# CS345 Lambda Calculus Project

By Levi Wiseman Heng Xiong

```
import Data.List
import System
data Expr = Var String
           App Expr Expr
           Lambda String Expr deriving (Eq. Read)
instance Show Expr where
        show (Var x) = x
        show (App x y) = "(" ++ show x ++ " " ++ show y ++ ")
        show (Lambda x e) = "(\\" ++ x ++ " -> " ++ show e ++ ")"
freeVars :: Expr -> [Expr]
free Vars x((Var) = [x]
freeVars (App t \bar{s}) = freeVars t ++ freeVars s
freeVars (Lambda x t) = [candidate | candidate <- freeVars t, candidate /= (Var x)]
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[M1M2][x := N] \equiv (M1[x := N]) (M2[x := N])
  (λγ.Μ)[x:="N] = λγ.(M[x:= N]), if x ≠ y and y ∉ FV(N)
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import Data.List
import System
-- abstract syntax tree
data Expr = Var String
App Expr E
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## Captures the idea of function application Corresponds to a computational step

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normalize :: Expr -> Expr
normalize = last . normalReduce
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#### Church Numerals

- Def: Church Numerals are functions of two arguments, where the first argument is applied a number of times to the second, the number of times it is applied is the integer representation.
- Example:
  - zero =  $\lambda fx.x$
  - one =  $\lambda fx.fx$
  - two =  $\lambda fx.f(fx)$
  - ...

### **Library Functions**

- We built the library functions based on Peano Axioms
- Started with inc(increment):
  - App inc zero = one
  - App inc one = two
  - ...
- plus:
  - App(App plus two) three = five
    - Currying: App plus two is a function of one argument
    - Binded two to plus first, then applied with three

### Library Function con.

- Booleans in Lambda Calculus
  - true =  $\lambda xy.x$
  - false =  $\lambda xy.y$
- Now we can do more functions with booleans
  - isZero
    - App isZero zero =  $\lambda xy.x$  (true)
    - App isZero one =  $\lambda xy.y$  (false)
  - ifThenElse
    - App (App ifThenElse true) (Var "42")) (Var "666")) = 42
    - App (App (App ifThenElse true) (Var "42")) (Var "666")) = 666

### **Big Cool Factorial**

- Before we do factorial, we need to handle recursion
  - A fixed point combinator (Y combinator) is a higherorder function that computes a fixed point of functions.
    - $(0)^2 = 0$ , therefore 0 is a fixed point of  $f(x) = x^2$
- factorial = fix (λfx.ifThenElse (isZero x) 1 (mult x (f (dec x))))
  - App factorial three = six
  - App factorial four = twenty-four
    - Without normalize, it will look like ...
    - Hence, thanks to normalization!

#### **Aftermath**

 We've shown that Lambda Calculus ≤ Haskell's type system, therefore we have trivially proved that Haskell's type system is Turing complete!

# FIN