

1) To measure the breaking strength of material Swedish physicist Waloddi Weibul proposed the Weibul distribution with density function $f(x) = \lambda x^{\lambda-1} \exp(-x^\lambda)$ for $x, \lambda > 0$. Suppose you have a random sample x_1, \dots, x_n from the Weibul distribution and would like to estimate the maximum likelihood of λ .

- a. Find the likelihood function $L(\lambda|x_1, \dots, x_n)$ and the log-likelihood function.
- b. Find the score functions $S(\lambda) = \frac{\partial \ln(L(\lambda|x_1, \dots, x_n))}{\partial \lambda}$ and its derivative.
- c. Suppose we observe 5 values (0.1, 0.25, 0.5, 1, 2). Plot the log-likelihood function.
- d. Write an R function that uses Newton's method to find the MLE of λ .

2) Insert keys (45, 24, 13, 0, 7, 23, 30, 43, 32) into a binary search tree where each node is a list of 4 fields: (left pointer, key, rank, right pointer). The rank of a node is one plus the number of nodes in its left subtree.

- a. Show the tree with nodes composed of value and rank field.
- b. Show the inorder traversal of the tree.
- c. Show the sequence of calls to find the median, minimum and maximum.
- d. Let $S[x]$ be the number of nodes in a binary search tree rooted at x . Find a recursive expression for $S[x]$.

3) Suppose the postorder traversal of a binary tree gives (A, B, G, K, H, F, Y, Z, X, R, M) whereas its inorder traversal gives (A, B, F, G, H, K, M, R, X, Y, Z). Draw the binary tree and give its preorder traversal.

4) Obtain the time complexity of the following code segments in terms of N .

- a. `for (i in 1 : N){for (j in 1 : i) print(x)}`
- b. `f2 = function(N){if(N == 0)return(1);x = 0;for (i in 1 : N) x = x + f2(N - 1)}`
- c. `f3 = function(N){if(N == 1)return(0);return(1 + f3(N/2))}`

5) When a data set is too large to be stored in primary memory, standard algorithms for computing sample quantiles are not directly applicable. You will need to write a function for estimating the sample median of n observations in m memory locations using n random numbers from a standard normal distribution. To describe the method of recursive medians, suppose that m is of the form b^k where b and k are integers. The **remedian algorithm** processes the observations sequentially in groups of size b . A median is computed for each group, yielding b^{k-1} medians in the first stage. This step is repeated recursively until a single estimate is found. A median can be computed using k arrays of size b requiring $m = kb$ storage locations. Let $b = 15$ and $k = 4$ so that $n = 50,625$ and $m = 60$. Replicate your result 100 times. Obtain univariate summary statistics and compare to the true median. See the paper on the Remedian.

6) We are given n objects in a knapsack. Object i has weight w_i and the knapsack has capacity M . If object i is placed into the knapsack, then a profit of p_i is earned. The objective is fill the knapsack to maximize profit. Hence we want to maximize $\sum_{i=1}^n p_i x_i$ subject to $\sum_{i=1}^n w_i x_i \leq M$ and $x_i = 0, 1$ for $i = 1, \dots, n$.

- a. Consider the following instance of the knapsack problem: $n = 3$, $M = 20$, $P = (p_1, p_2, p_3) = (25, 24, 15)$, and $W = (w_1, w_2, w_3) = (18, 15, 10)$. By hand calculations find the feasible solutions and obtain the optimal selection when $X = (x_1, x_2, x_3)$ is i) $(1, 0, 1)$, ii) $(1, 0, 0)$, iii) $(0, 0, 1)$, and iv) $(1, 1, 1)$.
- b. We would like to devise a randomized algorithm to find a suboptimal solution. Normally, P and W are given vectors, and we generate X to maximize the objective function. In this case, suppose we generate $n = 20$ observations from $P \sim \chi^2(1)$, and $W \sim U(0, 1)$ and fix them. Write an R function that finds a suboptimal solution by generating 1000 instances of the problem generating vector X with independent components $Bernoulli(0.5)$. Let $M = 2$.