CS 524: Introduction to Optimization Lecture 35: Two stage SPs

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The problem

A furniture maker can manufacture and sell four different dressers. Each dresser requires a certain number t_{cj} of man-hours for carpentry, and a certain number t_{fj} of man-hours for finishing, $j=1,\ldots,4$. In each period, there are d_c man-hours available for carpentry, and d_f available for finishing. There is a (unit) profit \bar{c}_j per dresser of type j that's manufactured. The owner's goal is to maximize total profit:

$$\max_{x \ge 0} \ 12x_1 + 25x_2 + 21x_3 + 40x_4$$
 (profit)

subject to

$$4x_1 + 9x_2 + 7x_3 + 10x_4 \le 6000$$
 (carpentry)
 $x_1 + x_2 + 3x_3 + 40x_4 \le 4000$ (finishing)

Succinctly:

$$\max_{x} c^T x \text{ s.t. } Tx \leq d, x \geq 0$$

Is your time estimate that good?

Solution is (1333, 0, 0, 67), value \$18,667

- The time for carpentry and finishing for each dresser cannot be known with certainty
- Each entry in T takes on four possible values with probability 1/4, independently
- 8 entries of T are random variables: s = 65,536 different T's each with same probability of occurring
- But decide "now" how many dressers x of each type to build
- Might have to pay for overtime (for carpentry and finishing)
- Can make different overtime decision y_s for each scenario s recourse!

Creating a Stochastic Model

Here is a general procedure for making a (scenario-based) 2-stage stochastic optimization problem

- For a "nominal" state of nature (scenario), formulate an appropriate LP model
- Decide which decisions are made before uncertainty is revealed (here and now), and which are decided after
- All second stage variables get "scenario" index
- Constraints with scenario indices must hold for all scenarios
- Second stage variables in the objective function should be weighted by the probability of the scenario occurring

Extensive Form (Deterministic Equivalent)

- Assume $\Omega = \{\omega_1, \omega_2, \dots \omega_S\} \subseteq \mathbb{R}^r$, $\mathsf{P}(\omega = \omega_s) = p_s, \forall s = 1, 2, \dots, S$
- $T_s \stackrel{\text{def}}{=} T(\omega_s), d_s \stackrel{\text{def}}{=} d(\omega_s), q_s \stackrel{\text{def}}{=} q(\omega_s), W_S = W(\omega_s)$
- Then can write extensive form:

$$c^{T}x + p_{1}q_{1}^{T}y_{1} + p_{2}q_{2}^{T}y_{2} + \cdots + p_{s}q_{s}^{T}y_{s}$$
s.t.
$$Ax = b$$

$$T_{1}x + W_{1}y_{1} = d_{1}$$

$$T_{2}x + W_{2}y_{2} = d_{2}$$

$$\vdots + \vdots$$

$$T_{S}x = b$$

$$\vdots + W_{S}y_{s} = d_{s}$$

$$x \in X = y_{1} \in Y = y_{2} \in Y = y_{s} \in Y$$

Extensive Form Problem (furnrun.gms)

$$\min_{x,y} -c^{T}x + \sum_{s=1}^{65,536} p_{s}q^{T}y_{s}$$

subject to

$$T^s x - y_s \le d, \quad s = 1, \dots, 65, 536$$

 $x, y_s \ge 0$

- Immediate costs + expected future costs
- Stochastic program with recourse
- Solution is (205, 0, 696, 33), value \$18, 128

About the extensive form

- $y_s \equiv y(\omega_s)$ is the recourse action to take if scenario ω_s occurs.
- Pro: It's a linear program.
- Con: It's a BIG linear program.
- Imagine the following (real) problem. A Telecom company wants to expand its network in a way in which to meet an unknown (random) demand.
- There are 86 unknown demands. Each demand is independent and may take on one of seven values.
- $S = |\Omega| = \prod_{k=1}^{86} (7) = 7^{86} = 4.77 \times 10^{72}$
 - ▶ The number of subatomic particles in the universe.
- How do we solve a problem that has more variables and more constraints than the number of subatomic particles in the universe?

But Its Even Worse!

- If Ω is not a countable set (say if it is made up of continuous-valued random variables, our "deterministic equivalent" would have ∞ variables and constraints. :-)
- The answer is we can't!
- We solve an approximating problem obtained through sampling.

Monte Carlo Methods

(*)
$$\min_{x \in X} \{ f(x) \equiv \mathbb{E}_P F(x, \xi) \equiv \int_{\Omega} F(x, \xi) dP(\xi) \}$$

- Most of the theory presented holds for (*)—A very general SP problem
- Naturally it holds for our favorite SP problem:
 - $X \stackrel{\text{def}}{=} \{ x \mid Ax = b, \ x \ge 0 \}$
 - $f(x) \equiv c^{T}x + \mathbb{E}\{Q(x,\omega)\}\$
 - $Q(x,\omega) \equiv \min_{y \ge 0} \{ q(\omega)^T y | Wy = d(\omega) T(\omega)x \}$

Sampling

- Instead of solving (*), we solve an approximating problem.
- Let ξ^1, \dots, ξ^N be N realizations of the random variable ξ :

$$\min_{\mathbf{x}\in X}\{\widehat{f}_{N}(\mathbf{x})\equiv N^{-1}\sum_{j=1}^{N}F(\mathbf{x},\xi^{j})\}.$$

- $\widehat{f}_N(x)$ is just the sample average function
- Since ξ^j drawn from P, $\widehat{f}_N(x)$ is an unbiased estimator of f(x)
 - $\blacktriangleright \ \mathbb{E}[\widehat{f}_N(x)] = f(x)$

What do we learn? (furnsaa.gms)

- Deterministic solution: $x^{mvs} = (1333, 0, 0, 67)$
- Expected profit using this solution: \$16,942
- Expected (averaged) overtime costs: \$1,725
- Extensive form solution: $x^s = (257, 0, 666, 34)$ with expected profit \$18,051
- Deterministic solution is not optimal for stochastic program, but more significantly it isn't getting us on the right track!
- Stochastic solution suggests large number of "type 3" dressers, while deterministic solution has none!

Stochastic Programming

A Stochastic Program

$$\min_{x \in X} f(x) \stackrel{\text{def}}{=} \mathbb{E}_{\omega}[F(x,\omega)]$$

2 Stage Stochastic LP w/Recourse

$$F(x,\omega) \stackrel{\text{def}}{=} c^T x + Q(x,\omega)$$

- c^Tx : Pay me now
- $Q(x,\omega)$: Pay me later

The Recourse Problem

$$Q(x,\omega) \stackrel{\text{def}}{=} \min q(\omega)^T y$$

$$W(\omega)y = d(\omega) - T(\omega)x$$
$$y \ge 0$$

Specialized techniques use this recourse function for decomposition algorithms