

第二章 知识表示方法

主要内容:

§ 2.1 状态空间表示

§ 2.2 问题归约法

§ 2.3 谓词逻辑表示

§ 2.4 语义网络法

§ 2.5 框架与本体

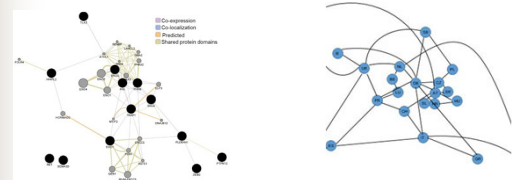
状态空间

§ 2.1 状态空间表示

■ State space method

A method for problem representation and problem-solving based on the *solution space*.

■ It is based on the *State* and *Operator*.



State & Operator

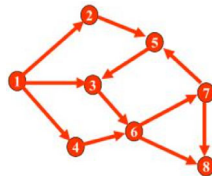
- **State:** Some variables, describing the difference among different things or events of some class

$$Q = [q_0, q_1, \dots, q_n]^T$$

- **Operator:** A means that is used to transform a problem from one state to another.

Nodes represent states,

Arcs represent occurring events.



状态空间

- **State space:** a graph that represents all possible states and their relationships with the problem.

Example: 3-element states $\{S, F, G\}$

S : set of all possible initial states of problem

F : set of operators

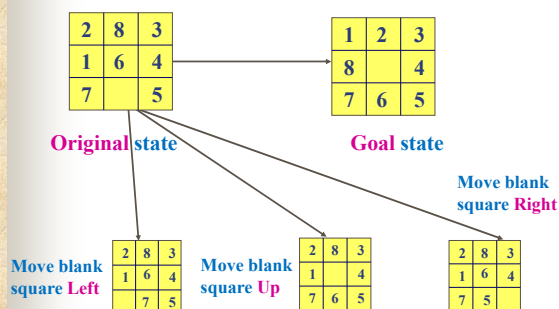
G : set of goal states



State space

八数码问题

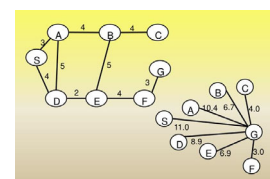
八数码问题的状态空间



基本算法

状态空间法的基本算法

- 规定一组操作(算子), 能够使状态从一个状态变为另一个状态。 $n_i \xrightarrow{\text{operator}} n_j$
- 决定一种搜索策略, 使得能够从初始状态出发, 沿某个路径达到目标状态。



Route Planning

Example: Route Planning

■ Initial state

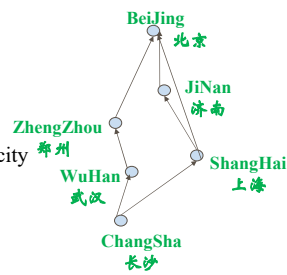
Starting from ChangSha

■ Operators

Driving from a city to another city

■ Goal state

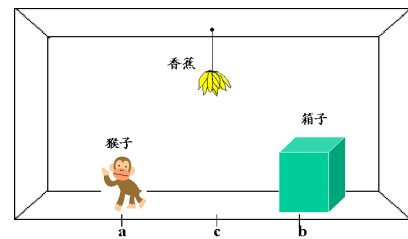
Destination city - BeiJing



Route Planning

Monkey & Banana

Example: Monkey & Banana



Monkey and Banana Problem

Monkey & Banana

■ Problem State Set : (W, x, Y, z)

➢ W Monkey's horizontal position

➢ x Monkey is above the box $x = 1$, else $x = 0$

➢ Y Box's horizontal position

➢ z Monkey get the banana $z = 1$, else $z = 0$

■ Operator:

➢ **Goto** (U) means monkey go to position U

$(W, 0, Y, z) \xrightarrow{\text{goto } (U)} (U, 0, Y, z)$

➢ **Pushbox** (V) monkey pushes the box to position V :

$(W, 0, W, z) \xrightarrow{\text{pushbox } (V)} (V, 0, V, z)$

Monkey & Banana

➢ **Climbbox**: monkey climbs the top of the box

$(W, 0, W, z) \xrightarrow{\text{Climbbox}} (W, 1, W, z)$

Attention: when used these two operators, the monkey and the box should be on the same horizontal position and the monkey is not on the top of the box.

➢ **Grasp**: monkey get the banana

$(c, 1, c, 0) \xrightarrow{\text{Grasp}} (c, 1, c, 1)$

c : the position that right below the banana

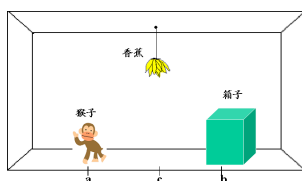
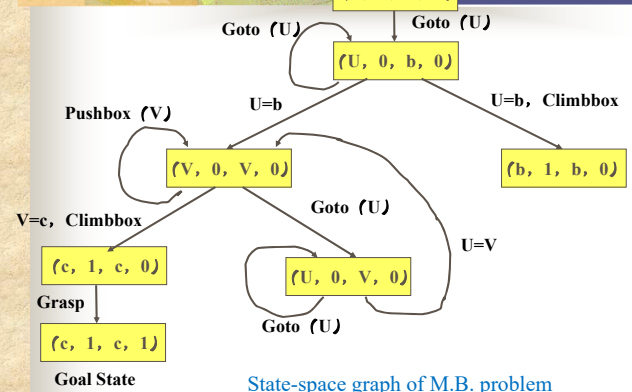
Monkey & Banana

■ Initial state is $(a, 0, b, 0)$

■ Operator sequence from the initial state to the goal state:

$\{\text{Goto } (b), \text{Pushbox } (c), \text{Climbbox}, \text{Grasp}\}$

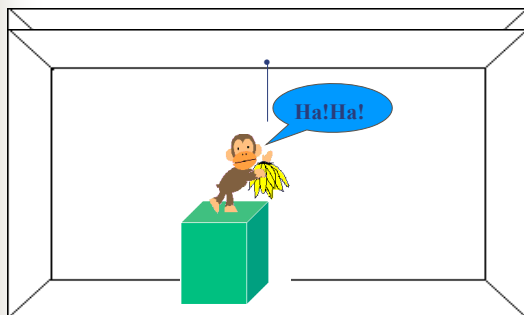
is a solution of the problem.

Initial State $(a, 0, b, 0)$ Monkey & Banana

State-space graph of M.B. problem

Monkey & Banana

Solving processes



传教士野人

举例：传教士野人问题

三个传教士M和三个野人C过河，只有一条能装下两个人的船。如果在河的一方（含船上），野人的人数大于传教士的人数，那么传教士就会有危险，提出一种安全的渡河方案！



传教士野人

状态可有多种表示方法

左岸传教士数,右岸传教士数,左岸野人数,右岸野人数,船的位置

简化表示

左岸传教士数,左岸野人数,船的位置



初始状态: (0, 0, 0) 0: 右岸

目标状态: (3, 3, 1) 1: 左岸

传教士野人

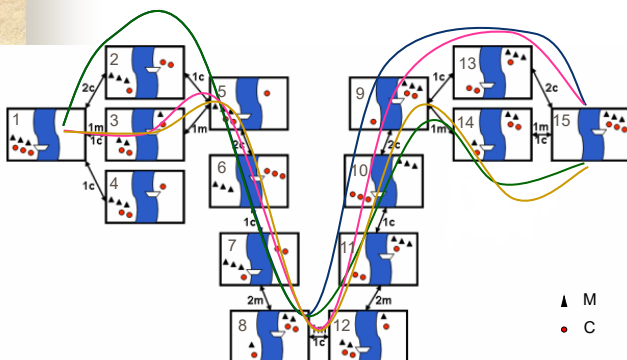
状态的转换 MC问题中的算子

- 将传教士或野人运到河对岸
- Move-1m1c-lr: 将一个传教士(m)和一个野人(c)从左岸(L)运到右岸(R)
- 所有可能操作

Move-1m1c-lr	Move-1m1c-rl	Move-2c-lr
Move-2c-rl	Move-2m-lr	Move-2m-rl
Move-1c-lr	Move-1c-rl	Move-1m-rl
Move-1m-rl		

传教士野人

传教士野人问题状态空间图

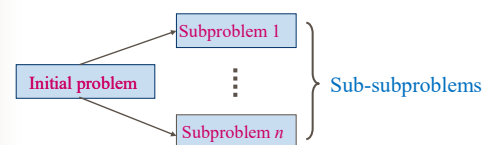


问题的规约

§ 2.2 问题规约法

■ 问题的规约

将一个大的问题变换成若干子问题,子问题又可分解成更小的子问题,这样一直分解到可以直接求解为止,全部子问题的解即为大的问题的解。



规约三要素

■ 归约方法求解问题三要素

- 初始问题的描述
- 一组将问题变换成子问题的变换规则
- 一组本原问题的描述

■ 问题规约的实质

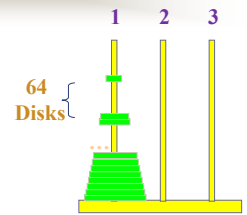
- 从初始问题出发, 建立子问题以及子问题的子问题, 直至把初始问题规约为一个本原问题的集合

Hanoi Puzzle

Working Mechanism-Reduction

1. Tower of Hanoi Puzzle (THP)

- Three pegs and sixty-four disks
 - Moving times: 64 disks
- $$2^{64}-1 \approx 2^{64} = 10^{19.27}$$

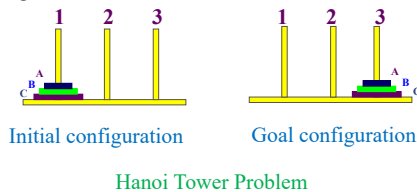


If one person moves 1 disk in one second, then to finish this problem needs more than 3000 billion years.

Hanoi Puzzle

■ Three Disk Hanoi-Tower

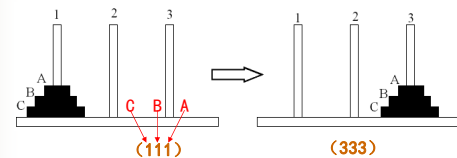
- **Question:** There are three pegs-1, 2 and 3 and three disks of different sizes A, B and C. The disks have holes in their centers so that they can be stacked on the pegs.



Hanoi Puzzle

■ Constrain conditions

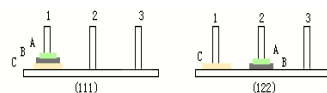
- Only one disk can be moved at a time;
- Only the top disk on a peg can be moved;
- Larger disk may never be placed on top of a smaller one.



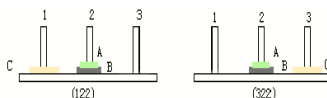
Hanoi Puzzle

原问题可以归约为3个子问题:

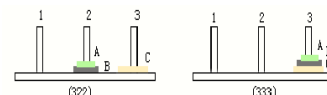
子问题1: 移圆盘A和B至
柱子2 (借助柱子3)



子问题2: 移动圆盘C至柱
子3

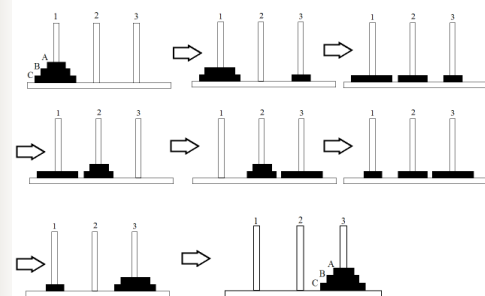


子问题3: 把圆盘A和B移至
柱子3 (借助柱子1)

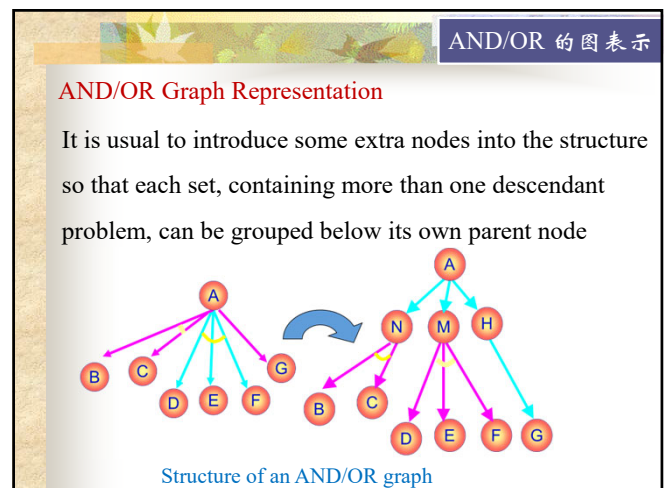
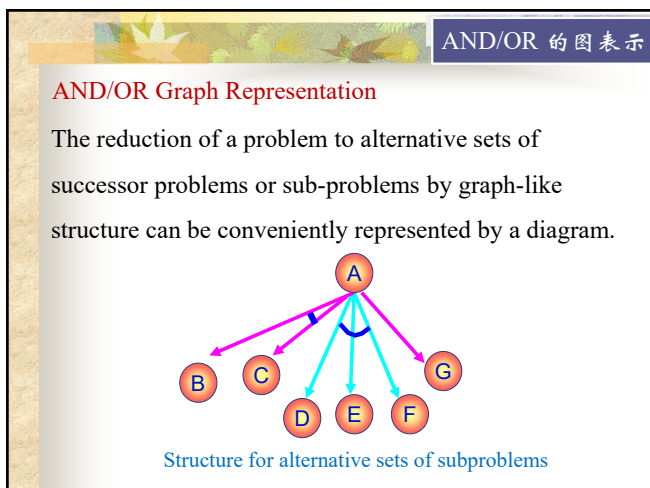
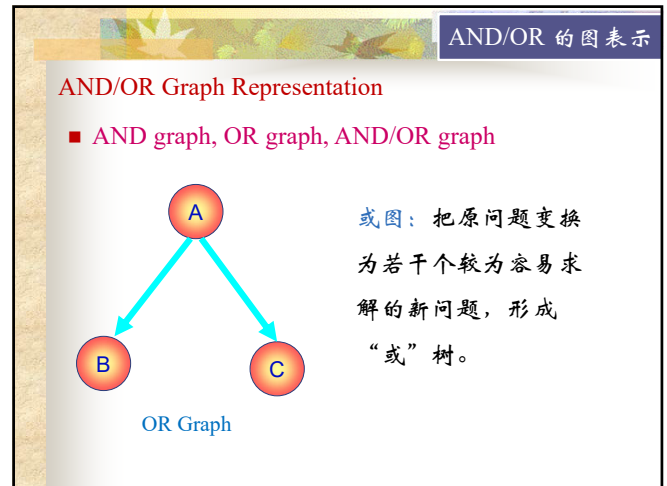
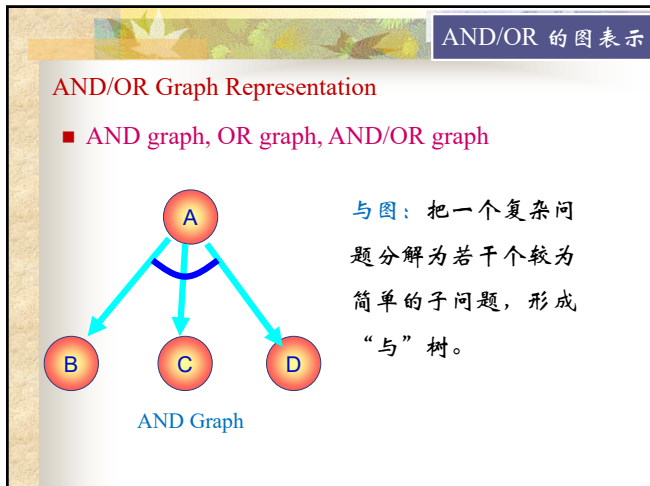
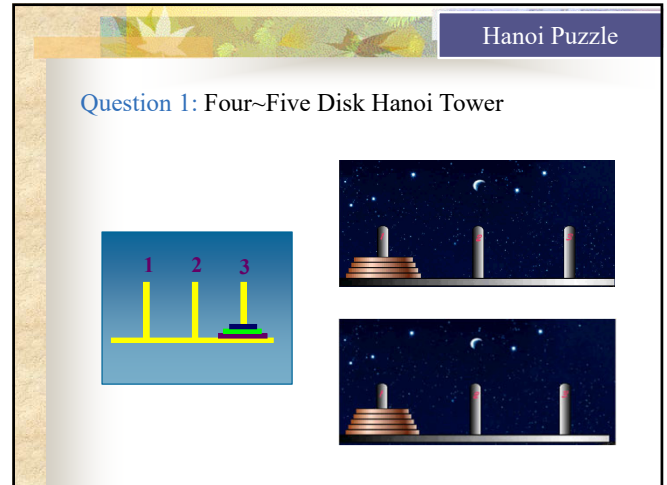
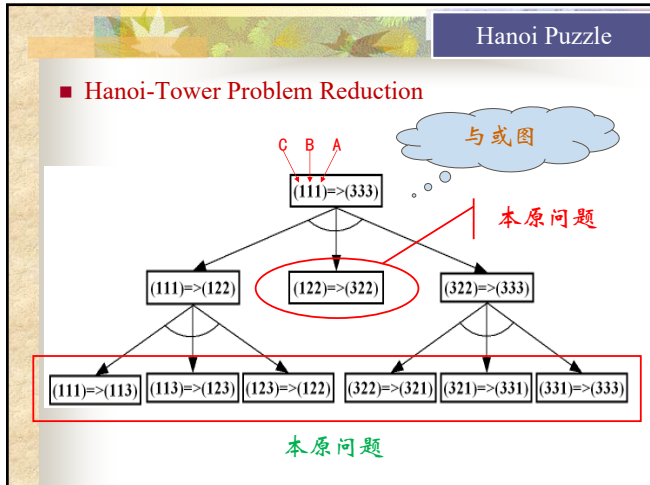


Hanoi Puzzle

归约过程 (3个圆盘)

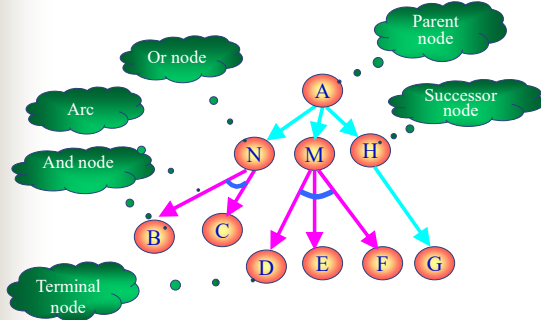


The solving process of Hanoi tower problem



AND/OR 的图表示

AND/OR Graph Representation

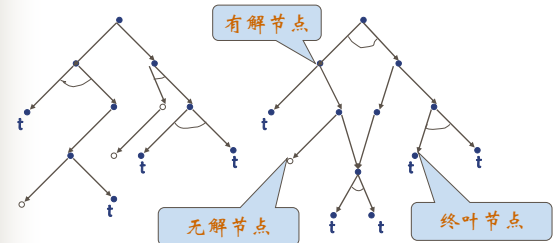


Structure of an AND/OR graph

AND/OR 的图表示

AND/OR Graph Representation

Solved node and unsolved node



The solved node and unsolved node in an AND/OR graph

AND/OR 的图表示

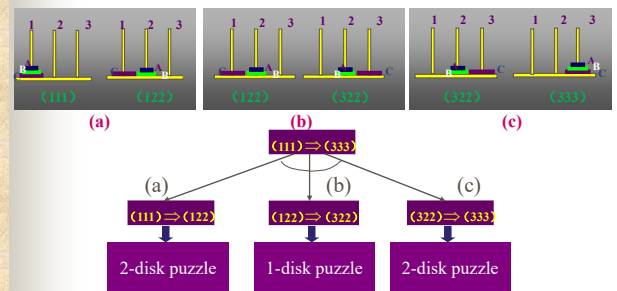
AND/OR Graph Representation

■ 与或图的构建规则

- Every node represents a **single problem** or a **set of problems**.
- Terminal nodes correspond to primitive problems.
- Operators transform a problem into a set of sub-problems.
- Directed links are towards every AND node of the set of sub-problems.

AND/OR 的图表示

AND/OR Graph Representation

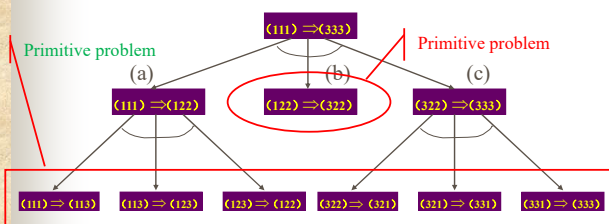


Reduction graph for the 3-disk Hanoi Tower

AND/OR 的图表示

AND/OR Graph Representation

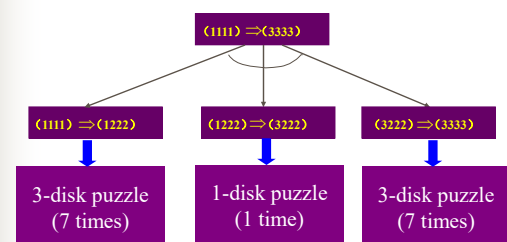
- Puzzle number (b): a primitive problem
- Puzzle number (a) and (c): 2-disk puzzle



And/or graph for the 3-disk Hanoi Tower

AND/OR 的图表示

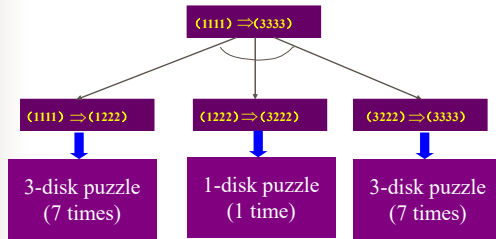
Four Disk Hanoi Tower Graph Representation



Reduction graph for the 4-disk Hanoi Tower

AND/OR 的图表示

Four Disk Hanoi Tower Graph Representation



Reduction graph for the 4-disk Hanoi Tower

Hanoi Tower

Hanoi Tower

Number of Dishes	Moving times
1	1
2	$2^2-1=3$
3	$2^3-1=7$
...	...
64	$2^{64}-1$

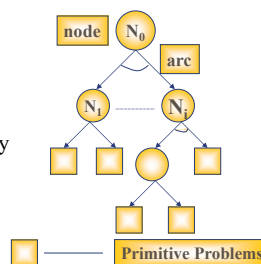
AND/OR 的图表示

AND/OR Graph Representation

■ Conclusions

- A node is used to represent a problem.
- Original problem is reduced into several subproblems and sub-subproblems.
- The solution can be obtained by searching the AND/OR graph.
- The solution is a solution tree.

Reduction graph



谓词逻辑法

§ 2.3 谓词逻辑法

一个陈述句由主语和谓语两个部分组成

在谓词逻辑中，为揭示命题内部结构及不同命题的内部结构关系，就按这两个部分进行分析

主语称为个体或客体

谓语称作谓词

个体和谓词

■ 个体和谓词

- 在原子命题中，所描述的对象，称为个体
- 用以描述个体性质，或个体间关系的部分，称为谓词

个体

- 个体，是指可以独立存在的事物，具体的或抽象的。

例如：李明、香蕉、电视机、精神、思想，等

- 表示特定的个体，称为个体常元

例如：李明、字母A、数字5

- 表示不确定的个体，叫作个体变元

例如：花朵、星星、男生、学校

谓词

- **谓词**:用以描述个体性质, 或个体间关系的部分, 称为谓词
- 当谓词与一个个体相关时, 刻划了该个体的性质
例如: 5是质数, A是字母 一元谓词
- 当谓词与两个或两个以上个体相关时, 刻划了个体间的关系
例如: 李明生于北京, 10比5大 二元谓词

命题的谓词形式

■ 命题的谓词形式

单独的个体和谓词不能构成命题

■ 命题的表示:

将表示个体的小写字母, 写在表示谓词的大写字母右侧的括号里。

谓词形式

命题的谓词形式

“李明是大学生”

其中: 李明是个体, 用小写 a 表示

“...是大学生”是谓词, 用大写 S 表示

则该命题表示为 $S(a)$: 李明是大学生

谓词形式

命题的谓词形式

“张华出生在北京”

其中: 张华和北京是个体, 用 a 表示张华, 用 b 表示北京

“...出生在...”是谓词, 用大写 B 表示

则该命题表示为 $B(a, b)$: 张华出生在北京

形式语言

- **一阶谓词演算**:是一种形式语言, 它把数学中的逻辑论证符号化。
- **形式语言**:一种有别于自然语言的符号语言。

基本符号

1.谓词演算语法和语义

■ 基本符号

- 谓词符号、变量符号、函数符号、常量符号、括号和逗号

$\text{MARRIED}(\text{father}(\text{LI}), \text{mother}(\text{LI}))$

predicate function constant

- 原子公式 = 谓词 + 项

$\text{Inroom}(\text{robot}, \text{r1})$

连词和量词

2. 连词和量词

■ 连词(Connectives)

- 与及合取 (AND, conjunction)

Like (I, Music) \wedge Like (I, Sport)

- 或及析取 (OR, disjunction)

Like (I, Music) \vee Like (I, Sport)

- 蕴涵 (Implication) 如果 \rightarrow 那么

Runs (Philip, Fastest, 100m) \rightarrow Wins (Philip, Champion)

- 非 (Not) \sim 或者 \neg

\sim IN (Ma, Shanghai, Sept-2022)

连词和量词

■ 量词

➢ 全称量词

一个原子公式 $P(x)$, 对于所有可能的变量 x 都具有真值 T 。
这个特性可由在 $P(x)$ 前面加上 **全称量词** $\forall(x)$ 来表示。例如,
句子“所有的机器人都是灰色的”可表示为:

$\forall(x)[\text{ROBOT}(x) \Rightarrow \text{COLOR}(x, \text{GRAY})]$

➢ 存在量词

如果至少有一个 x 值可使 $P(x)$ 具有真值 T , 那么这一特性可在 $P(x)$ 前面加上存在量词 $\exists(x)$ 来表示。句子“1号房间内有个物体”可表示为: $\exists(x) \text{ INROOM}(x, R1)$

全称和存在量词

■ Quantifier

Universal quantifier

An atomic formula $P(x)$ has the value T for **All** variables x :

$(\forall x) P(x)$

Universal Quantifier \rightarrow Bound variables
Domain of universal quantifier

Existential quantifier

If there **Exists** at least one variable of x that makes $P(x)$ has the value T : $(\exists x) P(x)$

E.G. $(\exists x) \text{ INROOM}(x, r1)$ (1号房间内有个物体)

Existential Quantifier \rightarrow Bound variables
Domain of existential quantifier

谓词公式

2. 谓词公式

■ 原子谓词公式

用 $P(x_1, x_2, \dots, x_n)$ 表示一个 n 元谓词公式, 其中 P 为 n 元谓词, x_1, x_2, \dots, x_n 为客体变量或变元。通常把 $P(x_1, x_2, \dots, x_n)$ 叫做谓词演算的原子公式, 或原子谓词公式。

■ 分子谓词公式

用连词把原子谓词公式组成复合谓词公式, 并把它叫做分子谓词公式。

谓词公式

■ WFF (well-formed formulas)

- Legal sentences = WFF

- Definition

- 原子公式是合适公式
- If A is a WFF, then $\sim A$ is also a WFF
- If A and B are WFF, then $(A \wedge B)$, $(A \vee B)$, $(A \Rightarrow B)$, $(A \longleftrightarrow B)$ are also WFF
- If A is WFF, and x is a free variable in A , then $\forall(x)A$ and $\exists(x)A$ are also WFF
- Only the formula by above 4 rules, would be WFF

真值表

合适公式的真值表

P	Q	$P \vee Q$	$P \wedge Q$	$P \Rightarrow Q$	$\sim P$
T	T	T	T	T	F
F	T	T	F	T	T
T	F	T	F	F	F
F	F	F	F	T	T

Equivalence

■ Equivalence

Two WWFs are **Equivalence** if they have the same truth value under every interpretation.

(1) **Double negative elimination** $\neg(\neg P) \equiv P$

(2) $(P \vee Q) \equiv \neg P \rightarrow Q$

(3) **De Morgan's law**

$$\neg(P \wedge Q) \equiv \neg P \vee \neg Q \quad \neg(P \vee Q) \equiv \neg P \wedge \neg Q$$

(4) **Distributive law**

$$P \wedge (Q \vee R) \equiv (P \wedge Q) \vee (P \wedge R)$$

$$P \vee (Q \wedge R) \equiv (P \vee Q) \wedge (P \vee R)$$

Equivalence

(5) **Commutative laws**

$$P \wedge Q \equiv Q \wedge P, \quad P \vee Q \equiv Q \vee P$$

(6) **Associative law**

$$(P \wedge Q) \wedge R \equiv P \wedge (Q \wedge R)$$

$$(P \vee Q) \vee R \equiv P \vee (Q \vee R)$$

(7) **Contra-negative law**

$$P \rightarrow Q \equiv \neg Q \rightarrow \neg P$$

Equivalence

合适公式的性质

■ $(P \wedge \neg Q) \rightarrow \neg R$

P: Jon likes cats, Q: Jon likes dogs, R: Jon likes all kinds of pets.

■ An expression is a sentence \longleftrightarrow It can be formed of legal symbols

Equivalence

合适公式的性质

■ $((P \wedge \neg Q) \rightarrow \neg R) \equiv \neg P \vee Q \vee R$

➢ $(P \wedge \neg Q) \rightarrow \neg R$ symbols conjunction implication

➢ $\neg P \vee Q \vee R$ symbols disjunction

➢ $((P \wedge \neg Q) \rightarrow \neg R) \equiv \neg P \vee Q \vee R$ equivalence

置换与合一

3. 置换与合一

■ 置换

➢ 假元推理

$$\left. \begin{array}{l} W_1 \Rightarrow W_2 \\ W_1 \end{array} \right\} \rightarrow \text{Generate } W_2$$

➢ 全称化推理

$$(\forall x) W(x) \rightarrow W(A) \quad \begin{array}{l} \text{Bound var.} \quad \text{Free const} \end{array}$$

➢ 综合推理

$$\left. \begin{array}{l} (\forall x) [W_1(x) \rightarrow W_2(x)] \\ W_1(A) \end{array} \right\} \rightarrow \text{Generate } W_2(A)$$

置换与合一

■ 合一

➢ 寻找项对变量的置换，以使两表达式一致。

➢ 如果一个置换 s 作用于表达式集 $\{E_i\}$ 的每个元素，则我们用 $\{E_i\} s$ 来表示置换例的集。我们称表达式集 $\{E_i\}$ 是可合一的。

$$\text{If } E_{1s} = E_{2s} = E_{3s} = \dots$$

Expression set $\{E_i\}$ is unifiable, s is unifier.

定义 & 组成

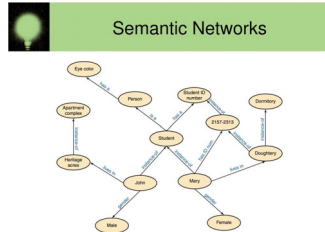
§ 2.4 语义网络法

■ 定义 & 组成

➤ 定义：语义网络是知识的图形表示，由节点和链接组成。

➤ 组成

- 词法
- 结构
- 过程
- 语义



二元语义网络

■ A collection of predicate calculus expression can be represented by a graph structure-Semantic Network.

1. Representation of Two-Element Semantic Network

(二元语义网络的表示)

-Representation of simple facts

举例：小燕是一只燕子，燕子是鸟；巢-1是小燕的巢，巢-1是巢中的一个

ISA (is-a) Relation

ISA (is-a) Relation

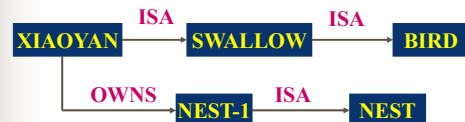
- The ISA (is-a) relation is often used to link instances to classes, classes to super-classes
 - Some links (e.g. isPart) are inherited along ISA paths
 - The semantics of a SN can be relatively very formal or informal.
- Often defined at the implementation level

Example

■ All swallow are birds



■ Xiaoyan is a swallow, swallow is birds; Nest-1 is xiaoyan's nest, Nest-1 is one of the nests.



语义基元

■ 选择语义基元

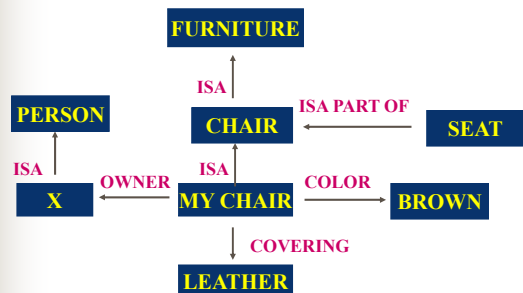
➤ Looking for basic concepts and basic arcs

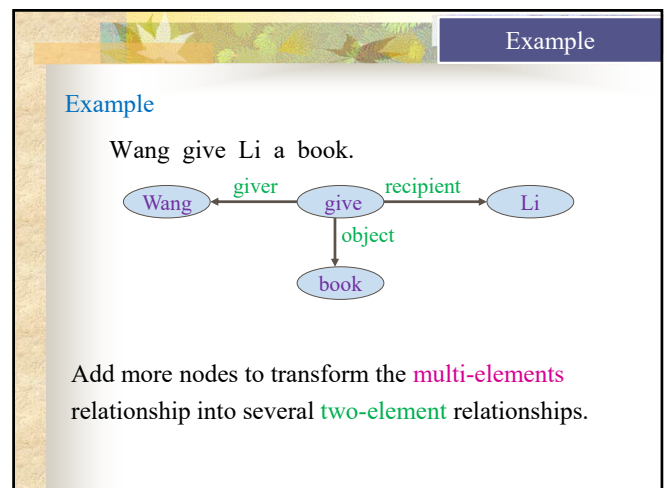
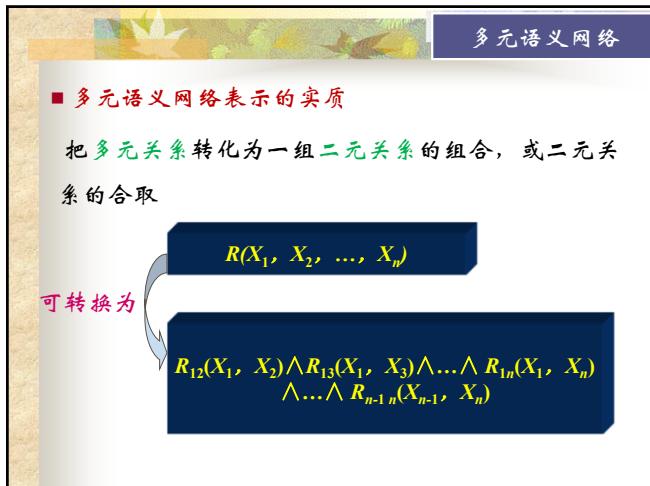
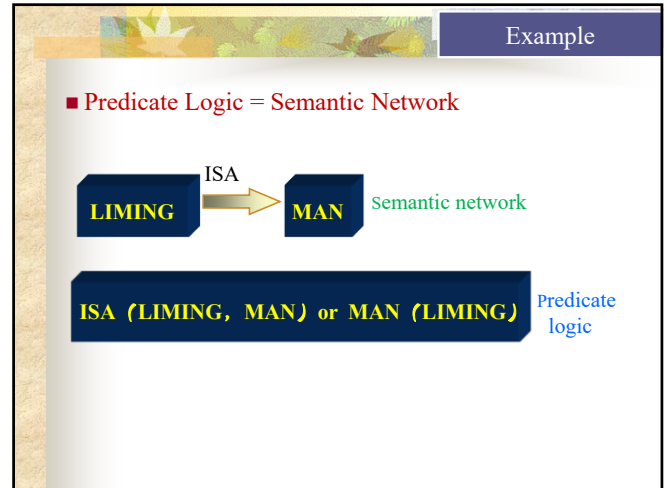
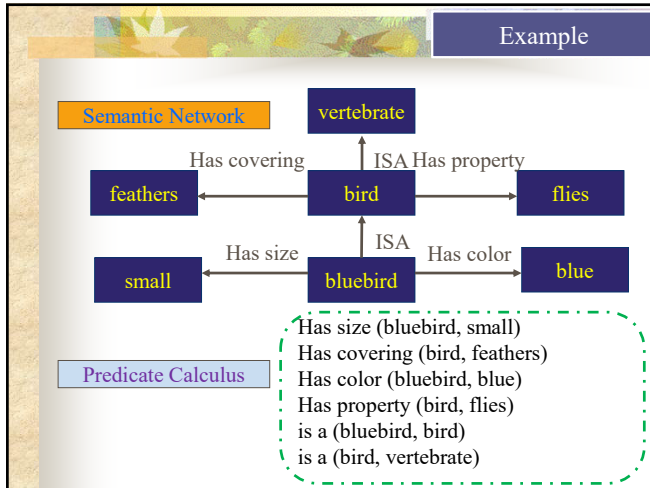
试图用一组基元来表示知识，以便简化表示，可用简单的知识来表示更复杂的知识。

Example:

The color of my chair is brown; the covering of the chair is leather; chair is one kind of furniture; chair is part of seat; chair's owner is X; X is a person

Example



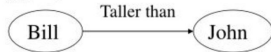


Example

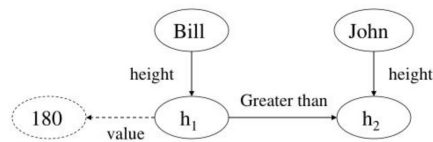
Example

“Bill is taller than John.”

■ Non appropriate scheme :



■ Appropriate scheme :

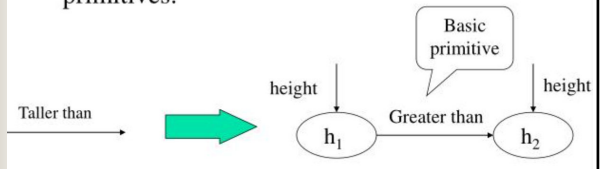


Example

Example

For an appropriate scheme:

- Draw relations on the basic of primitives.
- Represent complicated relations with this primitives.



Inheritance

■ Inheritance (继承)

把事物的描述从概念节点传递到实例节点

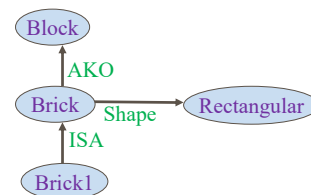
- Pass (遗留) — the descendants get the properties from the ancestors
- Add (添加) — the descendants expand the properties of the ancestors
- Exclude (排除) — stop inheriting the properties

Inheritance

Three methods for inheritance

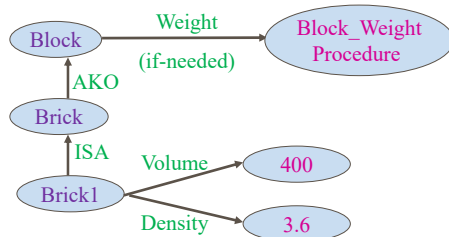
- 值继承 is-a, AKO (a kind of)

What is the shape of Brick1?



Inheritance

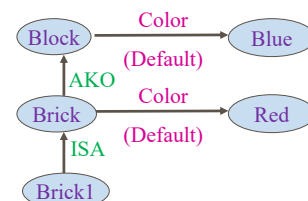
- 如果-需要“继承” If-needed (it can not inherit from the ancestors, we get it from other programs)



Inheritance

- 缺省“继承” Default—it is mostly the truth

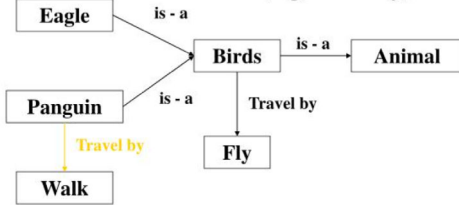
Example: the birds can fly;



Example

Example

Some times we had to override a relation for an inherited node (e.g travel by).



Example

Example

Use semantic networks to represent the below:

- Nellie is an elephant,
- He likes apples.
- Elephants are a kind of mammals,
- They live in Africa,
- And they are big animals.
- Mammals and reptiles are both animals,
- All animals have head.

Example

