Object-Oriented Programming Programming Project #1

郭建志

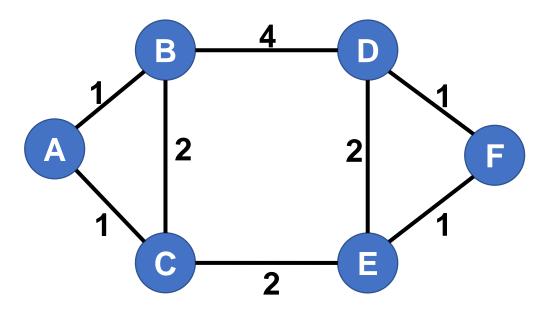
Background

- Local area network within a data center
- Routing efficiency is critical to system performance and quality of service



The Routing Problem

• Consider a scenario, where each request has a different size:

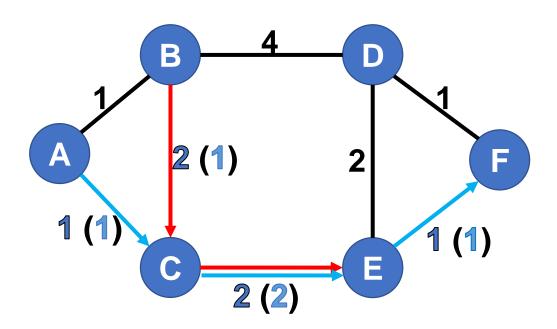


Requests

(src, dest; demand):

- (A, F; 1)
- (A, F; 2)
- (B, E; 1)
- (B, E; 3)
- (B, E; 4)
- (B, C; 1)
- (B, D; 3)
- (A, B; 4)
- (A, C; 2)
- (D, F; 4)
- (E, F; 5)
- (B, C; 6)
- (B, D; 2)
- (B, D; 1)
- (B, D; 2)
- (C, E; 3)
- (C, E; 1)

One Routing Solution



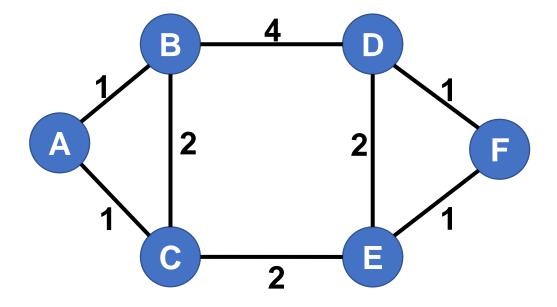
Requests

(src, dest; demand):

- (A, F; 1)
 - (A, F; 2)
- **(B, E; 1)**
 - (B, E; 3)
 - (B, E; 4)
 - (B, C; 1)
 - (B, D; 3)
 - (A, B; 4)
 - (A, C; 2)
 - (D, F; 4)
 - (E, F; 5)
 - (B, C; 6)
 - (B, D; 2)
 - (B, D; 1)
 - (B, D; 2)
 - (C, E; 3)
 - (C, E; 1)

The Online Routing Problem

• The requests are arrived in an online fashion.



Requests (src, dest; demand): (A, F; 1) (A, F; 2) (B, E; 1) (B, E; 3) (B, E; 4) (B, C; 1) (B, D; 3) (A, B; 4) (A, C; 2) (D, F; 4) (E, F; 5) (B, C; 6) (B, D; 2)

(B, D; 1)

(B, D; 2)

(C, E; 3)

(C, E; 1)

• Input:

- Numbers of nodes and undirected links
- Undirected links with non-negative link capacities
- Requested flows with their sources, destinations, and sizes

Procedure:

- Accept or reject the flow one by one
- If accepted, then assign a path

• Output:

- Number of accepted flows
- Total throughput
- The assigned paths of each accepted flow
- The grade is proportional to the total throughput

Input file:

5	10		
0	0	1	14
1	0	2	12
2	0	3	9
0	0	1	8

...

Output file:

15	124	
0	0	1
1	2	0

...

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Input file:

10

Э	10		
0	0	1	1
1	0	2	1
2	0	3	9
0	0	1	8
1	2	1	6

Output file:

15 0	124 0	1	
1	2	0	1
1	4	<u> </u>	1

怎麼辦

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- Numbers of nodes and undirected links
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Procedure:

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Input file:

10

3	10		
0	0	1	1
1	0	2	1
2	0	3	9
•••			
0	0	1	8
1	2	1	6

Output file:

15	124		
0	0	1	
1	2	0	

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Input:

- Numbers of nodes and undirected links
- Undirected links with non-negative link capacities
- Requested flows with their sources, destinations, and sizes

Procedure:

- Accept or reject the flow one by one
- If accepted, then assign a path

• Output:

- Number of accepted flows
- Total throughput
- The assigned paths of each accepted flow
- Implement a designated algorithm to select the flows

Input file:

5	10		
0	0	1	14
1	0	2	12
2	0	3	9
0	0	1	8

...

Output file:

15	124	
0	0	1
1	2	0

..

Designated Algorithm:

- Initially, cost(e) = the minimum of double for each edge e
- Examine the current requested flow *f* with size size(f) and find the shortest path P (i.e., the path with the smallest sum of edge costs)
- 3. If all edges on the path *P* can accommodate this flow (i.e., load(e) + $size(f) \leq capacity(e)$), then accept the flow *f*; otherwise, reject it and go to step 5
- 4. After accepting a flow, update load(e) += size(f) and then cost(e) = $\frac{load(e)}{capacity(e)-load(e)}$ for each edge $e \in P$
- Go to step 2 and examine the next requested flow until no next one exists

Input file:

10

0 14 12 9

8

Output file:

15 124

Discussion

- Why does the algorithm work?
- It chooses the path that avoids the link with a heavy load
- What can be added to improve the performance?
- Discussion & bonus

Discussion

- Why does the algorithm work?
- It chooses the path that avoids the link with a heavy load
- What can be added to improve the performance?
- Discussion & bonus
- Design a new edge cost formula?
- Add an acceptance condition?

Further Reading

- Section 9.1 in "The Design of Competitive Online Algorithm via a Primal-Dual Approach", Foundations and Trends in Theoretical Computer Science, 2009.
- Dynamic Routing for Network Throughput Maximization in Software-Defined Networks, in INFOCOM 2016
- ...
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Input Sample: request.txt

Format:

```
#nodes #undirectedLinks
#linkID #firstNode #secondNode #linkCapacity
...
#requestFlows
#flowID #sourceID #destinationID #flowSize
...
```

下次揭曉,先自己設計input

Output Sample: result.txt

Format:

```
#acceptedFlows totalThroughput
#flowID #firstNode #secondNode ... #lastNode
```

下次揭曉,用自己設計input產生答案

Input Sample: request.txt

9

#nodes #undirectedLinks #linkID #firstNode #secondNode #linkCapacity

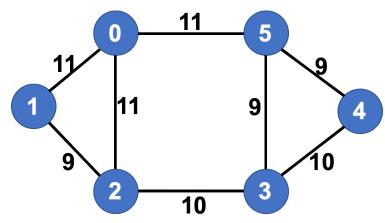
Format:

#requestFlows

#flowID #sourceID #destinationID #flowSize

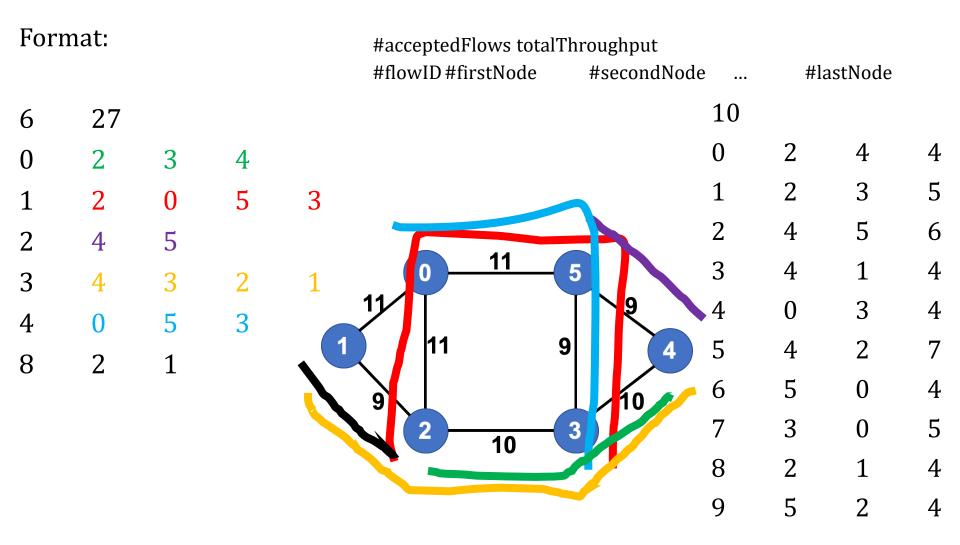
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6	8		
0	0	1	11
1	0	2	11
2 3	0	5	11
3	1	2	9
4	2	3	10
5	3	4	10
6	3	5	9
7	4	5	9



Output Sample: result.txt

Format:



Note

- Deadline:3/26 Tue
- E-course

• C++ Source code

Show a good programming style