

CS271 WEEK 5 TUES LECTURE + THURS

Modularization

- Breaking a program into segments that can be solved independent.
- Main procedure becomes just a caller of functions/Processes.

Processes

- Pushes next instruction offset to ~~EIP~~ when Call called, to stack
- Popped by "ret" into EIP.

Programming Process w/ Procedures

- Write main and procedure stubs
→ test!
- declare variables
→ test!
- Implement procedures 1 at a time
→ test each separately
→ Implement: Out, In, Proc...
- Clean up.

System Stack

- Stack is Data Structure
- LIFO \equiv FILO
- Operations: push, pop
- Managed by CPU,
- Uses registers ESP (top of stack) and SS (stack segment)

Pop and Push

- Push decrements ESP by $\overset{\text{for 32-bit}}{4}$,
copies value into mem location

where ESP points to. $\left\{ \begin{array}{l} [reg] = \text{value} \\ \text{of register} \\ @ \text{address} \end{array} \right.$

- Pop copies value pointed to by ESP into DEST.
→ increments ESP by 4

Documenting a Procedure

- Description
- Preconditions
- Relieves
- Registers modified
- Returns

Parameters : Arguments

- "Argument" (Actual Parameter) is the value that is being passed.
- "Parameter" (Formal Parameter) is the value received by the proc

Parameter Passing

- Input Parameter $\hat{=}$ "Value" or "pass by value"
- Output Parameter $\hat{=}$ "reference" or "pass by reference". Passed using an address (using offset). Used only as return space for output.
- I/O Parameter = pass by ref, something that is both used and returned by Process
→ edited/modified.

Ways of Passing Parameters

→ Global Vars → Registers → Stack.

- Global Variable Parameters
 - Bad Idea
- Register Passing Parameters
 - Bad Idea, but used quite often. Very simple.
- Passing on System Stack
 - push then pop to reg
 - Stack frames
 - Better for system resources