```
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//KDtree
                                                                                             void query(int k, int now) {
const int MAXN=200005, MAXD=2;
                                                                                                        double d, dI=INF, dr=INF;
                                                                                                        d=dis(t[k],T)\;;\;\;ans=min(ans,d)\;;\;\;
const double INF=1e10;
int n.m.D:
                                                                                                        if (t[k]. I) dl=get(t[k]. I. T):
                                                                                                        if (t[k].r) dr=get(t[k].r,T);
struct P{
          int d[MAXD], M[MAXD], m[MAXD], I, r;
                                                                                                        if (dl<dr) {
          P(int x=0, int y=0) { I=r=0; d[0]=x; d[1]=y; }
                                                                                                                   if (dl<ans) query(t[k].l.now^1):
          int& operator [](int x) { return d[x]; }
                                                                                                                   if (dr<ans) query(t[k].r, now^1);</pre>
          friend bool operator \langle (P a, P b) \{ return a[D] \langle b[D]; \} \}
                                                                                                        }
          friend bool operator == (P a, P b) { return
                                                                                                        else {
a. d[0]==b. d[0]&&a. d[1]==b. d[1];
                                                                                                                   if (dr<ans) query(t[k].r, now^1);</pre>
}p[MAXN];
                                                                                                                   if (dl<ans) query(t[k].1, now^1);
inline double dis(P a, P b)
{ return
sqrt(1.0*(a[0]-b[0])*(a[0]-b[0])+1.0*(a[1]-b[1])*(a[1]-b[1]));
                                                                                              inline double query(P p) { ans=INF; T=p; query(root, 0); return
inline bool in(int x1, int y1, int x2, int y2, int X1, int Y1, int X2, int Y2)
                                                                                   ans: }
{ return x1<=X1&&X2<=x2&&y1<=Y1&&Y2<=y2; }
                                                                                   }kd:
inline bool out(int x1, int y1, int x2, int y2, int X1, int Y1, int X2, int Y2)
{ return x1>X2||x2<X1||y1>Y2||y2<Y1; }
struct KDtree{
                                                                                   //Line intersect
          P t[MAXN], T;
                                                                                   inline bool intersect (P a, P b, P c, P d) {
          int root.n:
                                                                                              if (\max(a. x, b. x) \le \min(c. x, d. x) \mid \min(a. x, b. x) \ge \max(c. x, d. x))
          double ans:
                                                                                   return false:
          inline void update(int k) {
                                                                                             if (\max(a. y, b. y) \le \min(c. y, d. y) \mid \min(a. y, b. y) \ge \max(c. y, d. y))
                     int i, I=t[k]. I, r=t[k]. r;
                                                                                   return false;
                     for (i=0; i < MAXD; ++i) {
                                                                                             doub Le
                                if (I) {
                                                                                   x=cross(a-b, c-b), y=cross(a-b, d-b), z=cross(c-d, a-d), w=cross(c-d, b-d);
                                                                                              if (sgn(x)==0\&sgn(y)==0) return false;
          t[k].M[i]=max(t[k].M[i],t[l].M[i]);
                                                                                              if (sgn(z)==0) return true; else if (sgn(w)==0) return false;
                                                                                              if (sgn(x)*sgn(y) \le 0\&sgn(z)*sgn(w) \le 0) return true; else
          t[k].m[i]=min(t[k].m[i],t[l].m[i]);
                                                                                   return false;
                                if (r) {
                                                                                   inline P work (P a, P b, P c, P d) {
                                                                                             P ans:
          t[k]. M[i]=max(t[k]. M[i], t[r]. M[i]);
                                                                                             double a1=b. y-a. y, b1=-(b. x-a. x), c1=a. x*b. y-a. y*b. x;
                                                                                             double a2=d. y-c. y, b2=-(d. x-c. x), c2=c. x*d. y-c. y*d. x;
          t[k].m[i]=min(t[k].m[i],t[r].m[i]);
                                                                                             ans. x=(c1*b2-c2*b1)/(a1*b2-a2*b1);
                                                                                   ans. y=(a1*c2-a2*c1)/(a1*b2-a2*b1);
                                                                                             return ans:
          }
                                                                                   }
          int build(int I, int r, int now) {
                                                                                   //Point
                     D=now:
                     int i.mid=(l+r)>>1:
                                                                                   const int MAXN=300005, Mo=1000000007;
                     nth_element(p+I, p+mid, p+r+1); t[mid]=p[mid];
                                                                                   const double eps=1e-10, INF=1e20, PI=3. 141592653589793238462;
                     for (i=0; i<MAXD;++i)
                                                                                   inline int sgn(double x) { if (fabs(x) \leq eps) return 0; return x>0?1:-1; }
t[mid]. M[i]=t[mid]. m[i]=t[mid][i];
                                                                                   struct P{
                     if (I<mid) t[mid]. I=build(I, mid-1, now^1);</pre>
                                                                                             double x.v:
                     if (mid<r) t[mid].r=build(mid+1,r,now^1);</pre>
                                                                                             P(double xx=0, double yy=0):x(xx), y(yy) {}
                                                                                             double len() { return sqrt(x*x+y*y); }
                     update(mid); return mid;
                                                                                             P operator +(const P &E) const { return P(x+E. x, y+E. y); }
          inline void build(int nn) { n=nn; root=build(1, nn, 0); }
                                                                                             P operator -(const\ P\ \&E) const\{\ return\ P(x-E.\ x,\ y-E.\ y)\ ;\ \}
          void insert(int &k, int now) {
                                                                                             P operator *(const double &f)const{ return P(x*f, y*f); }
                     if (k==0) {
                                                                                             P operator /(const double &f)const{ return P(x/f, y/f); }
                                t[k].r=++n; t[n]=T;
                                                                                              double operator *(const P &E) const{ return x*E. y-E. x*y; }
                                for (int i=0; i \le MAXD; ++i)
                                                                                             double operator ^(const P &E)const{ return x*E.x+y*E.y; }
t[n].M[i]=t[n].m[i]=t[n][i];
                                                                                   } A. B. C. O:
                                return;
                                                                                   inline double dist(P a, P b) { return
                                                                                   sqrt((a. x-b. x)*(a. x-b. x)+(a. y-b. y)*(a. y-b. y)); 
                     if (T[now] \ge t[k][now]) insert(t[k].r, now^1);
                                                                                   inline double cross(P a, P b, P o) { return (a-o)*(b-o); }
                     else insert(t[k].l, now^1);
                                                                                   inline double dot(P a, P b, P o) { return (a-o)^(b-o); }
                     update(k);
                                                                                   inline double deg(P p) { double tmp=atan2(p.y,p.x); return
                                                                                   tmp<0?tmp+PI*2:tmp; }</pre>
          inline void insert(P p) { T=p; T. I=T. r=0; insert(root, 0); }
                                                                                   inline P rot(P p, double a) { return
          inline double get(int k, P p) {
                                                                                   P(\cos(a)*p. x-\sin(a)*p. y, \sin(a)*p. x+\cos(a)*p. y);}
                     double tmp[2] = \{0.0, 0.0\};
                                                                                   struct | {
                     for (int i=0:i<MAXD:++i) {</pre>
                                                                                             double a.b.c:
                                if (tmp[i] < t[k].m[i]-p[i])
                                                                                             L(P A, P B) \{ a=B. y-A. y; b=A. x-B. x; c=-(A*B); \}
tmp[i]=t[k].m[i]-p[i];
                                                                                   } | 11, | 12;
                                if (tmp[i]<p[i]-t[k].M[i])</pre>
                                                                                   inline P clammer(L I, L r) {
                                                                                             P ans (-1. c*r. b+1. b*r. c, -1. a*r. c+1. c*r. a);
tmp[i]=p[i]-t[k].M[i];
                                                                                             return ans/(I.a*r.b-r.a*I.b);
                                                                                   }
                     return sqrt(tmp[0]*tmp[0]+tmp[1]*tmp[1]);
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int ConvexHull(P p[], int n, P q[]) {
          sort(p+1, p+n+1);
                                                                              double vlen(P u) {
          int i.m=0:
                                                                                        return sqrt (u. x*u. x+u. y*u. y+u. z*u. z);
          for (i=1;i<=n;i++) {
                    while (m)1\&coss(q[m]-q[m-1], p[i]-q[m-1]) \le eps)
                                                                              P pvec(P a, P b, P c) {
                                                                                        return xmult(subt(a, b), subt(b, c));
m--:
                    q[++m]=p[i];
                                                                              double Dis(P a. P b. P c. P d) {
          }
          int k=m:
                                                                                        return fabs (dmult (pvec (a, b, c), subt (d, a))) / vlen (pvec (a, b, c));
          for (i=n-1; i>0; i--) {
                    while (m>k\&\&cross(q[m]-q[m-1], p[i]-q[m-1]) \le eps)
                                                                              struct T3dhull{
                                                                                        int n;//初始点数
m--:
                    q[++m]=p[i];
                                                                                        P ply[N];//初始点
                                                                                        int trianglecnt;//凸包上三角形数
          if (n>1) m--;
                                                                                        fac tri[N];//凸包三角形
                                                                                        int vis[N][N];//点i到点j是属于哪个面
          return m:
                                                                                        double dist(P a) \{ return \ sqrt(a. \ x*a. \ x+a. \ y*a. \ y*a. \ z*a. \ z); \} //两
}
                                                                              点长度
                                                                                        double area(Pa, Pb, Pc) {return dist((b-a)*(c-a));}//三角形
//旋转卡亭
                                                                              面积*2
double Rotating_calipers(P p[], int n, P q[], int m) {
                                                                                        double volume (P a, P b, P c, P d) {return (b-a)*(c-a)^(d-a);}//
          int v, u, i;
                                                                              四面体有向体积*6
          double ans=INF, tmp;
                                                                                        double ptoplane (P &p, fac &f) {//正: 点在面同向
          for (i=u=1; i \le n; ++i) if (p[u]. y-eps > p[i]. y) u=i;
          for (i=v=1; i \le m; ++i) if (q[v].y+eps \le q[i].y) v=i;
                                                                              m=ply[f.b]-ply[f.a], n=ply[f.c]-ply[f.a], t=p-ply[f.a];
          for (i=1; i \le n; ++i) {
                                                                                                   return (m*n) ^t;
                    while
((tmp=cross(p[u+1]-p[u],q[v+1]-p[u])-cross(p[u+1]-p[u],q[v]-p[u]))>ep
                                                                                        void deal(int p, int a, int b) {
                                                                                                  int f=vis[a][b];//与当前面(cnt)共边(ab)的那个面
                              { v++; if (v==m+1) v=1; }
                    ans=min(ans, work(p[u], p[u+1], q[v]));
                                                                                                  if(tri[f].ok){
                    if (fabs(tmp)<=eps) {</pre>
                                                                                                             if((ptoplane(ply[p], tri[f]))>PR)
                                                                              dfs(p,f);//如果p点能看到该面f,则继续深度探索f的3条边,以便更新新
                              ans=min(ans, work(p[u], p[u+1], q[v+1]);
                              ans=min(ans, work(q[v], q[v+1], p[u]));
                                                                              的凸包面
                              ans=min(ans, work(q[v], q[v+1], p[u+1]));
                                                                                                            else//否则因为p点只看到 cnt 面,看不到 f
                                                                              面,则p点和a、b点组成一个三角形。
                    u++; if (u==n+1) u=1;
          return ans:
                                                                              add. a=b, add. b=a, add. c=p, add. ok=1;
}
                                                                              vis[p][b]=vis[a][p]=vis[b][a]=trianglecnt;
                                                                                                                       tri[trianglecnt++]=add;
//三维凸句
                                                                                                            }
#include<stdio.h>
#include<string.h>
#include(math.h)
                                                                                        void dfs (int p, int cnt) {//维护凸包,如果点p在凸包外更新凸包
#include<algorithm>
                                                                                                  tri[cnt]. ok=0;//当前面需要删除,因为它在更大的凸包
#include <iostream>
using namespace std;
                                                                              //下面把边反过来(先b,后a),以便在 deal()中判断与当前面(cnt)共边(ab)的
#define PR 1e-9
                                                                              那个面。即判断与当头面(cnt)相邻的3个面(它们与当前面的共边是反向的,如
#define N 1100
                                                                              下图中(1)的法线朝外(即逆时针)的面 130 和 312,它们共边 13,但一个方向是
struct P{
                                                                              13, 另一个方向是 31)
          double x, v, z:
                                                                                                  deal(p, tri[cnt].b, tri[cnt].a);
          P(double x=0, double y=0, double z=0): x(x), y(y), z(z) {}
                                                                                                  deal(p, tri[cnt]. c, tri[cnt]. b);
          P operator-(const P p) {return P(x-p. x, y-p. y, z-p. z);}
                                                                                                  deal(p, tri[cnt]. a, tri[cnt]. c);
          P operator*(const P p) {return
P(y*p. z-z*p. y, z*p. x-x*p. z, x*p. y-y*p. x);}//叉积
                                                                                        bool same (int s. int e) {//判断两个面是否为同一面
          double operator (const P p) {return x*p. x+y*p. y+z*p. z;}//点
                                                                                                  P a=ply[tri[s].a],b=ply[tri[s].b],c=ply[tri[s].c];
和
                                                                                                  return fabs(volume(a, b, c, ply[tri[e]. a])) <PR</pre>
                                                                                                            &&fabs(volume(a, b, c, ply[tri[e].b]))\langle PR
};
P dd:
                                                                                                            &&fabs(volume(a, b, c, ply[tri[e].c])) <PR;
struct fac{
          int a, b, c;//凸包一个面上的三个点的编号
                                                                                        void construct() {//构建凸包
          bool ok;//该面是否是最终凸包中的面
                                                                                                  int i, j;
                                                                                                  trianglecnt=0;
}:
P xmult(P u, P v) {
                                                                                                  if(n<4) return;
                                                                                                  bool tmp=true:
          return P(u. y*v. z-v. y*u. z, u. z*v. x-u. x*v. z, u. x*v. y-u. y*v. x);
                                                                                                  for (i=1; i<n; i++) {//前两点不共点
double dmult(P u, P v) {
                                                                                                             if((\mathsf{dist}(\mathsf{ply}[0]\mathsf{-}\mathsf{ply}[i]))\mathsf{>}\mathsf{PR})
          return u. x*v. x+u. y*v. y+u. z*v. z;
                                                                                                             {
                                                                                                                      swap(ply[1], ply[i]); tmp=false;
P subt(P u, P v) {
                                                                              break;
                                                                                                            }
          \texttt{return} \ \mathsf{P}(\mathsf{u}.\ \mathsf{x-v}.\ \mathsf{x}, \mathsf{u}.\ \mathsf{y-v}.\ \mathsf{y}, \mathsf{u}.\ \mathsf{z-v}.\ \mathsf{z})\ ;
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}
                   if(tmp) return;
                                                                                     return 0;
                   tmp=true:
                   for (i=2; i<n; i++) {//前三点不共线
                                                                           //hdu4266
          if((dist((ply[0]-ply[1])*(ply[1]-ply[i])))>PR) {
                                      swap(ply[2],ply[i]); tmp=false;
                                                                           //半平面交
break:
                                                                           #include <cmath>
                                                                           #include <cstdio>
                   if(tmp) return ;
                                                                           #include <algorithm>
                   tmp=true;
                                                                           using namespace std;
                   for (i=3; i<n; i++) {//前四点不共面
                                                                           const int maxn=1505;
                                                                           const double eps=1e-8;
          if(fabs((ply[0]-ply[1])*(ply[1]-ply[2])^(ply[0]-ply[i]))>PR)
                                                                           int n, pn, dq[maxn], top, bot;//数组模拟双端队列
{
                                                                           struct Point{ double x, y;}p[maxn];
                                      swap(ply[3],ply[i]); tmp=false;
                                                                           struct Line{
break:
                                                                                    Point a.b.
                            }
                                                                                    double angle;//极角
                                                                                    Line& operator =(Line I) {
                   if(tmp) return ;
                                                                                              a. x=1. a. x, a. y=1. a. y;
                   fac add:
                                                                                              b. x=1.b. x, b. y=1.b. y;
                   for (i=0; i<4; i++) {//构建初始四面体(4 个点为
                                                                                              angle=1.angle;return *this;
ply[0], ply[1], ply[2], ply[3])
                                                                           } | [maxn];
          add. a=(i+1)\%4, add. b=(i+2)\%4, add. c=(i+3)\%4, add. ok=1;
                                                                           int dblcmp(double k) {//精度函数
                                                                                     if (fabs(k) <eps) return 0;
                             if((ptoplane(ply[i], add))>0)
swap (add. b, add. c);//保证逆时针,即法向量朝外,这样新点才可看到。
                                                                                     return k>0?1:-1;
         vis[add. a] [add. b]=vis[add. b] [add. c]=vis[add. c] [add. a]=trian
                                                                           double multi(Point p0, Point p1, Point p2) {//叉积
glecnt://逆向的有向边保存
                                                                                    return (p1. x-p0. x)*(p2. y-p0. y)-(p1. y-p0. y)*(p2. x-p0. x);
                             tri[trianglecnt++]=add;
                                                                           bool cmp (const Line& I1, const Line& I2) {
                   for (i=4; i<n; i++) {//构建更新凸包
                                                                                     int d=dblcmp(I1.angle-I2.angle);
                             for (j=0; j<trianglecnt; j++) {//对每个点判
                                                                                     if (!d) return dblcmp(multi(I1.a, I2.a, I2.b))<0;
断是否在当前3维凸包内或外(i表示当前点,j表示当前面)
                                                                                    //大于0取半平面的左半, 小于0取右半
                                                                                    return d<0:
          if(tri[j].ok&&(ptoplane(ply[i],tri[j]))>PR){//对当前凸包面
进行判断, 看是否点能否看到这个面
                                                                           void addLine(Line& I, double x1, double y1, double x2, double y2) {
                                                dfs(i, i): break://点
                                                                                     L.a. x=x1: L.a. v=v1:
能看到当前面,更新凸包的面(递归,可能不止更新一个面)。当前点更新完成后
                                                                                     I.b. x=x2; I.b. y=y2;
break 跳出循环
                                                                                     I. ang le=atan2(y2-y1, x2-x1);
                                      }
                                                                           void getIntersect(Line I1, Line I2, Point& p) {
                                                                                    double A1=I1. b. y-I1. a. y;
                   int cnt=trianglecnt;//这些面中有一些 tri[i]. ok=0,
                                                                                    double B1=I1. a. x-I1. b. x;
                                                                                    double C1=(I1. b. x-I1. a. x)*I1. a. y-(I1. b. y-I1. a. y)*I1. a. x;
它们属于开始建立但后来因为在更大凸包内故需删除的,所以下面几行代码的作
用是只保存最外层的凸包
                                                                                    double A2=12. b. y-12. a. y;
                   trianglecnt=0;
                                                                                    double B2=12. a. x-12. b. x;
                   for (i=0; i <cnt; i++)
                                                                                    double C2=(12. b. x-12. a. x)*12. a. y-(12. b. y-12. a. y)*12. a. x;
                   if(tri[i].ok) tri[trianglecnt++]=tri[i];
                                                                                    p. x=(C2*B1-C1*B2)/(A1*B2-A2*B1);
                                                                                    p. y=(C1*A2-C2*A1)/(A1*B2-A2*B1);
         double res() {
                   double _min=1e300;
                                                                           bool judge (Line 10, Line 11, Line 12) {
                   for(int i=0;i<trianglecnt;i++) {</pre>
                                                                                    Point p;getIntersect(I1, I2, p);
                                                                                     return dblcmp(multi(p, 10. a, 10. b))>0;
                             doub le
now=Dis(ply[tri[i].a], ply[tri[i].b], ply[tri[i].c], dd);
                                                                                    //与上面的注释处的大于小于符号相反,大于 0,是 p 在向量
                             if(_min>now) _min=now;
                                                                           10. a->10. b 的左边, 小于 0 是在右边, 当 p 不在半平面 10 内时, 返回 true
                                                                           void HalfPlaneIntersect() {
                   return _min;
                                                                                     int i, j;
}hull;
                                                                                     sort(I, I+n, cmp);//极角排序
                                                                                     for (i=0, j=0; i<n; i++)
int main() {
         while (scanf ("%d", &hull. n) !=E0F) {
                                                                                     if (dblcmp(|[i].angle-|[j].angle)>0) |[++j]=|[i];//排除极角
                   if(hull.n==0) break;
                                                                           相同(从了 | [1] 开始比较)
                   int i, j, q;
                                                                                    n=j+1;//个数
                   for (i=0: i<hull. n: i++)
                                                                                    da[0]=0: da[1]=1:
scanf("%|f%|f%|f", &hull.ply[i].x, &hull.ply[i].y, &hull.ply[i].z);
                                                                                     top=1; bot=0;
                   hull.construct();
                                                                                    for (i=2; i<n; i++) {
                   scanf("%d", &q);
                                                                                              while
                   for (j=0; j < q; j++) {
                                                                           (top>bot&&judge(|[i],|[dq[top]],|[dq[top-1]])) top--;
                             scanf("%|f%|f%|f", &dd. x, &dd. y, &dd. z);
                                                                                              while
                                                                           \label{locality} $$(top>bot\&judge(I[i],I[dq[bot]],I[dq[bot+1]]))$ bot++;
                            printf("%.4If\n", hull.res());
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da[++top]=i:
          while (top>bot\&\&judge(|[dq[bot]],|[dq[top]],|[dq[top-1]]))
top--;
                                                                                //原根
          while (top>bot&&judge(|[dq[top]],|[dq[bot]],|[dq[bot+1]]))
                                                                                LL pow_mod(LL a, LL b, LL c) {
                                                                                          LL ans=1;
bot++;
          dq[++top]=dq[bot];
                                                                                          while (b) {
          for (pn=0, i=bot; i<top; i++, pn++)</pre>
                                                                                                    if (b&1) ans=ans*a%c:
getIntersect(I[dq[i+1]], I[dq[i]], p[pn]);//更新重复利用 p 数组
                                                                                                    a=a*a%c; b>>=1;
double getArea() {
                                                                                          return ans;
          if (pn<3) return 0;
          double area=0;
                                                                                vector<LL> a;
          for (int i=1; i < pn-1; i++) area+=multi(p[0], p[i], p[i+1]); //利
                                                                                bool g_test(LL g, LL p) {
用p数组求面积
                                                                                          for (LL i=0; i < a. size(); ++i)
          return fabs(area)/2;
                                                                                          if (pow_mod(g, (p-1)/a[i], p) == 1) return 0;
}
int main() {
          int t, i;
                                                                                LL primitive_root(LL p) {
          scanf ("%d", &t);
                                                                                          LL tmp=p-1;
          while (t--) {
                                                                                          for (LL i=2; i*i <= tmp; ++i)
                    scanf ("%d", &n);
                                                                                          if (!(tmp%i)) {
                    for (i=0; i<n; i++) scanf ("%|f%|f", &p[i].x, &p[i].y);
                                                                                                    a.push_back(i);
                    for (i=0; i<n-1; i++)
                                                                                                    while (!(tmp\%i)) tmp/=i;
addLine(I[i],p[i].x,p[i].y,p[i+1].x,p[i+1].y);
                    addLine(I[i],p[i].x,p[i].y,p[0].x,p[0].y);
                                                                                          if (tmp>1) a.push_back(tmp);
                    HalfPlaneIntersect();
                                                                                          for (LL g=1; g++) if (g_test(g, p)) return g;
                    printf ("%. 2If\n", getArea());
                                                                               }
          return 0;
                                                                                //Miller_Rabin
}
                                                                                LL add(LL a, LL b, LL c) {
PKU 1279 题为顺时针方向
                                                                                          LL ans=0;
                                                                                          while (b) {
                                                                                                    if (b&1) {
                                                                                                              ans+=a;
                粉论
                                                                                                              if (ans>=c) ans-=c;
//exGCD
int extend_gcd(int a, int b, int &x, int &y) {
                                                                                                    a<<=1; b>>=1;
          if (b==0) { x=1; y=0; return a; }
                                                                                                    if (a>=c) a-=c;
                                                                                          }
          else {
                    int tmp=extend_gcd(b, a%b, y, x);
                                                                                          return ans;
                    y=x*(a/b); return tmp;
          }
                                                                                LL power (LL a, LL b, LL c) {
                                                                                          LL ans=1:
}
                                                                                          while (b) {
                                                                                                    if (b&1) ans=add(ans, a, c);
AX+BY=GCD (A. B)
B(Y+A/B*X)+A\%B*X=GCD(B, A\%B)
                                                                                                    a=add(a, a, c); b>>=1;
*/
                                                                                          }
                                                                                          return ans;
//中国剩余定理
                                                                                bool Miller_Rabin_test(LL a, LL b, LL n) {
int CRT(int a[], int m[], int n) {
                                                                                          if (n==2) return true;
          int i, M=1, ans=0;
                                                                                          else if (n==a||!(n&1)) return false;
          for (i=1; i \le n; i++) M*=m[i];
                                                                                          while (!(b&1)) b>>=1;
          for (i=1; i \le n; i++) {
                                                                                          LL tmp=power(a, b, n);
                    int x, y, Mi=M/m[i];
                                                                                          while (b!=n-1\&\&tmp!=n-1\&\&tmp!=1) \{ tmp=add(tmp, tmp, n); \}
                                                                                b<<=1; }
                    extend_Euclid(Mi, m[i], x, y);
                    ans=(ans+Mi*x*a[i])%M;
                                                                                          return (tmp==n-1) | | ((b&1) && (tmp==1));
          if (ans<0) ans+=M;
                                                                                bool isprime(LL n) {
          return ans;
                                                                                          int a[]={2, 3, 7, 61, 24251}; //10<sup>16</sup>
                                                                                          if (n==46856248255981LL) return false;
}
                                                                                          for (int i=0; i<5; ++i) {
Problem:
                                                                                                    if (n==a[i]) return true;
x=ai(%mi)[0<=i<n]
                                                                                                    if (!Miller_Rabin_test(a[i], n-1, n)) return false;
m0, m1, m2, ..., mn-1 两两互质
                                                                                          }
已知 mi, ai, 求 x
                                                                                          return true;
Solution:
                                                                               }
Mi=pi(mj) [i!=j]
gcd(Mi, mi)=1 --> Mipi+miqi=1
Let ei=Mipi, then ei=\{0 (\%mj) [i!=j], 1 (\%mj) [i==j]\}
                                                                                //exBSGS
ans0=(e0a0+e1a1+e2a2+...+en-1an-1)%pi(mi)
                                                                                #include <cmath>
```

```
#include <cstdio>
                                                                                             memset(flag, 0, sizeof(flag));
#include <cstring>
                                                                                             LL tmp=BabyStep_GiantStep(A, B, C);
#include <iostream>
                                                                                             if (tmp==-1) puts("No Solution\n"); else
#include <algorithm>
                                                                          printf("%|Id\n", tmp);
using namespace std:
                                                                                             scanf ("%||d%||d%||d", &A, &C, &B):
typedef long long LL;
                                                                                   }
const int Mo 131071;
                                                                                   return 0;
                                                                         }
II ton:
bool flag[Mo*2];
                                                                          扩展 BSGS
struct HashNode{ LL data, id, next; }hash[Mo*2];
void Insert(LL a, LL b) {
                                                                          简单的来说,就是加了一个把不互质的数通过去除公因数变为互质的,再进行
         LL k=b&Mo;
                                                                          BSGS
         if (flag[k]==false) {
                                                                          具体的就是
                   flag[k] = true; hash[k].next = -1; hash[k].id = a;
                                                                          考虑 a 与 p 不互质的情况: 1. 对于 a x=b mod p, 我们可以考虑从 x 个 a 中拿
hash[k].data=b;
                                                                          出c个a与b和p消去公因子,直到a和p,互质为止。
                                                                          2. 一旦互质了, 那么方程就是 v*a^(x-c)=b' (mod p'), v 是拿出 c 个 a 消去
                   return;
                                                                          公因子后剩下的东西, b', p' 是消去公因子的 b, p。
         }
         while (hash[k].next!=-1) {
                                                                          这个时候还要求 v 的逆元, 方程变为 a^(x-c)=b'*v^(-1) \pmod{p'}。
                   if (hash[k].data==b) return;
                                                                          此时就可以用 baby-step-giant-step 做了,答案为 BSGS 的答案+c。
                  k=hash[k].next;
                                                                          注意:有可能 c 会大于 x, 所以必须约去一次就特判一次方程两边(就是 v 和 b')
                                                                          是不是相等了。如果两边相等那么直接返回 c。
         if (hash[k].data==b) return;
         hash[k].next=++top; hash[top].next=-1; hash[top].id=a;
hash[top].data=b;
LL Find(LL b) {
                                                                          const int Mo=1000000007, rev=500000004;
         LL k=b&Mo:
                                                                          void FWT(int a[], int n) {
         if(flag[k]==false) return -1;
                                                                                   for (int d=1; d<n; d<<=1)
         while (k!=-1) {
                                                                                             for (int m=d <<1, i=0; i < n; i+=m)
                   if (hash[k].data==b) return hash[k].id;
                                                                                                      for (int j=0; j< d; j++) {
                  k=hash[k].next:
                                                                                                                int x=a[i+j], y=a[i+j+d];
         }
         return −1;
                                                                                   a[i+j]=(x+y) \text{ mod}, a[i+j+d]=(x-y+mod) \text{ mod};
LL gcd(LL a, LL b) { return b?gcd(b, a%b):a; }
                                                                                   //xor:a[i+j]=x+y, a[i+j+d]=(x-y+mod) mod;
LL ext_gcd (LL a, LL b, LL& x, LL& y ) {
                                                                                   //and:a[i+j]=x+y, a[i+j+d]=a[i+j+d];
         LL t.ret:
         if (b==0) { x=1, y=0; return a; }
         ret=ext\_gcd(b,a\%b,x,y)\;;\;\;t=x\;;\;\;x=y\;;\;\;y=t-a/b*y\;;
                                                                                   //or:a[i+j]=a[i+j],a[i+j+d]=x+y;
         return ret:
                                                                                                     }
LL mod_exp(LL a, LL b, LL n) {
                                                                          void UFWT(int a[], int n) {
                                                                                   for (int d=1; d \le n; d \le = 1)
         LL ret=1; a=a%n;
         while (b>=1) {
                                                                                             for (int i=0, m=d<<1: i<n: i+=m)
                   if (b&1) ret=ret*a%n;
                                                                                                      for (int j=0; j<d; j++) {
                  a=a*a%n; b>>=1;
                                                                                                                int x=a[i+j], y=a[i+j+d];
         }
         return ret;
                                                                                   a[i+j]=1LL*(x+y)*rev\%mod, a[i+j+d]=(1LL*(x-y)*rev\%mod+mod)\%m
                                                                          od;
LL BabyStep_GiantStep(LL A, LL B, LL C) {
         top=Mo; B%=C;
                                                                                   //xor:a[i+j]=(x+y)/2, a[i+j+d]=(x-y)/2;
         LL i, tmp=1;
         for (i=0; i \le 100; tmp=tmp*A%C, i++) if (tmp==B%C) return i;
                                                                                   //and:a[i+j]=x-y, a[i+j+d]=a[i+j+d];
         LL D=1, cnt=0:
         while ((tmp=gcd(A, C))!=1) {
                                                                                   //or:a[i+j]=a[i+j],a[i+j+d]=y-x;
                   if (B\%tmp) return -1;
                   C/=tmp; B/=tmp; D=D*A/tmp%C; cnt++;
                                                                          }
                                                                          void solve(int a[], int b[], int n) {
         LL M=(LL) \operatorname{sqrt}(C+0.0);
                                                                                   FWT(a,n); FWT(b,n);
         for (tmp=1, i=0; i<=M; tmp=tmp*A%C, i++) Insert(i, tmp);</pre>
                                                                                   for(int i=0; i<n; i++) a[i]=1LL*a[i]*b[i]%mod;
         LL x, y, K=mod_exp(A, M, C);
         for (i=0; i \le M; i++) {
                                                                          //当指异或非运算、与非运算、或非运算时, 我们可以将 直接用异或运算、与运
                  ext\_gcd(D, C, x, y); //D*X=1 (mod C)
                                                                          算、或运算的方法求出来,然后将互反的两位交换即可
                   tmp=((B*x)%C+C)%C;
                   if((y=Find(tmp))!=-1) return i*M+y+cnt;
                  D=D*K%C:
         }
         return −1;
}
int main() {
         LL A. B. C:
         scanf("%||d%||d%||d", &A, &C, &B);
```

while (!(A==0&&B==0&&C==0)) {