CellSDN: Taking control of cellular core networks



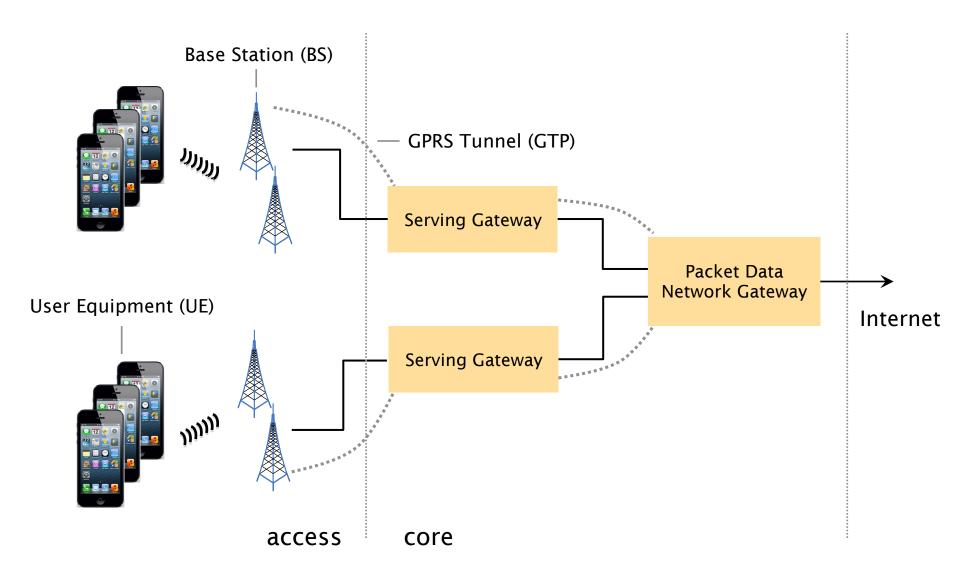
Laurent Vanbever

Princeton University

IBM Watson Research Center

Tue 9 April 2013

Joint work with Xin Jin, Li Erran Li, and Jennifer Rexford



Most functionalities are implemented at the PGW:

- Stateful firewall, DPI, lawfull intercept
- Application optimization
- Paging
- Monitoring and Billing

Packet Data Network Gateway

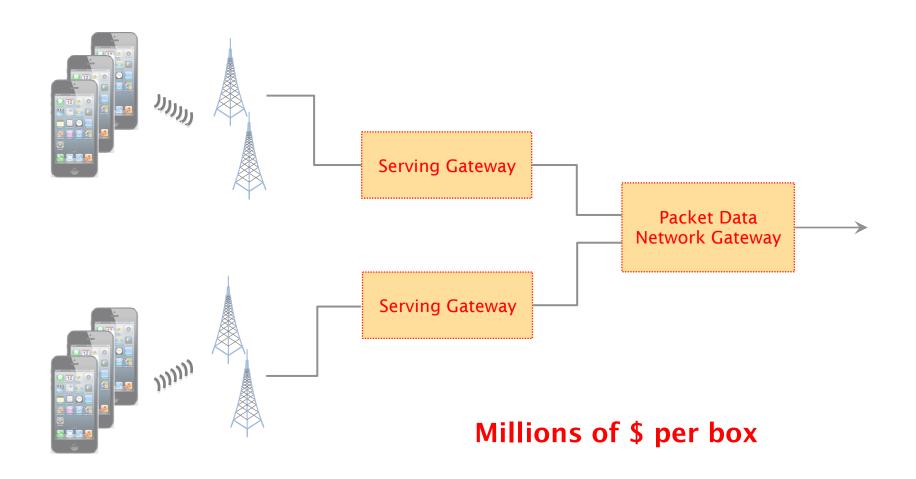
Most functionalities are implemented at the PGW:

- Stateful firewall, DPI, lawfull intercept
- Application optimization
- Paging
- Monitoring and Billing

Packet Data Network Gateway

- Hard to mix-and-match
- Slow innovation

Cellular core networks are not flexible and costly!



Key idea: separate Software from Hardware

Packet Data Network Gateway

Packet Data Network Gateway Software

Hardware

Packet Data Network Gateway Virtualized Software

Hardware

Packet Data Network Gateway Virtualized Software

Hardware

easy deployment up-/down-scale

Packet Data Network Gateway Virtualized Software

Commodity Hardware

easy deployment up-/down-scale

Packet Data Network Gateway Virtualized Software

Commodity Hardware

easy deployment up-/down-scale

cheap

Cellular policies are fine-grained, based on

- customer-attributes smartphone model, OS, billing plan, ...
- applications video, web, voice, ...
- a combination of both video from iPhone 5 gold subscribers, ...

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- customer-attributes smartphone model, OS, billing plan, ...
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 video, web, voice, ...
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can lead to *million* of different flows

Challenges

Data-plane Forwarding table size

Control-plane Overhead

Slow reactivity

Challenges

Solutions

Data-plane

Forwarding table size

Hierarchical tagging

Control-plane

Overhead

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Challenges

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Forwarding table size

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Slow reactivity

Hierarchical controller

CellSDN: Taking control of cellular core networks



Architecture

software-defined network

Scaling the data-plane multi-dimensional tagging

Scaling the control-plane tasks delegation

CellSDN: Taking control of cellular core networks



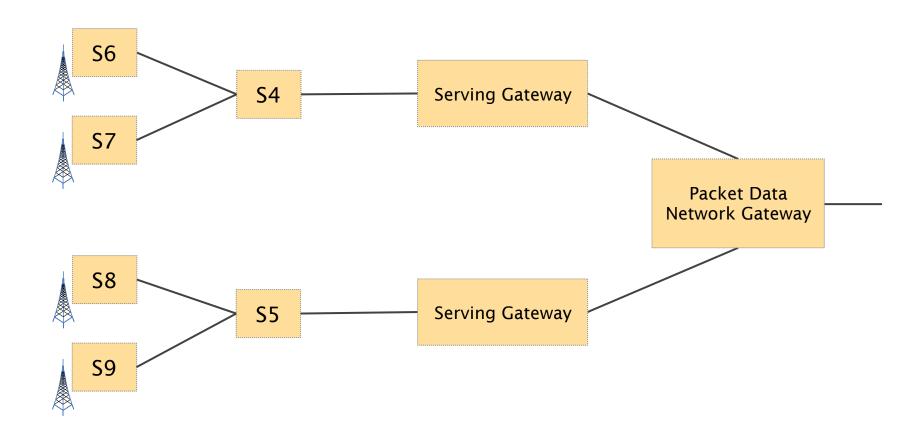
1 Architecture

software-defined network

Scaling the data-plane multi-dimensional tagging

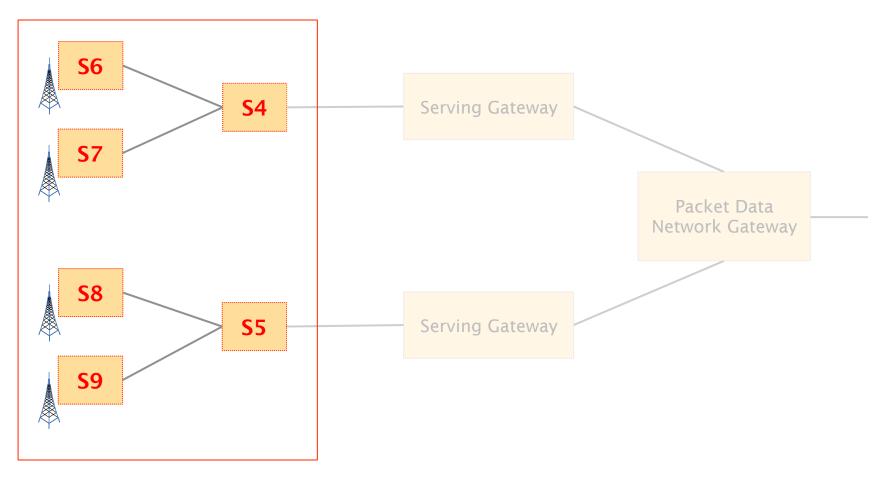
Scaling the control-plane tasks delegation

CellSDN relies on commodity, SDN-enabled switches

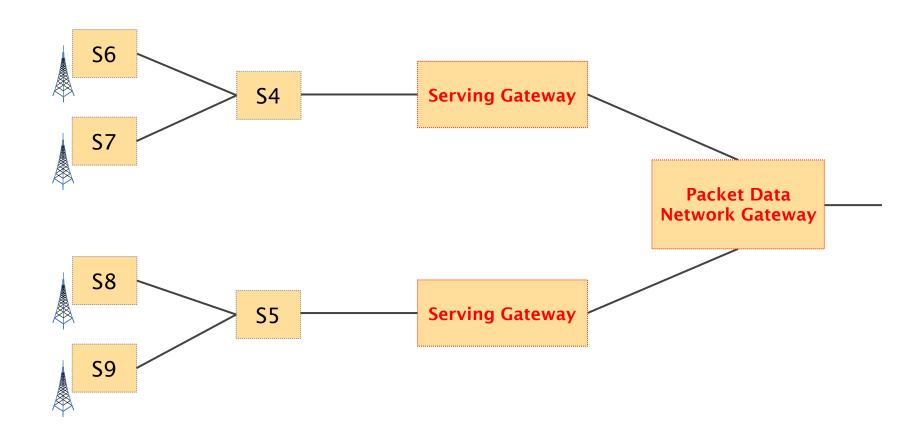


CellSDN relies on commodity, SDN-enabled switches

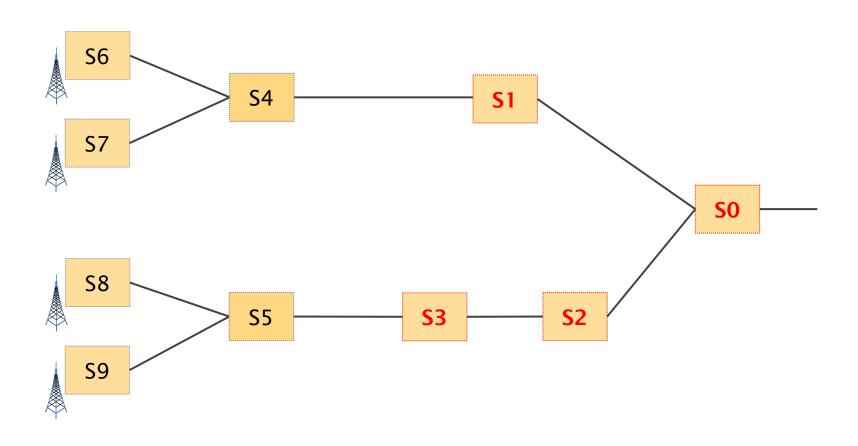
OpenFlow switches



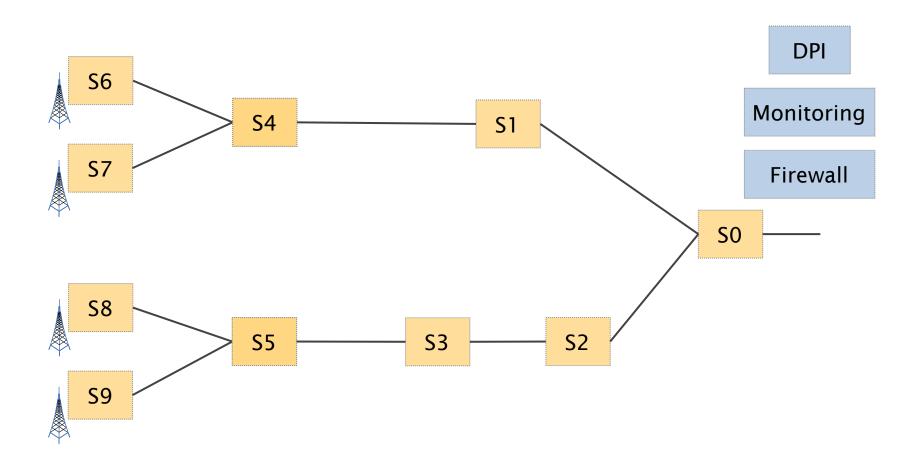
CellSDN also relies on commodity SDN switches for specialized appliances

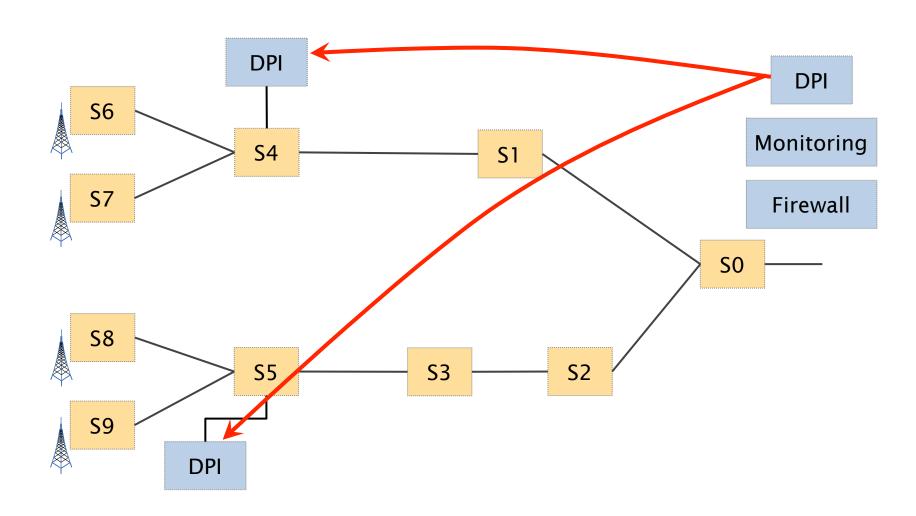


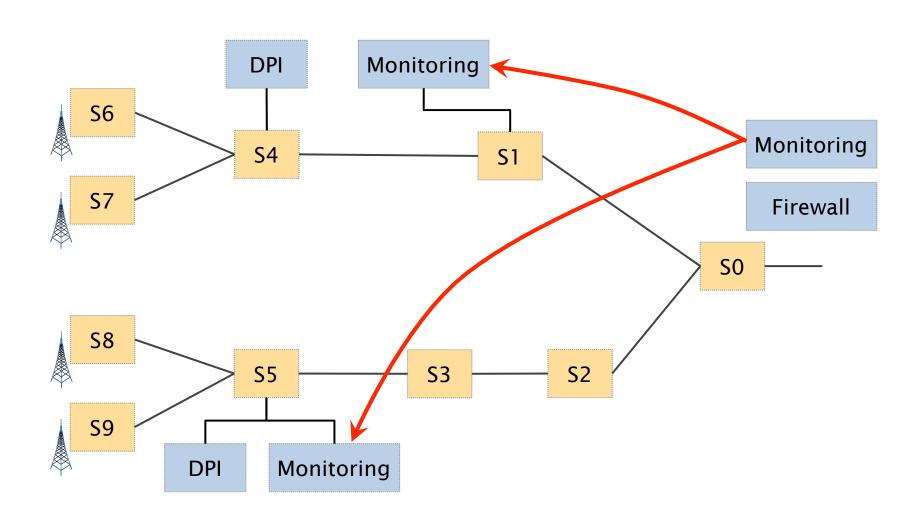
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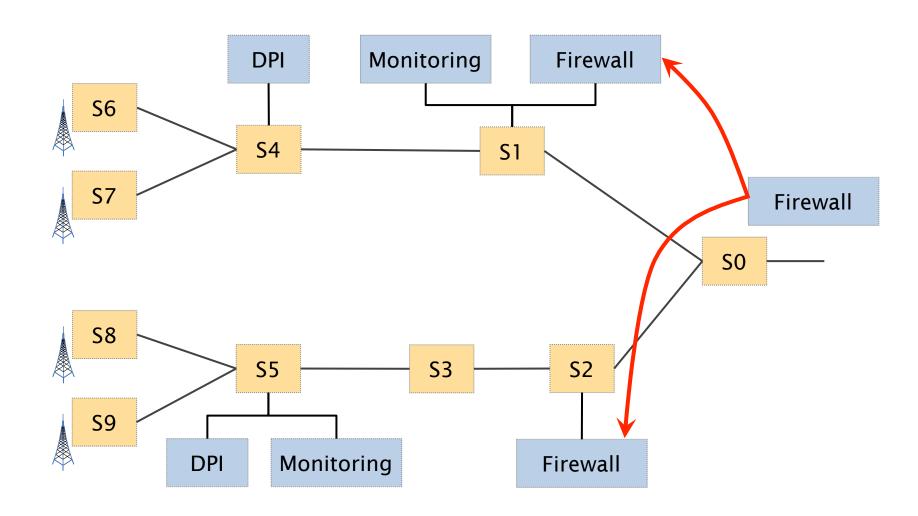


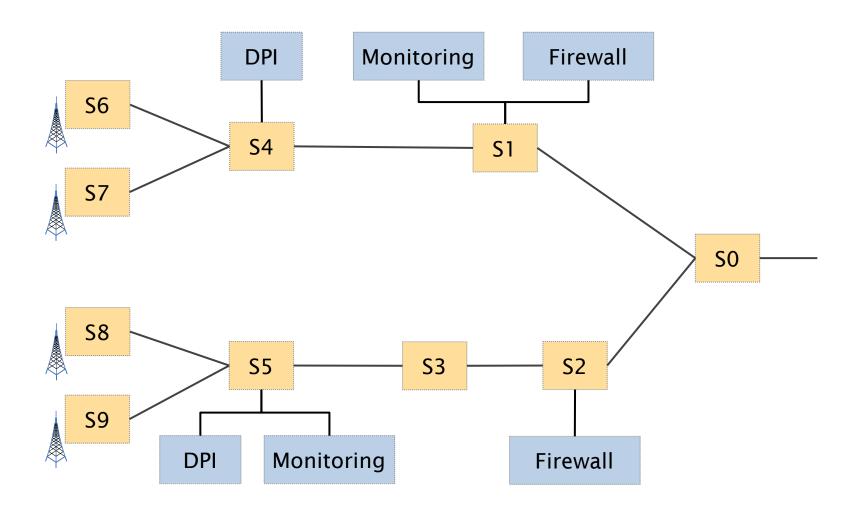
CellSDN distributes network processing











CellSDN route traffic based on high-level service policies

A service policy is composed of a

predicate A boolean expression, e.g. on subscriber attributes,

applications attributes, cell properties ...

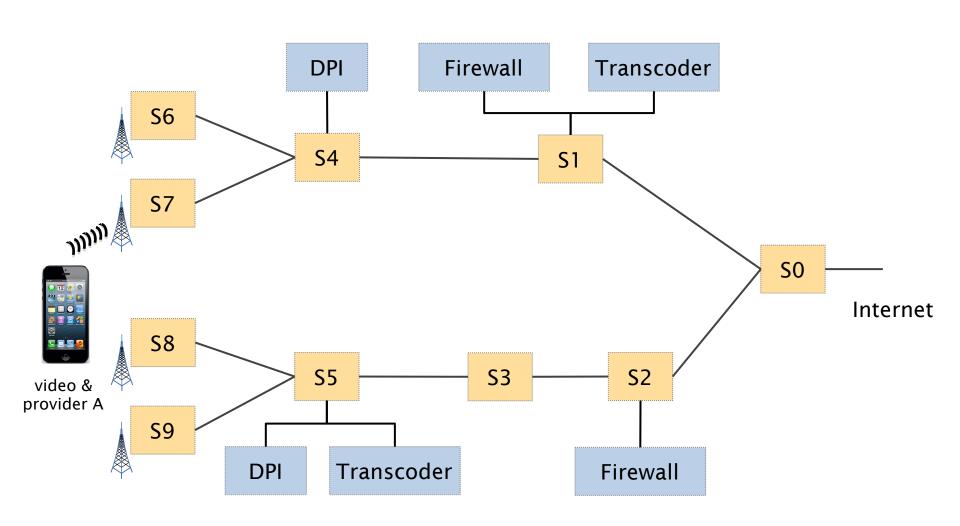
action A list of middleboxes, access-control

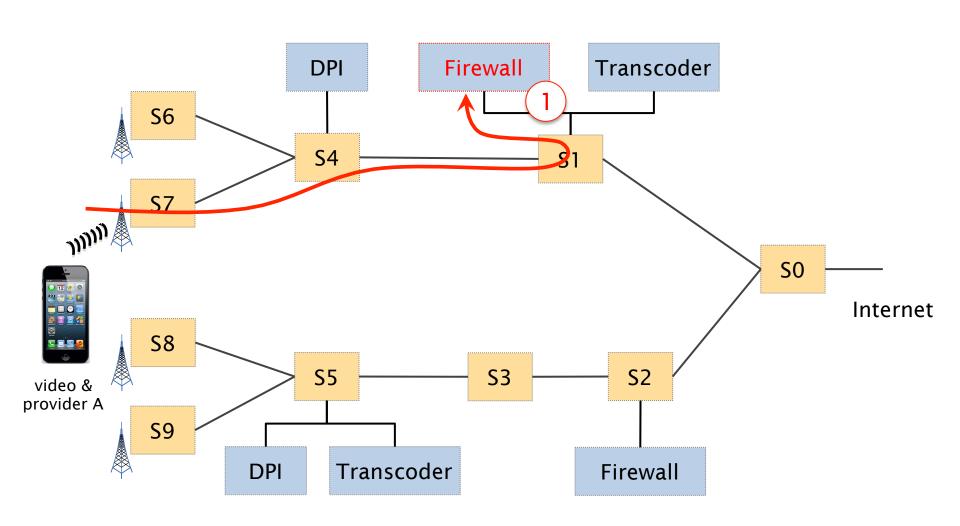
specifications (i.e., allow/deny) ...

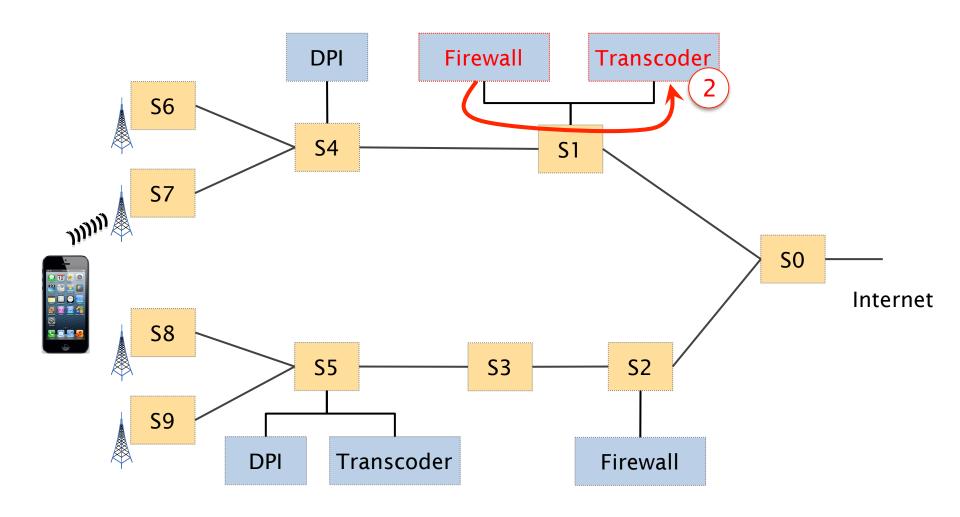
priority An integer. To disambiguate overlapping policies

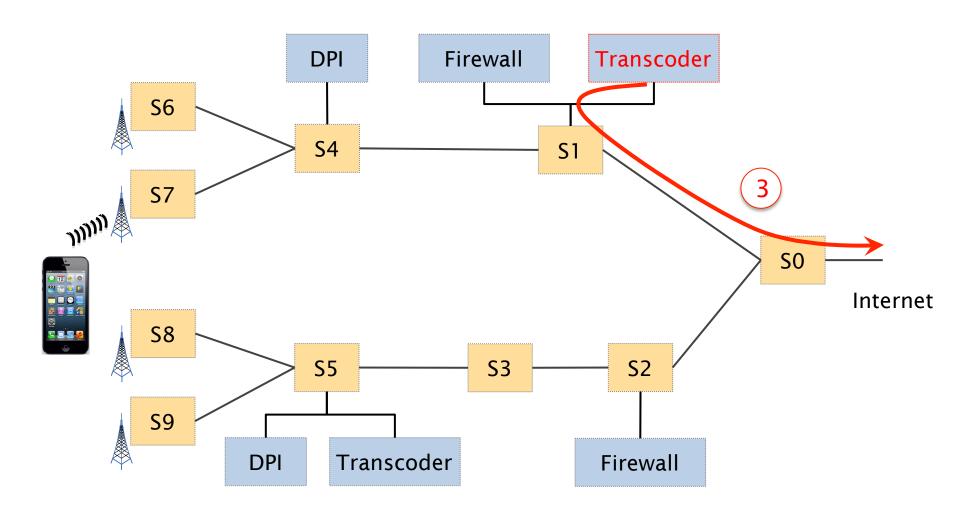
Example of service policy

priority	predicate	action
1.	provider is <i>B</i>	Firewall
2.	provider is not A	Drop
3.	traffic is video and plan is Silver	Firewall> Transcoder
4.	*	Firewall

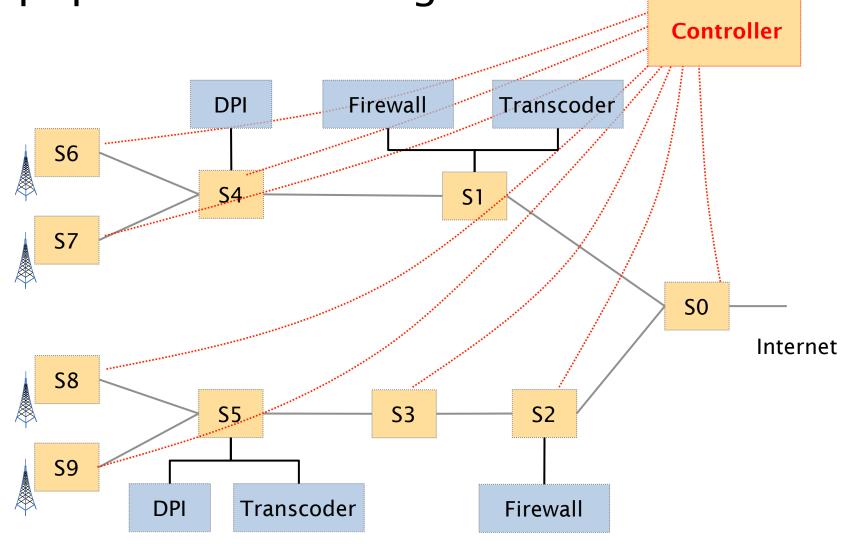








CellSDN uses a centralized controller to populate forwarding tables



CellSDN: Taking control of cellular core networks

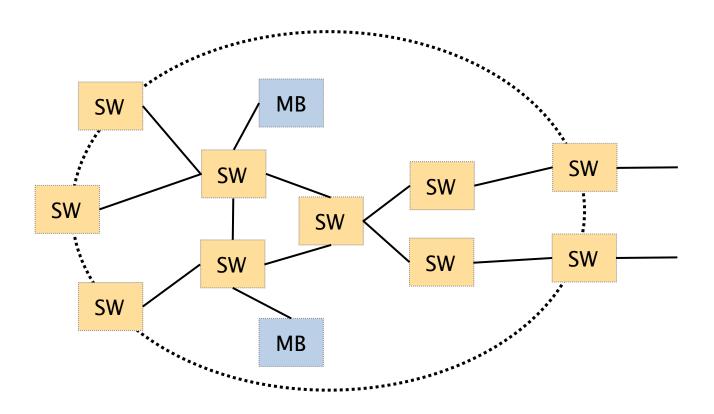


Architecture software-defined network

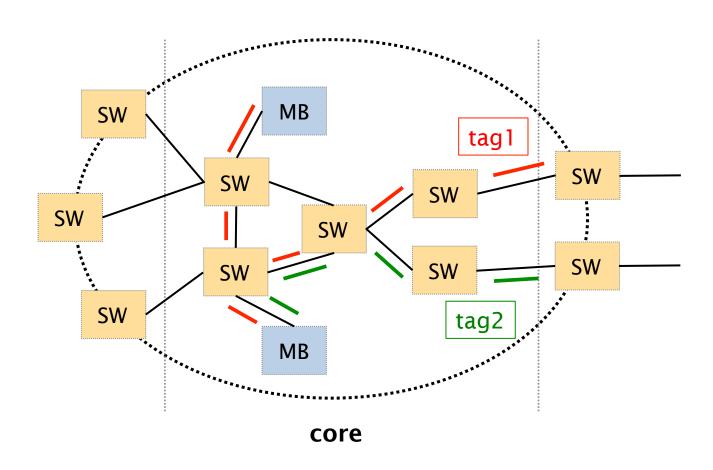
Scaling the data-plane multi-dimensional tagging

Scaling the control-plane tasks delegation

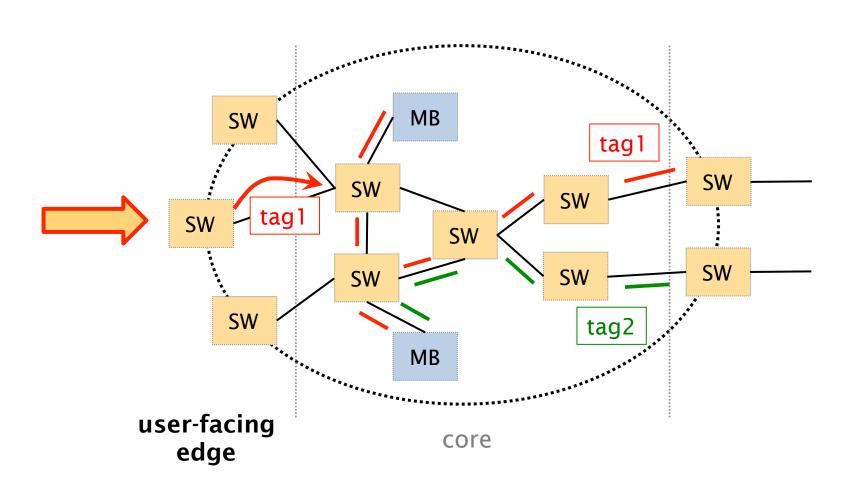
CellSDN relies on tagging to implement forwarding policies



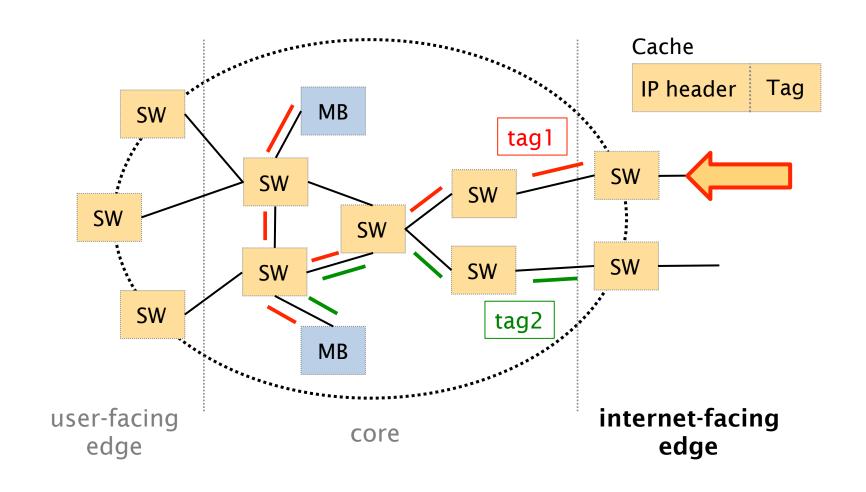
CellSDN installs stable policy path in the core



CellSDN classifies and tags traffic at the user-facing edge



CellSDN caches tags at the Internetfacing edge for the return traffic



If each path gets one tag, millions of tag are needed



If each path gets one tag, millions of tag are needed



At best, switches support a few tens of thousands entries

[Stephen, Conext12]

CellSDN relies on multi-dimensional tag

CellSDN tags are composed of three parts:

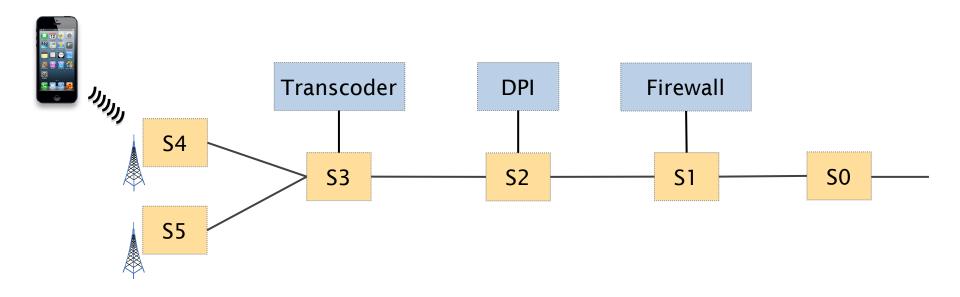
- Policy
- Location
- User Equipment Identifier

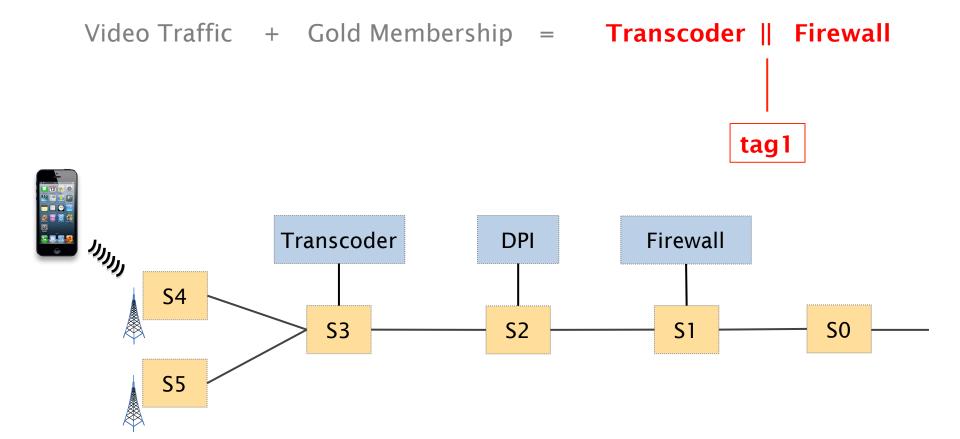
switches can selectively match on any part

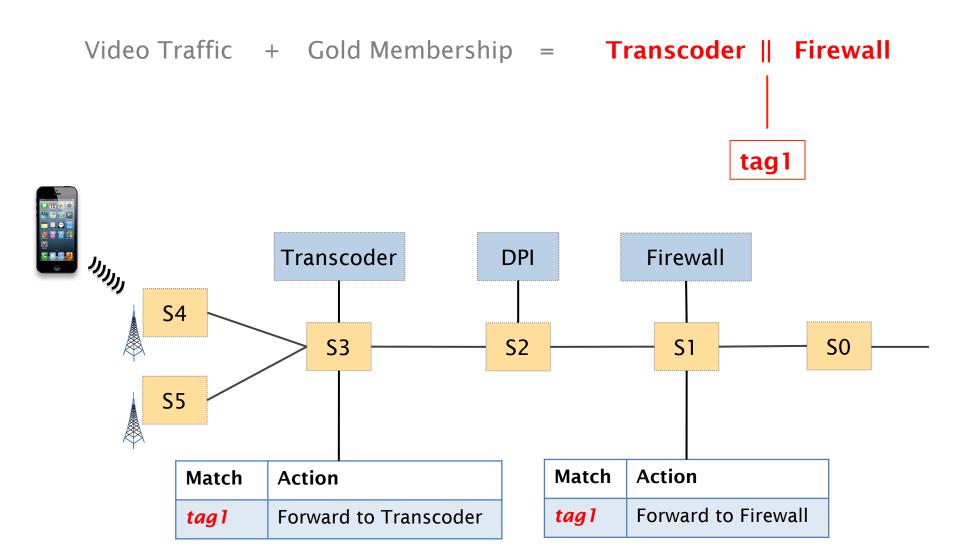
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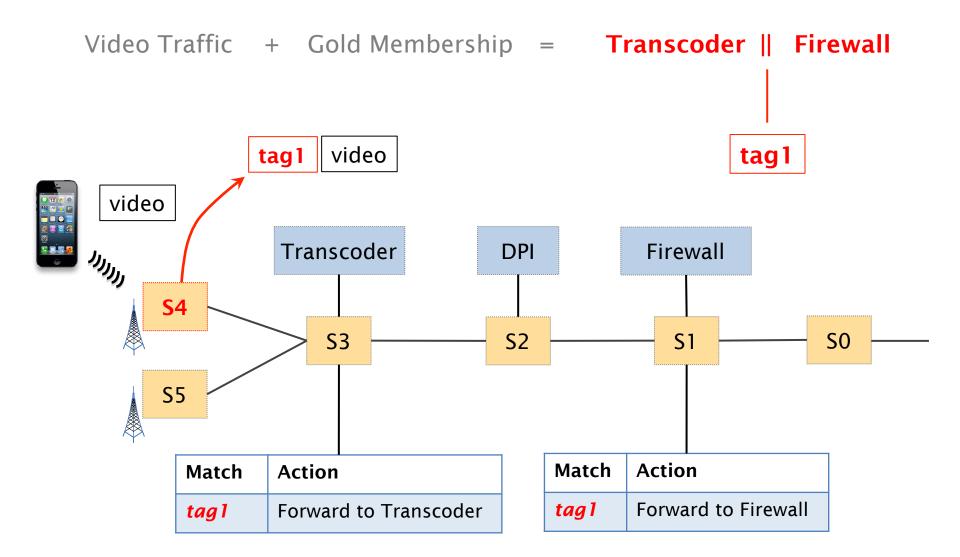
- Policy
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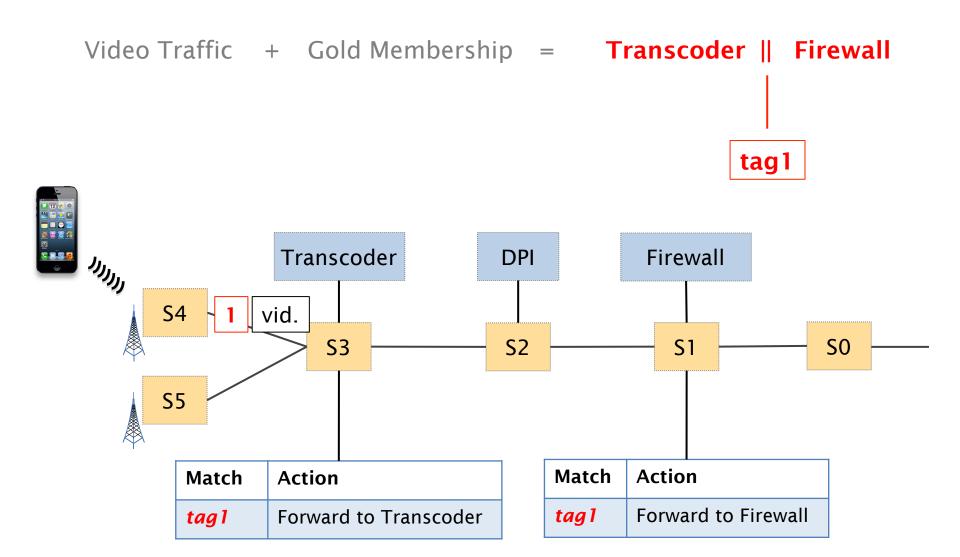
Video Traffic + Gold Membership = Transcoder || Firewall



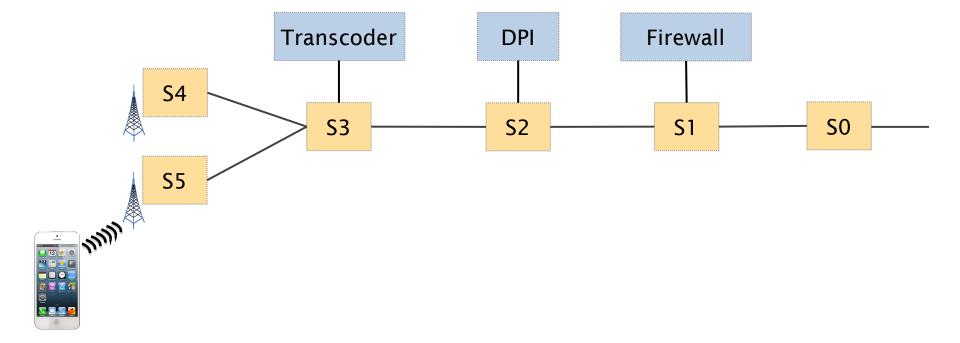


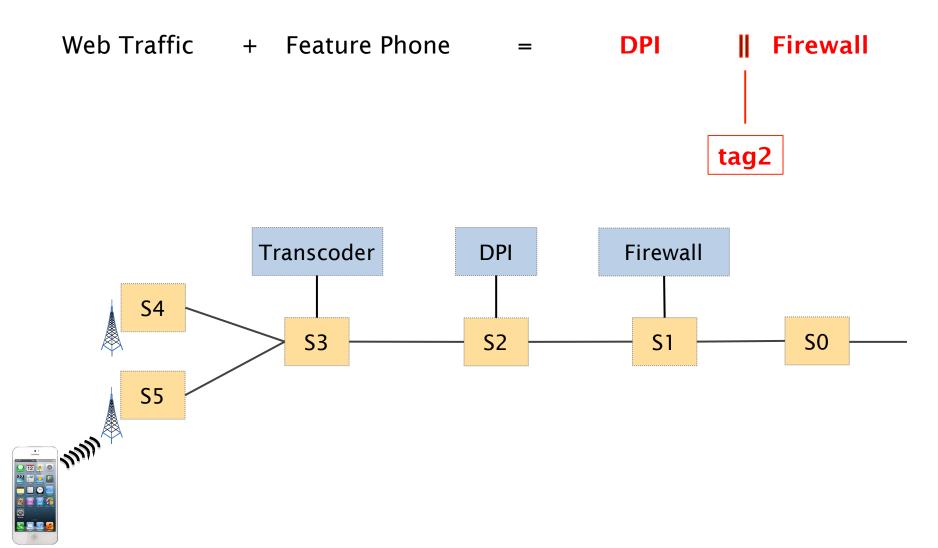


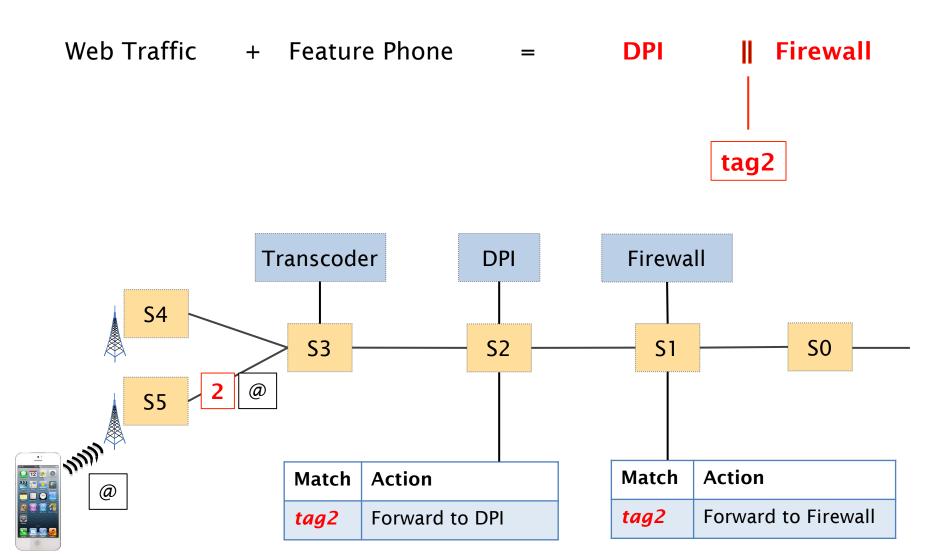




Web Traffic + Feature Phone = DPI || Firewall





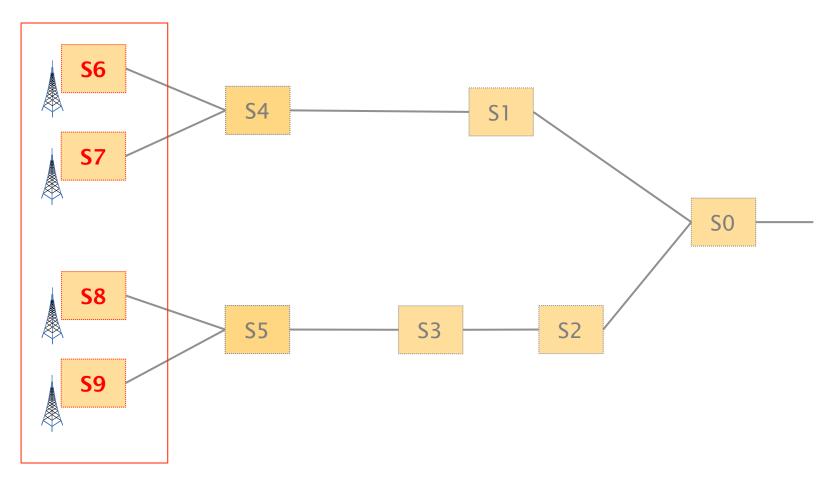


CellSDN tags are composed of three parts:

- Policy
- Location
- User Equipment Identifier

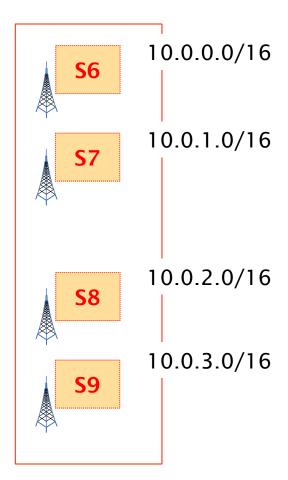
In CellSDN, each base station is associated with an IP prefix

Access switches



In CellSDN, each base station is associated with an IP prefix

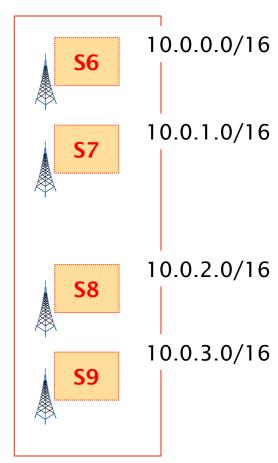
Access switches



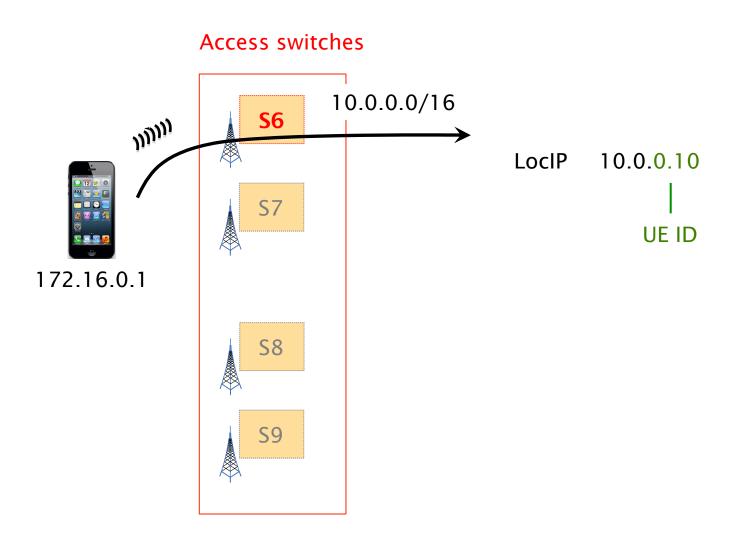
User Equipment receive a unique IP address

Access switches

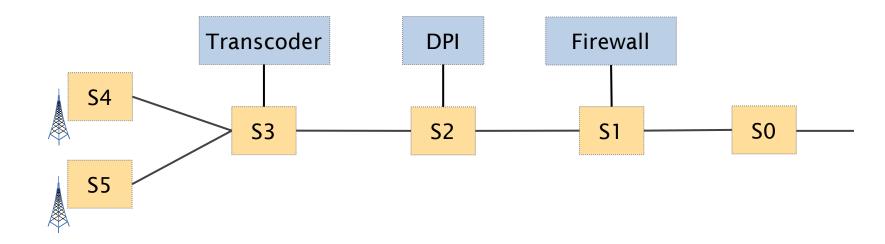




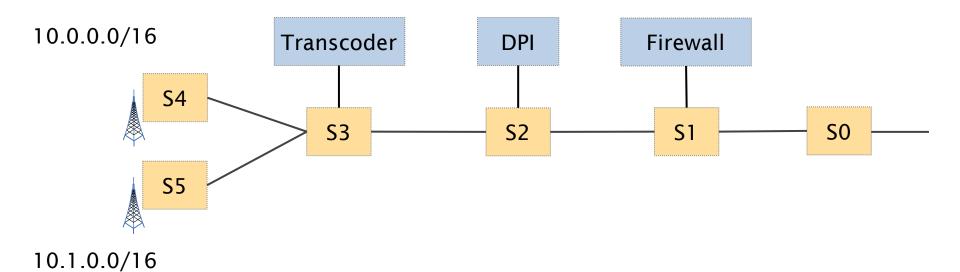
Access switches rewrite the UE IP into a location-dependent address



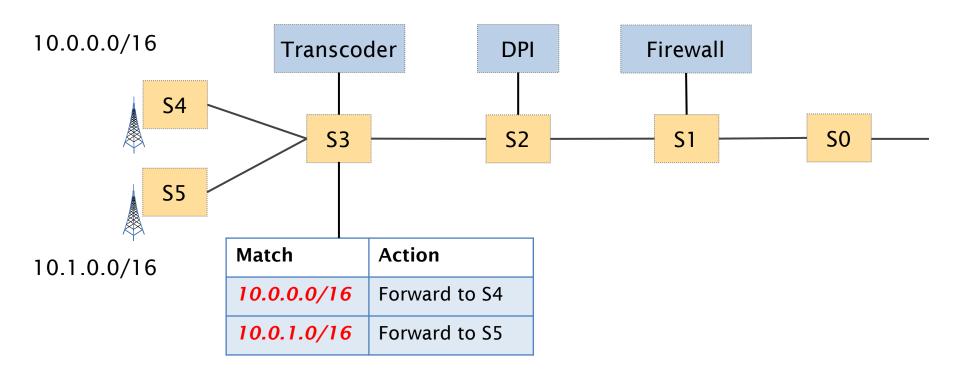
CellSDN can route based on Base Station prefixes



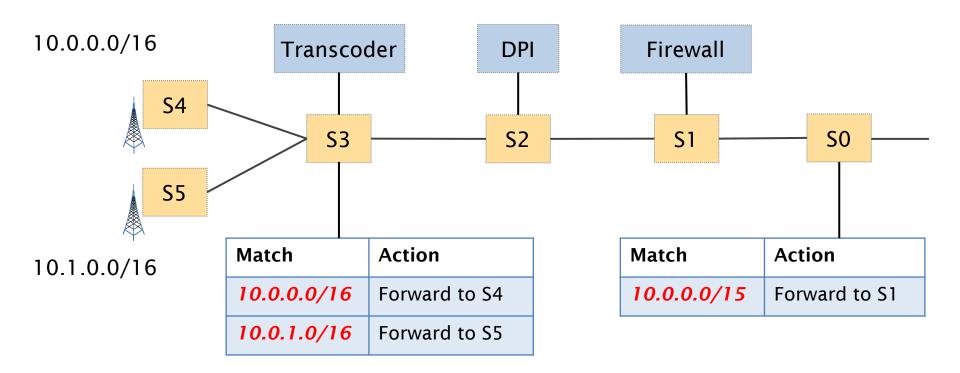
CellSDN can route based on Base Station prefixes



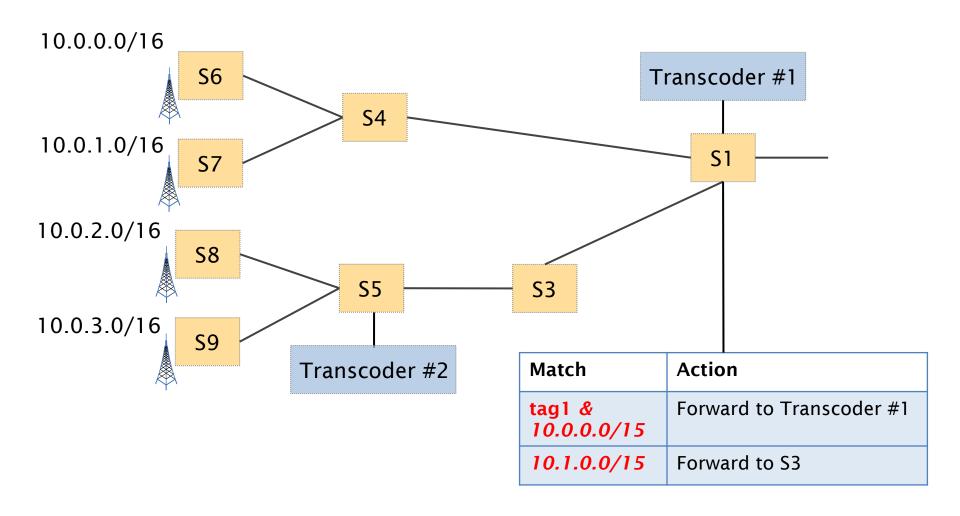
CellSDN can route based on Base Station prefixes



CellSDN automatically aggregates adjacent prefixes



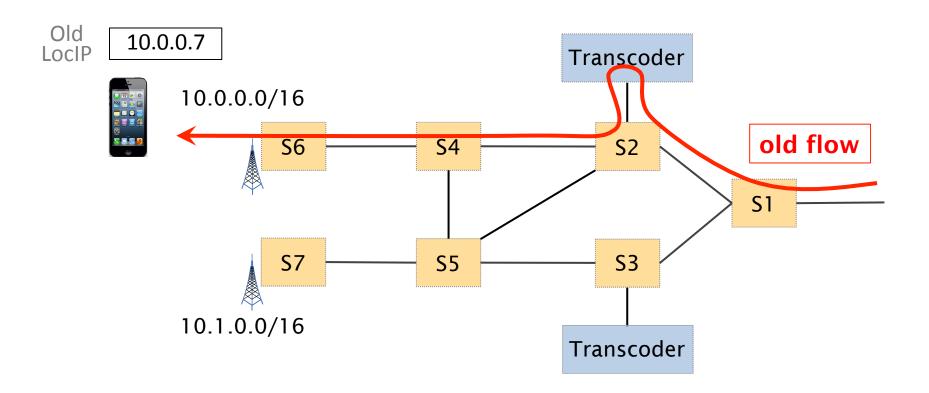
CellSDN can selectively match on tag and BS ID for load-balancing



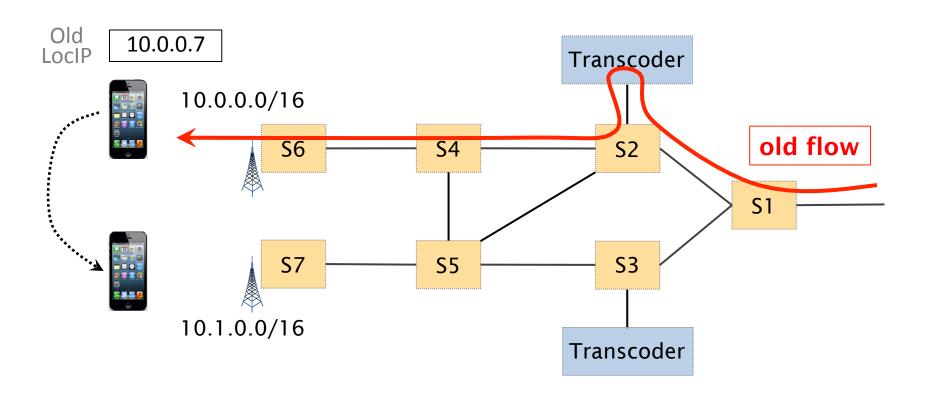
CellSDN tags are composed of three parts:

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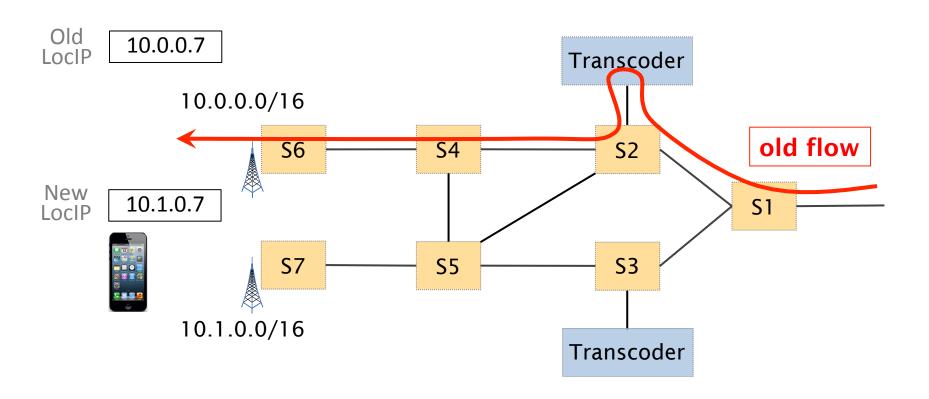
CellSDN UE tag enables mobility



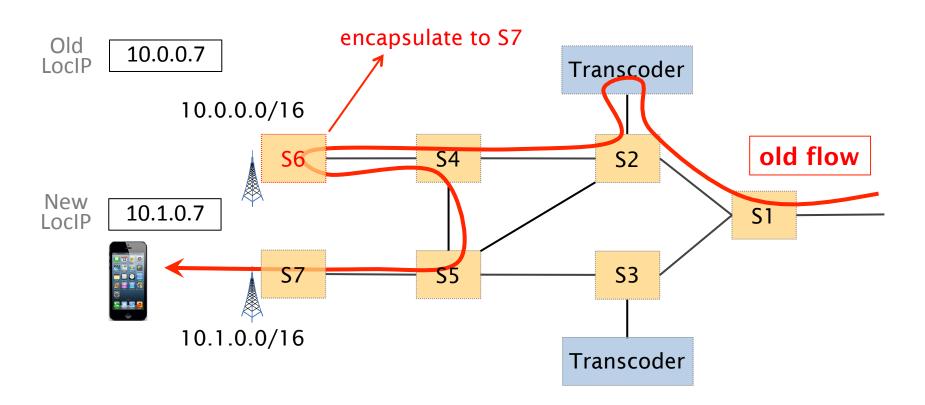
When an user moves, she receives a new LocIP



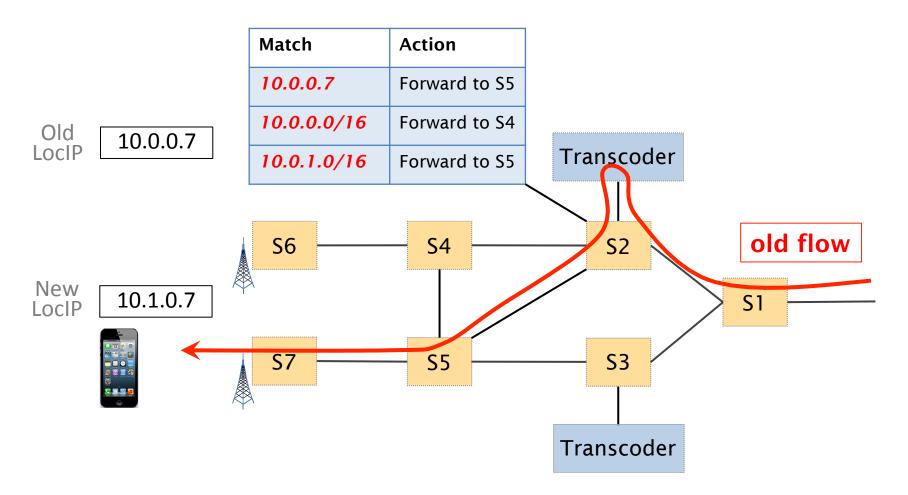
When an user moves, she receives a new LocIP



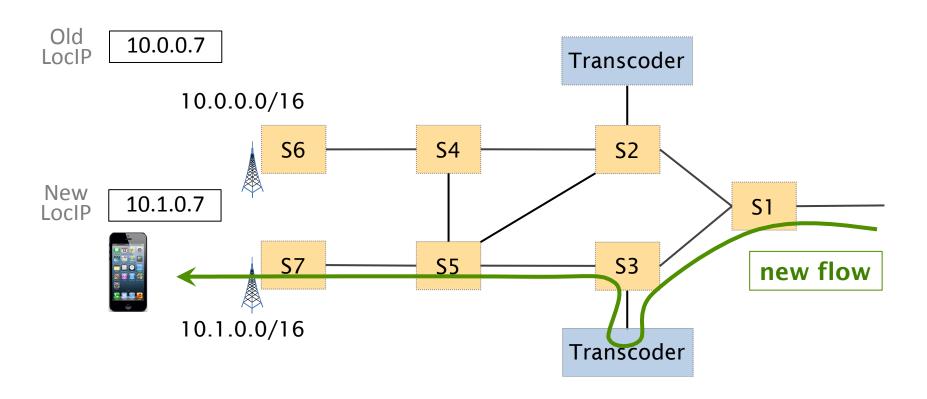
Old flows reach the previous base station, ensuring policy consistency



To avoid triangle routing, CellSDN can install shortcut path after the last MB

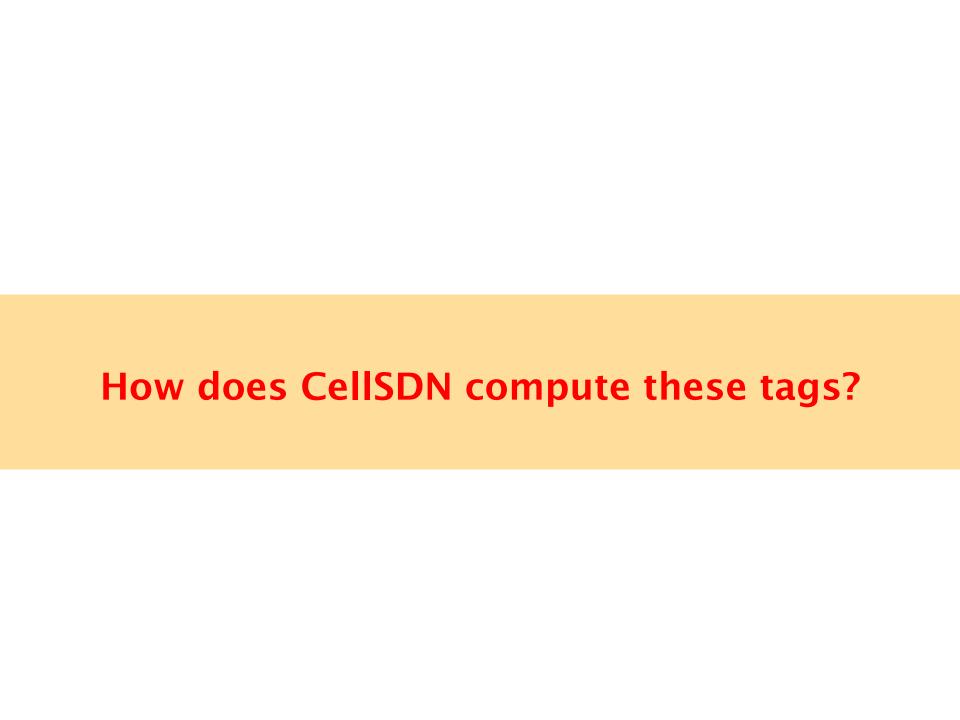


New flows automatically flow along new policy paths



CellSDN tags are composed of three parts:

- Policy
- Location
- User Equipment Identifier

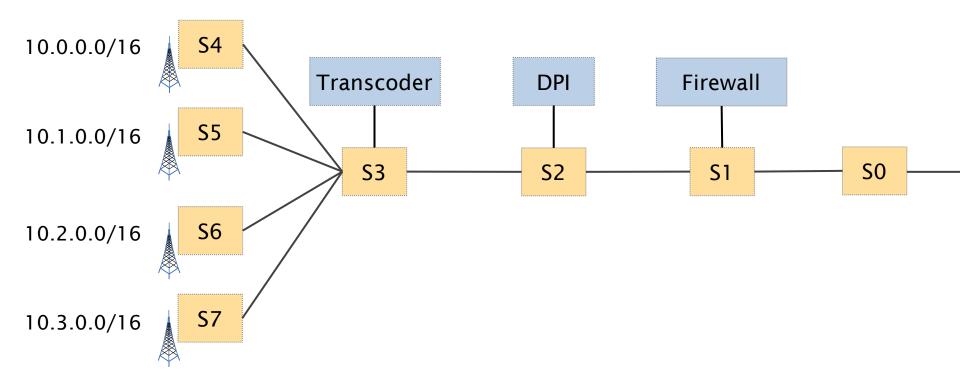


CellSDN minimizes the forwarding tables size by reusing tags

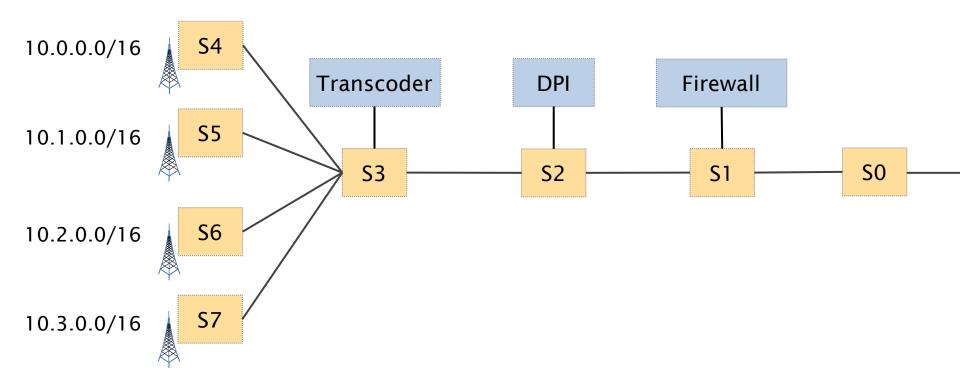
Given a service policy-path,

- 1. Compute the actual path *P* used in the network
- 2. For each candidate tag *t* used on the path *P*,

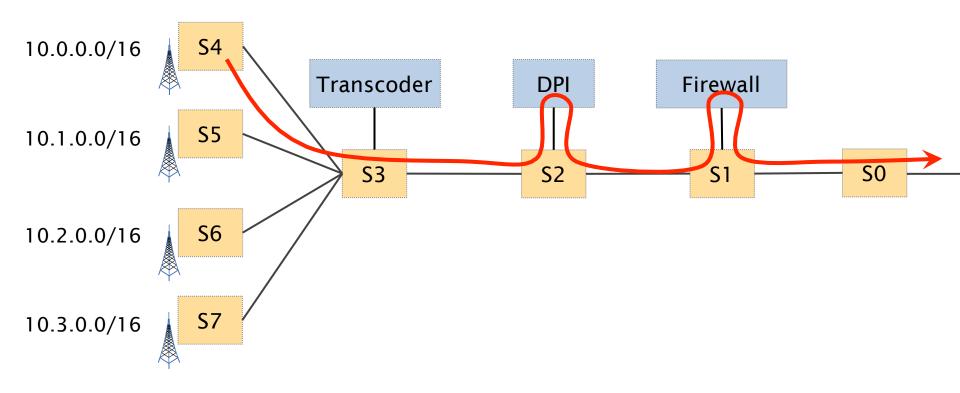
 Compute the # of new rules needed if *t* is used
- Select the candidate tag minimizing the # of new rules, Create new tag if none available
- 4. Install the forwarding entries in the network



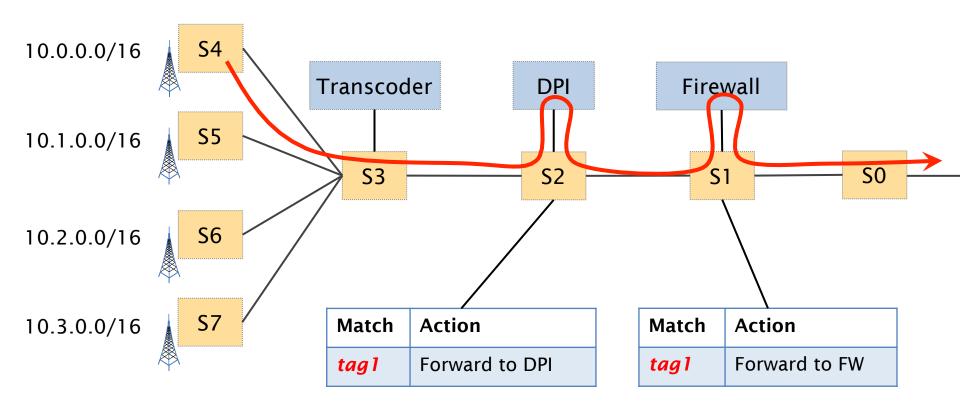




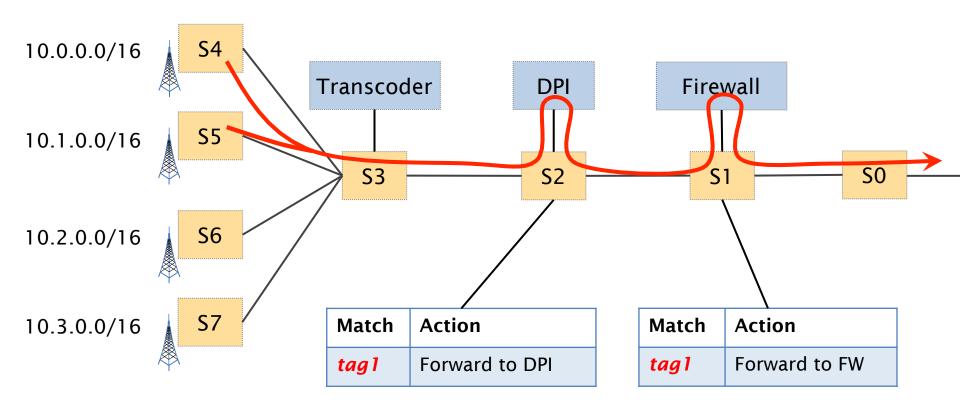




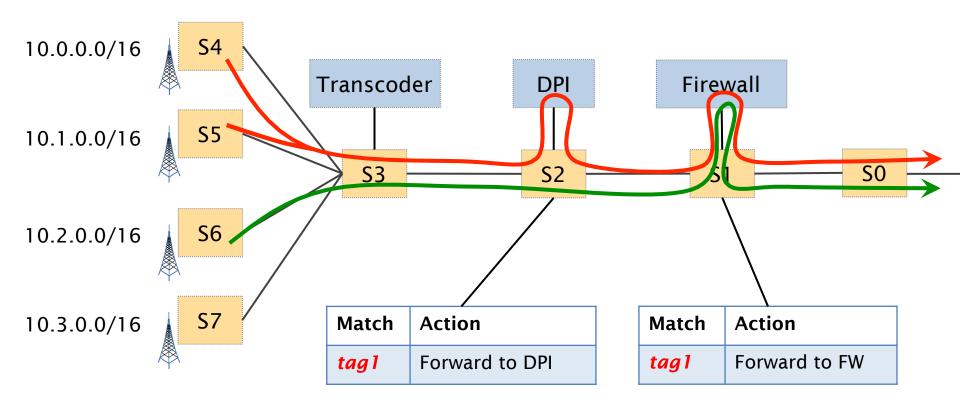




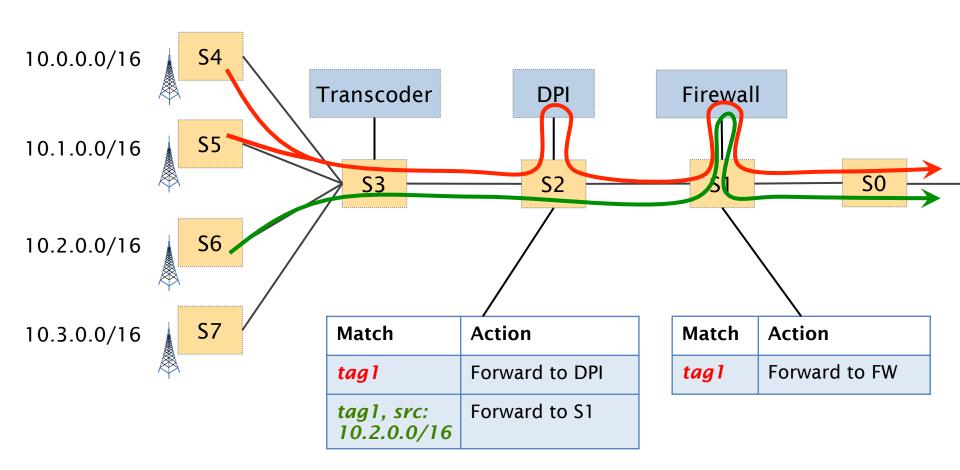




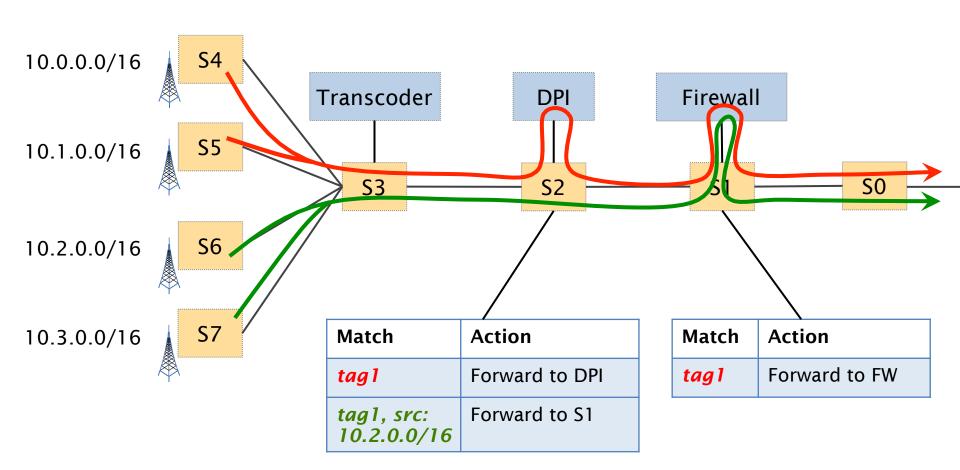




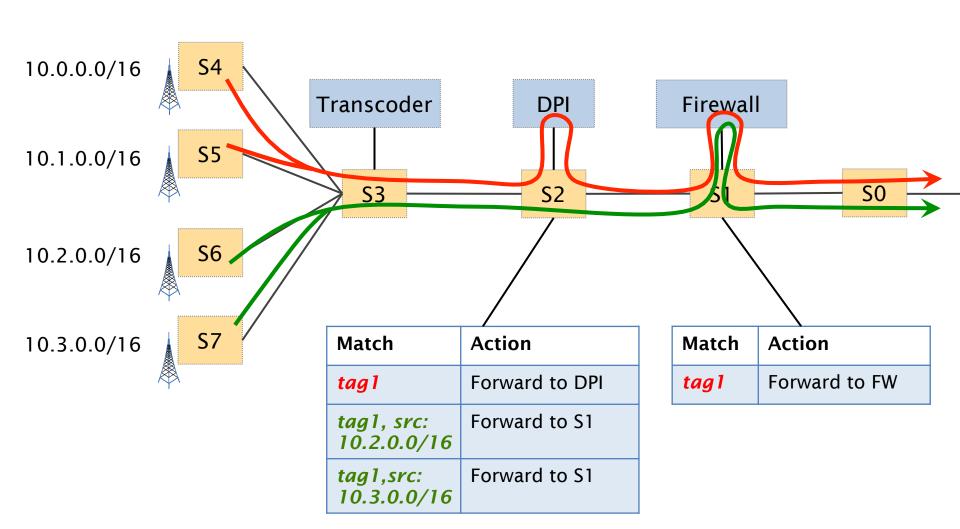




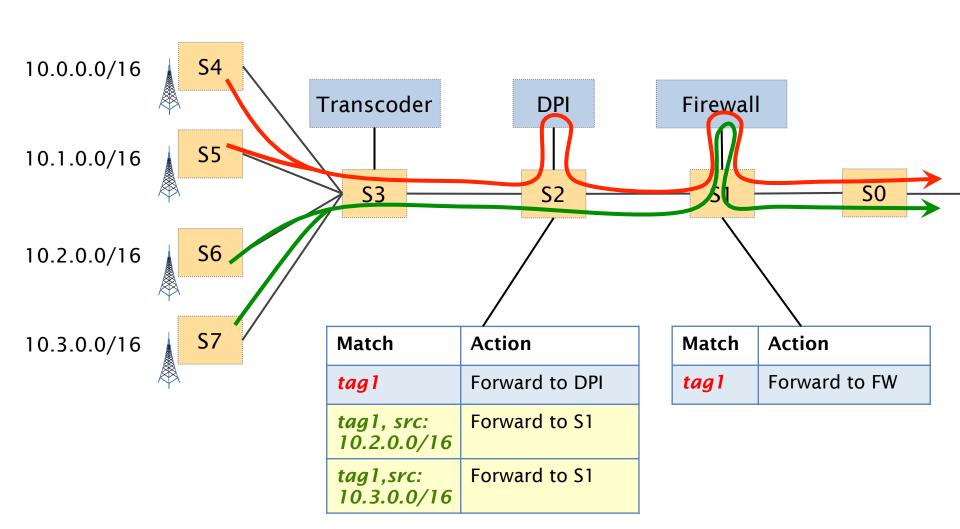
policy path S7 --> Firewall



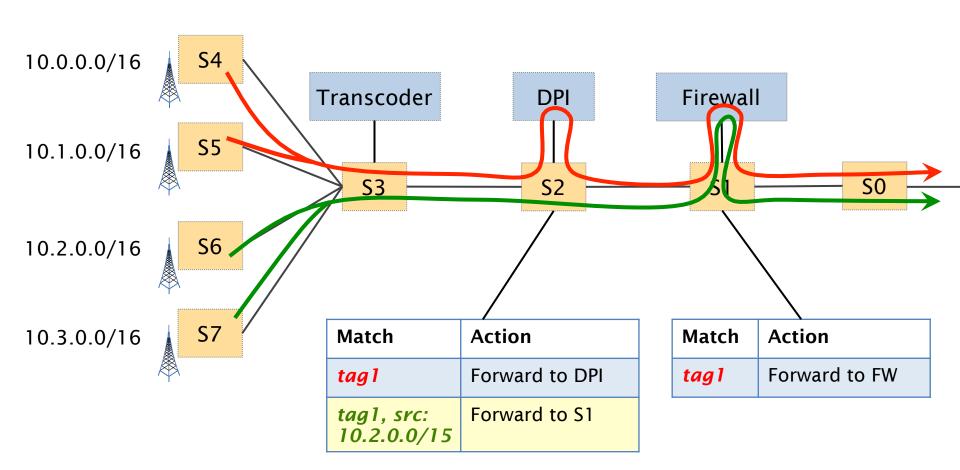
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policy path S7 → Firewall



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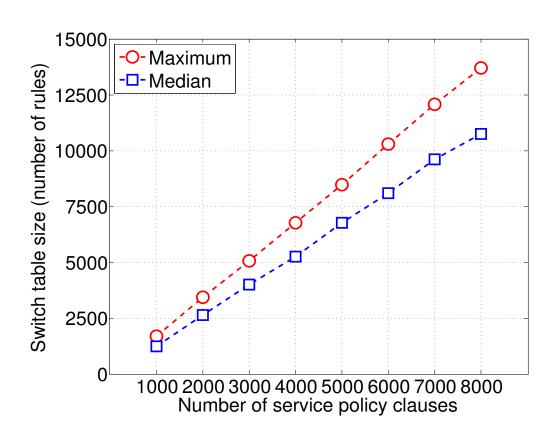
CellSDN dataplane can support a large number of service policies

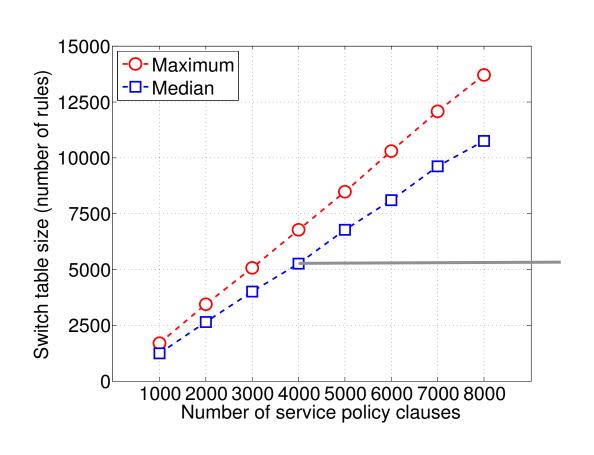
Simulation

- FatTree topology
- 128 switches && 1280 base stations
- 8 types of middleboxes
- 5 middleboxes long service policy

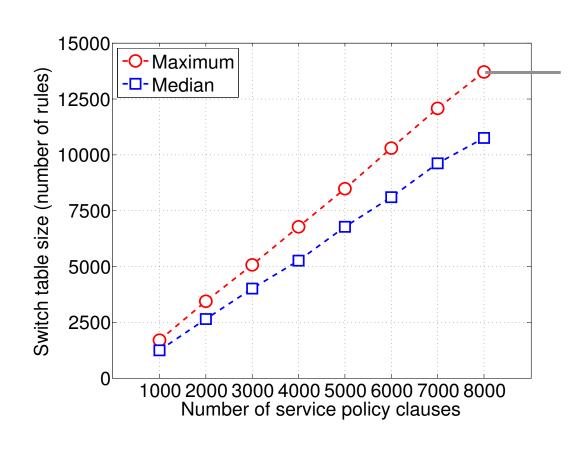
Evaluation

table size in function of the # of policies

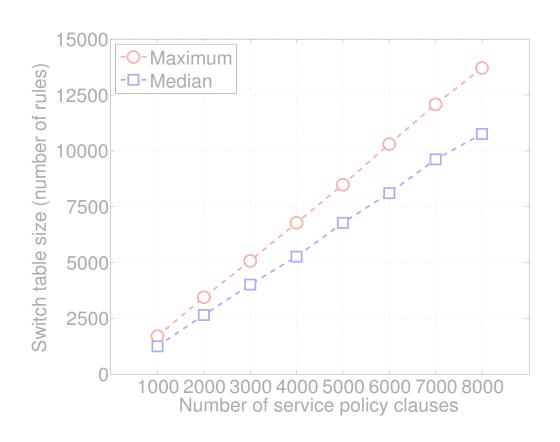




Only 5k entries are required to support 4k policy paths

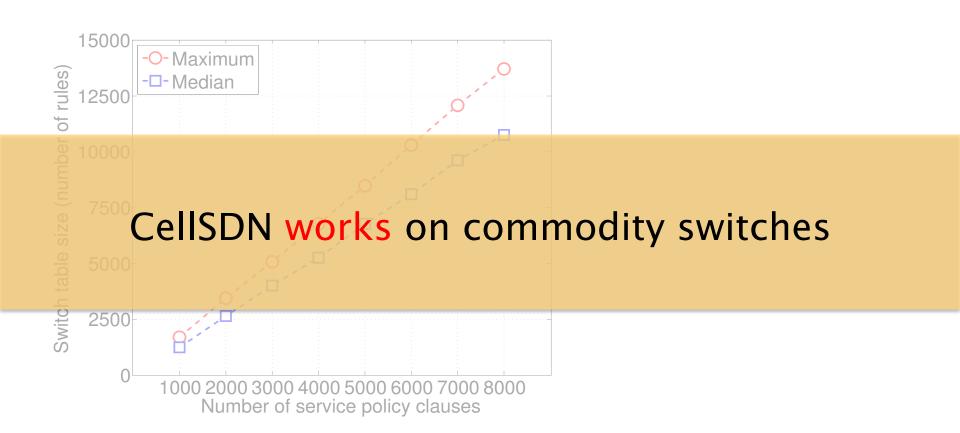


In the worst case, 8k policy paths require 13.6k entries



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Today, operators require a few hundreds policies



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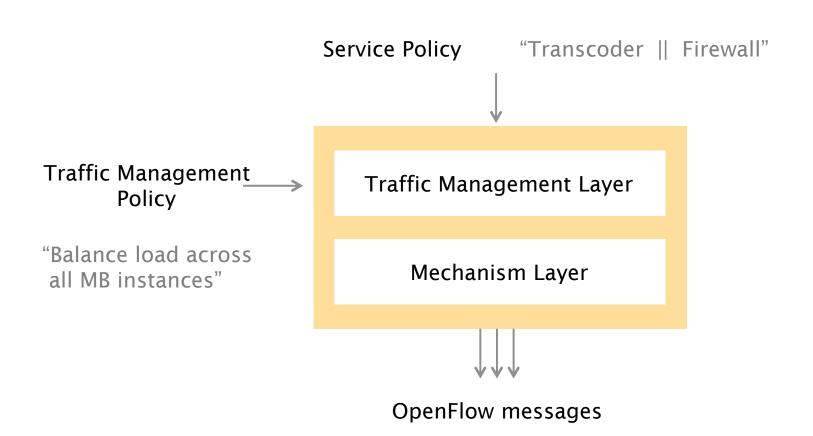


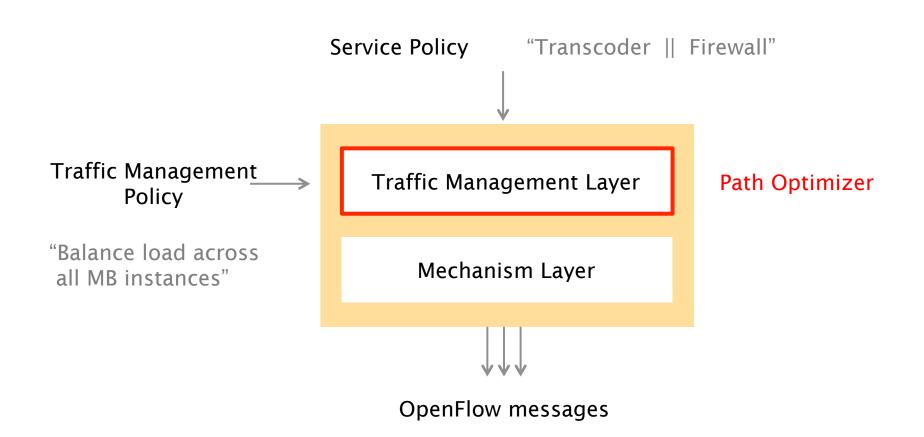
Architecture software-defined network

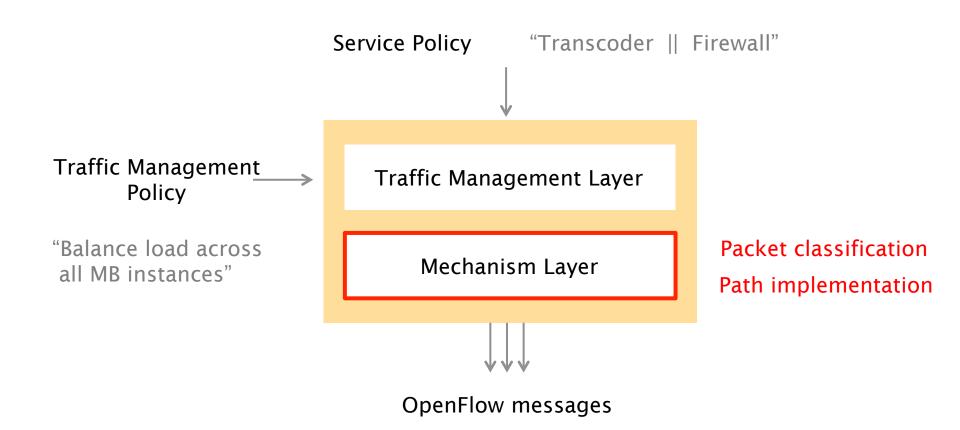
Scaling the data-plane multi-dimensional tagging

Scaling the control-plane tasks delegation

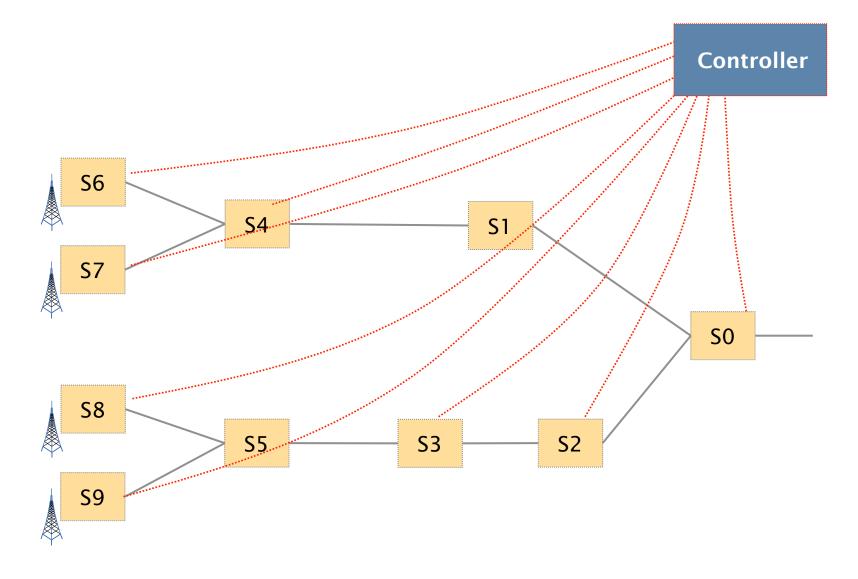
CellSDN controller separates traffic management from rules installation



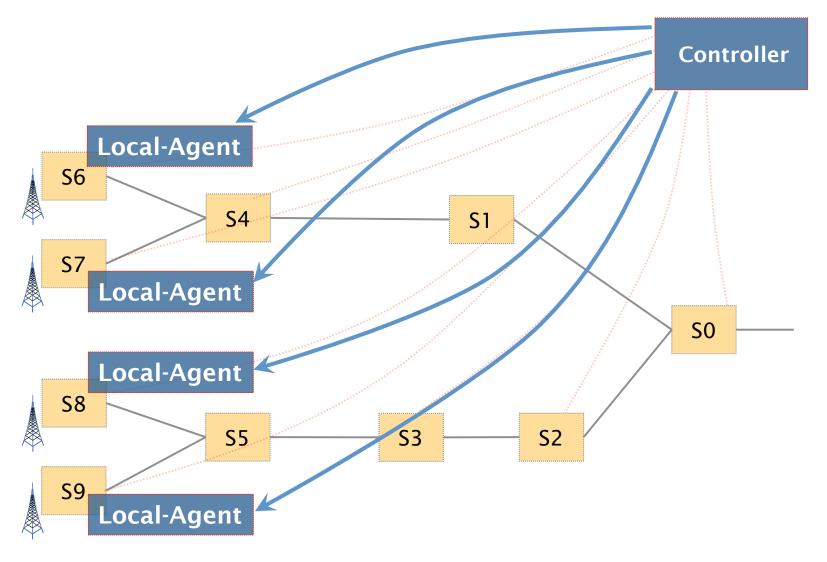




To scale, CellSDN uses a hierarchical controller



Most of the tasks are delegated to local-agents



Local-agents act as cache, reducing the load on the main controller

Local agents handle *locally* most frequent events

- cache a list of packet classifiers
- contact central controller upon cache miss
- tag flows

The central controller deals with less frequent, but more complex events

Local agents handle *locally* most frequent events

- cache a list of packet classifiers
- contact central controller upon cache miss
- tag flows

Central controller *globally* handle less frequent events

- UE arrival, handoff
- topology changes
- dynamic policies

CellSDN control-plane can handle the load of large cellular networks

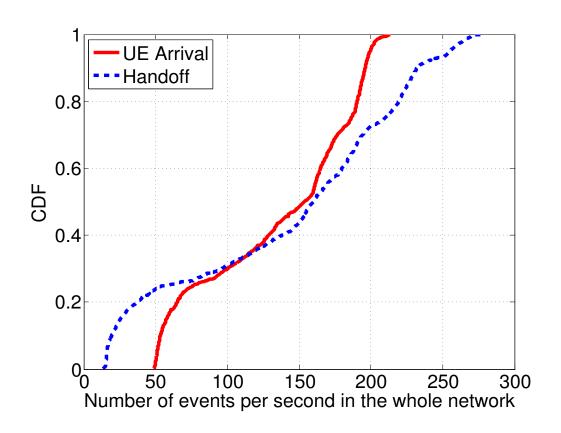
Dataset

- 1 week of traces from a large LTE network
- 1500 base stations
- 1 million users

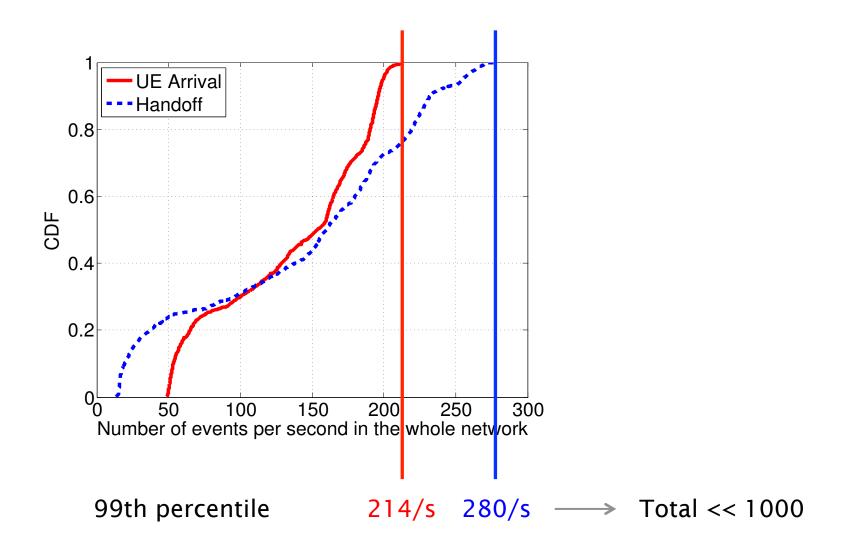
Evaluation

- # of events per second
- # of active users per second

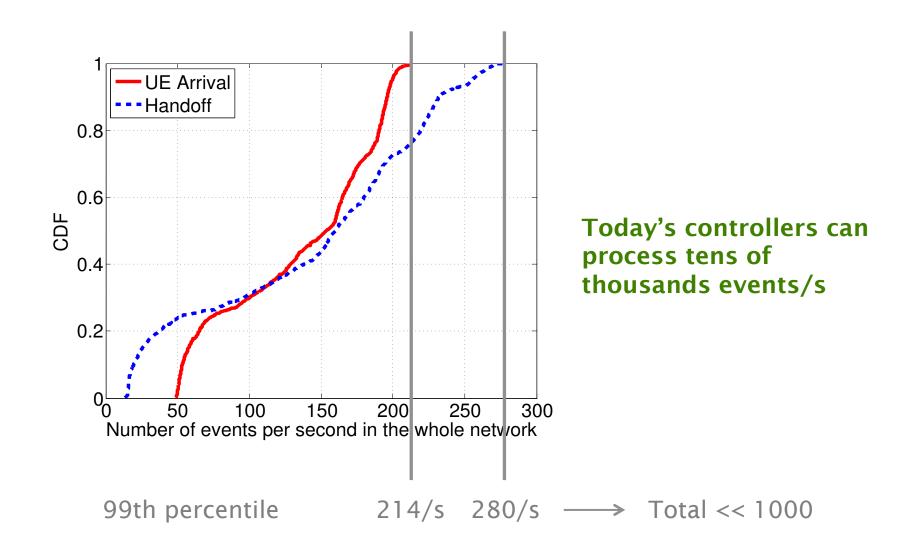
The number of events going to the main controller is small



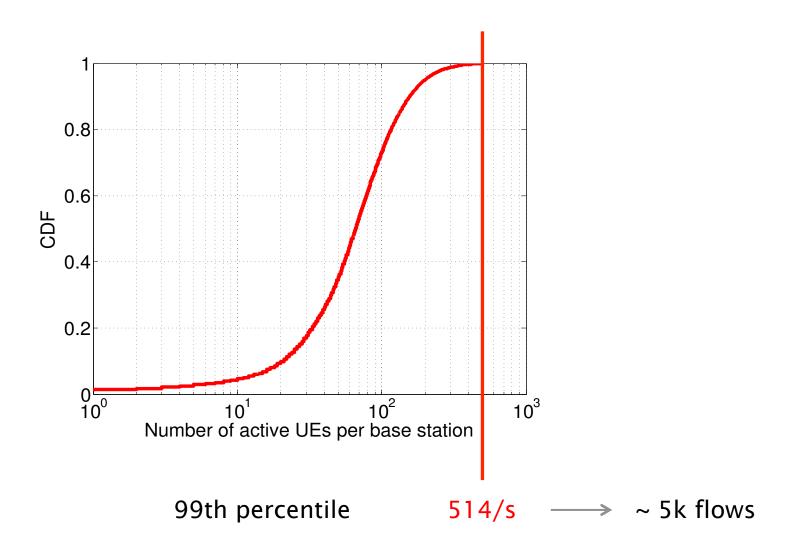
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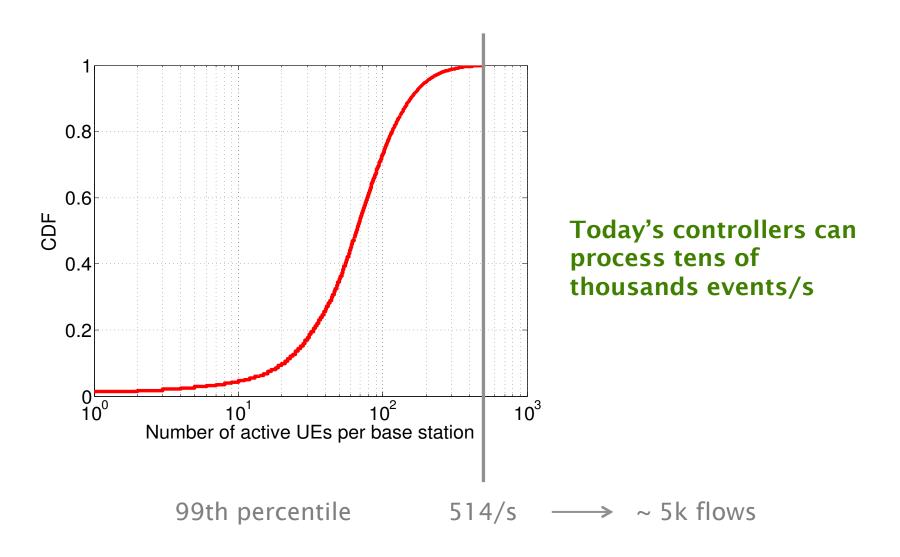
The number of events going to the main controller is small



The number of flows handled by the local-agent is small



The number of flows handled by the local-agent is small



CellSDN: Taking control of cellular core networks



Architecture software-defined network

Scaling the data-plane multi-dimensional tagging

Scaling the control-plane tasks delegation

CellSDN enables flexible and cost-effective cellular networks

CellSDN supports flexible fine-grained policies

CellSDN achieves scalability with

- Multi-dimensional aggregation
- Asymmetric edge design
- Hierarchical controller

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Laurent Vanbever

www.vanbever.eu

IBM Watson Research Center

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