

Converse: $p \rightarrow q \Rightarrow q \rightarrow p$ Inverse: $p \rightarrow q \Rightarrow \neg p \rightarrow \neg q$
Contrapositive: $p \rightarrow q \Rightarrow \neg q \rightarrow \neg p \equiv p \rightarrow q$ DeMorgan's:
 $\neg(p \wedge q) = \neg p \vee \neg q$ Set Equality: $A = B$ iff $A \subseteq B \wedge B \subseteq A$
Subset: $A \subseteq B$ if $x \in A \rightarrow x \in B$ Proper Subsets: $A \subseteq B \wedge A \neq B$ Power Set $|P(A)| = 2^{|A|}$ Set Union: combine both sets
Set Intersection: Things in both sets Disjoint

Sets: $A \cap B = \emptyset$ Difference of sets: $A - B$ or A / B remove all elements of B in A . Not for sets: \bar{A} = all not in A . req: $D \subseteq D$. DeMorgan Set: $A \cup B = \overline{A \cap B}$ Distributive Laws $A \cap (B \cup C)$

$= (A \cap B) \cup (A \cap C)$ Contradict $p \rightarrow q$ Contrapos $p \rightarrow q$ S.1] $\neg p \rightarrow \neg q$

PMI 1] PS PMI 2] Suppose $p \wedge \neg q$ 3] Suppose $\neg q \equiv T$ S] $p \rightarrow q$

2] Base Case 3] Unpack Definition 4] show $\neg p$ follows

5] I. H. 1] see ex 2] win contra

Possible Subgraphs for

1] PMI: $\forall n \geq 1 \quad B \cap \left(\bigcup_{k=1}^n A_k \right) = \bigcup_{k=1}^n (B \cap A_k)$

$= \bigcup_{k=1}^n (B \cap A_k)$ // Iter 4: // eval $n+1$

2] B.C. $n=1 \quad B \cap A_k = B \cap A_k$ and $S = \emptyset \Rightarrow L(z) = \text{path len}$

3] I.H: $B \cap \left(\bigcup_{k=1}^n A_k \right) = \bigcup_{k=1}^n (B \cap A_k)$ while $z \notin S$

4] $B \cap \left(\bigcup_{k=1}^{n+1} A_k \right)$ // Pop off top $u = \alpha$ vertex not in S with $L(u)$ minimal

$= B \cap \left(\bigcup_{k=1}^n A_k \cup A_{n+1} \right)$ // Apply // $S := S \cup \{u\}$

$= \left(B \cap \bigcup_{k=1}^n A_k \right) \cup (B \cap A_{n+1})$ for all vertices not in S :

$= \left(\bigcup_{k=1}^n (B \cap A_k) \right) \cup (B \cap A_{n+1})$ Add a vertex to S with minimal label, and update Lb 's of vertices not in S

$= \left(\bigcup_{k=1}^{n+1} (B \cap A_k) \right)$ 5] Boom

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Arrange n distinct items: $n!$ Arrange n with r repetition: n^r
Permutations: $P(n, r) = \frac{n!}{(n-r)!}$, ordered. Combination: $\frac{n!}{r!(n-r)!}$, unordered. The number of subsets of length k formable. Coefficients: $\binom{n}{k}$ Power choose term st. $\in \mathbb{Q}$. Pascal's Identity $\binom{n+1}{k} = \binom{n}{k-1} + \binom{n}{k}$
Multisets: How many of length n = $\binom{n+k-1}{k}$ Graphs: (V, E)
Select k from n distinct

Graph Types and Descriptions			Graph	Edges	Mult. E. allow	Loop
ordered	Unordered	Repeats	simple	UNDIR	N	N
n^k	$\binom{n+k-1}{k}$	not dist	multigr.	UNDIR	Y	N
$P(n, k)$	$\binom{n}{k}$	distinct no rep	Pseudogr	UNDIR	Y	Y
			Sim. Dir	DIR	Y indiff dir	N
			Mult. dir	DIR	Y	Y
			Mixte	*	*	*

Neighborhood: all vertices conn'd to vertex. Denoted $N(v)$.
 $N(\text{set}) = \bigcup_{v \in A} N(v)$

Deg: number of edges leading into a vertex. In degree: $\deg^-(v)$ out degree: $\deg^+(v)$

Note: $|E| = \sum \deg^-(v \in V) = \sum \deg^+(v \in V)$

Complete Graph: K_n , all vertices connected, $\sum_{k=0}^{n-1} \binom{n-1}{k}$. Cycle Graph: n vertices, n edges. Wheel: put a vertex in center of C_n , connect to all vertices. Handshake:

$\sum_{v \in V} \deg(v) = 2|E|$ Bipartite: color red/blu so no red touch red or blue touch blue

Subgraph: $G = (V, E)$, s.g. $= (W, F) \mid F \subseteq E \wedge W \subseteq V$ Induced: only incl edges between selected subg such that only edges

Adjacency Matrix: make matrix, and put 1 if line connects vertices.

Ring Permutations: $\frac{P(n, k)}{k}$. 4 people, 4 seats: $4!/4$ Dijkstra: 1) start at point, for all connected points, add weight to current point, put in box. 2) take least valued point not already visited. 3) rep till end.

There is 1 pers every loves: $L(x, y) = x \perp y : \exists x (A_y(L(x, y)) \wedge \forall z (A_z(L(x, z)) \rightarrow z = y))$