1-	Separate Ochaining:
	Separate Chaining: 100 days 8
	0 (11) (2 3 4 5 6 7 8 900 10
	Linear (Probing 10 < 3 Tobo with 10 12) 100 0 minus
D V	3 0 [98] (98 9 42 70
the cont	0 1 2 3 4 5 76 7 8 99 940 mm
	Quadratic Probing : , and who a managed book wood
	w on 42 and One 12 3 and 20 9 20 page 98 tong
	00 de 10 2 3 4 5 6 5 70 8 H 9 0 16 (0) 2 co
	I will choose 500, Because I don't know
	the total entries in the hash table so maximum
4	1913'22 to worthen bestook choice. The hours of subsen
	(1) 53491/106963 = 0.5 entries per bucket 30000 01
	(2) Yes, hashtable should rehash when the total
	humber of entires is 75% or less of your total size.
	(3) NO, it will be better to rehash the when loadfactor
	=0.78.
4.	Insert (x) Olly Olly
	(Rehesh (@) OCH) 41 0) (1)
	Remove (X) OCI) SI 41 01 01
	Contains (x) O (1) 1 + 2 1
7-	I think the problem is the Arraylist (HashItem(7>)
	old Browny = array; Berause It it will make a
4	deep capy of away and that will be expensive
	when table cells get filled over 50% when you
	rehash the table.

Separate (D Quanty push Cxl (1) 0 (1) COD () Olpopa V de de La Octoya) priority Queue Collection sientants E>C) OCh) 9. Suppose people charge for machine services. Each usor pay the same fee every time for using the machine, but different user need different service time. In order to obtain the maximum profit, suppose that as long as there are user, the machine will be ide, the users maiting to use one machine can be organized in to a minimum priority queue, and the priority is the service time required by the user. When a new user needs to use the machine - Regnest will be added to the priority quene 10 parent - 12 and may 20 mono 7.0 = 208801 (18482 (1) Heft schild szor lander blook eldelson set co a right childizer the NES & some to morning 2) NO. if will be better to rehish the Opt of 10 12 12 (0 (4 M) 0 (10) dulas (0 Remove (X) 0 (1) 21 4) 10 12 10 11 12 10 15 57 6 12 10 15 14 von 16 12 (0 (5) 14) 11 (0) 6

<u></u>		
	(2	1 3 5 6 12 10 15 14 11
	(3,	3 6 5 11 12 to 15 14
		5 6 10 11 12 14 15
		6 11 10 05 (5 812 14. JO) P 131 AS
	14	Bubble sort O(u2) Yes
		Insertion sort o(n) 8
		Iteap Sort O(nlogn) No
		Merge Sort O(nlogn) (1 4es) [1] [FS]
		Radix Sort OCH Yes
		Quick Sort O Caloga) Mo
	15.	Quick sort does not require any additional momory
		while sorting but, merge - sor requires extra metra
	, i	memory in order of O(n). Quick sort \$5 not a
	1	stable sorting algorithm while Merge is merge sort
		preserves the relative order of the elements with
		equal values but, quick sort does not guarantee that
		when compared with running times quick sort is a
		bit faster than merge sort and uses no extra
~		memory. Thus onfluencing languages to dose punck
8		50rt over mergesort. Is
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	The state of the s	

