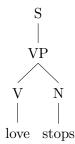
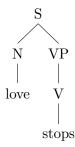
## CS4248 Assignment 3

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1. (Refer to algorithm at the end)

```
2. Input: "love stops"
     Chart[0]:
      \checkmark
                     S0
                              \gamma \to \bullet S
                                                    [0,0]
                                                              Dummy start state
                              S \rightarrow \bullet N VP
      ✓
                     S1
                                                    [0,0]
                                                              Predictor
                              S \to \bullet VP
                     S2
                                                     [0,0]
                                                              Predictor
      ✓
                     S3
                              VP \rightarrow \bullet V
                                                    [0,0]
                                                              Predictor
                     S4
                              \mathrm{VP} \to \bullet \ \mathrm{V} \ \mathrm{N}
                                                              Predictor
                                                    [0,0]
     Chart[1]:
                              N \to love \bullet
                      S5
                                                    [0,1]
                                                              Scanner
                     S6
                              V \rightarrow love \bullet
                                                     [0,1]
                                                              Scanner
     S7
                              S \to N \bullet VP
                                                              Completer (S5)
                                                     [0,1]
                     S8
                              \mathrm{VP} \to \mathrm{V} \, \bullet
                                                     [0,1]
                                                              Completer (S6)
                              VP \rightarrow V \bullet N
                     S9
                                                    [0,1]
                                                              Completer (S6)
                              VP \rightarrow \bullet V
                     S10
                                                              Predictor
                                                     [1,1]
                     S11
                              VP \rightarrow \bullet V N
                                                    [1,1]
                                                              Predictor
                              S \to VP \bullet
                     S12
                                                    [1,1]
                                                              Completer (S8)
                              \gamma \to S \bullet
                     S13
                                                              Completer (S12)
                                                    [1,1]
     Chart[2]:
                     S14
                              N \to stops \bullet
                                                    [1,2]
                                                              Scanner
                     S15
                              V \rightarrow stops \bullet
                                                    [1,2]
                                                              Scanner
                              VP \rightarrow V N \bullet
                                                    [1,2]
                     S16
                                                              Completer (S14)
                              VP \rightarrow V \bullet
                     S17
                                                     [1,2]
                                                              Completer (S15)
                     S18
                              VP \rightarrow V \bullet N
                                                    [1,2]
                                                              Completer (S15)
                              S \to VP \bullet
                     S19
                                                     [1,2]
                                                              Completer (S16)
                     S20
                              S \to N VP \bullet
                                                    [1,2]
                                                              Completer (S16)
                              \gamma \to S \bullet
                     S21
                                                    [1,2]
                                                              Completer (S19)
                              \gamma \to S \bullet
                     S22
                                                    [1,2]
                                                              Completer (S20)
```





- 3. The conventional CKY algorithm has time complexity of  $O(n^3)$ . To go through each item in the table takes  $n^2$ , and for each item checking through all the sub-trees takes n. Now that we consider probabilistic CKY, for each sub-tree there are multiple ways of parsing them, each with its own probability value. If we consider the worst case, there is an additional |G| steps, where G is the grammar. Thus, time complexity is  $O(|G| \cdot n^3)$ .
- 4. Sentence in consideration: "Flying helicopters can be hazardous."



In this interpretation, "Flying" modifies "helicopter". It describes a third-person perspective and means helicopters that are flying are hazardous.



In this interpretation, "Flying" is a verb. It describes a first-person perspective and means the act of piloting helicopters can be hazardous

```
Algorithm 1 convertRegexToCFG(R, G_t, A_{prev}) return G
```

```
R: Regex
G_t: Grammar
A_{prev}: Previous symbol or S (First call)
if R is terminal symbol then
  G_t.add(A_{prev} \to R)
  return G_t
else if R in form of (r) then
  Generate new non-terminal symbol, A_{new}
  G_t.add(A_{prev} \to A_{new})
  return convertRegexToCFG(r, G_t, A_{new})
else if R in form of r+ then
  return convertRegexToCFG(rr^*, G_t, A_{new})
else if R in form of r^* then
  Generate new non-terminal symbol, A_{new}
  G_t.add(A_{prev} \to A_{prev}A_{new})
  G_t.add(A_{prev} \to \epsilon)
  return convertRegexToCFG(r, G_t, A_{new})
else if R in form of r? then
  Generate new non-terminal symbol, A_{new}
  G_t.add(A_{prev} \to A_{new})
  G_t.add(A_{prev} \to \epsilon)
  return convertRegexToCFG(r, G_t, A_{new})
else if R in form of r\{d\} then
  Generate new non-terminal symbol, A_{new}
  A_{temp} \leftarrow null
  for 1 to d do
     A_{temp}.append(A_{new})
  end for
  G_t.add(A_{prev} \to A_{temp})
  return convertRegexToCFG(r, G_t, A_{new})
else if R in form of r\{d,\} then
  Generate new non-terminal symbol, A_{new}
  A_{temp} \leftarrow null
  for 1 to d do
     A_{temp}.append(A_{new})
  end for
  G_t.add(A_{prev} \to A_{temp})
  return convertRegexToCFG(r^*, G_t, A_{new})
else if R in form of r\{d_1,d_2\} then
  Generate new non-terminal symbol, A_{new}
  for d_i in d_1 to d_2 do
    convertRegexToCFG(r\{d_i\}, G_t, A_{new})
  end for
else if R in form of r_1 \mid r_2 then
  Generate 2 new non-terminal symbols, A and B
  G_t.add(A_{prev} \to A)
  G_t.add(A_{prev} \to B)
  return convertRegexToCFG(r_2, convertRegexToCFG(r_1, G_t, A), B)
else if R in form of r_1r_2 then
  Generate 2 new non-terminal symbols, A and B
  G_t.add(A_{prev} \to A)
  G_t.add(A_{prev} \to B)
  return convertRegexToCFG(r_2, convertRegexToCFG(r_1, G_t, A), B)
end if
```