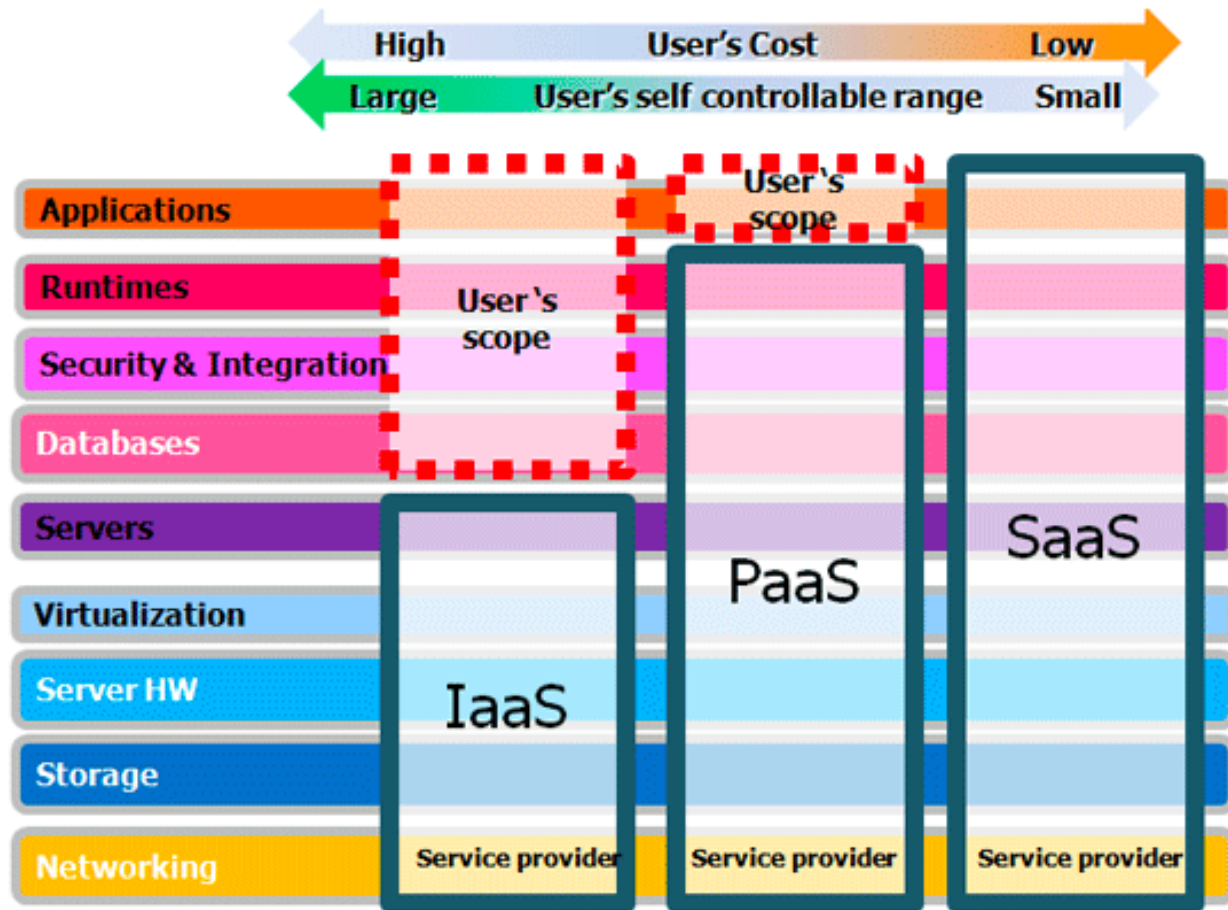


# Object-Oriented Programming Programming Project #2

郭建志

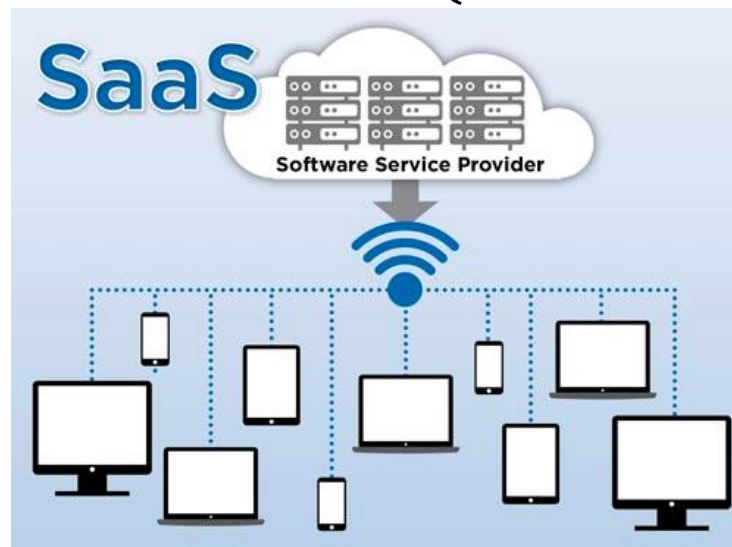
# Background

- Cloud services can be broadly divided into main three models
- Infrastructure-as-a-Service (**IaaS**)
- Platform-as-a-Service (**PaaS**)
- Software-as-a-Service (**SaaS**)



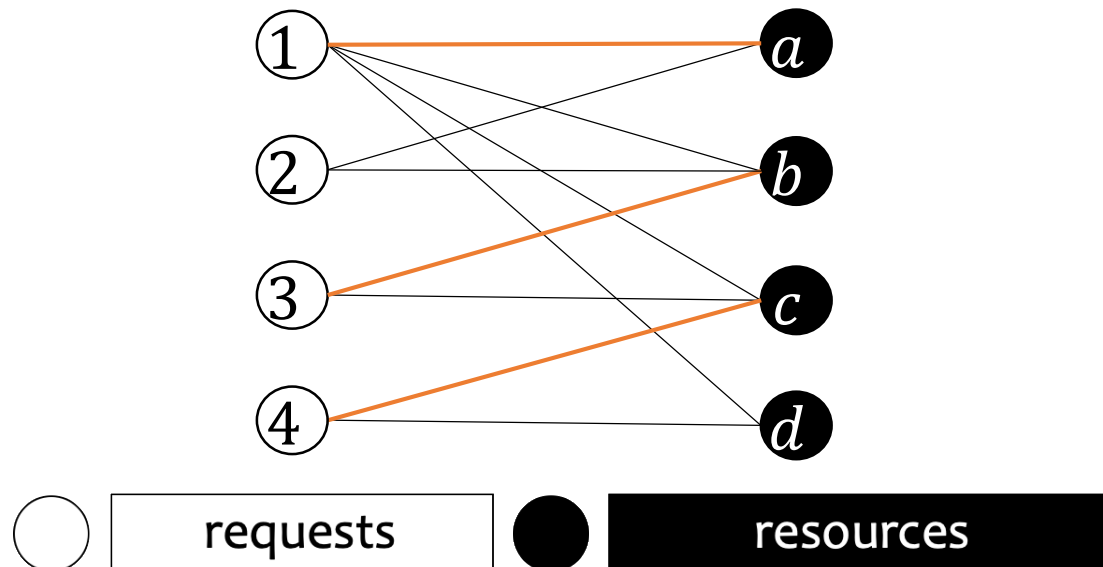
# Background

- Service providers have resources (such as virtual infrastructures, virtual platform, service software) and users request the resources for services
- For example, Google constructs data centers and maintain the hardware and software to provide the mail service for users (i.e., Gmail)



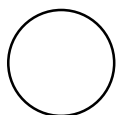
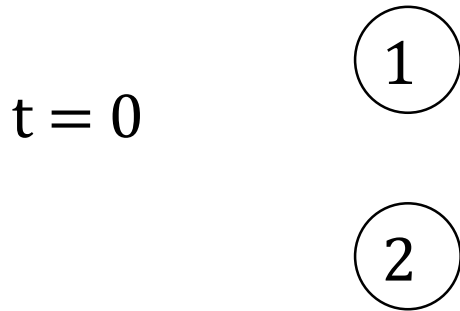
# Background

- Resource allocation can be modeled as a bipartite matching problem
- But resource allocation is usually **dynamic**
- Some request arrives and some resources are released at a time

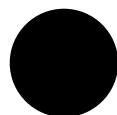


# Online Bipartite Matching Problem (BMP)

- Consider a scenario, where some request arrives and some resources are released at a time:



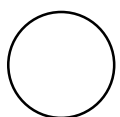
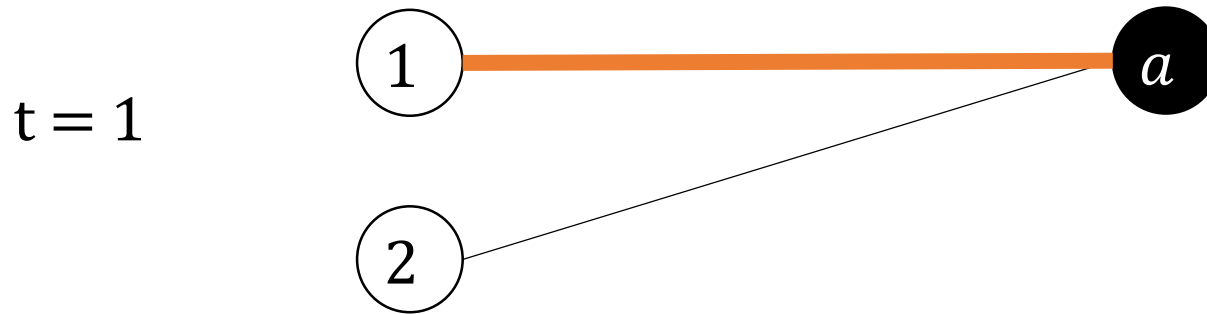
requests



resources

# Online Bipartite Matching Problem (BMP)

- Consider a scenario, where some request arrives and some resources are released at a time:



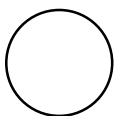
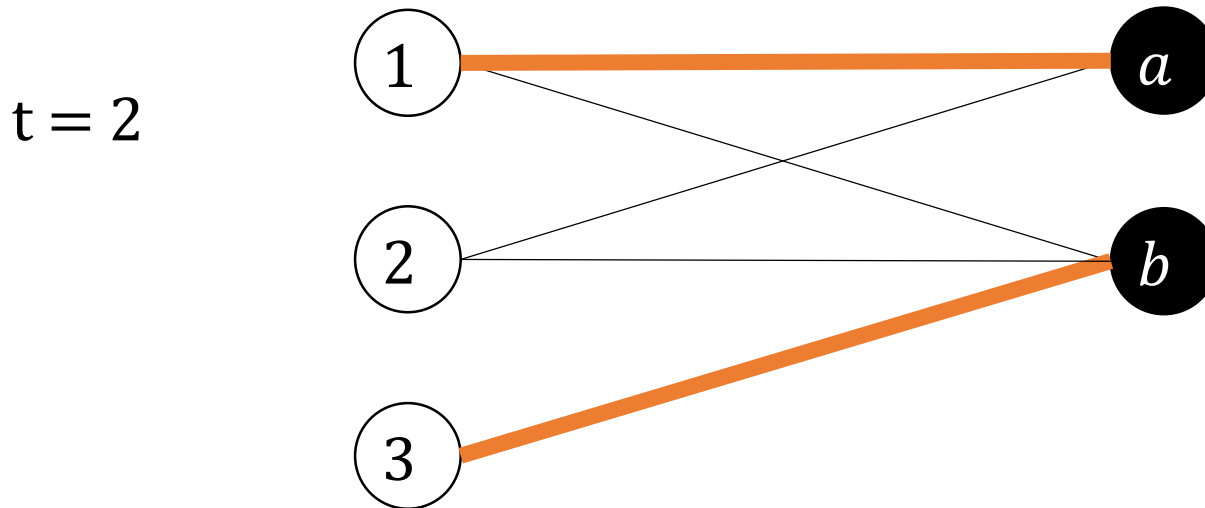
requests



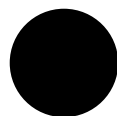
resources

# Online Bipartite Matching Problem (BMP)

- Consider a scenario, where some request arrives and some resources are released at a time:



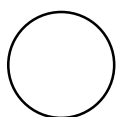
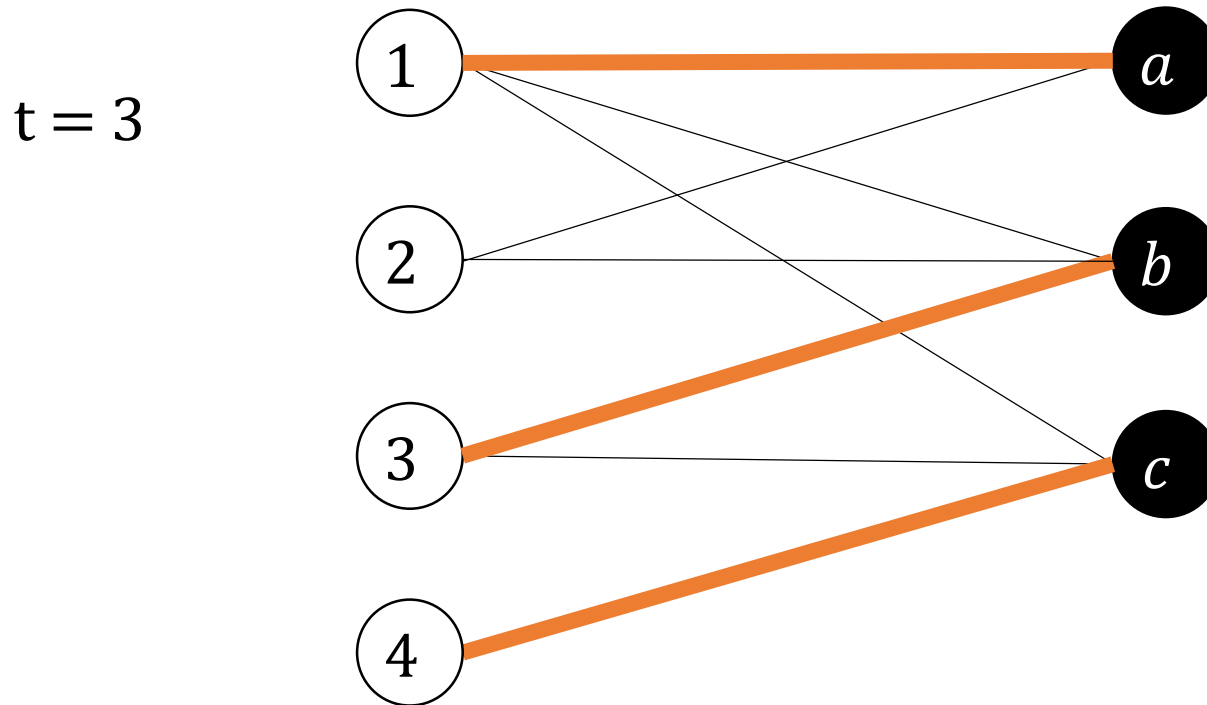
requests



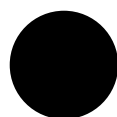
resources

# Online Bipartite Matching Problem (BMP)

- Consider a scenario, where some request arrives and some resources are released at a time:



requests

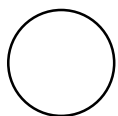
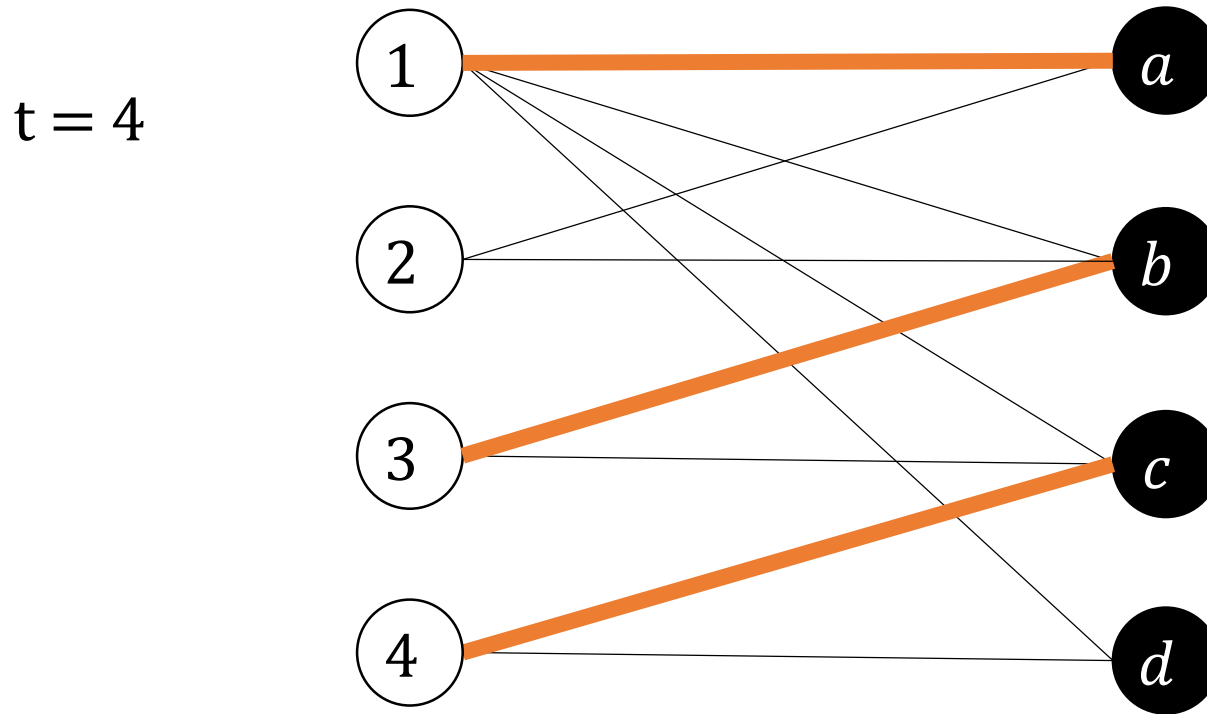


resources

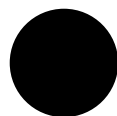


# Online Bipartite Matching Problem (BMP)

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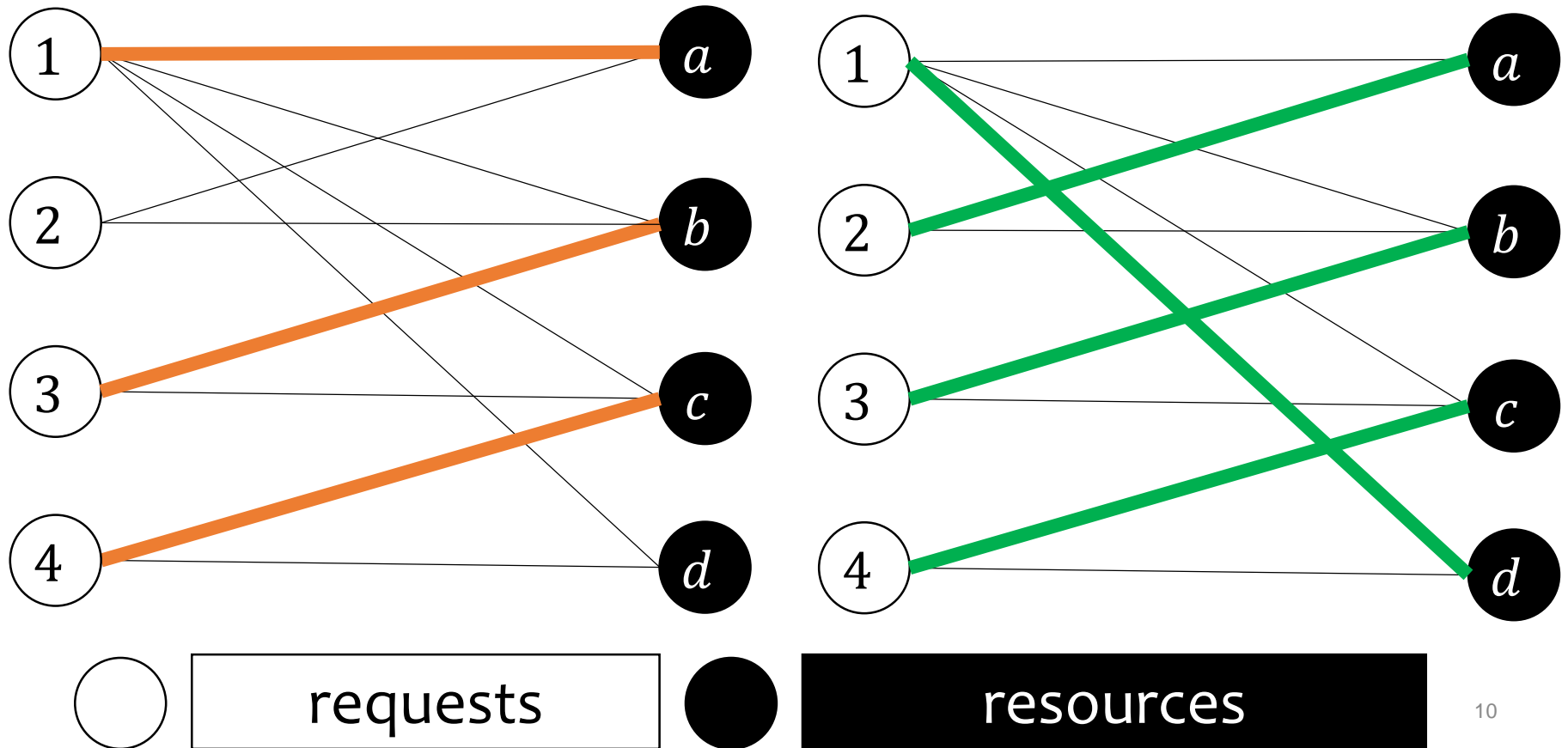
requests



resources

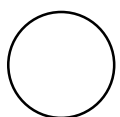
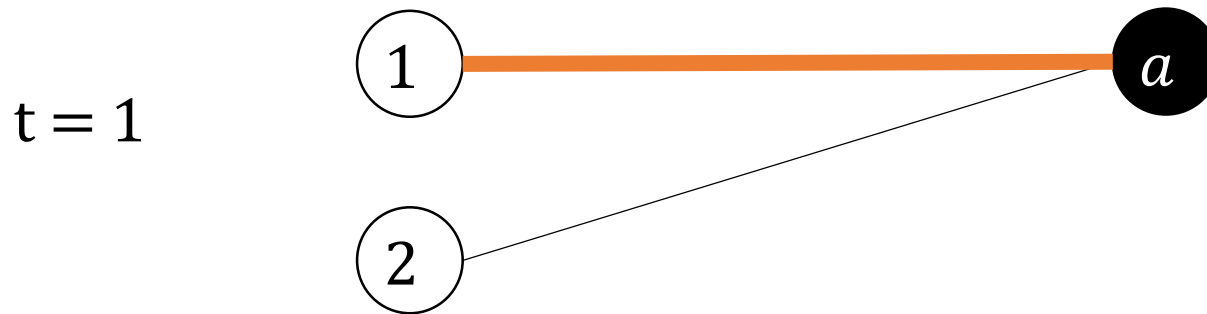
# Online Bipartite Matching Problem (BMP)

- Compared with the offline optimal solution



# Hardness of the Online BMP

- If your algorithm is deterministic



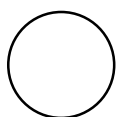
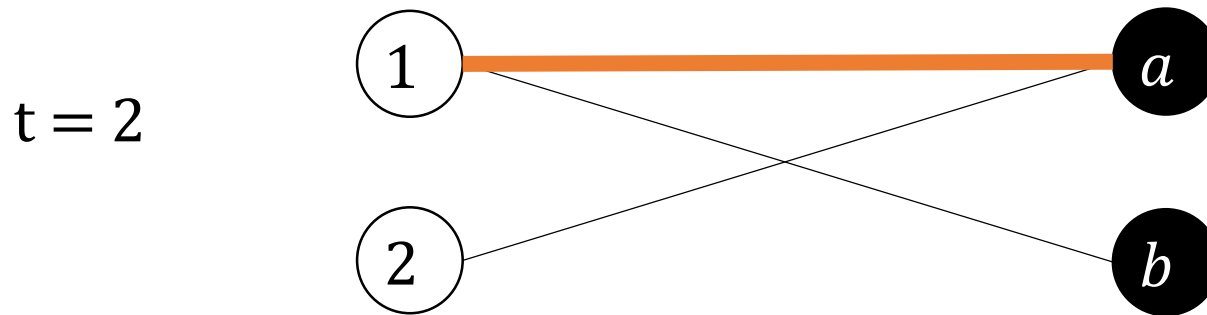
requests



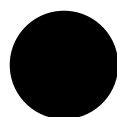
resources

# Hardness of the Online BMP

- If your algorithm is deterministic



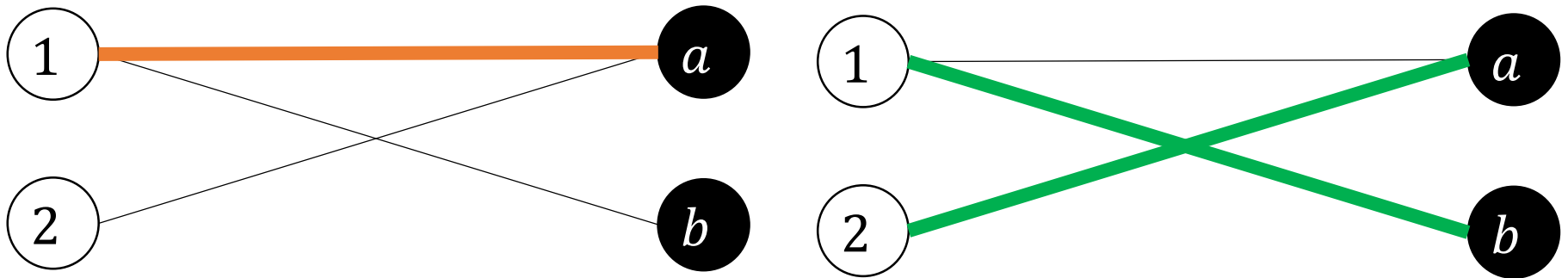
requests



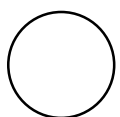
resources

# Hardness of the Online BMP

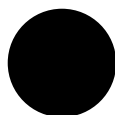
- Compared with the offline optimal solution



- No deterministic algorithm can bound the ratio larger than  $1/2$



requests



resources

# Programming Project #2:

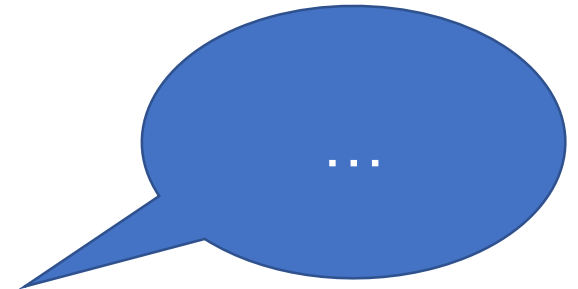
## Online resource allocation

- Input:
  - # time slots
  - Time slot ID, # requests, # resources, request IDs, and resource IDs at the time slot
  - Number of edges and edges between a request ID and a resource ID at the time slot
- Procedure:
  - Assign the resource to a request at the time slot when the resource is released
- Output:
  - # assigned pairs
  - The time slot
- The grade is proportional to # assigned pairs

# Programming Project #2:

## Online resource allocation

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老師又要  
騙我了嗎



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沒錯!

# Programming Project #2:

## Online resource allocation

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- Procedure:
  - Assign the resource to a request at the time slot when the resource is released
- Output:
  - # assigned pairs
  - The time slot
- Implement a designated randomized algorithm

# Programming Project #2:

## Online resource allocation

- **Designated Algorithm:**

1. For each time slot:
2. Uniform-randomly assign a value in  $[0, 1]$  for each arrival requests
3. Assign the available resource with the smallest ID to the unsatisfied request with the maximum random value
4. Repeat step 3 until there is no available resource
5. Go to step 1 for the next time slot

- **Note:**

1. Use default\_random\_engine and uniform\_real\_distribution
2. Let the random seed be a parameter of main function
3. Define classes Resource and Request, and use vector to store the objects of class Resource and Request
4. Generate new resource and request by push\_back
5. Design static member variables and static member functions to count and get the number of satisfied requests

# Programming Project #2:

## Online resource allocation

```
// You need forward declaration
class Resource{
    int id;
    bool matched;
    int requestId;
public:
    // constructors...
    bool operator->* (Request &b){...}
    // ->* check whether both of them are un-matched.
    // If they are, then match them and return true;
    // otherwise, return false
    // accessor
};
class Request{
    int id;
    bool matched;
    int resourceId;
    double weight; // a random value, indicates the priority
public:
    // constructors...
    friend class Resource;
    // accessor
};
```

# Discussion

- Why does the randomized algorithm work?
- What can be added to improve the performance?
- Discussion & bonus

# Discussion

- Why does the randomized algorithm work?
- It can avoid **being stuck in the trap of adversary**
- Adversary: deliberately choose difficult data to maximize the gap between the generated solution and the offline optimum
- What can be added to improve the performance?
- Distribution of the input
- Discussion & bonus

## Further Reading

- Randomized primal-dual analysis of RANKING for online bipartite matching, in SODA 2013
- Online stochastic matching beating  $1-1/e$ , in FOCS 2009
- ...
- 演算法頂尖研討會: STOC, FOCS, SODA...
- 一個字，神

# Input Sample: request.txt

Format:

#timeSlots

|             |           |            |
|-------------|-----------|------------|
| timeSlotID1 | #requests | #resources |
|-------------|-----------|------------|

|            |            |     |
|------------|------------|-----|
| requestID1 | requestID2 | ... |
|------------|------------|-----|

|             |             |     |
|-------------|-------------|-----|
| resourceID1 | resourceID2 | ... |
|-------------|-------------|-----|

#edges

|         |           |            |
|---------|-----------|------------|
| edgeID1 | requestID | resourceID |
|---------|-----------|------------|

|         |     |
|---------|-----|
| edgeID2 | ... |
|---------|-----|

timeSlotID2...

之後揭曉，先自己設計input



# Output Sample: result.txt

Format:

#satisfiedRequests

requestID      resourceID

...

下次揭曉，用自己設計input產生答案

# Input Sample: request.txt

Format:

3  
0 5 2  
0 1 2 3 4  
0 1  
3  
0 3 0  
1 0 1  
2 4 1  
1 1 1  
5  
2  
2  
3 0 2  
4 4 2

2 3 3  
6 7 8  
3 4 5  
9  
5 0 3  
6 1 3  
7 3 3  
8 6 3  
9 8 3  
10 5 4  
11 2 5  
12 4 5  
13 5 5

Format:

#timeSlots  
timeSlotID1 #requests #resources  
requestID1 requestID2 ...  
resourceID1 resourceID2 ...  
#edges  
edgeID1 requestID resourceID  
edgeID2 ...  
timeSlotID2...

# Output Sample: result.txt

Format:

6  
0 2  
2 5  
3 0  
4 1  
5 4  
8 3

Format:

#satisfiedRequests  
requestID resourceID  
...

# Note

- Deadline: 4/16 Tue
- 小老師deadline: 4/2 Tue (暫定)
- E-course
- **C++ Source code**
- Show a good programming style