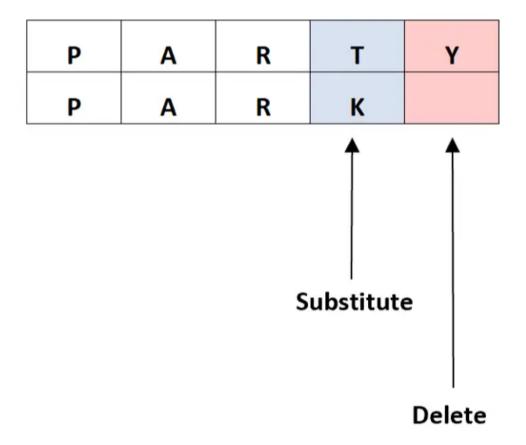


### The Levenshtein distance algorithm

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The **Levenshtein distance** (a.k.a edit distance) is a measure of similarity between two strings. It is defined as the minimum number of changes required to convert string a into string b (this is e by inserting, deleting or replacing a character in string a). The smaller the Levenshtein distance, the

# **Explanation**

# Using sub-problems

A pre-requisite for applying Dynamic Programming, to solve a problem, is to demonstrate that the solution to the original problem can be found by using the solutions to the sub-problems.

The approach here is somewhat simple and intuitive. Consider the strings a and b up to their last character:

- 1. If the last characters of both strings are the same, then the edit distance is equal to the edit distance of the same two strings, up to their second-to-last character.
- 2. If the last character is different, then the edit distance is equal to the *minimum* of the cost of inserting, deleting, or replacing the last character of string a.

### Understanding the array

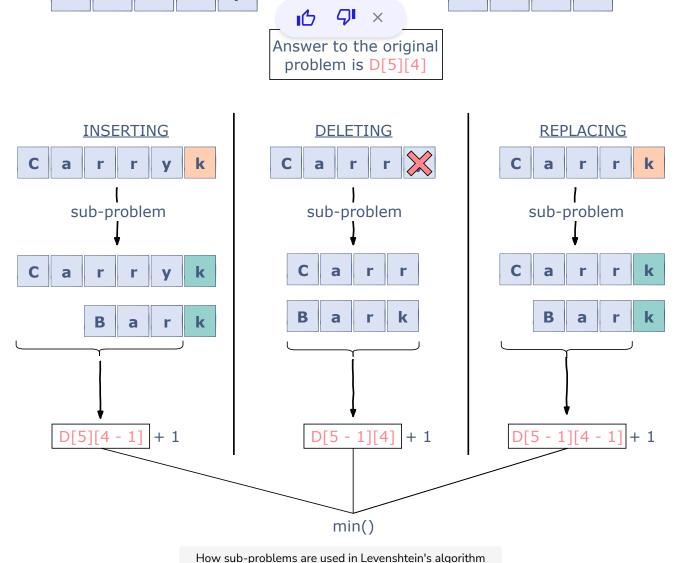
Since DP utilises memory for storing the solutions to sub-problems, we will use a 2D array (D) for this purpose.

The value in <code>D[i][j]</code> represents the edit distance *if* we consider the strings <code>a</code> and <code>b</code>, till their <code>ith</code> and <code>jth</code> character, respectively. Therefore, the answer to the original problem will be found in <code>D[length</code> of 'a'][length of 'b'].



C

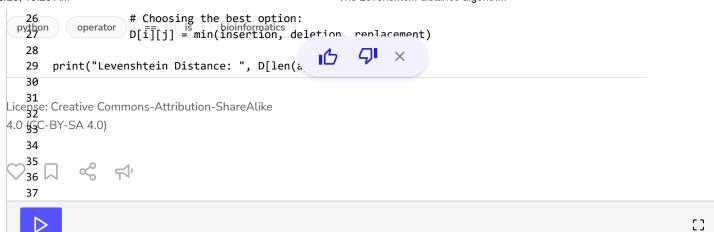
has to be changed into



#### How sub-problems are used in Levenshtein's algorithm

## Code

```
# Initialising first row:
   8
   9
       for i in range(len(a) + 1):
           D[i][0] = i
   10
   11
       # Initialising first column:
   12
       for j in range(len(b) + 1):
   13
           D[0][j] = j
   14
   15
   16
       for i in range(1, len(a) + 1):
   17
           for j in range(1, len(b) + 1):
               if a[i - 1] == b[j - 1]:
   18
                   D[i][j] = D[i - 1][j - 1]
   19
               else:
   20
   21
                   # Adding 1 to account for the cost of operation
   22
                   insertion = 1 + D[i][j - 1]
                                                                                                            6
   23
                   deletion = 1 + D[i - 1][j]
                   replacement = 1 + D[i - 1][j - 1]
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