

# Computer Network Security

ECE 4112/6612  
CS 4262/6262

Prof. Frank Li

\* Welcome to CityPower Grid Rerouting \*  
Authorised users only!  
New users MUST notify Sys/ops.  
login:

```
80/tcp      open   http          host<2_nc
81/tcp      open
100/tcp     open
113/tcp     open   nmap -v -SS -O 10.2.2.2
139/tcp     open
143/tcp     open
145/tcp     open
1539/tcp    open
22/tcp      open   ssh           Service
587/tcp     open
687/tcp     open
2432/tcp    open
50000/tcp   open
# nmap run completed -- 1 IP address (1 host up) scanned
# SSHNUKE 10.2.2.2 -rootpw:"210H0101" -successful.
# Connecting to 10.2.2.2:ssh ... successful.
# Attempting to exploit SSHv1 CRC32 IP Resetting root password to "210H0101"; successful.
# System open: Access Level <9>
# ssh 10.2.2.2 -l root
# root@10.2.2.2's password: [REDACTED]
```

EDITU1 SSHNUKE  
rcr ebx. 1  
bsr ecx. ecx  
shrd ebx. edi. CL  
shrd eax. edx. CL  
[mobile]  
[mobile]

RTF CONTROL  
ACCESS GRANTED

# Logistics

Tue, Oct 10	No Class (Fall Break)
Thu, Oct 12	Web security Part 1: Web attacks and defenses
Tue, Oct 17	Web security Part 2: Web attacks and defenses
Thu, Oct 19	Web security Part 3: Web attacks and defenses
Tue, Oct 24*	Quiz 2
Thu, Oct 26*	Authentication

HW2 due Tuesday, Oct 17 midnight

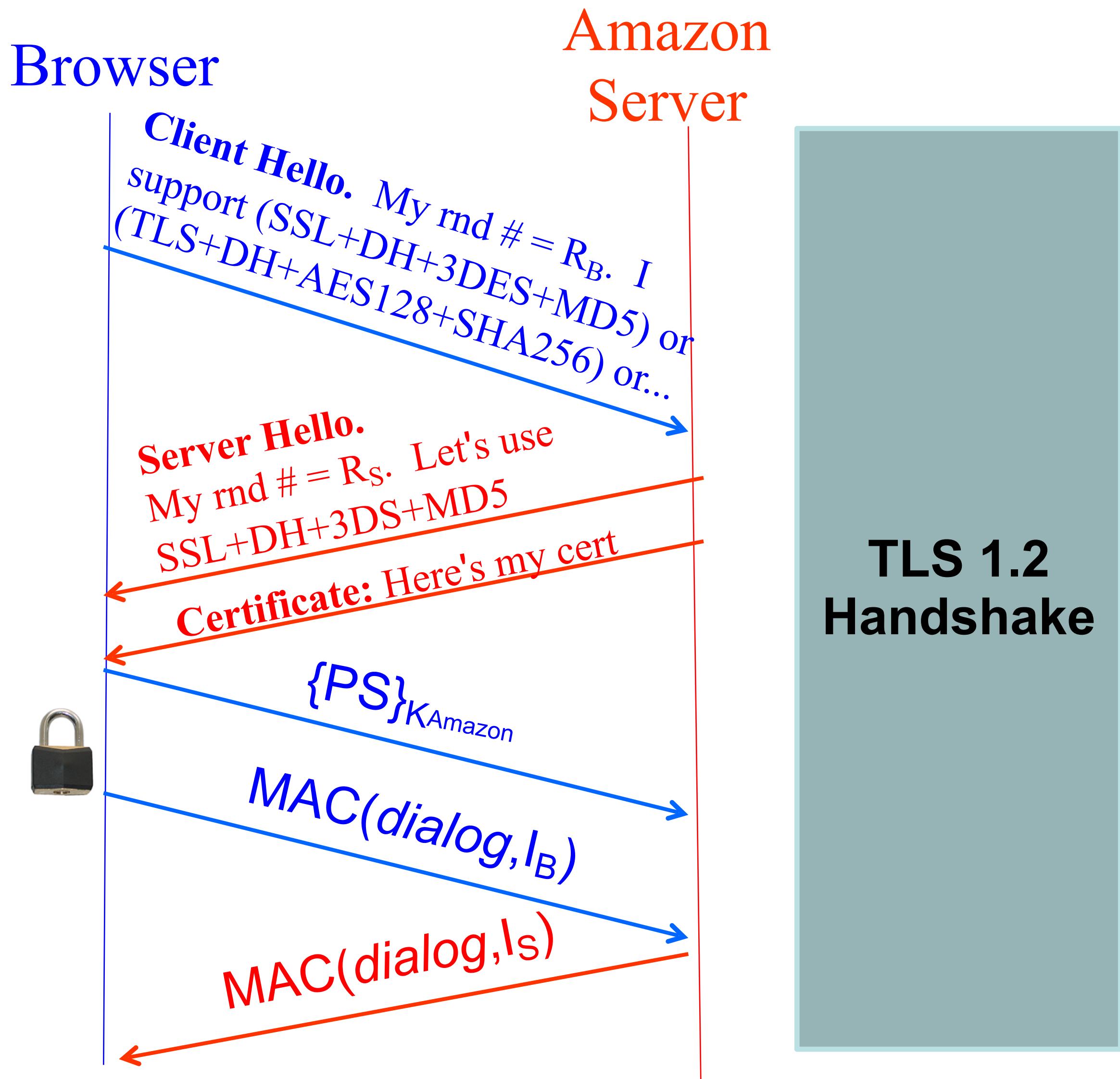
Quiz 2 in 1.5 weeks (Tue, Oct 24) -> No web security, but includes IP + BGP

Project presentation + report assignment to be released (including presentation scheduling)

*Note: I'll be away the week of Oct 24-26. TAs will proctor Quiz 2, guest lecture for Thu lecture.*

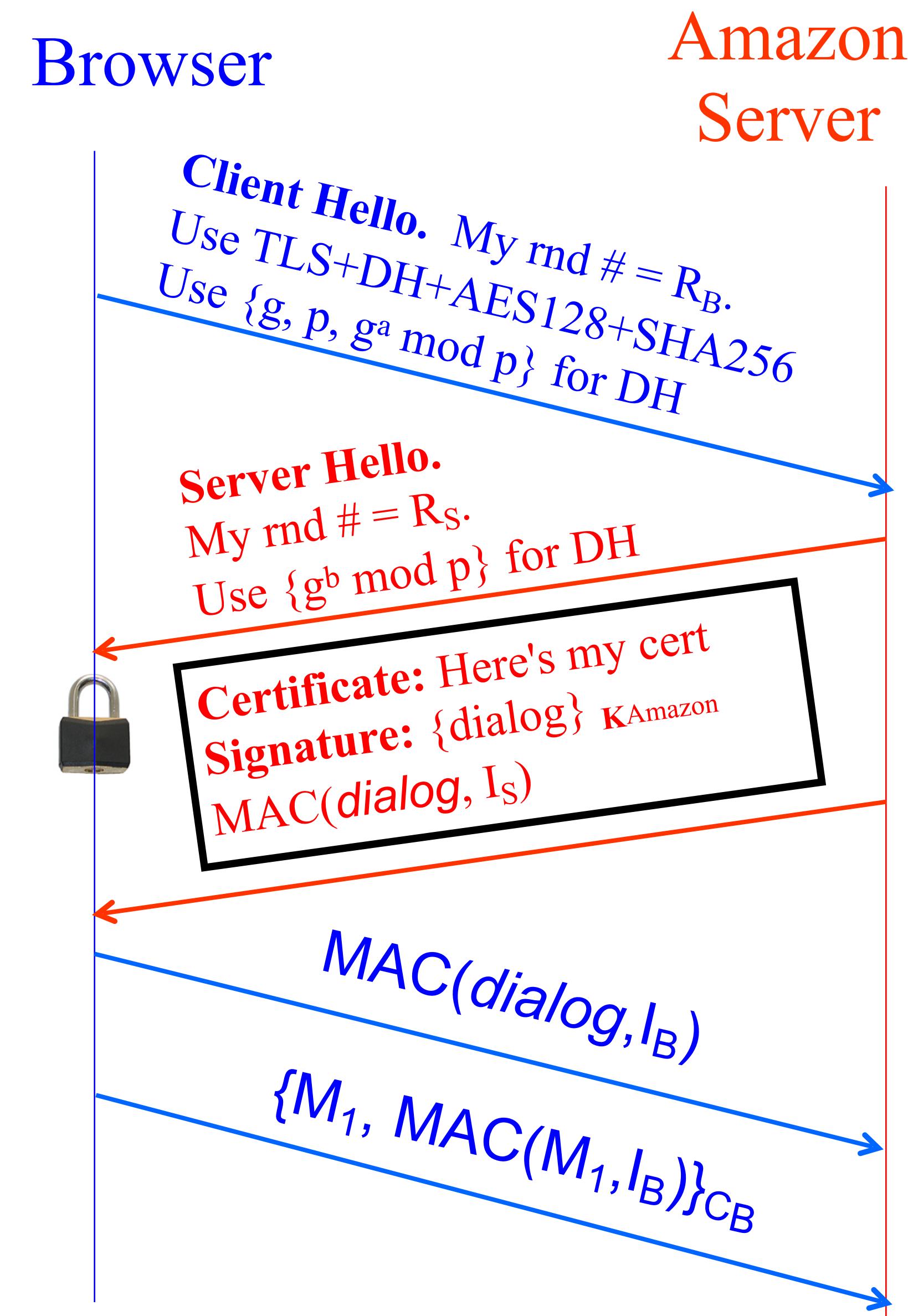
# Finishing up TLS

# TLS 1.2



**TLS 1.2  
Handshake**

# TLS 1.3



**TLS 1.3 Handshake**

# Validating Amazon's Identity

- Browser compares domain *name* in cert w/ URL
- Browser accesses separate cert belonging to **issuer**
  - Might be **hardwired into the browser** - **trusted root CAs!**
  - There could be a *chain* of these ...
- Browser applies issuer's public key to verify certificate's signature
  - Validates signature on its own **SHA-256** hash of Amazon's cert
- Assuming signature validates, now have high confidence it's indeed Amazon ...
  - ***assuming signatory is trustworthy***

# End-to-End $\Rightarrow$ Powerful Protections

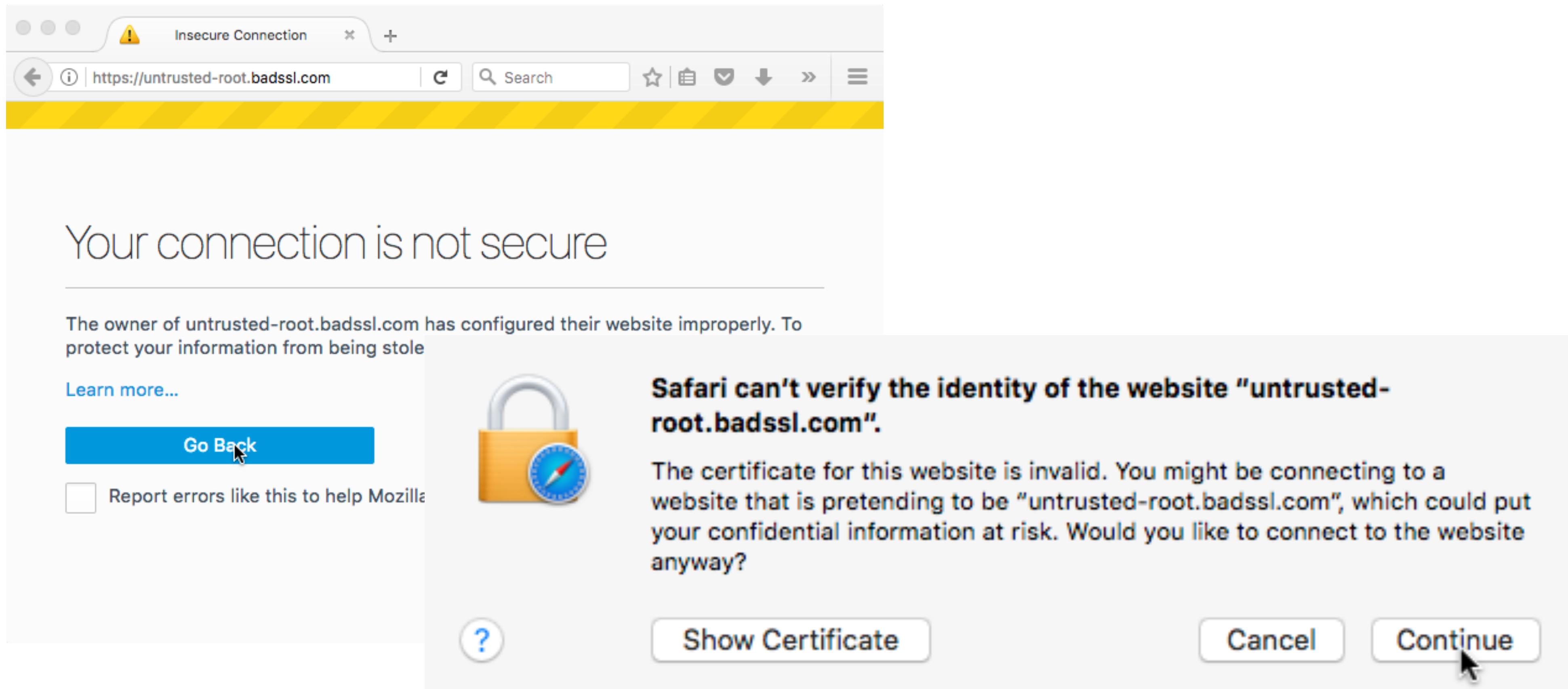
- Attacker runs a sniffer to capture our WiFi session?
  - (maybe by buying a cup of coffee to get the password)
  - **But:** encrypted application data is unreadable
    - No problem!
- DNS cache poisoning?
  - Client goes to wrong server
  - **But:** detects impersonation since attacker lacks valid cert
    - No problem!
- Attacker hijacks our connection, injects new traffic
  - **But:** data receiver rejects it due to failed integrity check
    - No problem!

# Powerful Protections, con't

- DHCP spoofing?
  - Client goes to wrong server
  - **But:** they can't read; we detect impersonation
    - No problem!
- Attacker manipulates BGP routing to run us by an eavesdropper or take us to the wrong server?
  - **But:** they can't read; we detect impersonation
    - No problem!
- Attacker slips in as a Man In The Middle?
  - **But:** they can't read, they can't inject
  - They can't even replay previous encrypted traffic
  - **No problem!**

# Validating Amazon's Identity, con't

- What if browser can't find a cert for the issuer or the cert doesn't pass?



# Validating Amazon's Identity, con't

- What if browser can't find a cert for the issuer or the cert doesn't pass?
  - Then warns the user that site has not been verified
    - Note, can still proceed, just **without authentication**
  - Q: Which end-to-end security properties do we lose if we incorrectly trust that the site is whom we think?
  - A: **All of them!**
    - Goodbye confidentiality, integrity, authentication
    - Attacker can read everything, modify, impersonate

# How do we know to use TLS?

- How do we know if example.com is over HTTP or HTTPS?
  - Method 1: Could try HTTPS, and if we don't get any response or an error, switch to HTTP.
  - Method 2: Could try HTTP, and if HTTPS is available, the server's HTTP response can redirect us to HTTPS

# How do we know to use TLS?

- What if there's a network MITM attacker??
  - W/ Method 1: MITM can block/drop HTTPS attempt, causing us to switch to HTTP
  - W/ Method 2: Block the HTTP redirection response so we don't know HTTPS is available
- In both cases, then MITM attacker can pretend to be the HTTP server. We will send our data to the attacker over HTTP and the attacker can relay data to the real HTTPS website.
- **SSL-Stripping attack**

# How do we know to use TLS?

- Solutions?
  - **HTTP Strict Transport Security (HSTS)**: Server tells client that it always uses HTTPS, so client knows to use HTTPS moving forward.
    - Done using an HTTP header

```
Strict-Transport-Security: max-age=31536000;
```
    - Assumes secure first connection though...
  - **(Proposed) HTTPS DNS Resource Records**: Add a new DNS RR indicating HTTPS is available, which clients can query for.
    - Unless DNSSEC used, may not help in MITM case

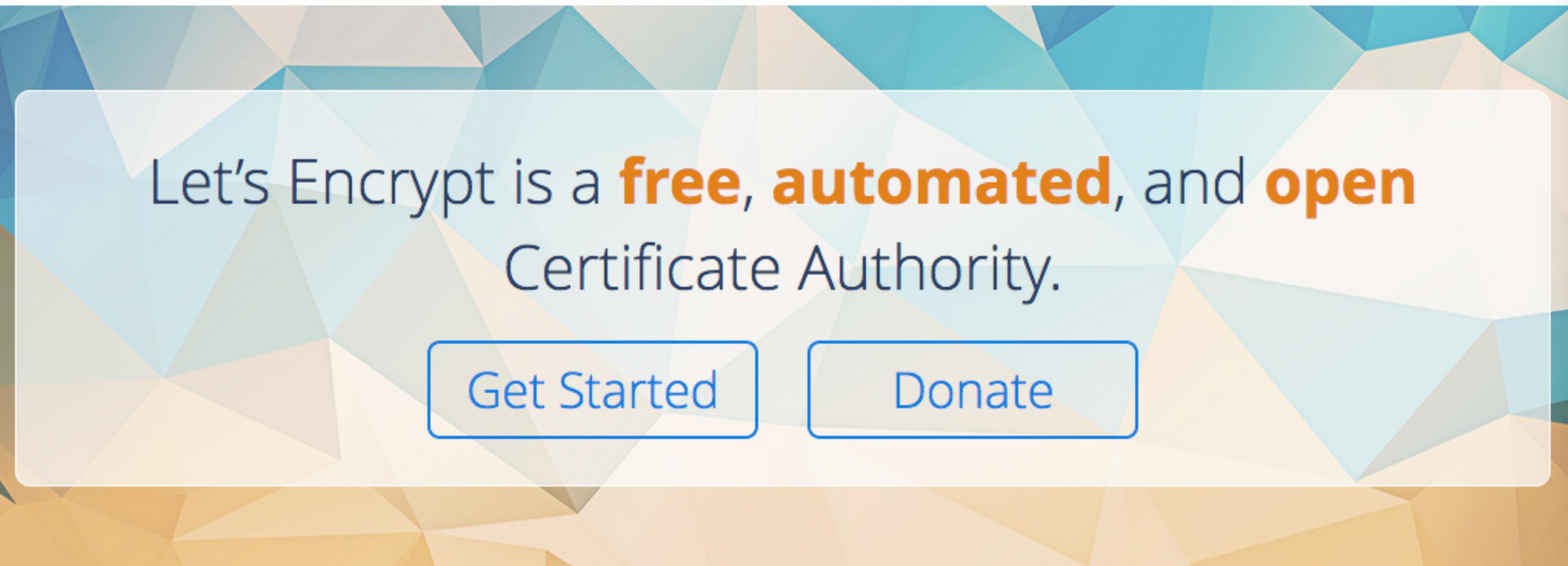
# TLS Limitations

- Properly used, TLS provides powerful end-to-end protections
- So why not use it for *everything*??
- Issues:
  - Cost of public-key crypto
    - Takes non-trivial CPU processing (but today a minor issue)
    - Note: *symmetric* key crypto on modern hardware is non-issue
  - Hassle of buying/maintaining certs (fairly minor)



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Mar 23, 2017

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We're pleased to announce that OVH has renewed their support for Let's Encrypt as a Platinum sponsor for the next three years.

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# TLS Limitations

- Properly used, TLS provides powerful end-to-end protections
- So why not use it for *everything*??
- Issues:
  - Cost of public-key crypto
    - Takes non-trivial CPU processing (but today a minor issue)
    - Note: *symmetric* key crypto on modern hardware is non-issue
  - Hassle of buying/maintaining certs (fairly minor)
  - **Latency**: extra round trips ⇒ pages take longer to load

# TLS Limitations, con't

- Problems that TLS does **not** take care of ?
- TCP-level attacks
  - RST injection
    - (but does protect against data injection!)
- Denial of Service Attacks
- Application-level vulnerabilities:
  - SQL injection / XSS / server-side coding/logic flaws
  - Browser coding/logic flaws
  - User flaws
    - Weak passwords
    - Phishing
  - HTTP server vulnerabilities

# **Web Security**

## **(Part 1)**

# What is the Web?

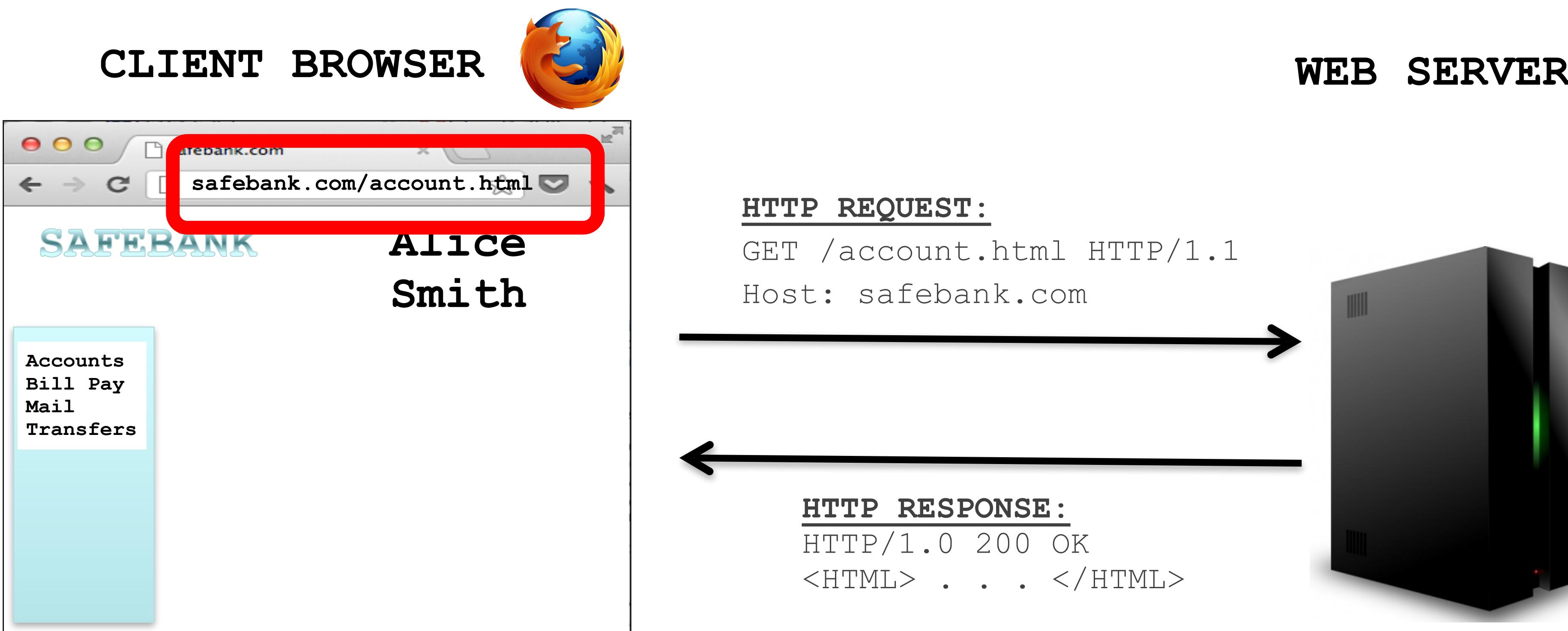
A platform for deploying applications and sharing information,  
*portably and securely (?)*



# HTTP

(Hypertext Transfer Protocol)

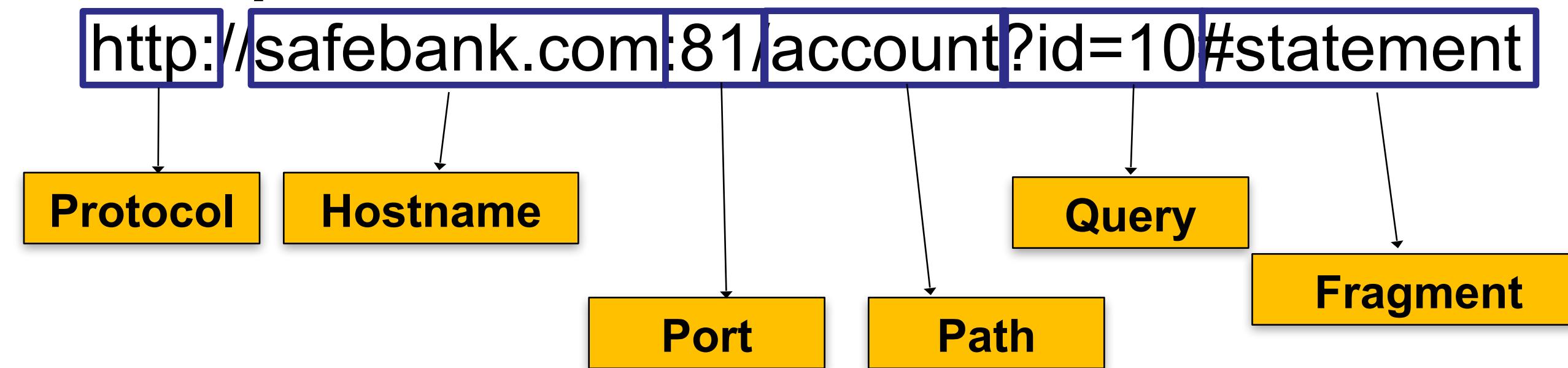
A common data communication protocol on the web



# URLs

Global identifiers of network-retrievable resources

**Example:**



# HTTP



# HTTP Request

GET: no side effect (supposedly)

POST: possible side effect, includes additional data

Method	Path	HTTP version	Headers
GET	/index.html	HTTP/1.1	Accept: image/gif, image/x-bitmap, image/jpeg, */* Accept-Language: en Connection: Keep-Alive User-Agent: Chrome/21.0.1180.75 (Macintosh; Intel Mac OS X 10_7_4) Host: safebank.com Referer: http://www.google.com?q=dingbats

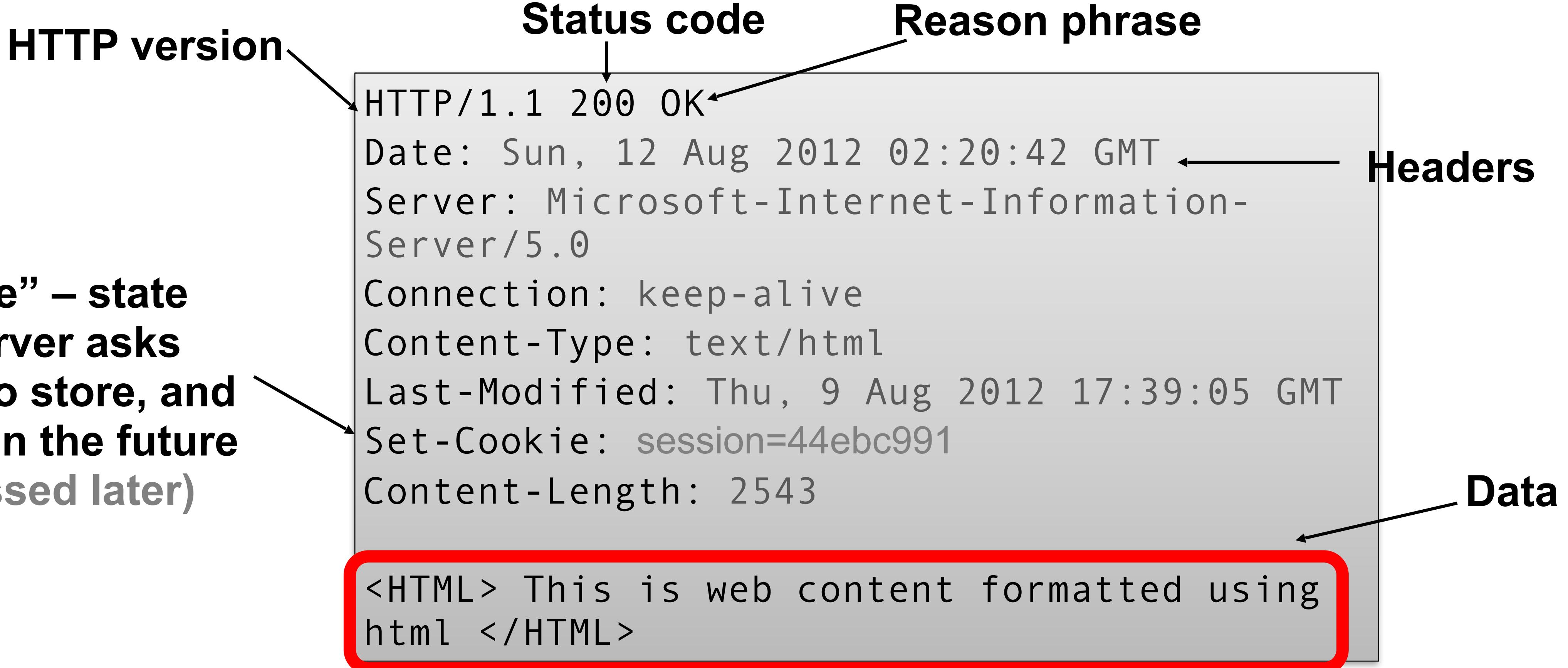
Blank line

Data – none for GET

# HTTP

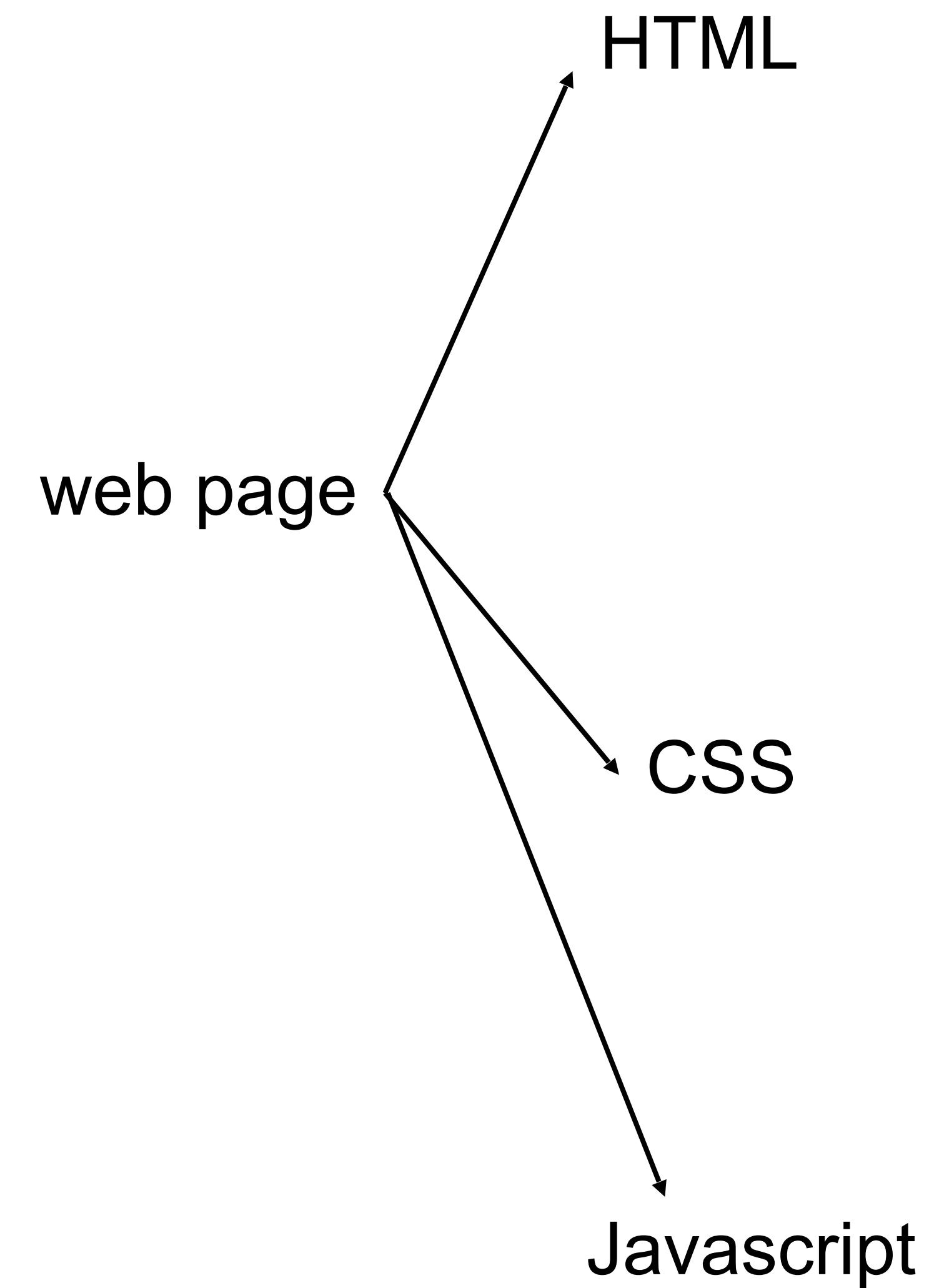


# HTTP Response



Can be a webpage, image, audio, executable ...

# Web page



# HTML

A language to create structured documents  
One can embed images, objects, or create interactive forms

## index.html

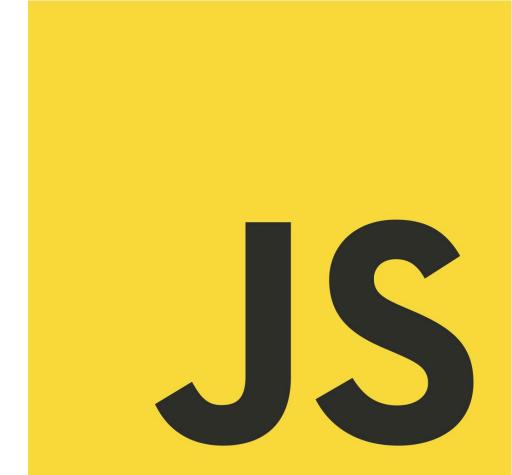
```
<html>
  <body>
    <div>
      foo
      <a href="http://google.com">Go to Google!</a>
    </div>
    <form>
      <input type="text" />
      <input type="radio" />
      <input type="checkbox" />
    </form>
  </body>
</html>
```

# CSS (Cascading Style Sheets)

Language used for describing the presentation of a document

## **index.css**

```
p.serif {  
    font-family: "Times New Roman", Times, serif;  
}  
p.sansserif {  
    font-family: Arial, Helvetica, sans-serif;  
}
```



JS

# Javascript

Programming language used to manipulate web pages. It is a high-level, untyped and interpreted language with support for objects.

Supported by all web browsers

```
<script>
function myFunction()
{
    document.getElementById("demo").innerHTML = "Text
changed.";
}
</script>
```

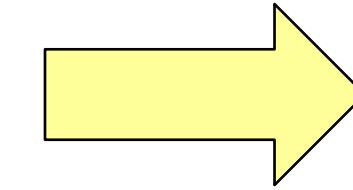
**Very powerful!**

# DOM (Document Object Model)

Cross-platform model for representing and interacting with objects in HTML

## HTML

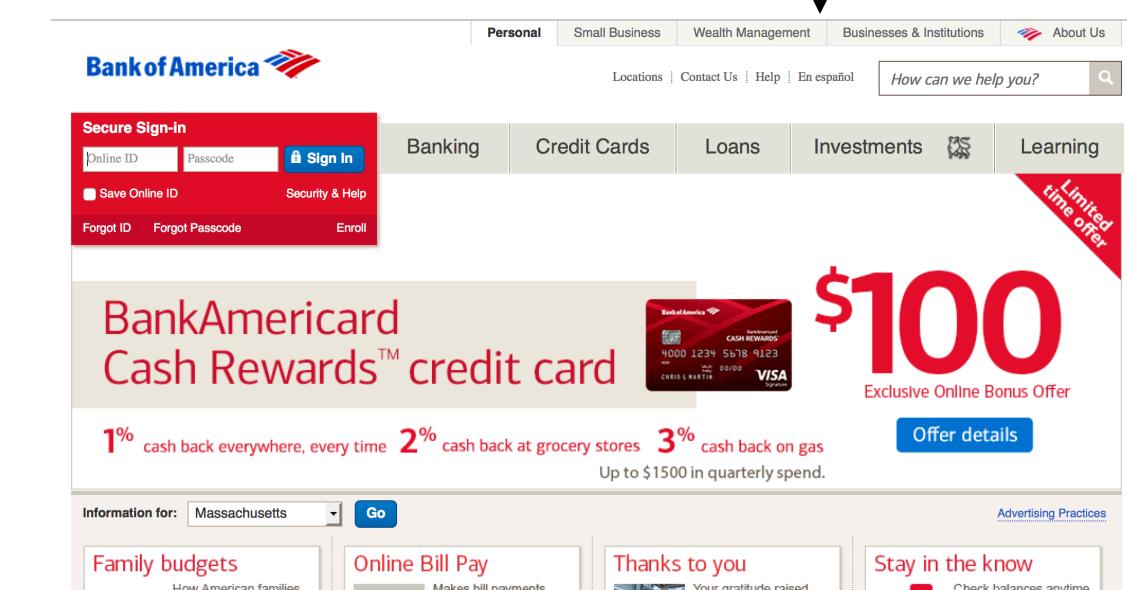
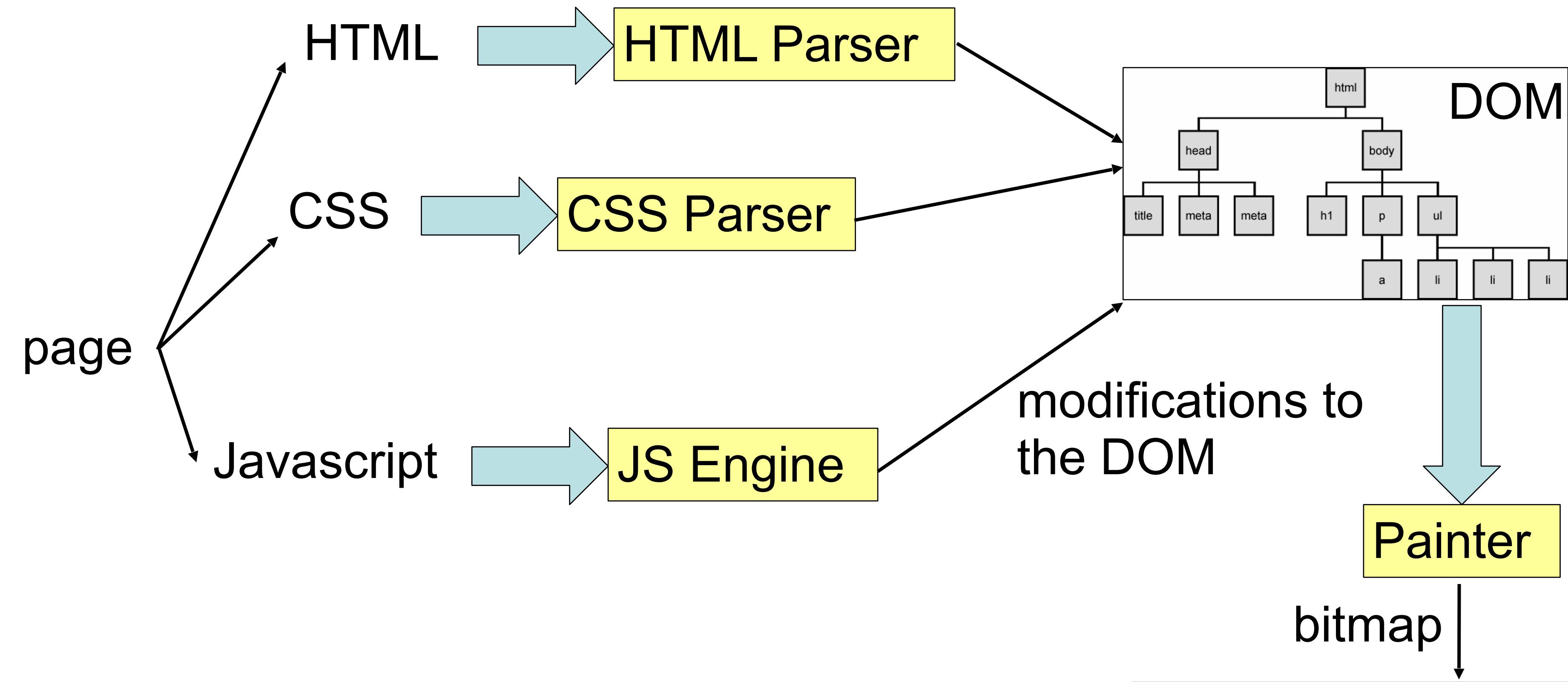
```
<html>
  <body>
    <div>
      foo
    </div>
    <form>
      <input type="text" />
      <input type="radio" />
      <input type="checkbox" /
    >
    </form>
  </body>
</html>
```



## DOM Tree

```
| -> Document
  | -> Element (<html>)
  | -> Element (<body>)
  | -> Element (<div>)
  | -> text node
  | -> Form
    | -> Text-box
    | -> Radio Button
    | -> Check Box
```

# Page rendering



# What can you do with Javascript?

Almost anything you want to the DOM!

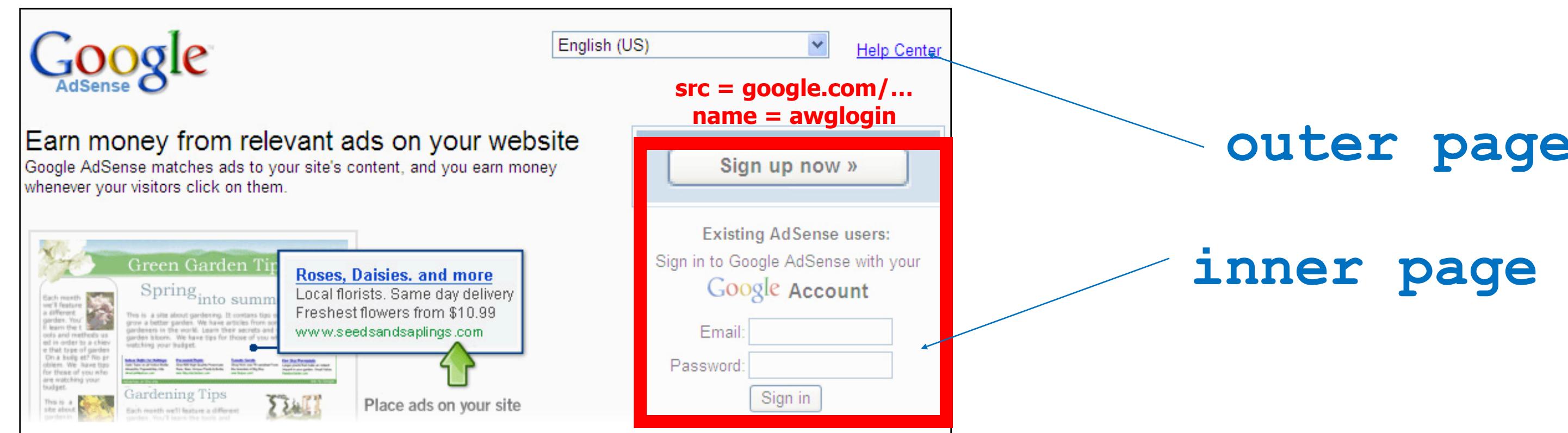
A JS script embedded on a page can modify  
in almost arbitrary ways the DOM of the page.

The same happens if an attacker manages to  
get you load a script into your page.

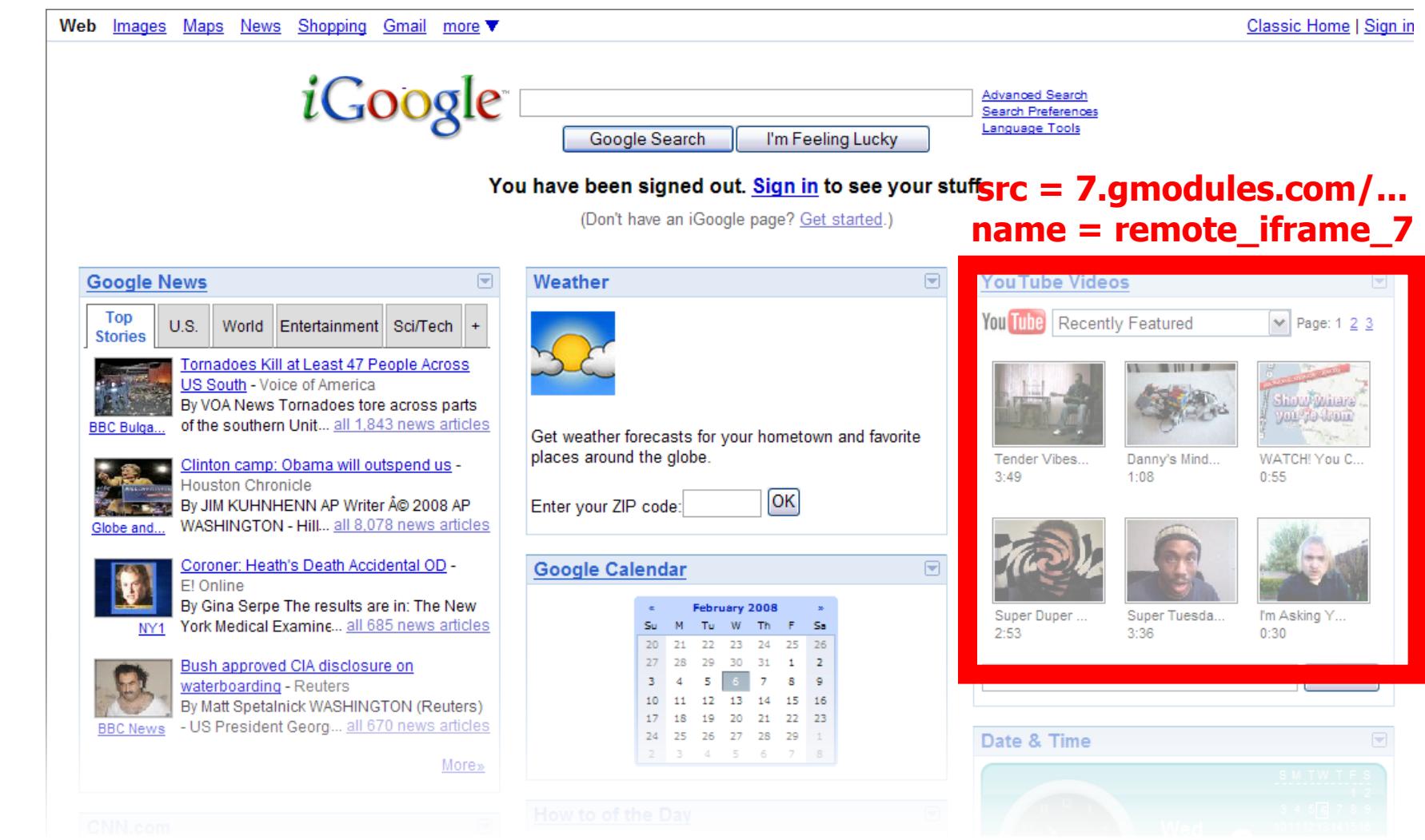
# Frames

- Enable embedding a page within a page

```
<iframe src="URL"></iframe>
```



# Frames



- Modularity
  - Brings together content from multiple sources
  - Client-side aggregation
- Delegation
  - Frame can draw only inside its own rectangle

# Frames

- Outer page can specify only sizing and placement of the frame in the outer page
- Frame isolation: Outer page cannot access contents of inner page; inner page cannot access contents of outer page (unless both pages are from the same site\*)

## Outrageous Chocolate Chip Cookies

★★★★★ 1676 reviews

Made 321 times

Recipe by: Joan

"A great combination of chocolate chips, oatmeal, and peanut butter."



# HTTP cookies

Print

## Ingredients

25 m 18 servings 207 cals

+ 1/2 cup butter

+ 1 cup all-purpose flour

On Sale

On

What's on sale near you.

+ 1/2 cup white sugar

+ 1 teaspoon baking soda

Target  
**TARGET**  
1057 Eastshore Hwy  
ALBANY, CA 94710  
Sponsored

Market Pantry Granulated  
Sugar - 4lbs



\$2.59

[SEE DETAILS](#)

ADVERTISEMENT

+ 1/4 teaspoon salt

+ 1/2 cup rolled oats

+ 1 cup semisweet chocolate chips

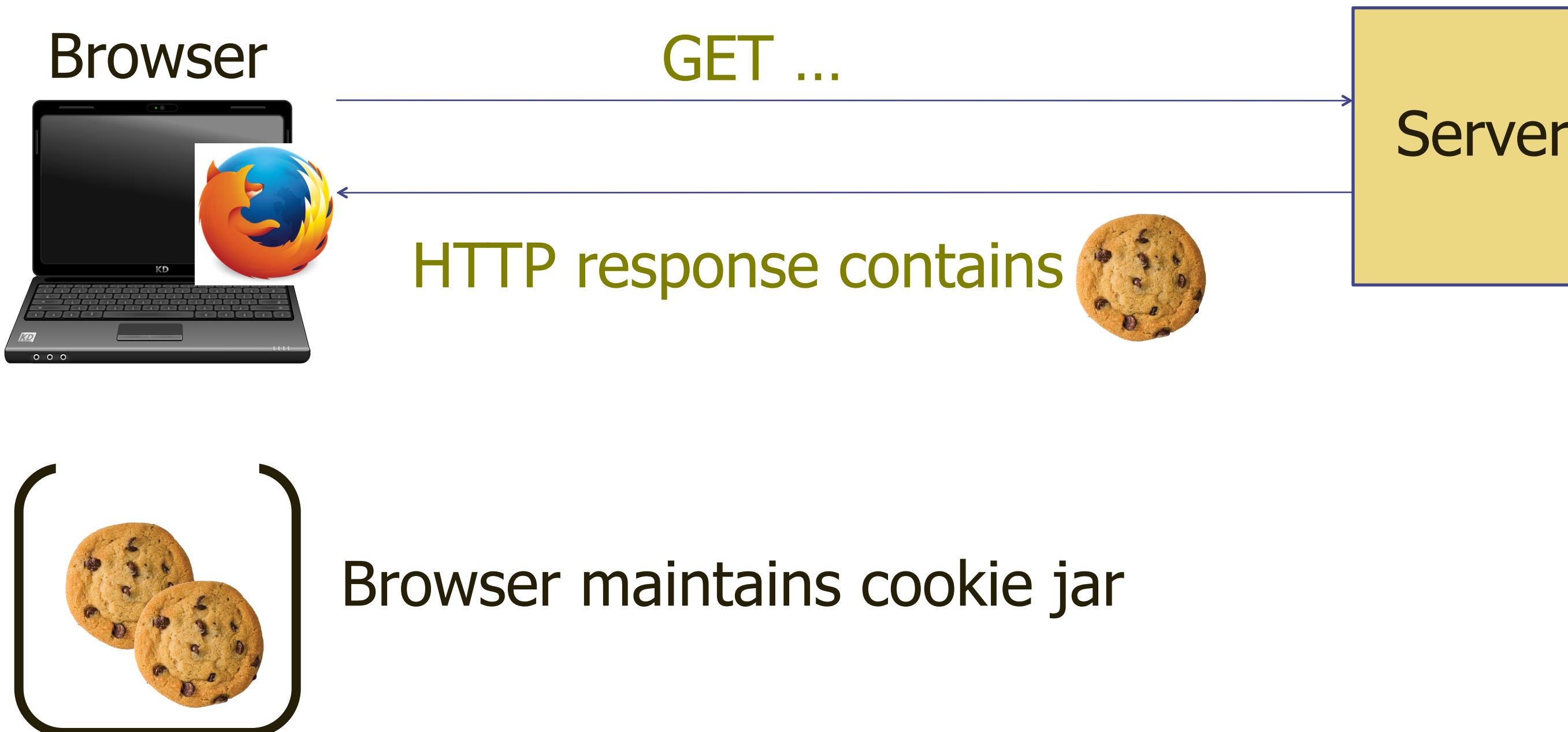
+ 1/3 cup packed brown sugar

May we suggest

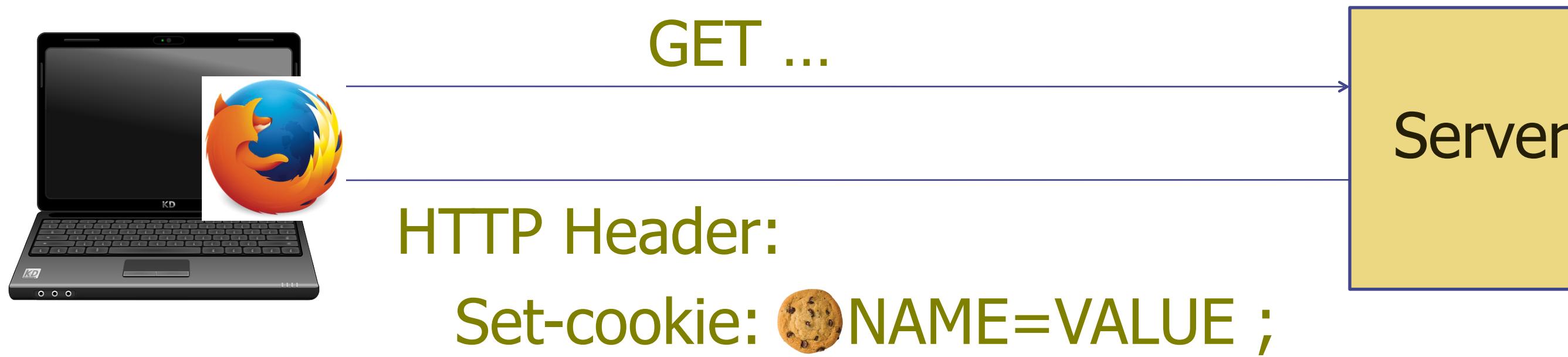
These nearby stores have  
ingredients on sale!

# Cookies

- ◆ A way of maintaining state

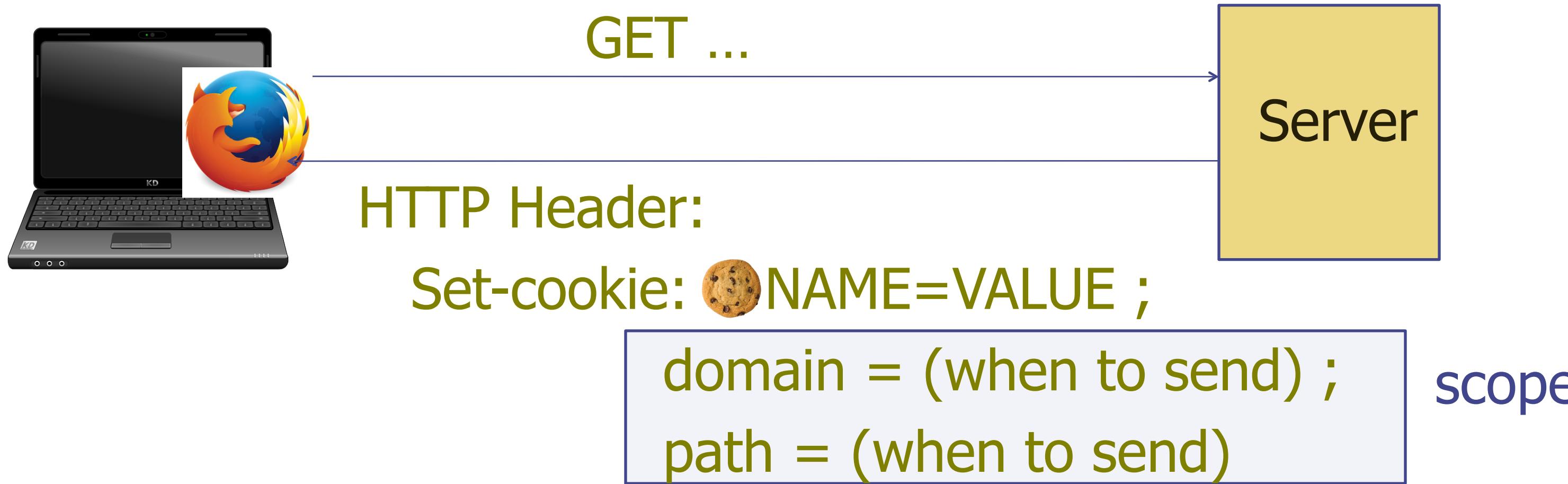


# Setting/deleting cookies by server



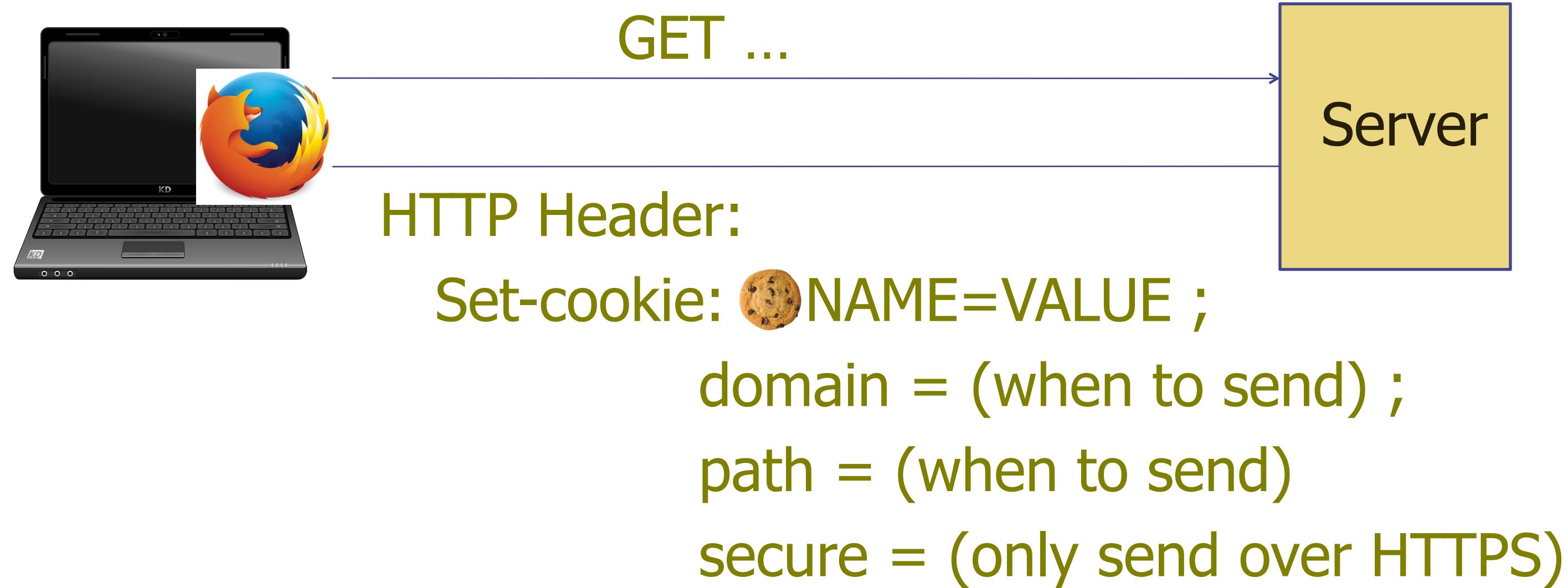
- ◆ The first time a browser connects to a particular web server, it has no cookies for that web server
- ◆ When the web server responds, it includes a **Set-Cookie:** header that defines a cookie
- ◆ Each cookie is just a name-value pair

# Cookie scope



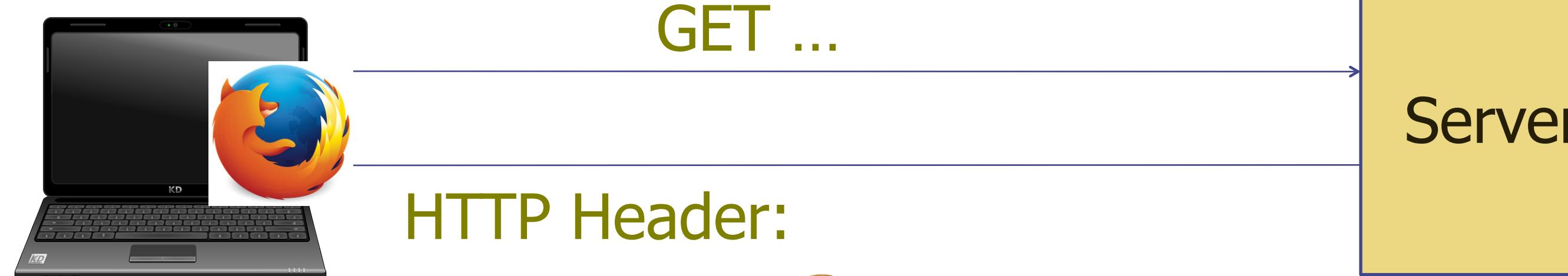
- ◆ When the browser connects to the same server later, it includes a **Cookie:** header containing the name and value, which the server can use to connect related requests.
- ◆ Domain and path inform the browser about which sites to send this cookie to

# Cookie scope



- Secure: sent over HTTPS only
  - HTTPS provides secure communication (privacy, authentication, integrity)

# Cookie scope



- Expires is expiration date
- HttpOnly: cookie cannot be accessed by Javascript, but only sent by browser

## Client side read/write: document.cookie

- ◆ Setting a cookie in Javascript:

`document.cookie = "name=value; expires=...;"`

- ◆ Reading a cookie:

`alert(document.cookie)`

prints string containing all cookies available for  
document (based on [protocol], domain, path)

- ◆ Deleting a cookie:

`document.cookie = "name=; expires= Thu, 01-Jan-70"`

`document.cookie` often used to customize page in Javascript

# Cookie scope

- ◆ Scope of cookie might not be the same as the URL-host name of the web server setting it

Rules on:

1. What scopes a URL-hostname (server) is allowed to set
2. When a cookie is sent to a URL

# What scope a server may set for a cookie

The browser checks if the server may set the cookie, and if not, it will not accept the cookie.

domain: any domain-suffix of URL-hostname, except TLD

[top-level domains,  
e.g. '.com']

example: host = "login.site.com"

allowed domains

**login.site.com**  
**.site.com**

disallowed domains

**user.site.com**  
**othersite.com**  
**.com**

- ⇒ Cookies are sent to domain + subdomains
- ⇒ **login.site.com** can set cookies for all of **.site.com**  
but not for another site or TLD

Problematic for sites like **.gatech.edu**

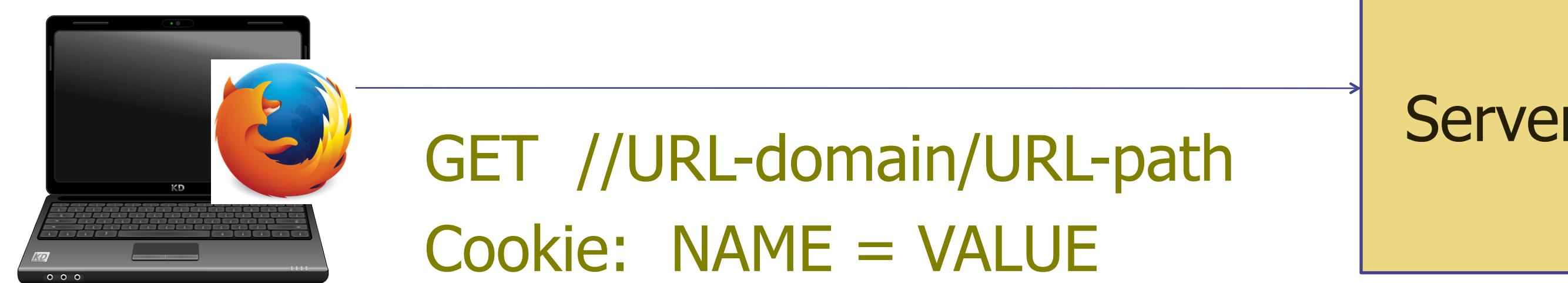
path: can be set to anything

# Examples

Web server at `foo.example.com` wants to set cookie with domain:

Domain	<b>Whether it will be set, and if so, where it will be sent to</b>
[value omitted]	<code>foo.example.com</code>
<code>bar.foo.example.com</code>	Cookie not set; more specific than original domain
<code>foo.example.com</code>	<code>*. foo.example.com</code>
<code>baz.example.com</code>	Cookie not set; domain mismatch
<code>example.com</code>	<code>*.example.com</code>
<code>ample.com</code>	Cookie not set; domain mismatch
<code>.com</code>	Cookie not set; TLD is too broad

# When browser sends cookie

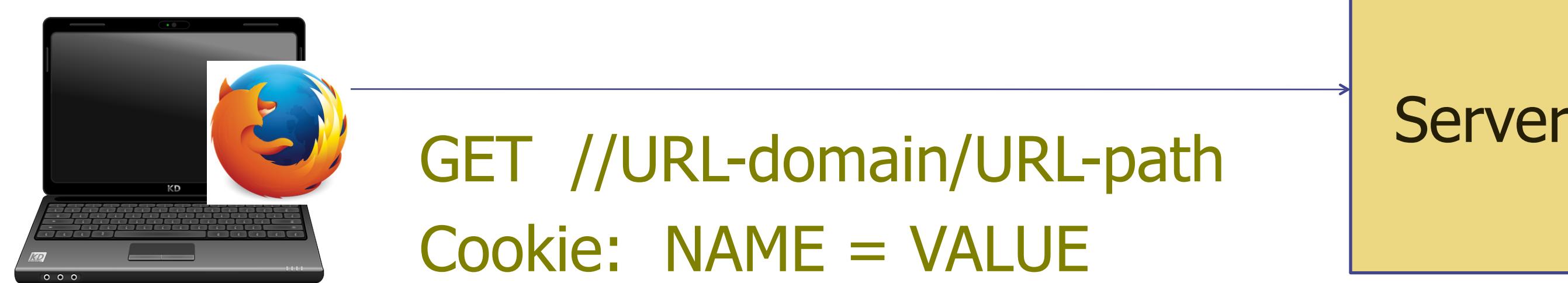


Goal: server only sees cookies in its scope

Browser sends all cookies in URL scope:

- cookie-domain is domain-suffix of URL-domain, and
- cookie-path is prefix of URL-path, and
- [protocol=HTTPS if cookie is “secure”]

# When browser sends cookie



A cookie with  
domain = **example.com**, and  
path = **/some/path/**  
will be included on a request to  
**http://foo.example.com/some/path/subdirectory/hello.txt**

## Examples: Which cookie will be sent?

cookie 1

name = **userid**

value = u1

domain = **login.site.com**

path = /

non-secure

cookie 2

name = **userid**

value = u2

domain = **.site.com**

path = /

non-secure

<http://checkout.site.com/>

cookie: userid=u2

<http://login.site.com/>

cookie: userid=u1, userid=u2

<http://othersite.com/>

cookie: none

# Examples

cookie 1

name = **userid**

value = u1

domain = **login.site.com**

path = /

**secure**

cookie 2

name = **userid**

value = u2

domain = **.site.com**

path = /

non-secure

<http://checkout.site.com/>

<http://login.site.com/>

<https://login.site.com/>

cookie: userid=u2

cookie: userid=u2

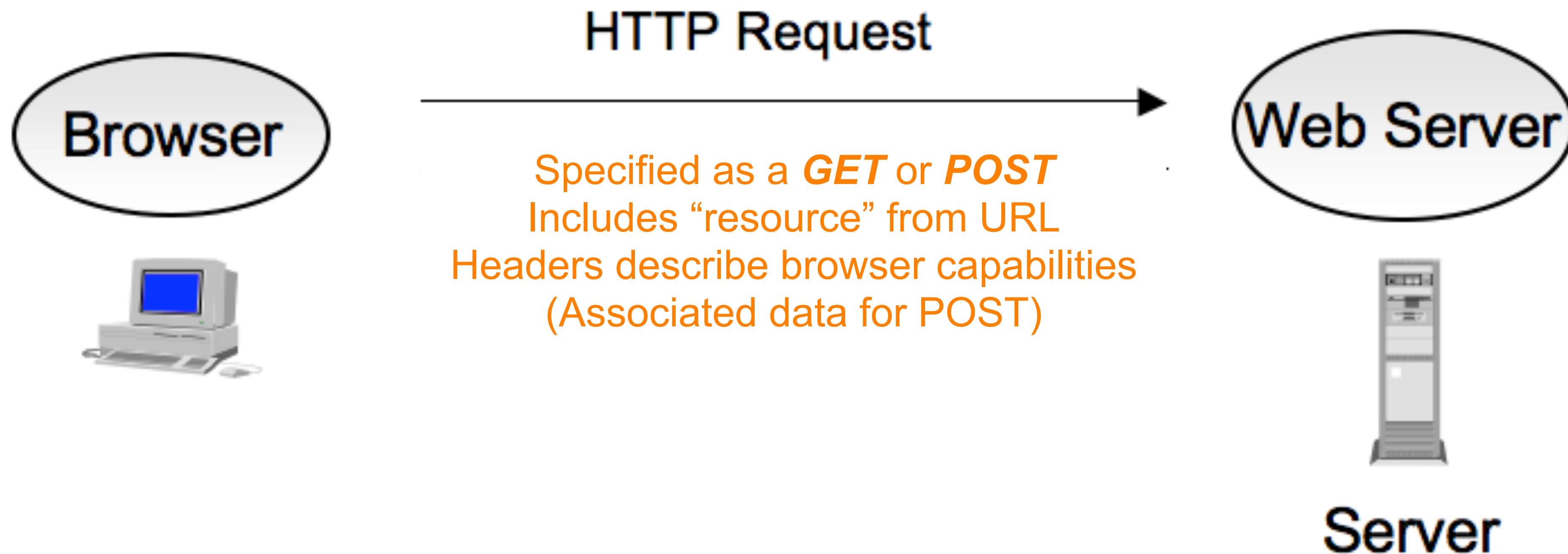
**cookie: userid=u1; userid=u2**

(arbitrary order)

# Cookies & Web Authentication

- ◆ One very widespread use of cookies is for web sites to **track users who have authenticated**
- ◆ E.g., once browser fetched *http://mybank.com/login.html?user=alice&pass=bigsecret* with a correct password, server associates value of “**session**” cookie with logged-in user’s info
  - An “authenticator”

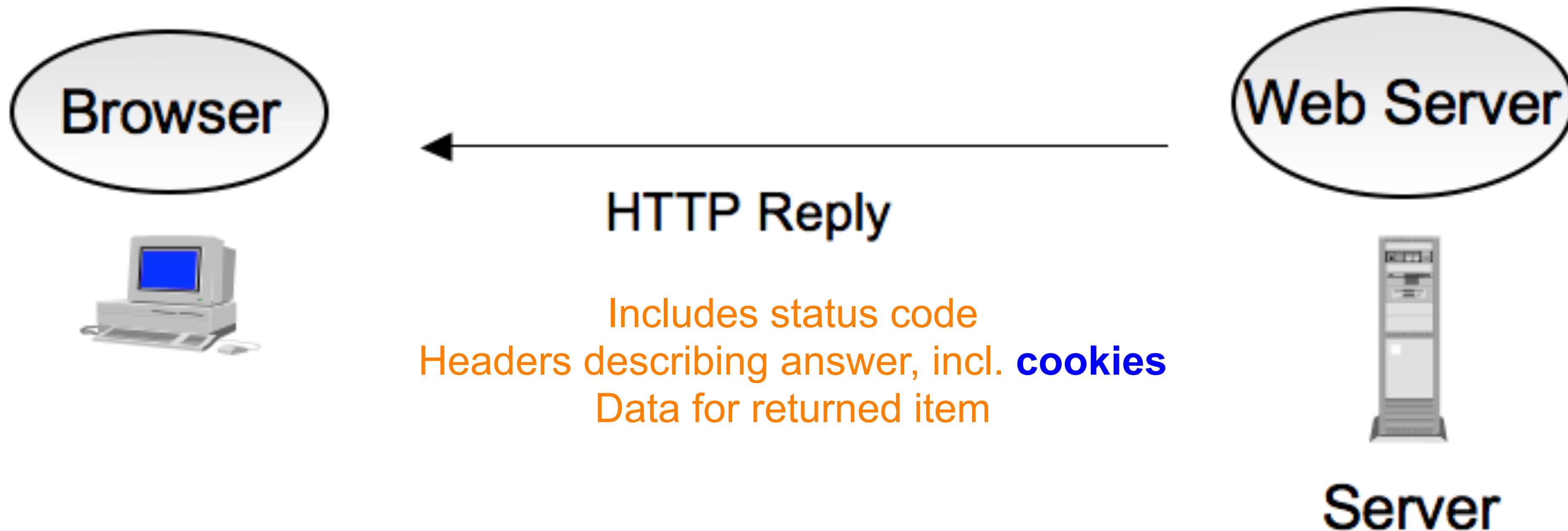
# Basic Structure of Web Traffic



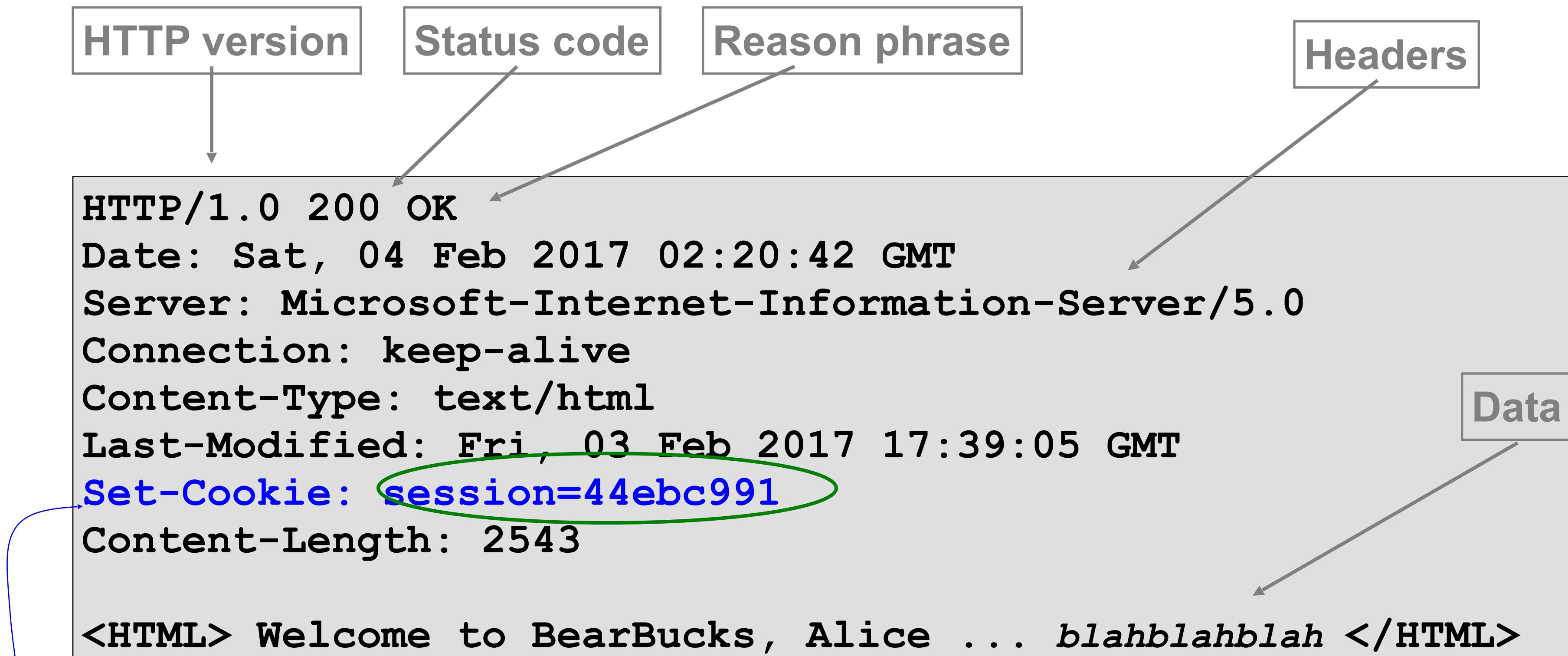
E.g., user clicks on URL:

***http://mybank.com/login.html?user=alice&pass=bigsecret***

# HTTP Cookies



# HTTP Response



**Cookie**

Here the server instructs the browser to remember the cookie "session" so it & its value will be included in subsequent requests

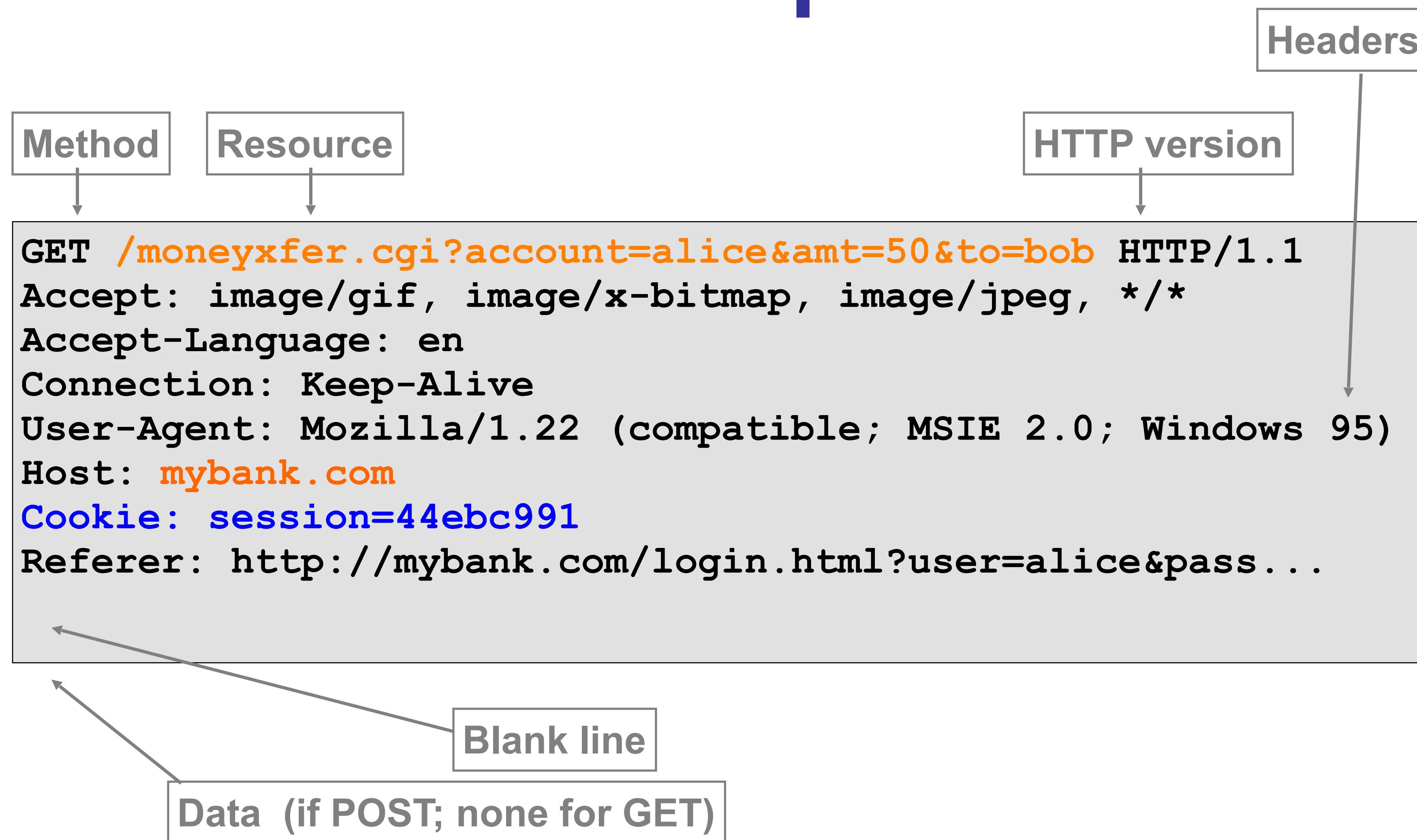
# Cookies & Follow-On Requests



E.g., Alice clicks on URL:

`http://mybank.com/moneyxfer.cgi?account=alice&amt=50&to=bob`

# HTTP Request



# Cookies & Web Authentication

- One very widespread use of cookies is for web sites to track users who have authenticated
- E.g., once browser fetched `http://mybank.com/login.html?user=alice&pass=bigsecret` with a correct password, server associates value of “session” cookie with logged-in user’s info
  - An “authentication token”
- Now server subsequently can tell: “I’m talking to same browser that authenticated as Alice earlier”  
⇒ *An attacker who can get a copy of Alice’s cookie can access the server impersonating Alice!*

# Thinking about Web Security

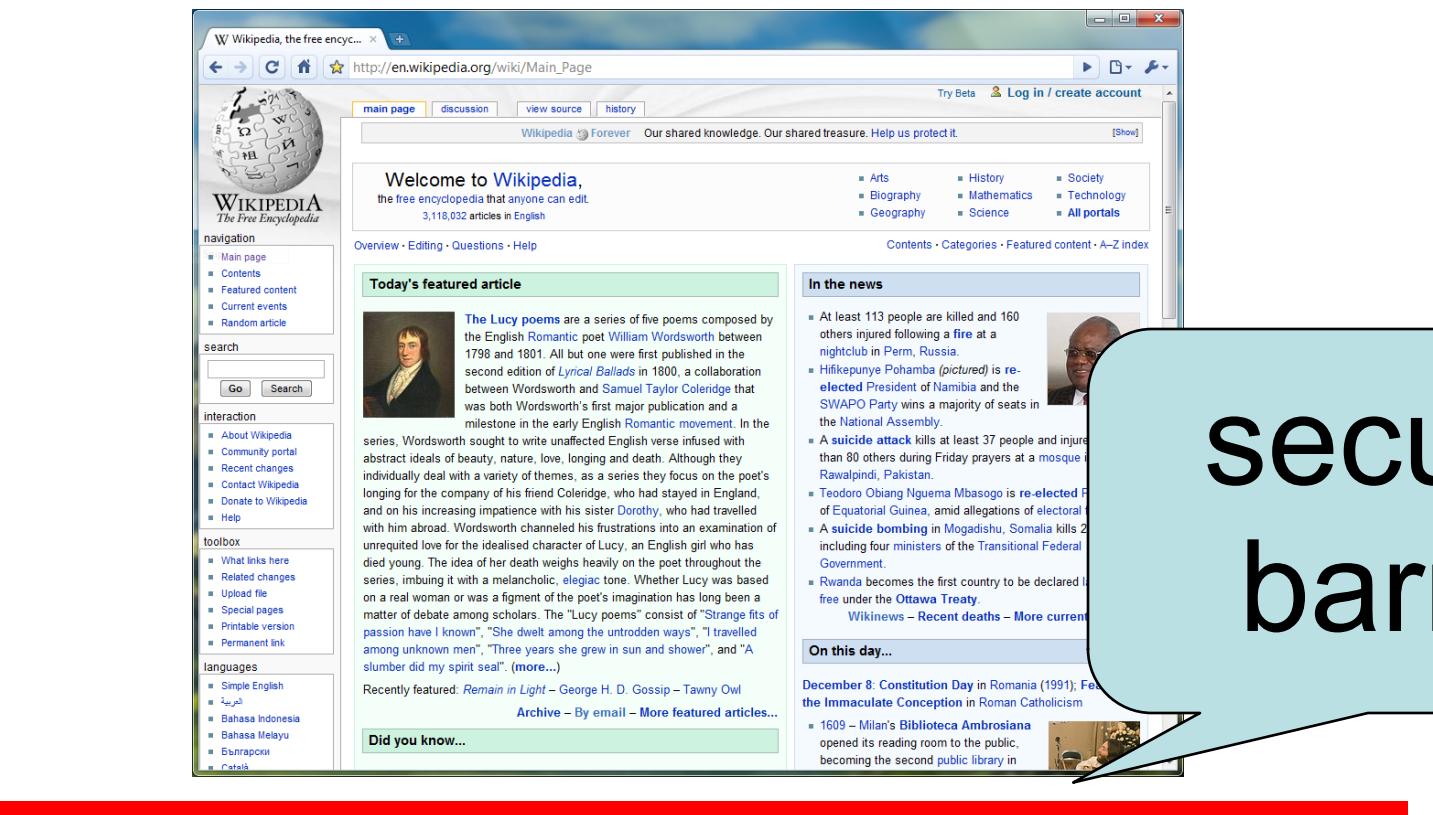
# Desirable security goals

- **Integrity:** malicious web sites should not be able to tamper with integrity of our computers or our information on other web sites
- **Confidentiality/Privacy:** malicious web sites should not be able to learn confidential information from our computers or other web sites
- **Availability:** malicious websites should not be able to keep us from accessing other web resources

# Same-origin policy

- Each site in the browser is isolated from all others

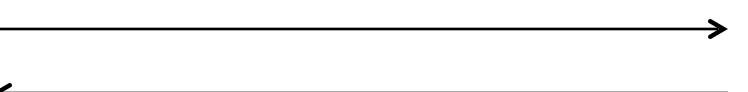
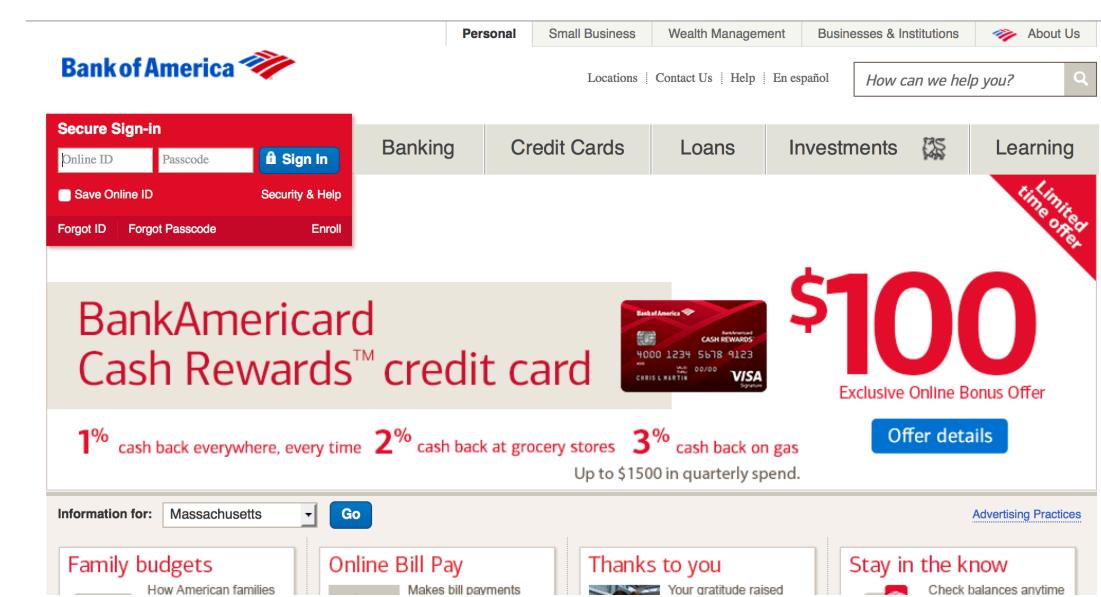
browser:



security  
barrier



wikipedia.org



bankofamerica.com

# Same-origin policy

- Multiple pages from the same site are not isolated

browser:



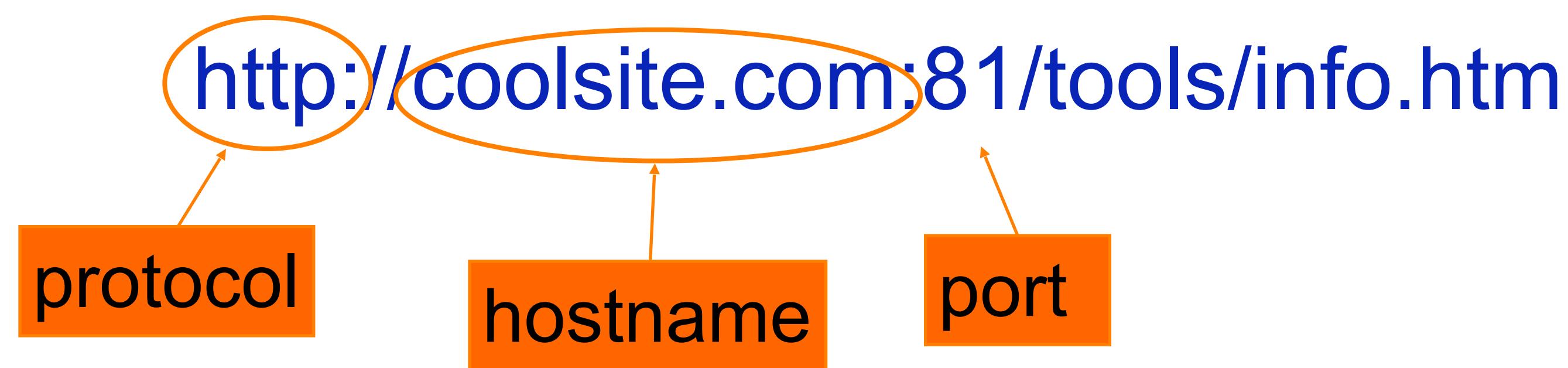
wikipedia.org



wikipedia.org

# Origin

- Granularity of protection for same origin policy
- Origin = protocol + hostname + port



- Hostname compared using ***string matching!*** If these match (along with protocol + port), it is same origin; else it is not. Even though in some cases, it is logically the same origin (e.g., google.com vs google.uk), if there is no string match, it is not.

# Same-origin policy

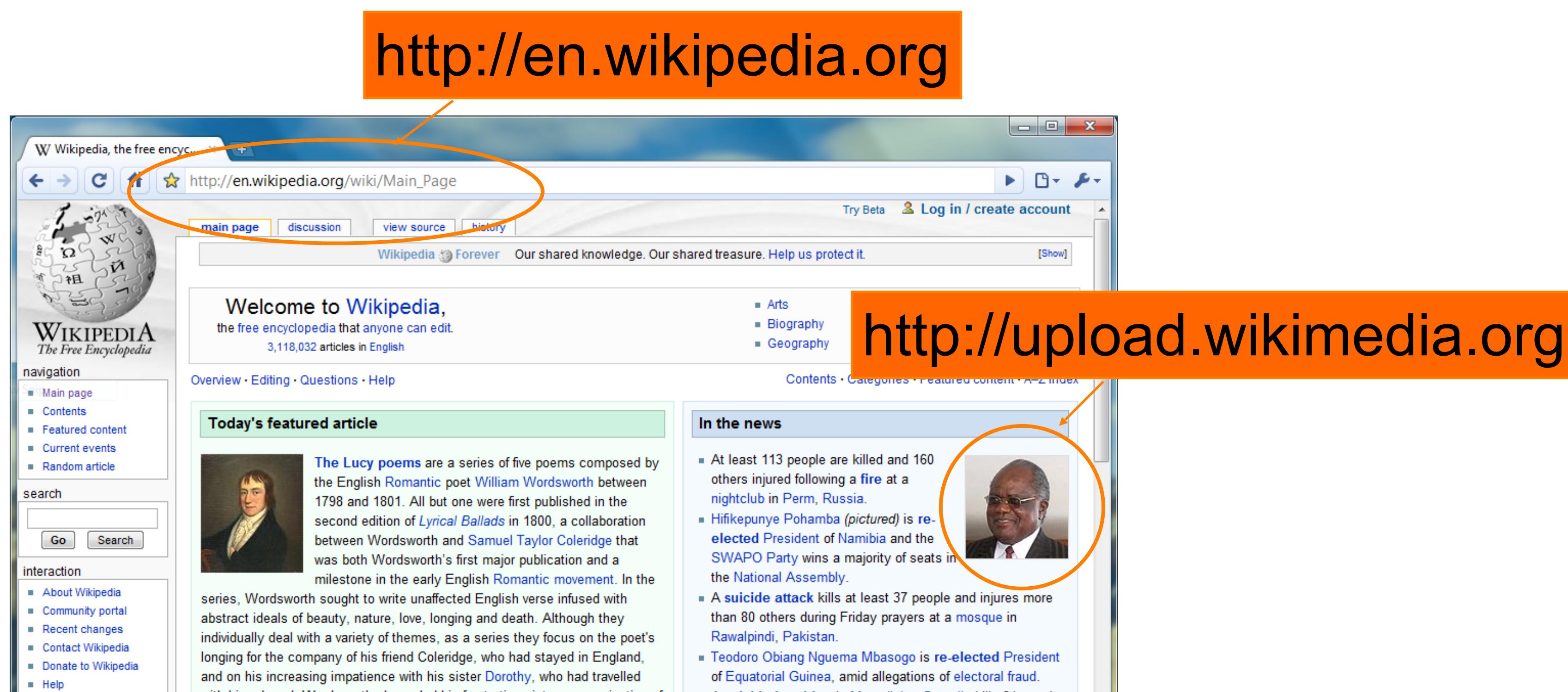
One origin should not be able to access the resources of another origin

Javascript on one page cannot read or modify pages from different origins.

The contents of an *iframe* have the origin of the URL from which the iframe is served; *not* the loading website.

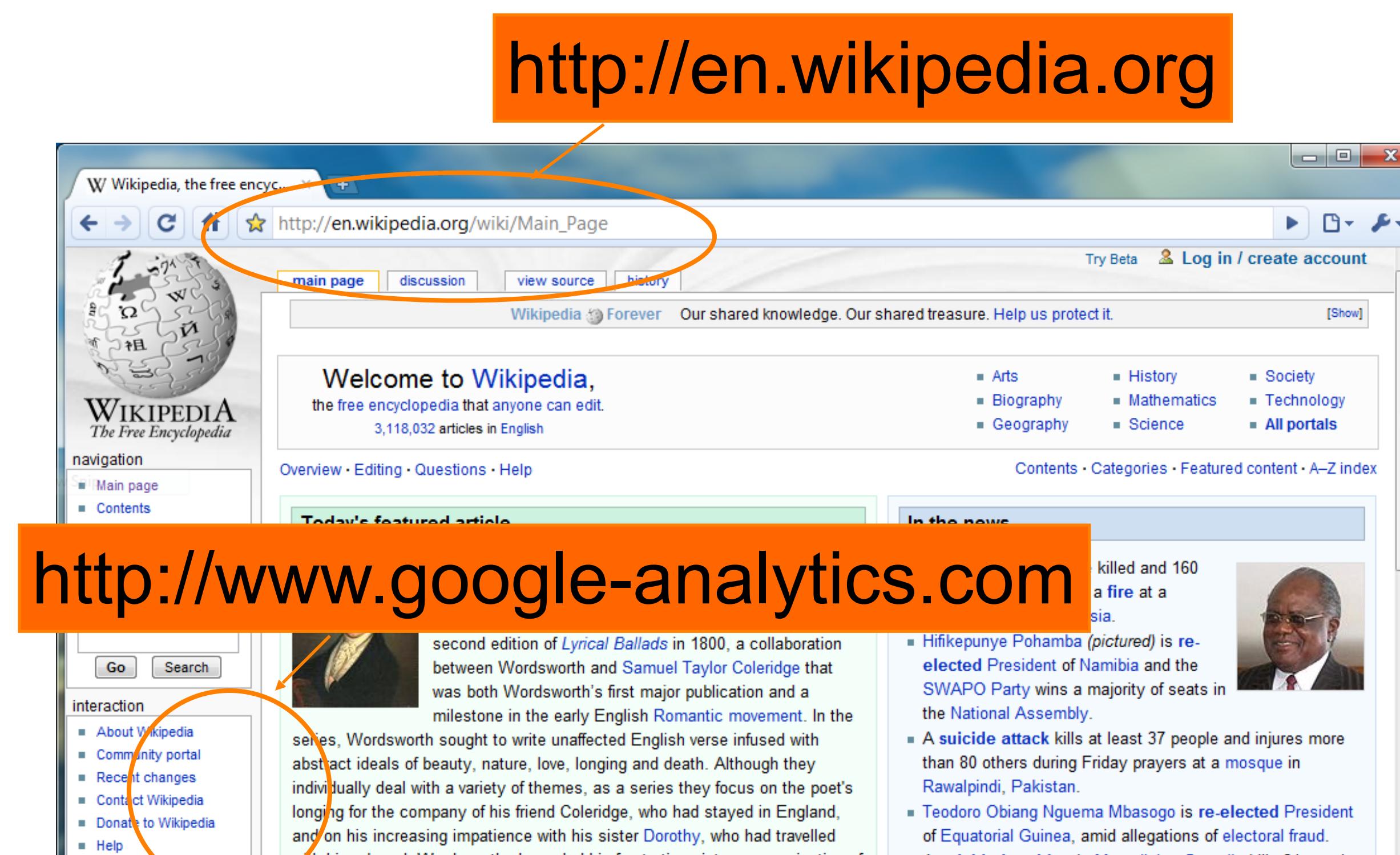
# Same-origin policy

- The origin of a page is derived from the URL it was loaded from



# Same-origin policy

- The origin of a page is derived from the URL it was loaded from
- **Special case:** Javascript runs with the origin of the page that loaded it



# Assessing SOP

Originating document	Accessed document	
http://wikipedia.org/a/	http://wikipedia.org/b/	
http://wikipedia.org/	http://www.wikipedia.org/	
<b>http://wikipedia.org/</b>	<b>https://wikipedia.org/</b>	
http://wikipedia.org:81/	http://wikipedia.org:82/	
http://wikipedia.org:81/	http://wikipedia.org/	

except

