MLFQ scheduler design

Policy

This policy employs a **round-robin** scheduling policy with **multilevel feedback** and **priority aging (usage decay)**, where:

Time unit: clock tick (10ms, based on 4.3BSD UNIX)

Number of priority queues: 128 (0, 1, 2···126,127), where 0-49 are kernel mode priorities and 50-127 are user mode priorities. (Consider to reduce it to 32 total queues later)

Quantum: 10 clock ticks based on 4.3BSD UNIX (not sure whether to have different quantums among different priority levels)

Process fields:

- 1. p_nice: base priority ranging from 0-19.
- 2. p_cpu: the accumulated and decayed CPU usage, it increments by 1 per running clock tick, every second divided by decay factor D.
- 3. p_usrpri: the priority a process belongs to, the higher p_usrpri, the lower the priority level.

The following methods of re-computation of process fields is modeled after UNIX System V: **After every running clock tick:**

```
p cpu = p cpu + 1
p usrpri := PUSER + 0.5 * p cpu + p nice
```

After every second:

```
p cpu = p cpu / D
p usrpri := PUSER + 0.5 * p cpu + p_nice
```

Here PUSER (to adjust the process to user mode priority) is 50, D (decay factor) is 2.

I'm not sure whether rule 4 and 5 of MLFQ must by followed under usage decay method, but rule 1-3 will be followed in this design.

```
Rule 1: If Priority(A) > Priority(B) run A.
```

Rule 2: If Priority(A) == Priority(B), run them in Round-Robin.

Rule 3: When a job enters the system, it is placed at the highest priority.

Rule 4: Once a job uses its time allotment at a given level, regardless of how many times it has given up CPU, it moves down a queue and its priority is reduced.

Rule 5: After some time period S, move all of the jobs in the system to the topmost queue.

Evaluation

The evaluation will focus on the aspects of

- 1. CPU utilization efficiency
- 2. Fairness (no starving processes)
- 3. Turnaround time
- 4. Response time
- 5. Throughput (I don't quite understand this metric yet)

The jobs to be tested can be classified in (assume all jobs enters in the same time)

- 1. All jobs are batch jobs (cpu bound)
- 2. All jobs are interactive jobs (IO bound)
- 3. A mix of the above jobs

Note

All underlined terms are better to be discussed before implementation

References

Decay-Usage Scheduling in Multiprocessors
https://www.researchgate.net/publication/234805812_Decay-Usage_Scheduling_in_Multiprocessors

Chapter 8 - Operating Systems: Three Easy Pieces

https://pages.cs.wisc.edu/~remzi/OSTEP/cpu-sched-mlfq.pdf