# CS208 Lab9 Practice

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# A Task-scheduling Problem

### **Desciption**

A unit-time task is a job, such as a program to be run on a computer, that requires exactly one unit of time to complete. Given a finite set S of unit-time tasks, a schedule for S is a permutation of S specifying the order in which to perform these tasks. The first task in the schedule begins at time S and finishes at time S, and so on.

The problem of scheduling unit-time tasks with deadlines and penalties for a single processor has the following inputs:

```
A set S = \{a_1, a_2, \dots, a_n\} of n unit-time tasks;
```

A set of n integer deadlines  $d_1, d_2, \ldots, d_n$ , such that each  $d_i$  satisfies  $1 \le d_i \le n$  and task  $a_i$  is supposed to finish by time  $d_i$ ;

A set of n nonnegative weights or penalties  $w_1, w_2, \ldots, w_n$ , such that we incur a penalty of  $w_i$  if task  $a_i$  is not finished by time  $d_i$ , and we incur no penalty if a task finishes by its deadline.

We wish to find a schedule for S that minimizes the total penalty incurred for missed deadlines.

### **Analysis**

It's 01 backpack. Suppose dp[i][j] represents in the first i tasks, we choose j tasks, the max sum of w. So if we do not choose i-th task, dp[i][j] = dp[i-1][j]; if we choose, dp[i][j] = dp[i-1][j-1] + w[i].

Time complexity:  $O(n^2)$ .

#### C++ Code

# **Another Task-scheduling Problem**

# **Description**

The problem of scheduling non-unit time tasks with deadlines and penalties for a single processor has the following inputs:

A set  $S = \{a_1, a_2, \dots, a_n\}$  of n non-unit time tasks;

A set of n integer time  $t_1, t_2, \ldots, t_n$ , to finish  $a_i$  need  $t_i$  time;

A set of n integer deadlines  $d_1, d_2, \ldots, d_n$ , such that each  $d_i$  satisfies  $1 \le d_i \le n$  and task  $a_i$  is supposed to finish by time  $d_i$ :

A set of n nonnegative weights or penalties  $w_1, w_2, \ldots, w_n$ , such that we incur a penalty of  $w_i$  if task  $a_i$  is not finished by time  $d_i$ , and we incur no penalty if a task finishes by its deadline.

We wish to find a schedule for S that minimizes the total penalty incurred for missed deadlines.

# Analysis & C++ Code

Same, just change the value from 1 to  $t_i$ .

$$dp[i][j] = egin{cases} dp[i-1][j], & ext{do not choose } i \ dp[i-1][j-t[i]] + w[i], & ext{choose } i \end{cases}$$

Time complexity:  $O(n \sum t_i)$ .

```
#include <iostream>
#include <algorithm>
using namespace std;
int main()
{
    int n;
    cin >> n;
    int d[n+1], t[n+1], w[n+1];
    int maxx = 0, total = 0, sumt = 0;
    for (int i = 1; i <= n; i++)
        cin >> d[i] >> t[i] >> w[i];
        sumt += t[i];
        total += w[i];
    int dp[n+1][n+1];
    for (int i = 0; i \leftarrow n; i++)
        for (int j = 0; j <= n; j++)
            dp[i][j] = 0;
    for (int i = 1; i <= n; i++)
        for (int j = 0; j \leftarrow sumt; j++)
             dp[i][j] = dp[i-1][j];
            if (j > 0 \text{ and } j \leftarrow d[i]) dp[i][j] = \max(dp[i][j], dp[i-1][j-t[i]] + w[i]);
    }
    for (int j = 0; j \leftarrow sumt; j++)
        maxx = max(maxx, dp[n][j]);
    cout << total - maxx;</pre>
    return 0;
```