
CS161 WEEK3 DISCUSSION 1C

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Contributed by Yewen W' and Shirley C's previous course materials

Agenda

UNIFORM-COST SEARCH

INFORMED SEARCH

A* SEARCH

(CSP NEXT WEEK)

UNIFORM-COST SEARCH

- Search the node(state) that has the smallest cost
- Update the distance record for neighboring nodes
- Keep searching even after a goal is found

Algorithm 4: Uniform-Cost Algorithm

Data: A weighted directed graph $G = (V, E)$ and a source vertex s

Result: A distance map where $dist[v]$ is the shortest path weight between s and v

Function UniformCost(G, s):

```
     $dist[s] = 0$ ;  
    Create an empty vertex set  $Q$ ;  
    Add  $s$  into  $Q$  with key  $dist[s]$ ;  
    while  $Q \neq \emptyset$  do  
        Find vertex  $u$  in  $Q$  whose  $dist[u]$  is the smallest;  
        Remove  $u$  from  $Q$ ;  
        foreach  $v \in Adj(u)$  do  
             $dist = dist[u] + w(u, v)$ ;  
            if  $v \in Q$  then  
                if  $dist < dist[v]$  then  
                     $dist[v] = dist$ ;  
                    Update  $Q$  on vertex  $v$  with new  $dist[v]$  value;  
                end  
            else  
                 $dist[v] = dist$ ;  
                Add  $v$  into  $Q$  with key  $dist[v]$ ;  
            end  
        end  
    end  
    return  $dist$ ;
```

UNIFORM-COST SEARCH

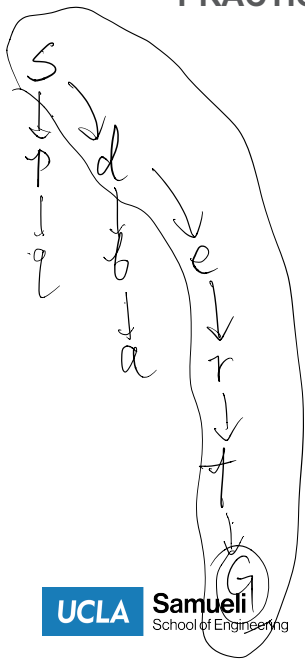
EVALUATION

- Optimal? T
- Complete? T
- Time $O(b^{C/\epsilon})$
- Space $O(b^{C/\epsilon})$

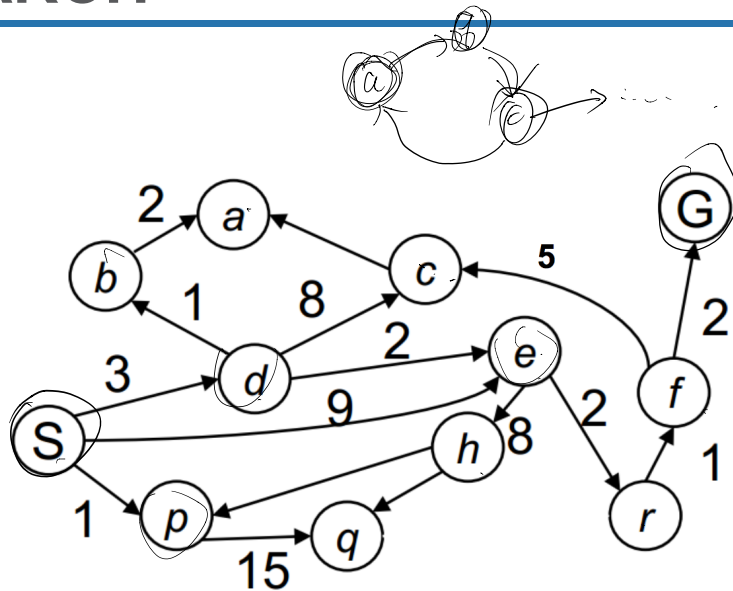
C : Total cost

ϵ : minimum cost at each step

PRACTICE

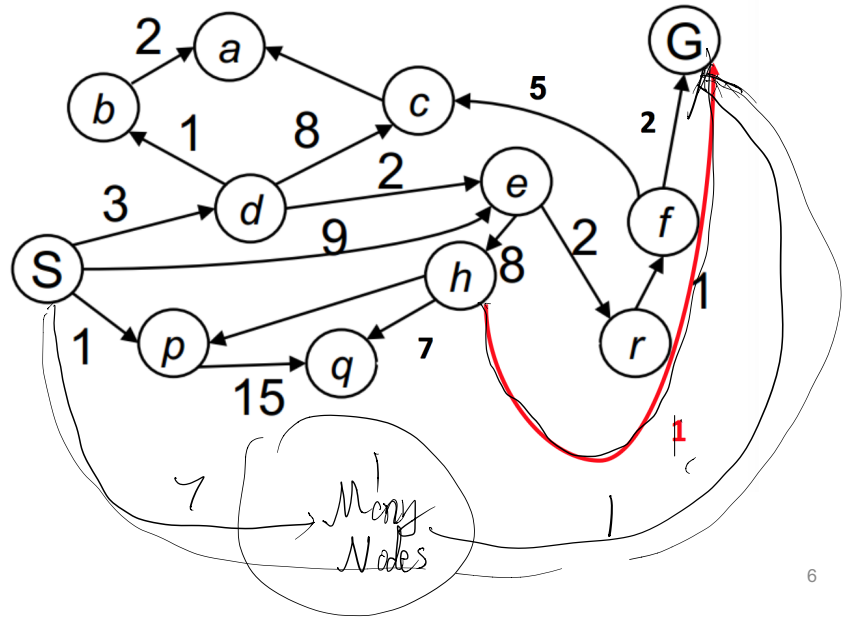


\sqrt{ab}
 $\sqrt{b^4}$
 $\sqrt{c^{11}}$
 $\sqrt{d^3}$
 $\sqrt{e^9} \rightarrow 5$
 $\sqrt{f^8}$
 $\sqrt{h^{13}}$
 $\sqrt{p^{11}}$
 $\sqrt{q^{15}}$
 $\sqrt{r^7}$
 G



UNIFORM-COST SEARCH

PRACTICE



UNIFORM-COST SEARCH

Uniform-cost search is similar to Dijkstra Algorithm

*We must know
the graph*

The difference is that Dijkstra keeps the record of all nodes in the graph.

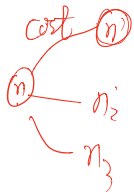
CA → DC

GREEDY BEST-FIRST SEARCH

- Expand the node that is the closest to the goal
- 'closest' is defined by the heuristic function
- Complete? \mathcal{N}
- Optimal? \mathcal{N}
- Time $\mathcal{O}(b^m)$
- Space $\mathcal{O}(b^m)$

A* SEARCH

- Define the total cost function $f(n) = g(n) + h(n)$
 - $g(n)$: path cost so far
 - $h(n)$: estimated cost to goal (heuristic)
- Properties of heuristic
 - Admissibility $h(n) \leq h^*(n)$
 - ($h^*(n)$ is the true cost from n to goal)
 - Consistency $h(n) \leq \text{cost}(n, n') + h(n')$
 - n' is the successor of n ; c is the cost of path from n to n' by choosing an action a



A* SEARCH

OPTIMALITY

If $h(n)$ is admissible, A* using tree-search is optimal

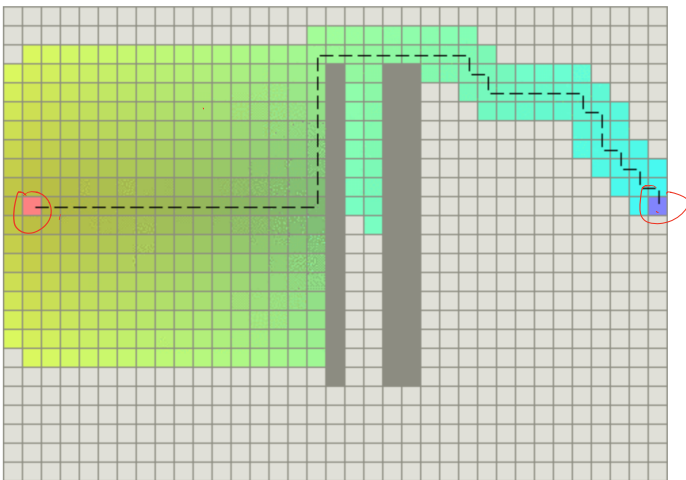
If $h(n)$ is consistent, A* using graph-search is optimal

→ whether we keep a record
of state/node visited before

A* SEARCH

PRACTICE

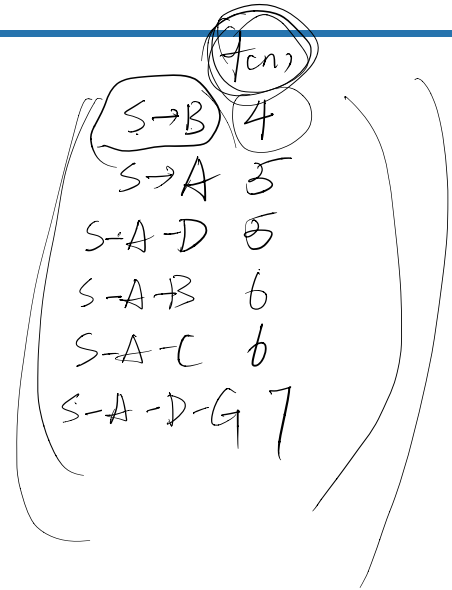
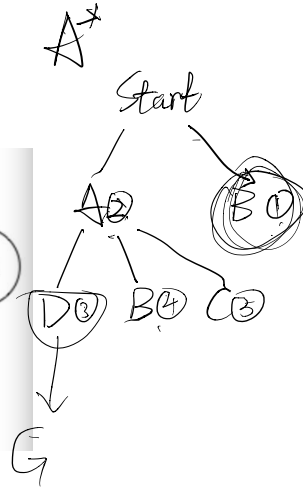
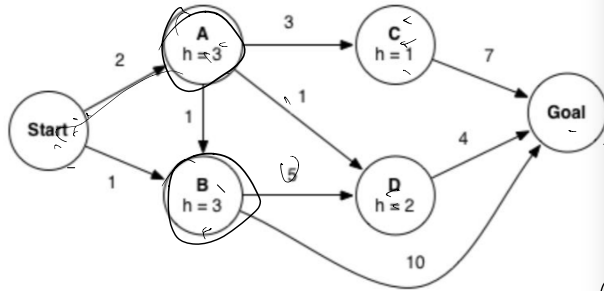
What are the possible heuristics for this question?



- ① Euclidean dist
- ② Manhattan dist
- ③ Drag

A* SEARCH

PRACTICE



A* SEARCH

EVALUATION

Complete? -Y

Time: $O(b^{\delta})$ where $\delta = \underbrace{h^* - h}$

Space: $O(b^d)$

HOMEWORK2

DFID

Three parts

DFS1: Perform DFS with depth d

DFS2: Start from a depth. Add 1 to depth each step and stop until it reaches d.

DFS3: Call DFS2 with start step = 0

```
(defun dfs1 (tree depth)
  (cond ((null tree) NIL)
        ((atom tree) (list tree))
        ((= depth 0) NIL)
        (t (append (dfs1 (car tree) (- depth 1))
                    (dfs1 (cdr tree) depth))))))
```

```
(defun dfs2 (tree depth maxdepth)
  (cond ((> depth maxdepth) NIL)
        (t (append (dfs tree depth) (dfs2 tree
                                             (+ depth 1) maxdepth))))))
```

HOMEWORK2

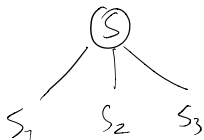
MISSIONARY-CANNIBAL PROBLEM

MC-DFS: Given a state

- If it is a goal, add it
- If it has been visited, ignore it
- Otherwise, explore it

MULT-DFS: Given expanded states

- If nothing to expand, ignore it
- Otherwise, call MC-DFS for each state



$(S_n S_{n-1} \dots S_0)$

$(S_0 \dots S_n)$

```

(defun mc-dfs (s path)
  (cond ((final-state s) (cons s path))
        ((on-path s path) nil)
        (t (mult-dfs (succ-fn s) (cons s path)))))
  
```

```

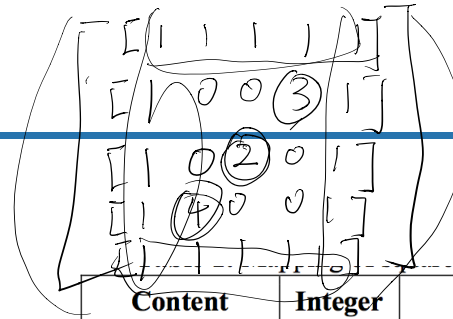
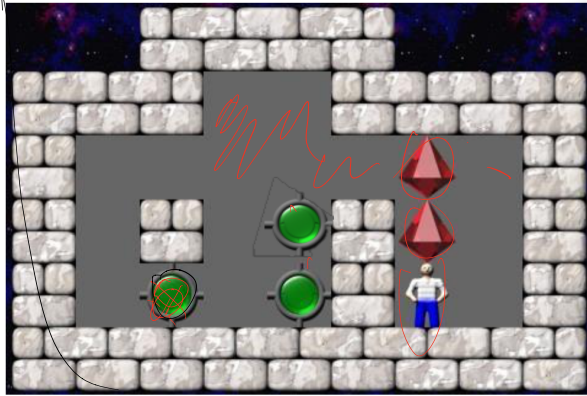
(defun mult-dfs (states path)
  (cond ((null states) nil)
        (t (or (mc-dfs (first states) path)
                 (mult-dfs (rest states) path)))))
  
```

nil
(search result of S_1) or nil result of S_2 or nil result of S_3 or nil

$(S_0 \rightarrow S_n)$
result of S_3

HOMWORK3

SOKOBAN



Content	Integer	ASCII
Blank	0	' ' (white space)
Wall	1	'#'
Box	2	'\$'
Keeper	3	'@'
Goal	4	'.'
Box + goal	5	'*'
Keeper + goal	6	'+'

HOMEWORK3

SOKOBAN

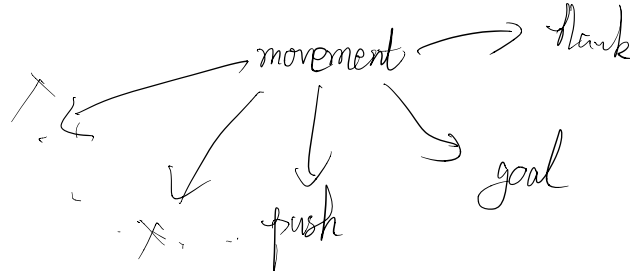
Break into parts

- Goal test
- Heuristic
- Action
 - Given current state
 - Move along a direction
 - Get the new state

Key function: try-move

You should consider all possible cases resulting from a move.

defun next-state ...



Q&A
