

CS 430 Lecture 23 Activities

Disjoint-set Data Structure

It is useful in many applications to have a structure that handles groups of disjoint sets. In particular, we support the following three operations:

- **MAKE-SET**(x): Creates a new set consisting of a single element x . Since the sets are disjoint, we assume that this element is not already contained in any set.
- **UNION**(x, y): Given elements in two different sets, forms the union of two existing sets into one new set.
- **FIND-SET**(x): Returns a pointer to the representative of the set containing element x .

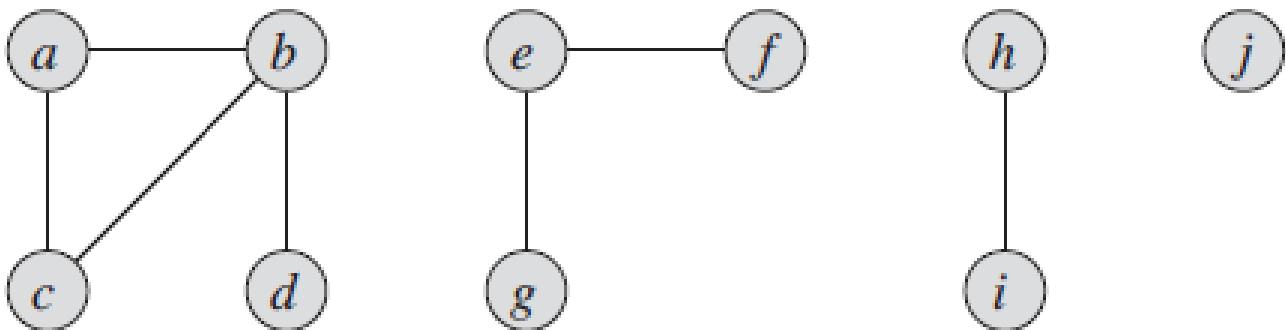
To identify a set, we return a pointer to any element in the set. The only constraint we make on the element chosen is that if the set does not change between calls to **FIND-SET**, we must return the same representative.

We analyze the algorithms implementing these operations in terms of n , the number of **MAKE-SET** operations (all **MAKE-SET**s are usually assumed to run first), and m , the total number of **MAKE-SET**, **UNION**, and **FIND-SET** operations ($n \leq m$).

1. Given the above, how many possible Union operations might there be?

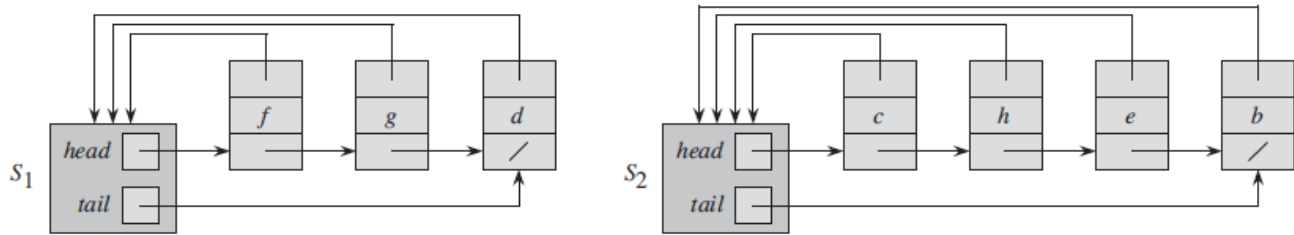
Disjoint-set Application

A graph data structure is a set of vertices and edges between those vertices, and supports problems where there can be relationships between any two items (vertices).



1. It is easy to see the disjoint sets (connected components) of the graph. Given a graph $G = (V, E)$ write an algorithm using **MAKE-SET**, **UNION**, and **FIND-SET** to find the connected components of any undirected graph.

Linked-List Representation



One simple approach is to represent a set as an unordered linked list. Each element contains two pointers, one to the next element (as in a simple linked list) and one to the head of the list. The head serves as the set representative.

- Describe the algorithms for `MAKE-SET(x)` and `FIND-SET(x)` including runtime. For `FIND-SET(x)`, assume you already have a reference to x .
- How do you think `UNION` is implemented?
- For n total elements, what is the maximum number of `MAKE-SET` and `UNION` operators that would need to be called in the worst case to get all the elements in one set? What is the maximum amortized runtime of each union?
- For n total elements, what is the minimum number of `MAKE-SET` and `UNION` operators that would need to be called in the best case to get all the elements in one set? What is the lower bound, worst case amortized runtime of each union?