STAR: Scaling Transactions through Asymmetric Replication

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Slides: https://tiny.cc/star_slides

Background

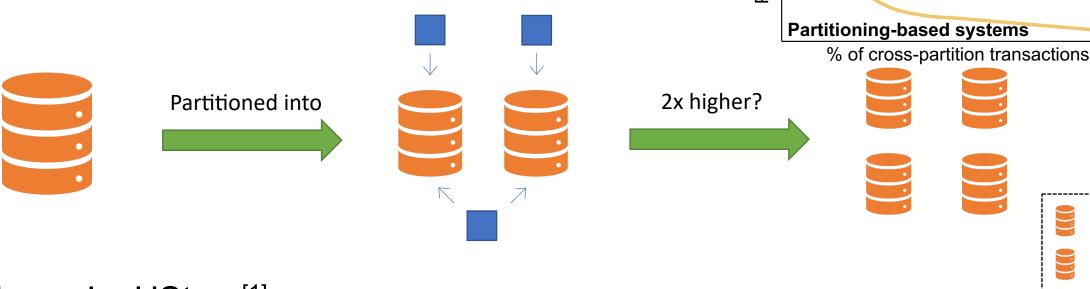
Transactions make programing easier and are used everywhere



High availability is crucial in modern OLTP applications

» Replication

Partitioning-based systems



Performance

Example: HStore^[1]

✓ Good fit for workloads with single-partition transactions

➤ Network communication and 2PC in cross-partition transactions

[1] Michael Stonebraker, Samuel Madden, Daniel J. Abadi, Stavros Harizopoulos, Nabil Hachem, Pat Helland

The End of an Architectural Era (It's Time for a Complete Rewrite). VLDB 2007: 1150-1160

Secondary replica

Performance Non-partitioned systems **Non-partitioned systems** % of cross-partition transactions 2x higher?

Example: Silo^[2]

- ✓ Good fit for workloads with cross-partition transactions
- X Cannot employ multiple nodes for parallel transaction execution

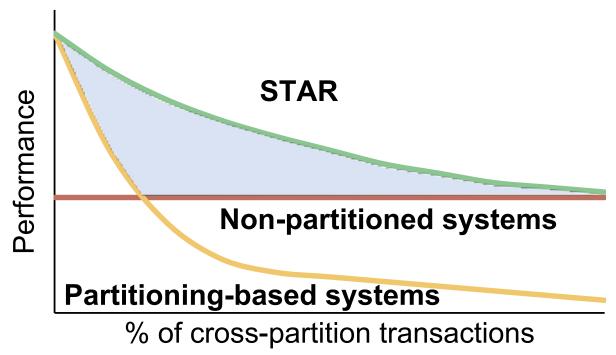
[2] Stephen Tu, Wenting Zheng, Eddie Kohler, Barbara Liskov, Samuel Madden

Speedy transactions in multicore in-memory databases. SOSP 2013: 18-32

Secondary replica

Our System: STAR

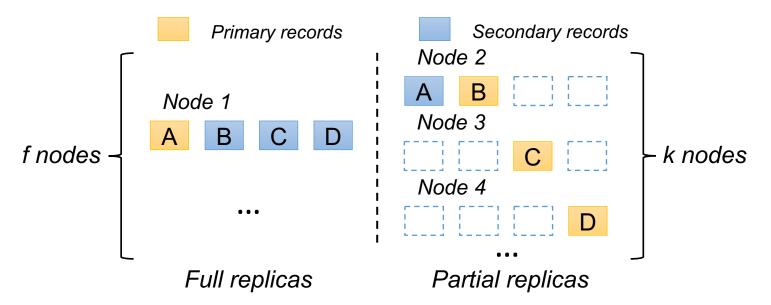
STAR uses partitioned and non-partitioned replicas to achieve the best of both worlds



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Asymmetric replication

- 1. One of these replicas is complete
- 2. One of these replicas is partitioned across several nodes



Amazon EC2 and Google Cloud now provide high memory instances with 12 TB RAM, and 24 TB instances are coming in the fall of 2019.

Partitioned phase and Single-master phase



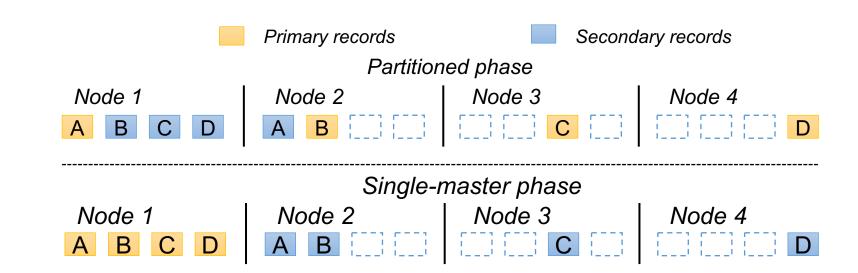
$$T_2: B = B + 1$$

$$T_3: C = C + 1$$

$$T_4: D = D + 1$$



$$T_2: B = B + C + D$$



Transactions only run over primary records.

The phase switching algorithm

Partitioned phase



Start the **partitioned phase** execution Sleep τ_p seconds

--- Replication fence ---

Start the **single-master phase** execution Sleep τ_s seconds

--- Replication fence ---

single-threaded execution per partition

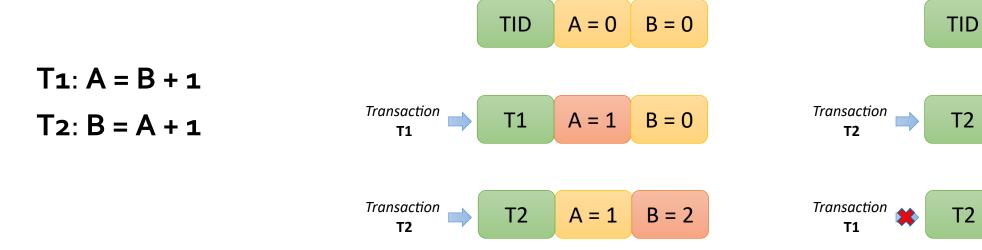
multi-threaded execution

Replication fence ensures replicas are consistent with one another before phase switching

Replication between replicas

Value replication:

✓ each write, tagged with a TID, has the value of a whole record



Secondary replica

A = 1

A = 0

A = 1

B = 0

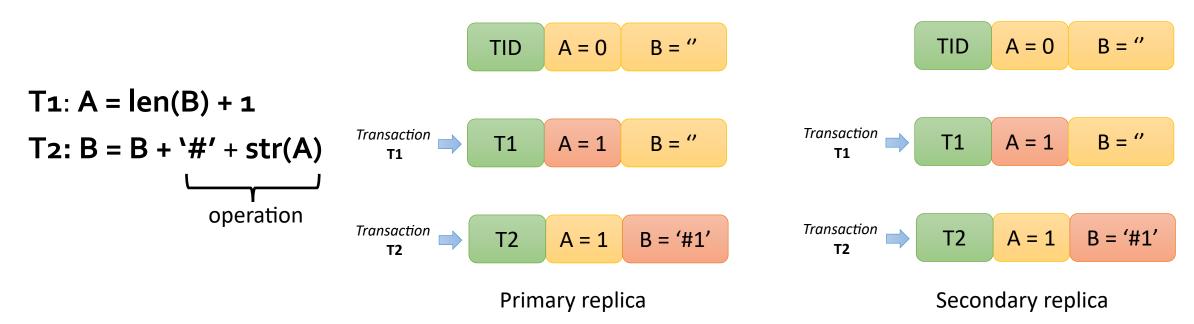
B = 2

Primary replica

Optimization: replicating operations

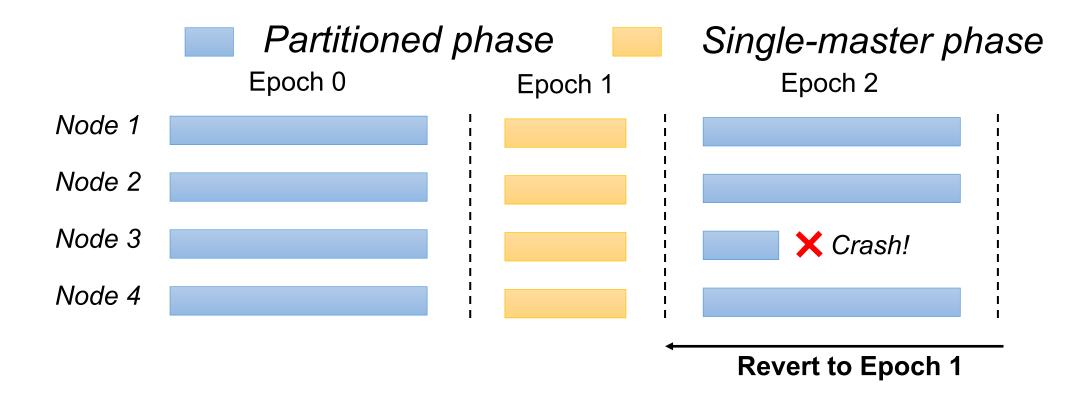
Operation replication:

✓ The replication strictly follows the commit order in the singlemaster phase



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Fault tolerance



Failure detection happens in replication fence

Experiments

A cluster of four m5.4xlarge nodes running on Amazon EC2

Benchmarks:

- » YCSB
- » TPC-C

Concurrency control algorithms:

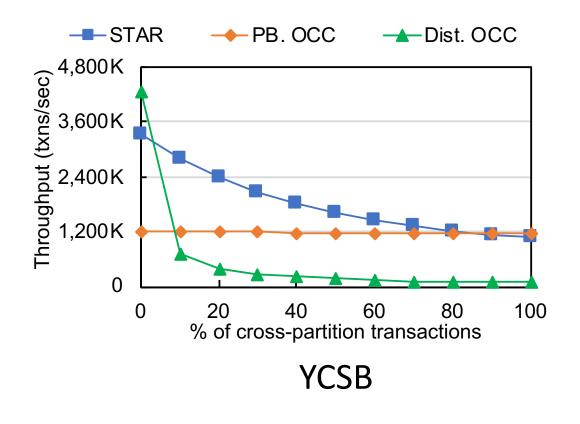
- » PB. OCC
- » Dist. OCC

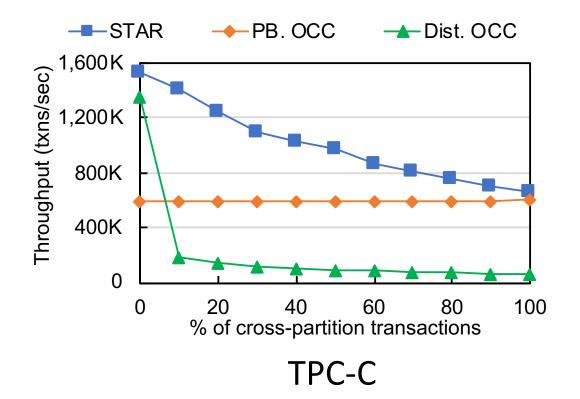
Synchronous replication
Asynchronous replication with epoch-based group commit

See our paper for comparison with Dist. S2PL and Calvin

Throughput comparison

Asynchronous replication and epoch-based group commit





Latency comparison

They all have similar latency due to epoch-based group commit

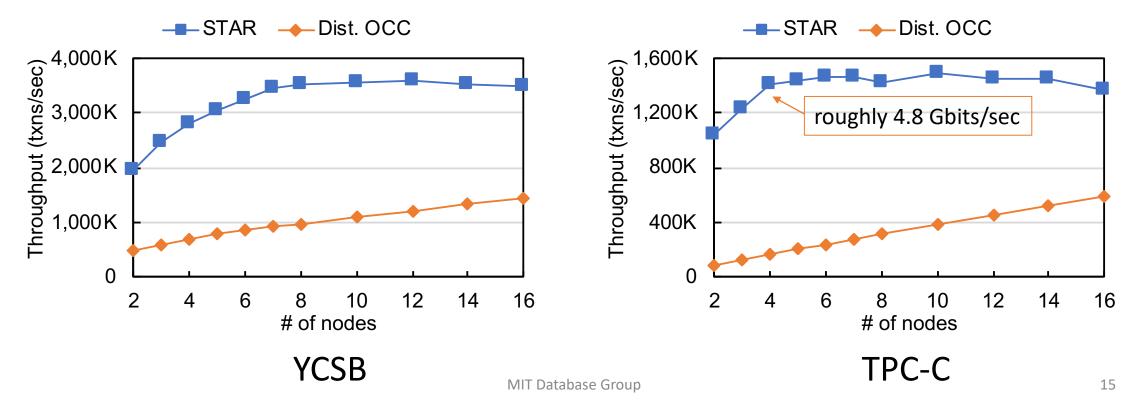
STAR	PB. OCC	Dist. OCC
6.2/9.4	5.5/11.3	6.4/11.4

50th percentile/99th percentile

Scalability experiment

STAR scales out until the network saturates.

Other systems achieve much lower throughput with the same number of nodes.

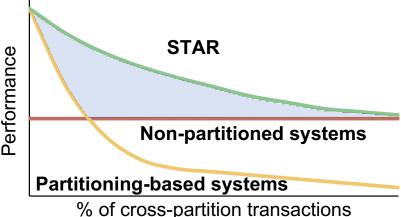


Conclusions

STAR employs a new phase-switching scheme

- » single-partition transactions are run on multiple machines in parallel
- » cross-partition transactions are run on a single machine by re-mastering records on the fly

STAR avoids cross-node communication and 2PC for distributed transactions.



Thank you

Scan the QR code to access our paper.



Paper: https://tiny.cc/star_paper

Slides: https://tiny.cc/star_slides

Code: https://tiny.cc/star_git