Exact Search-to-Decision Reductions for Time-Bounded Kolmogorov Complexity

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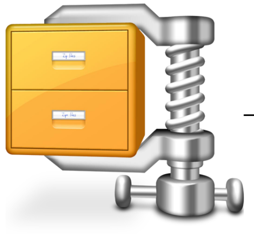
Overview

string x

01001101...0111010111



shortest encoding of x

010001110

Suppose given a string x, we can efficiently compute the length of an optimal compression of x.

Can we also efficiently find such a compression?

Kolmogorov Complexity

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Conditional Kolmogorov Complexity:

$\mathrm{K}(x \mid y) =$ "minimum length of a program $M \in \{0,1\}^*$ such that M outputs x given oracle (query) access to y"

Time-Bounded Kolmogorov Complexity

t-time-bounded Kolmogorov complexity:

$\mathrm{K}^t(x) =$ "minimum length of a program $M \in \{0,1\}^*$ such that M outputs x within time t"

Decision MINKT

Definition (MINKT):

- **Input**: $(x, 1^t, 1^s)$, where $x \in \{0,1\}^*$ and $t, s \in \mathbb{N}$
- **Task**: Decide whether $\mathrm{K}^t(x) \leq s$

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By trying $s = 1,2,\ldots,|x| + O(1)$, solving MINKT allows us to compute $\mathrm{K}^t(x)$, i.e.,
the length of a shortest t-time program that x

Computing K^t

<u>Conjecture</u>: $\boxed{\text{MINKT is NP-complete.}}$

Serach-MINKT

Definition (Search-MINKT):

- **Input**: $(x, 1^t)$, where $x \in \{0,1\}^*$ and $t \in \mathbb{N}$
- **Task**: Find a shortest t-time program that outputs x, i.e.,
  - A program M such that $|M| = \mathrm{K}^t(x)$
  - M outputs x within time t

Search-to-Decision

Question:

MINKT is easy	$\xrightarrow{?}$	Search-MINKT is easy

Search-to-Decision

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MINKT is easy

?

Search-MINKT is easy

A worst-case search-to-decision reduction

Average-Case Search-to-Decision

Question:

MINKT is easy on average $\overset{?}{\Rightarrow}$ Search-MINKT is easy on average

Average-Case Search-to-Decision

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Easy on average: For every poly-time samplable distribution D, there is an efficient algorithm that succeeds with high probability over a string $x \sim D$.

Average-Case Search-to-Decision

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Errorless	Error-Prone
• The algorithm outputs a correct answer for almost all $x \sim D$. • For the other x, the algorithm outputs $\perp$.	• The algorithm outputs a correct answer for almost all $x \sim D$. • For the other x, the algorithm can output a wrong answer.

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Prior Work

Theorem [Liu-Pass'20]:

MINKT is easy on average over the uniform distribution in the error-prone setting	→	Search-MINKT is easy on average over the uniform distribution in the error-prone setting

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Theorem [Liu-Pass'23]: Assume $\mathbf{E} \not\subseteq \mathbf{i.o.NSIZE}\left[2^{o(n)}\right]$

MINKT is easy on average over P-samplable distributions in the error-prone setting	$\Rightarrow$	Search-MINKT is easy on average over P-samplable distributions in the error-prone setting

Conditional Search-to-Decision

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Theorem [This work]: Assume $\mathbf{E} \not\subseteq \mathbf{i.o.SIZE}\left[2^{o(n)}\right]$

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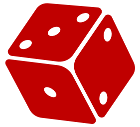
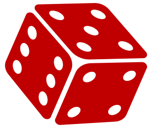
Can we get rid of the derandomization assumption?

Randomized Kolmogorov Compelxity



Randomized t-time-bounded Kolmogorov complexity:



$\mathrm{rK}^t_\lambda(x) =$ "minimum length of a randomized program $M \in \{0,1\}^*$ such that M outputs x within time t with probability $\geq \lambda$"

Decision MINrKT

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Definition (λ-MINrKT; First Attempt):

- **Input**: $(x, 1^t, 1^s)$, where $x \in \{0,1\}^*$ and $t, s \in \mathbb{N}$
- **Task**: Decide whether
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This problem is not very natraul and can only be placed in $\exists\cdot\mathbf{PP}$

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Definition (λ-MINrKT):

- **Input**: $(x, 1^t, 1^s, 1^k)$, where $x \in \{0,1\}^*$ and $t, s, k \in \mathbb{N}$
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This problem is in (promise) **MA**

Search MINrKT

Randomized t-time-bounded Kolmogorov complexity:

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Definition (λ-Search-MINrKT):

- **Input**: $(x, 1^t, 1^k)$, where $x \in \{0,1\}^*$ and $t, s, k \in \mathbb{N}$
- **Task**: Find an $(1/k)$-rK^t_λ witness of x, i.e.,
  - A randomized program M such that $|M| \leq \mathrm{rK}^t_\lambda(x)$
  - M outputs x with probability at least $\lambda - 1/k$

Average-Case Search-to-Decision for rK^t

Theorem [This work]:

λ-MINrKT is easy on average over P-samplable distributions in the erroless setting

$\Rightarrow$

λ-Search-MINrKT is easy on average over P-samplable distributions in the erroless setting

Proof Ideas

Proof Overview

Theorem [This work]: Assume $\mathbf{E} \not\subseteq \mathbf{i.o.SIZE}\left[2^{o(n)}\right]$

MINKT is easy on average over P-samplable distributions in the erroless setting	$\Rightarrow$	Search-MINKT is easy on average over P-samplable distributions in the erroless setting

Proof Overview

Theorem [This work]: Assume $\mathbf{E} \not\subseteq \mathbf{i.o.SIZE}\left[2^{o(n)}\right]$

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→

Search-MINKT is easy on average over P-samplable distributions in the erroless setting

High-Level Idea:

- Assume MINKT is easy on average.
- For a typical x from a **P**-samp distribution, there is an optimal t-time program $M \in \{0,1\}^*$ for x that admits a short encoding.
- We can then enumerate all such short encodings (and decode them) to find such an M.

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- Consider the lexicographically-first t-time program M for x.
- We know that x has short description given M.
- Here, we want that M has short description given x, so we need some kind of "symmetry of information".

Symmetry of Information

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Symmetry of information for time-unbounded Kolmogorov complexity:

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Does symmetry of information hold in the time-bounded setting, for K^t?

$$\mathrm{K}^t(x, y) \gtrsim \mathrm{K}^{\mathrm{poly}(t)}(x) + \mathrm{K}^{\mathrm{poly}(t)}(y \mid x)$$

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YES, assuming MINKT is easy on average and $\mathbf{E} \not\subseteq \mathbf{i.o.SIZE}\left[2^{o(n)}\right]$

Proof Overview

If MINKT is easy on average, then we have:	$\mathrm{K}^t(x, y) \gtrsim \mathrm{K}^{\mathrm{poly}(t)}(x) + \mathrm{K}^{\mathrm{poly}(t)}(y \mid x)$

- Fix x and t, let y_t be a shortest t-time program that outputs x.

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- Fix x and t, let y_t be a shortest t-time program that outputs x.

  - $\mathrm{K}^{\mathrm{poly}(2t)}(y_t \mid x) \lesssim \mathrm{K}^{2t}(x, y_t) - \mathrm{K}^{\mathrm{poly}(2t)}(x)$ By SoI for K^t
  - $\lesssim |y_t| - \mathrm{K}^{\mathrm{poly}(2t)}(x)$ Since given y_t, we can also recover x
  - $= \mathrm{K}^t(x) - \mathrm{K}^{\mathrm{poly}(2t)}(x)$ Since y_t is a shortest t-time program for x

Proof Overview

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Proof Overview

If MINKT is easy on average, then we have:

- Fix x and t, let y_t be a shortest t-time program that outputs x.

$$\mathrm{K}^{\mathrm{poly}(t)}(y_t \mid x) \lesssim \mathrm{K}^{t}(x) - \mathrm{K}^{\mathrm{poly}(t)}(x)$$

If $\mathrm{K}^{t}(x) - \mathrm{K}^{\mathrm{poly}(t)}(x)$ is small, then y_t admits a short and efficient encoding given x!

Proof Overview

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If $\mathrm{K}^{t}(x) - \mathrm{K}^{\mathrm{poly}(t)}(x)$ is small, then y_t admits a short and efficient encoding given x!

Claim: If MINKT is easy on average, then $\mathrm{K}^{t}(x) - \mathrm{K}^{\mathrm{poly}(t)}(x)$ is at most $O(\log t)$ for an average $x \sim D$

Conding Theorem

$\boxed{\text{If MINKT is easy on average, then we have coding theorem for } \mathrm{K}^t \text{ [Hir18]}}$

Conding Theorem

If MINKT is easy on average, then we have coding theorem for K^t [Hir18]

Coding theorem for time-unbounded Kolmogorov complexity: For every computable distribution D

$$\mathrm{K}(x) \lesssim \log\left(\frac{1}{D(x)}\right)$$

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If MINKT is easy on average, then we have coding theorem for K^t [Hir18]

Coding theorem for time-unbounded Kolmogorov complexity: For every computable distribution D

$$\boxed{\mathrm{K}(x) \lesssim \log\left(\frac{1}{D(x)}\right)}$$

if MINKT is easy on average, then for every P-samplable dist D and large enough polynomial t

$$\boxed{\mathrm{K}^t(x) \lesssim \log\left(\frac{1}{D(x)}\right)}$$

Proof Overview

Claim: If MINKT is easy on average, then $\mathrm{K}^t(x) - \mathrm{K}^{\mathrm{poly}(t)}(x)$ is small for an average $x \sim D$

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By the coding theorem for K^t, we have

$$\mathrm{K}^t(x) \lesssim \log\left(\frac{1}{D(x)}\right)$$

Fact: For every distribution D, with high probability over $x{\sim}D$,

$$\mathrm{K}^{\mathrm{poly}(t)}(x) \geq \mathrm{K}(x) \gtrsim \log\left(\frac{1}{D(x)}\right)$$

What about rK^t?

Proof Overview

If MINKT is easy on average, and $\mathbf{E} \not\subseteq \mathbf{i.o.SIZE}\left[2^{o(n)}\right]$, then we have

- symmetry of information for K^t
- coding theorem for K^t

Proof Overview

If MINKT is easy on average, and $\mathbf{E} \not\subseteq \mathbf{i.o.SIZE}\left[2^{o(n)}\right]$, then we have

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We want to say that if MINrKT is easy on average, then we have

- symmetry of information for rK^t
- coding theorem for rK^t

Proof Overview

If MINKT is easy on average, and $\mathbf{E} \not\subseteq \mathbf{i.o.SIZE}\left[2^{o(n)}\right]$, then we have

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We want to say that if MINrKT is easy on average, then we have

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Yes, but...

Proof Overview

If MINKT is easy on average, and $\mathbf{E} \not\subseteq \mathbf{i.o.SIZE}\left[2^{o(n)}\right]$, then we have

- symmetry of information for K^t $\boxed{\mathrm{K}^t(x,y) \geq \mathrm{K}^{\mathrm{poly}(t)}(x) + \mathrm{K}^{\mathrm{poly}(t)}(y \mid x) - \mathbf{log}(t)}$
- coding theorem for K^t $\boxed{\mathrm{K}^t(x) \leq \log(1/D(x)) + \mathbf{log}(t)}$

If MINrKT is easy on average, then we have

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Proof Overview

If MINrKT is easy on average, then we have

- symmetry of information for rK^t with **polylog** overhead
- coding theorem for rK^t with **polylog** overhead

This will give a quasi-polynomial-time search-to-decision reduction for rK^t

Proof Overview

If MINrKT is easy on average, then we have

- symmetry of information for rK^t with **polylog** overhead
- coding theorem for rK^t with **polylog** overhead

If MINrKT is easy on average, then we have

- symmetry of information for pK^t with **log** overhead [Goldberg-Kabanets-L.-Oliveira'22]
- coding theorem for pK^t with **log** overhead [L.-Oliveira-Zimand'22]

Proof Overview

If MINrKT is easy on average, then we have

- symmetry of information for pK^t $\boxed{\mathrm{pK}^t(x, y) \geq \mathrm{pK}^{\mathrm{poly}(t)}(x) + \mathrm{pK}^{\mathrm{poly}(t)}(y \mid x) - \mathbf{log}(t)}$
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Fix x and t, let y_t be a shortest t-time randomized program that outputs x.

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- $\lesssim |y_t| - \mathrm{pK}^{\mathrm{poly}(2t)}(x)$
- $= \mathrm{rK}^t(x) - \mathrm{pK}^{\mathrm{poly}(2t)}(x)$

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- $\mathrm{pK}^{\mathrm{poly}(t)}(y_t \mid x) \lesssim \mathrm{rK}^t(x) - \mathrm{pK}^{\mathrm{poly}(t)}(x)$

We want $\mathrm{rK}^t(x) - \mathrm{pK}^{\mathrm{poly}(t)}(x)$ to be small for an average x.

Proof Overview

If MINrKT is easy on average, then we have

- symmetry of information for pK^t $\boxed{\mathrm{pK}^t(x,y) \geq \mathrm{pK}^{\mathrm{poly}(t)}(x) + \mathrm{pK}^{\mathrm{poly}(t)}(y \mid x) - \mathbf{log}(t)}$
- coding theorem for pK^t $\boxed{\mathrm{pK}^t(x) \leq \log(1/D(x)) + \mathbf{log}(t)}$

Fix x and t, let y_t be a shortest t-time randomized program that outputs x.

- $\mathrm{pK}^{\mathrm{poly}(t)}(y_t \mid x) \lesssim \mathrm{rK}^t(x) - \mathrm{pK}^{\mathrm{poly}(t)}(x)$

We want $\mathrm{rK}^t(x) - \mathrm{pK}^{\mathrm{poly}(t)}(x)$ to be small for an average x.

But this requires coding for rK^t...

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via a magical lemma that we proved!

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This is small for an average x, by the coding theorem for pK^t

Open Problems

- Can we get worst-case search-to-decision reductions?

Theorem [This work]

MINrKT is easy on average $\Rightarrow$ An algorithm **A** that, given x, runs in $2^{O(n/\log n)}$ time and outputs an $o(1)$-rK^t witness of x, for some $\mathrm{poly}(n) \leq t \leq 2^{n^{\varepsilon}}$

Thank you!