

Plaintext P : 110101100101000

Key k : 010010101110101

$$w_0 = 01001010$$

$$w_1 = 11110101$$

$$\begin{aligned}w_2 &= w_0 \text{ XOR } 10000000 \text{ XOR SubNib (rot.Nib}(w_1)) \\&= 01001010 \text{ XOR } 10000000 \text{ XOR SubNib}(01011111) \\&= 11001010 \text{ XOR SubNib}(01011111) \\&= 11011101\end{aligned}$$

$$\begin{aligned}w_3 &= w_2 \text{ XOR } w_1 \\&= 11011101 \text{ XOR } 11110101 \\&= 00101000\end{aligned}$$

$$\begin{aligned}w_4 &= w_2 \text{ XOR } 00110000 \text{ XOR SubNib (Rot Nib}(w_2)) \\&= 11011101 \text{ XOR } 00110000 \text{ XOR SubNib}(11000010) \\&= 10000111\end{aligned}$$

$$\begin{aligned}w_5 &= w_4 \text{ XOR } w_3 \\&= 10000111 \text{ XOR } 00101000 \\&= 10101111\end{aligned}$$

Now the subkey are :

$$\begin{aligned}\text{key 0} &= w_0 w_1 \\&= 010010101110101\end{aligned}$$

$$\begin{aligned}\text{key 1} &= w_2 w_3 \\&= 1101110100101000\end{aligned}$$

$$\begin{aligned}\text{key 2} &= w_4 w_5 \\&= 1000011110101111\end{aligned}$$



## Encryption

Plaintext XOR Keys

1101 0111 0010 1000

0100 1010 1111 0101

$\oplus$  1001 1101 1101 1101

Round 1:

Input = 1001 1101 1101 1101

output = 0010 1110 1110 1110

= 0010 1110 1110 1110

$M_e = \begin{matrix} 0010 & 1110 \\ 1110 & 1110 \end{matrix} = \begin{matrix} S_{00} & S_{01} \\ S_{10} & S_{11} \end{matrix}$

$S = M_e \times S$

$$\begin{aligned} S_{00} &= 0010 \oplus (4 \times 1110) \\ &= 0010 \oplus (4 \times E) \\ &= 0010 \oplus 1101 \\ &= 1111 \end{aligned}$$

$$\begin{aligned} S_{10} &= (4 \times 0010) \text{ XOR } 1110 \\ &= 1000 \text{ XOR } 1110 \\ &= 0110 \end{aligned}$$

$$\begin{aligned} S_{01} &= 1110 \oplus (4 \times 1110) \\ &= 1110 \oplus (4 \times E) \\ &= 0011 \end{aligned}$$



$$\begin{aligned}
 S_{11} &= (4 \times 1110) \oplus 1110 \\
 &= 1101 \oplus 1110 \\
 &= 0011
 \end{aligned}$$

$$\begin{aligned}
 \text{output} &= S_{00} \ S_{01} \ S_{02} \ S_{11} \\
 &= 1111 \ 0110 \ 0011 \ 0011
 \end{aligned}$$

Add round 1 key:

$$\begin{aligned}
 &= 1111 \ 0110 \ 0011 \ 0011 \\
 &1101 \ 1101 \ 0010 \ 1000 \\
 &\oplus 0010 \ 1011 \ 0001 \ 1011
 \end{aligned}$$

Final round:

Nibble substitution (S-boxes)

$$= 1010 \ 0011 \ 0100 \ 0011$$

Shift row (2<sup>nd</sup> and 4<sup>th</sup>)

$$= 1010 \ 0011 \ 0100 \ 0011$$

Add Round 2 key:

$$\begin{aligned}
 &1010 \ 0011 \ 0100 \ 0011 \\
 &\oplus 1000 \ 0111 \ 1010 \ 1111 \\
 &= 0010 \ 0100 \ 1110 \ 1100
 \end{aligned}$$

Now we can the final cipher text

$$\text{cipher text} = 0010 \ 0100 \ 1110 \ 1100$$



## Decryption

Add Round 2 key

$$\begin{array}{cccc} 0010 & 0100 & 1110 & 1100 \\ \oplus 1000 & 0111 & 1016 & 1111 \\ \hline = 1010 & 0011 & 0100 & 0011 \end{array}$$

Inverse shift row (Same as normal)

$$= 1010 \ 0011 \ 0100 \ 0011$$

Inverse Nibble Sub (use the inverse of decryption)

$$= 0016 \ 1011 \ 0001 \ 1011$$

Add Round 1 key

$$\begin{array}{cccc} = 0010 & 1011 & 0001 & 1011 \\ \oplus 1101 & 1101 & 0010 & 1000 \\ \hline 1111 & 0110 & 0011 & 0011 \end{array}$$

Inverse mix columns

$$S = \begin{array}{cc} S_{00} & S_{01} \\ S_{10} & S_{11} \end{array}$$

$$= 1110 \ 0011$$

$$0110 \ 0011$$

$$S' = \begin{array}{cc} S_{00}' & S_{01}' \\ S_{10}' & S_{11}' \end{array}$$

$$\begin{array}{cc} = 9 \times S_{00} \text{ XOR } 2 \times S_{01} & 9 \times S_{01} \text{ XOR } 2 \times S_{10} \\ 2 \times S_{00} \text{ XOR } 4 \times S_{10} & 2 \times S_{01} \text{ XOR } 4 \times S_{11} \end{array}$$



$$\begin{aligned}
 S_{00}' &= (9 \times 1111) \text{ XOR } (2 \times 0110) \\
 &= 9 \times F \text{ XOR } 2 \times 6 \\
 &= F \text{ XOR } 6 \\
 &= 1110 \oplus 1100 \\
 &= 0010
 \end{aligned}$$

$$\begin{aligned}
 S_{10}' &= 2 \times 1111 \text{ XOR } 9 \times 0110 \\
 &= 2 \times F \text{ XOR } 9 \times 6 \\
 &= 1010 \text{ XOR } 1001 \\
 &= 1110
 \end{aligned}$$

$$\begin{aligned}
 S_{01}' &= 9 \times 0011 \text{ XOR } 2 \times 1001 \\
 &= 9 \times 3 \text{ XOR } 2 \times 9 \\
 &= 8 \text{ XOR } 6 \\
 &= 1000 \text{ XOR } 0110 \\
 &= 1110
 \end{aligned}$$

$$\begin{aligned}
 S_{11}' &= 2 \times 0011 \text{ XOR } 9 \times 0011 \\
 &= 1110
 \end{aligned}$$

$$\text{output} = 0010 \ 1110 \ 1110 \ 1110$$

Inverse shift row

$$= 0010 \ 1110 \ 1110 \ 1110$$

Inverse Nibble sub

$$= 1001 \ 1101 \ 1101 \ 1101$$

Add Round 1 Key:

$$= 1001 \ 1101 \ 1101 \ 1101$$

$$\oplus 0100 \ 1010 \ 1111 \ 0101$$

$$= 1101 \ 0110 \ 0010 \ 1000$$



Plaintext = 1101 0111 0010 1000

original = 1101 0111 0010 1000

The decryption worked