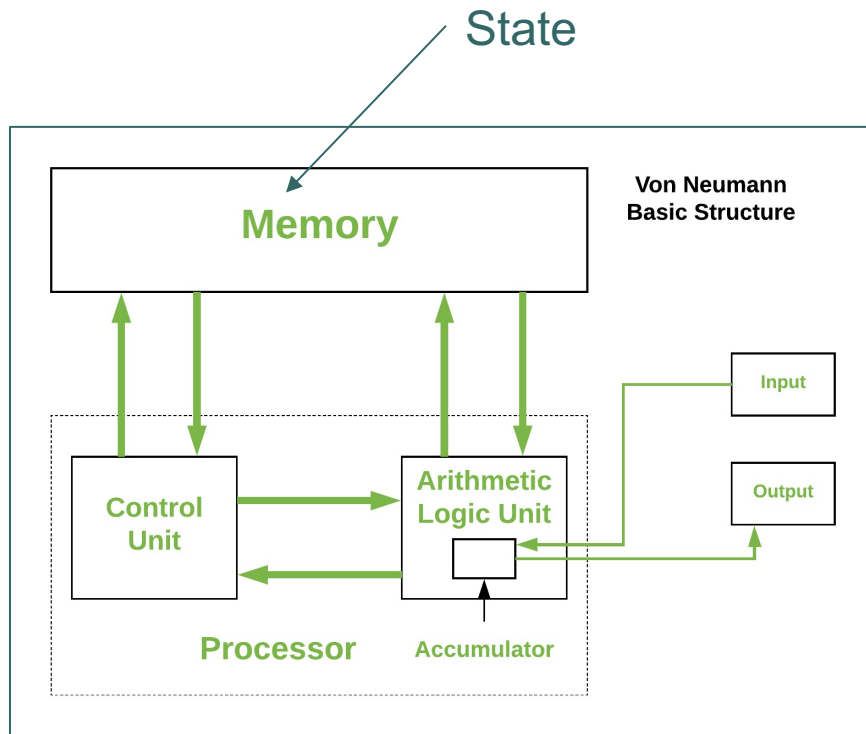
- The origins of imperative programming
- Types



The von Newman Architecture



John von Newman, Hungarian mathematician, 1903-1957.



- John von Newman's computing model gave rise to the notion of imperative programming
- Assembly/machine instructions directly manipulate processor memory
 - Imperative in the sense that each instruction states what memory will look like after it executes
- The contents of the memory defines the state of the computation at any particular point in time



The von Newman Architecture

```
section .data
    x dw 1
    y dw 2

section .bss
    z resw 1

section .text
    global _start

_start:
    mov ax, [x]      ; fetch x
    add ax, [y]      ; fetch and add y
    mov [z], ax      ; store result in z

    ; exit the program
    mov eax, 1
    xor ebx, ebx
    int 0x80
```

- Memory state is defined by three three memory **locations**
 - x,y,z
- The program changes the state by storing the sum of locations x and y into location z
- Here [<location name>] means reading/writing the value stored at that location



Imperative Programming – Foundations

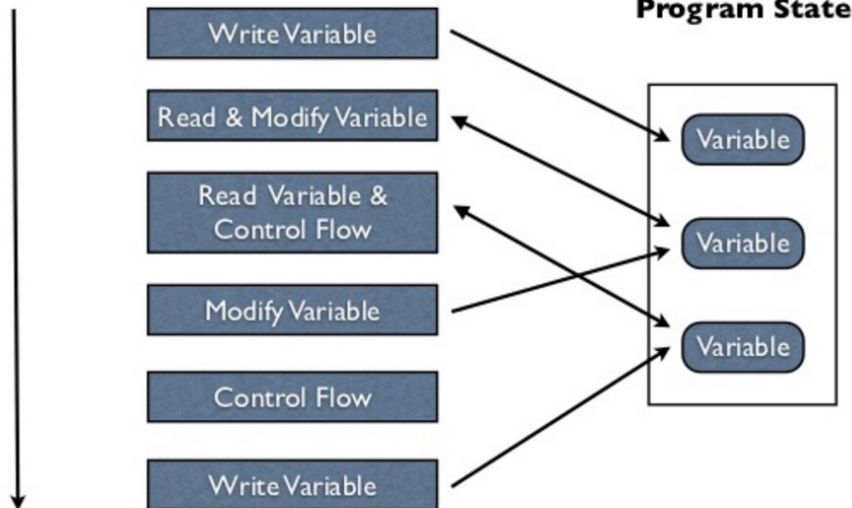
- In higher-level languages memory locations are abstracted into variables
 - This includes array/list variables
- Assembly/machine instructions are abstracted into programming language syntax
 - BUT, the assignment statement is still imperative, it tells us exactly what memory looks like after it executes.

```
let x = 1.  
let y = 2.  
let z = x + y.
```



Imperative Programming – Foundations

Program Flow



Imperative programming –

- Explicit statements that change the program state
- The program state is defined by the values assigned to the variables in a program
- The most common way to change the state in imperative programming is through an **explicit assignment of a new value to an existing variable**



Imperative Programming – Foundations

- Another example of an imperative program

```
-- sum the elements of a list
load system io.

-- initialize state
let lst = [1,2,3].
let sum = 0.

-- modify state each time around the loop by
--   (1) assigning a new value to x from the list
--   (2) incrementing sum by x
for x in lst do
|   let sum = sum + x.
end

io @println sum.
```



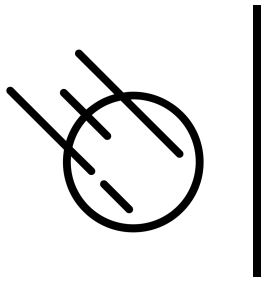
Imperative Programming – Foundations

- Let's review basic type theory for programming languages
- This is important in order to understand
 - Type hierarchies
 - Type checking
 - Type promotion



Reading

- Section 1 of the paper “Type Systems” by Luca Cardelli, Microsoft Research
 - lutzhamel.github.io/CSC493/docs/typesystems.pdf



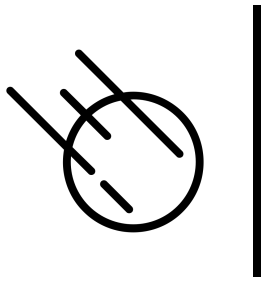
Types

A Type is a Set of Values

Consider the Rust statement:

```
let n : i32 = 3;
```

Here we constrain `n` to take on any value from the set of all 32bit integer values.



Types

Def: A type is a set of values.

Def: A primitive type is a type that is built into the language, e.g., integer, string.

Def: A constructed type is a user defined type, e.g., any type introduced by the user.
In Asteroid this is done through the 'structure' statement.

Example: Asteroid, primitive type

$q:\%real = 1.1;$	}	q is of type real, only a value that is a member of the set of all real values can be assigned to q.
$\underbrace{\hspace{1.5cm}}$		
type real \Rightarrow set of all possible real values		

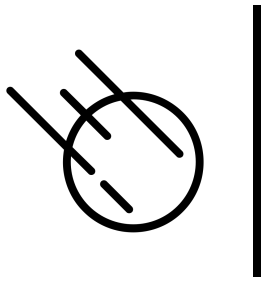


Types

Example: Rust, constructed type

```
struct Rectangle {  
    xdim: i32,  
    ydim: i32,  
}  
  
fn main() {  
    let r:Rectangle = Rectangle { xdim: 3, ydim: 4 };  
}
```

Now the variable `r` only accepts values that are members of type `Rectangle`;
☞ object instantiations of struct `Rectangle`.

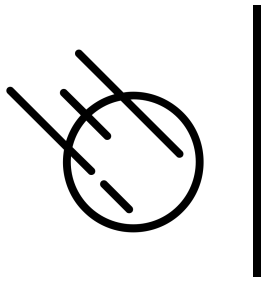


Types

Example: Asteroid, constructed type

```
structure Rectangle with  
  data xdim.  
  data ydim.  
end  
  
let r:%Rectangle = Rectangle(4,2).
```

an element of
type Rectangle.



Subtypes

Def: a subtype is a subset of the elements of a type.

Example: C

Short is a subtype of int: `short < int`

The notation $A < B$ means
A is a subtype of B.

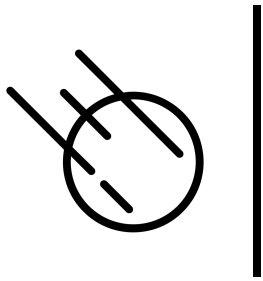
Observations:

- (1) converting a value of a subtype to a values of the super-type is called widening type conversion. (safe)
- (2) converting a value of a supertype to a value of a subtype is called narrowing type conversion. (not safe)

Example: C, partial type hierarchy

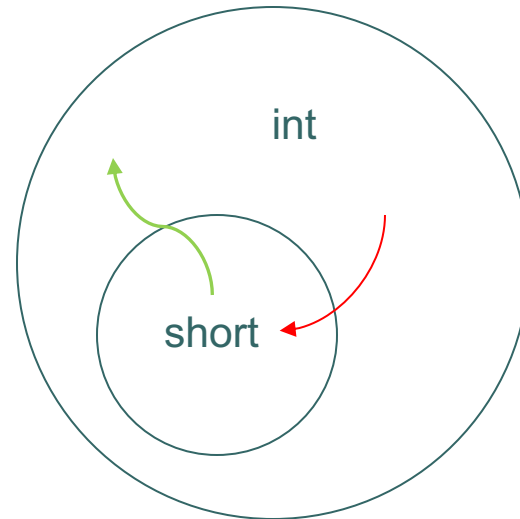
`char < short < int < float < double`

Subtypes give rise to type hierarchies and
type hierarchies allow for automatic type
coercion – widening conversions!

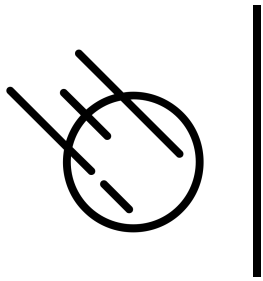


Subtypes

- A convenient way to visualize subtypes is using Venn diagrams
- Consider,
 $\text{short} < \text{int}$
- It is easy to see that the shorts are a subset of the integer values
- The green arrow represents a widening type conversion is always safe
- The red arrow represents a narrowing type conversion and is never safe



e.g. In Rust we have $i16 < i32$



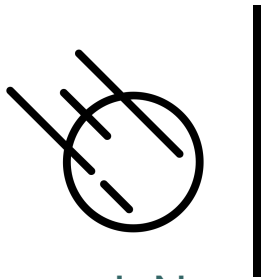
Why do we use types?

- Types allow the language system to assist the developer write better programs. Type mismatches in a program usually indicate some sort of programming error.
 - Static type checking – check the types of all statements and expressions at compile time.
 - Rust
 - Dynamic type checking – check the types at runtime.
 - Asteroid
 - Python



Type Equivalence

- Fundamental to type checking is the notion of type equivalence:
 - Figuring out whether two type description are equivalent or not
 - This is trivial for primitive types
 - But not so straight forward for constructed types like class/struct objects.



Type Equivalence

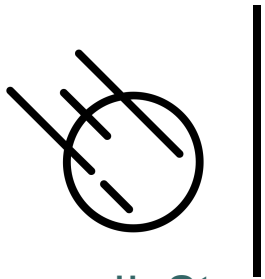
- I. Name (nominal) Equivalence – two objects are of the same type if and only if they share the same type name.

Example: Rust – constructed type

```
1 struct Type1 {x:i64, y:i64}
2 struct Type2 {x:i64, y:i64}
3
4 fn main () {
5     let x: Type1 = Type1{x:1,y:2};
6     let y: Type2 = x;
7     println!("{:?}",y);
8 }
```

Error; even though the types look the same, their names are different, therefore, Rust will not compile.

☞ Rust uses name equivalence



Type Equivalence

II. Structural Equivalence – two objects are of the same type if and only if they share the same type structure.

Example: Haskell

```
1  type Type1 = (Integer, Integer)
2  type Type2 = (Integer, Integer)
3
4  x :: Type1
5  y :: Type2
6
7  x = (1, 2)
8  y = x
```

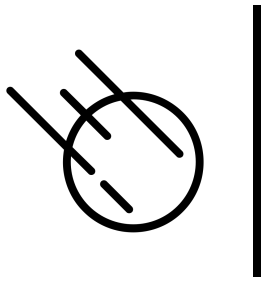
Even though the type names are different, Haskell correctly recognizes this statement.

👉 Haskell uses structural equivalence.



Type checking

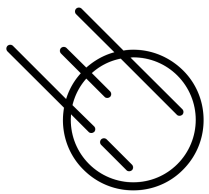
- Type checking refers to the process of making sure that all expressions and statements are properly typed.



Type Checking

- Here is the Python type checker in action
 - int and str are not part of a common type hierarchy.

```
Python 3.8.10 (default, Nov 14 2022, 12:59:47)
[GCC 9.4.0] on linux
Type "help", "copyright", "credits" or "license" for more information.
>>> "my string" + 1
Traceback (most recent call last):
  File "<stdin>", line 1, in <module>
TypeError: can only concatenate str (not "int") to str
>>> █
```

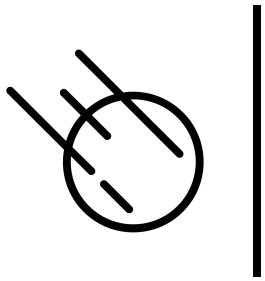


Type Checking

- Here is the type checker of the Rust compiler in action
 - i16 < i32

```
fn main () {  
    let x:i32 = 3;  
    let y:i16 = 2*x;  
    print!("{}",y);  
}
```

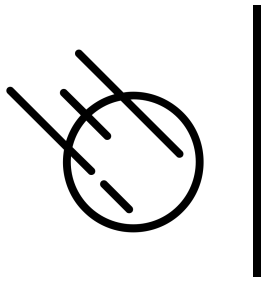
```
ubuntu$ rustc assign.rs  
error[E0308]: mismatched types  
--> assign.rs:3:16  
  
3 |     let y:i16 = 2*x;  
  |               ^^^ expected `i16`, found `i32`  
  |               |  
  |               expected due to this  
  
help: you can convert an `i32` to `i16` and panic if the converted value wouldn't fit  
  
3 |     let y:i16 = (2*x).try_into().unwrap();  
  |                  ^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^^  
  
error: aborting due to previous error  
  
For more information about this error, try `rustc --explain E0308`.  
ubuntu$
```



Type Checking in Asteroid

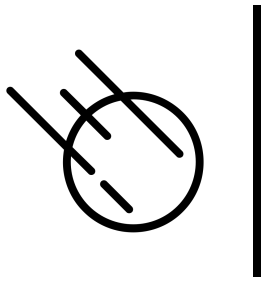
- The Asteroid type checker in action
 - Integer < real

```
Asteroid Version 1.1.4
(c) University of Rhode Island
Type "asteroid -h" for help
Press CTRL-D to exit
[ast> let x:%real = 3.1.
[ast> let y:%integer = x.
error: pattern match failed: expected type 'integer' got a term of type 'real'
ast> █
```



Type Promotion

- Convert a subtype to a supertype (automatically)
 - Widening conversion
- This usually happens at the operator level



Type Promotion - Python

- The addition operation is only defined for operands of the same type
- In order to apply the operator in a mixed-type situation one of the operands needs to be promoted
 - If promotion is not possible then flag a type error

```
Python 3.8.10 (default, Nov 14 2022, 12:59:47)
[GCC 9.4.0] on linux
Type "help", "copyright", "credits" or "license" for more information.
>>> isinstance(3.5 + 1, float)
True
>>> █
```

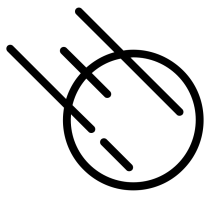
Promotion int → float



Type Promotion - Asteroid

```
Asteroid Version 1.1.4
(c) University of Rhode Island
Type "asteroid -h" for help
Press CTRL-D to exit
ast> load system type.
ast> type @gettype (3.5 + 1).
real
ast> █
```

Promotion integer → real



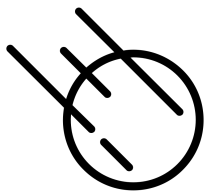
Imperative Programming – Asteroid

- Let's take a closer look at the imperative aspects of Asteroid
- We start with the type system



Primitive Types & Constants in Asteroid

- Constants are available for all the primitive data types,
 - integer, e.g. 1024
 - real, e.g. 1.75
 - string, e.g. "Hello, World!"
 - boolean, e.g. true



Type Hierarchies

- Asteroid arranges primitive data types in a type hierarchy,
 - $\text{boolean} < \text{integer} < \text{real} < \text{string}$
- As we have seen, type hierarchies facilitate automatic type promotion

```
let x:%string = "value: " + 1.
```

In002/let2.ast

Type promotion: plus as string concatenate op



Structured Data Types

- Asteroid also supports the built-in data types:
 - list
 - tuple
- These are structured data types in that they can contain entities that belong to other data types.
- Lists are mutable objects whereas tuples are immutable.
- Some examples,

Note: $(1,) \neq (1)$

```
let l = [1,2,3]. -- this is a list
let t = (1,2,3). -- this is a tuple
let one_tuple = (1,). -- this is a 1-tuple
```

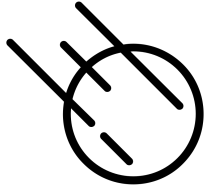
In002/let1.ast



Structured Data Types

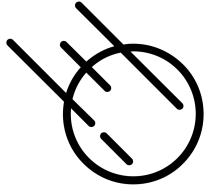
- Lists and tuples themselves are also embedded in type hierarchies, although very simple ones:
 - `list < string`
 - `tuple < string`
- That is, any list or tuple can be viewed as a string. This is very convenient for printing lists and tuples,

```
Asteroid Version 1.1.4
(c) University of Rhode Island
Type "asteroid -h" for help
Press CTRL-D to exit
ast> load system io.
ast> io @println ("this is my list: " + [1,2,3]).
this is my list: [1,2,3]
ast> █
```



The None Type

- Asteroid supports the `none` type.
- The `none` type has only one member
 - A constant named `none`.
 - The empty set of parentheses `()` can be used as a shorthand for the `none` constant.
 - That is: `none = ()`

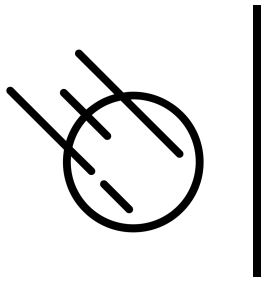


Other Data Types

- In Asteroid we also have additional data types:
 - function
 - pattern
 - user defined data types via structures

```
load system type.  
  
-- define a function  
function inc with x do  
|   return x+1.  
end  
  
-- show that 'inc' is of type 'function'  
assert (type @gettype(inc) == "function").
```

In002/ftype.ast



Reading

- The Basics

- asteroid-lang.readthedocs.io/en/latest/User%20Guide.html#the-basics