



**NANYANG  
TECHNOLOGICAL  
UNIVERSITY**  
SINGAPORE

# Discrete Mathematics

## MH1812

### Topic 2.2 - Propositional Logic II

Dr. Gary Greaves

# Topic Overview

# What's in store...

**L**

ogical Equivalence Laws

**C**

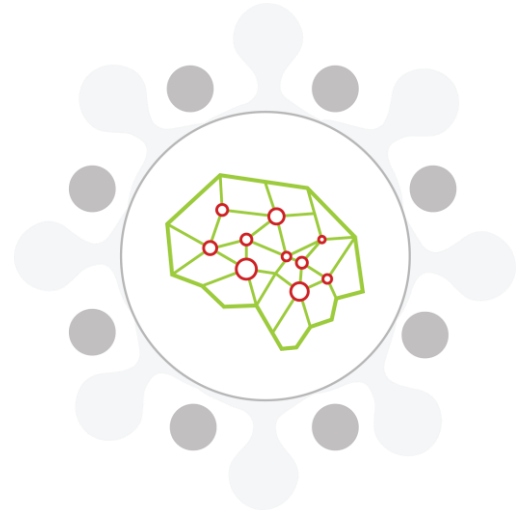
onditional Operator

**B**

iconditional Operator

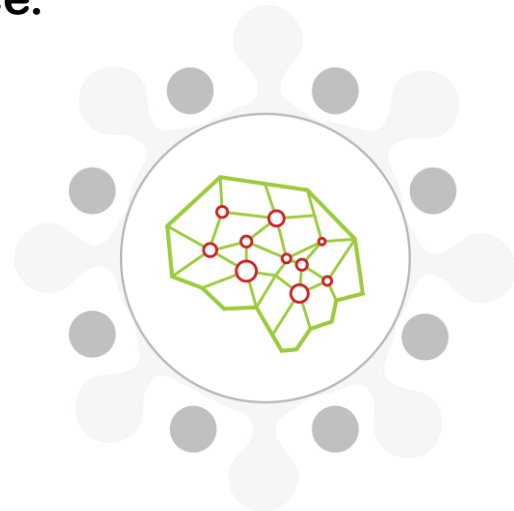
**O**

perator Precedence



# By the end of this lesson, you should be able to...

- Prove equivalence using logical equivalence laws.
- Use the conditional operator to combine propositions.
- Use the biconditional operator to combine propositions.
- Evaluate logical expressions using operator precedence.





# Logical Equivalence Laws

# Logical Equivalence Laws: Already Seen

- Useful laws to **transform** one logical expression to an equivalent one

## Axioms

$T \equiv$  Tautology

$C \equiv$  Contradiction

$$\neg T \equiv F$$

$$\neg F \equiv T$$

$$\neg T \equiv C \equiv F$$

$$\neg C \equiv T \equiv T$$

## De Morgan

$$\neg(p \wedge q) \equiv \neg p \vee \neg q$$

$$\neg(p \vee q) \equiv \neg p \wedge \neg q$$

## Commutativity

$$p \wedge q \equiv q \wedge p$$

$$p \vee q \equiv q \vee p$$

# Logical Equivalence Laws: More Laws

## Double Negation

$$\neg(\neg p) \equiv p$$

## Absorption

$$p \vee (p \wedge q) \equiv p$$

$$p \wedge (p \vee q) \equiv p$$

## Idempotent

$$p \wedge p \equiv p$$

$$p \vee p \equiv p$$

# Logical Equivalence Laws: Distributive Law

## Distributivity

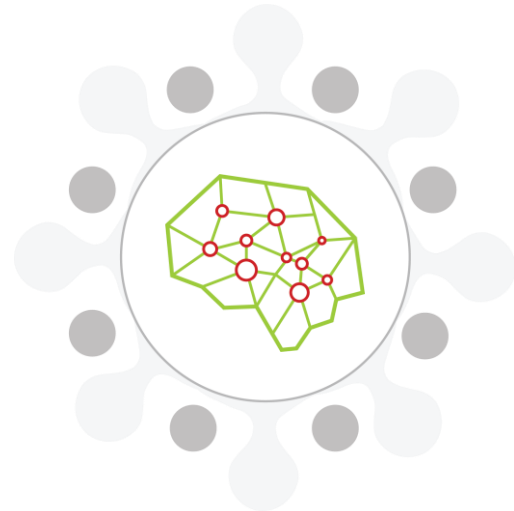
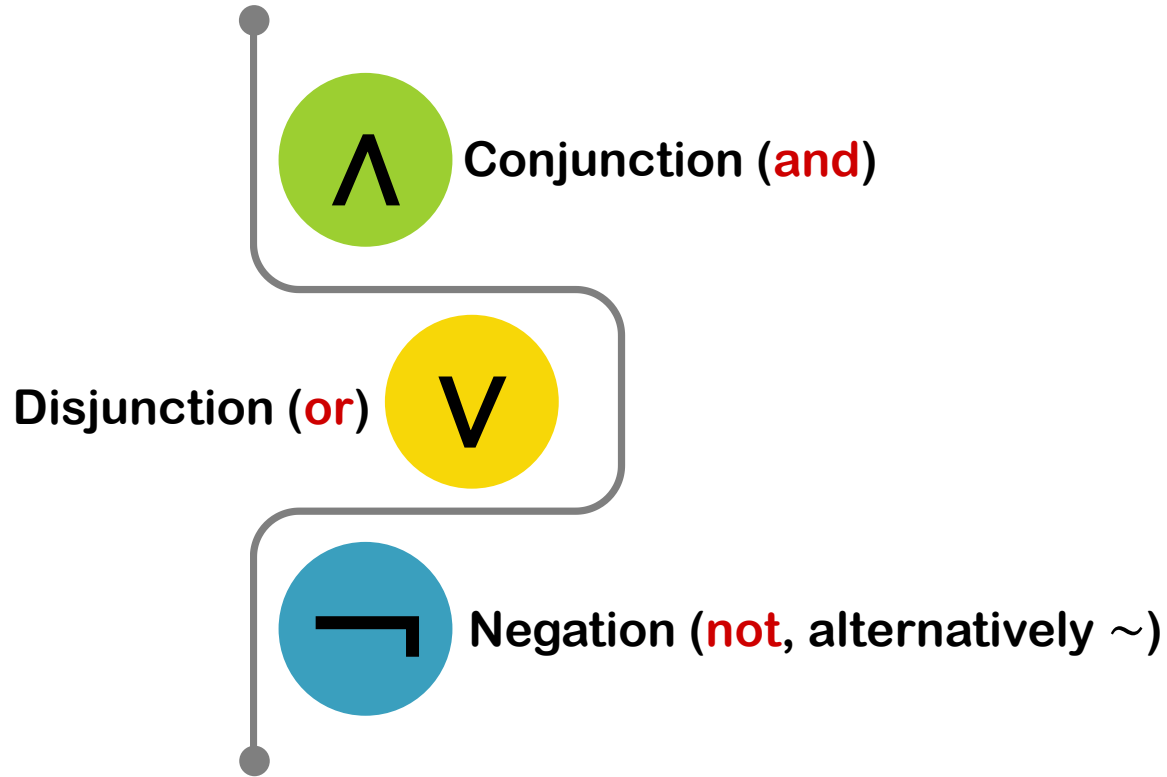
$$p \vee (q \wedge r) \equiv (p \vee q) \wedge (p \vee r)$$

$$p \wedge (q \vee r) \equiv (p \wedge q) \vee (p \wedge r)$$



# Conditional Operator

# Conditional Operator: Known Operators



# Conditional Operator: If Then



If  $p$  then  $q$ :  $p \rightarrow q$ .

By definition, when  $p$  is false,  $p \rightarrow q$  is true.

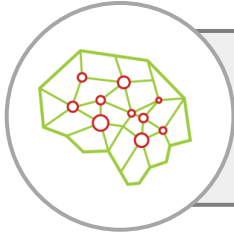
This is called **vacuously true** or **true by default**. →

$p$	$q$	$p \rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

```
gap>
gap> a:=10;; if (a>5) then Print("yes"); fi;
yes
gap> a:=1;; if (a>5) then Print("yes"); fi;
gap>
gap>
```

Not really the same!

# Conditional Operator: Conversion Theorem



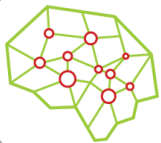
Theorem:  $p \rightarrow q \equiv \neg p \vee q$

## Proof

$p \ q$	$p \rightarrow q$	$\neg p$	$\neg p \vee q$
T T	T	F	T
T F	F	F	F
F T	T	T	T
F F	T	T	T

# Conditional Operator: Converse, Inverse, Contrapositive

Statement	$p \rightarrow q$
Converse	$q \rightarrow p$
Inverse	$\neg p \rightarrow \neg q$
Contrapositive	$\neg q \rightarrow \neg p$



Theorem:  $\neg q \rightarrow \neg p \equiv p \rightarrow q$

## Proof

$$\begin{aligned}\neg q \rightarrow \neg p \\ &\equiv \neg(\neg q) \vee \neg p \\ &\equiv q \vee \neg p \\ &\equiv \neg p \vee q \\ &\equiv p \rightarrow q\end{aligned}$$

# Conditional Operator: Only If

- $p$  **only if**  $q \triangleq \neg q \rightarrow \neg p$
- $\neg q \rightarrow \neg p$  is the **contrapositive** of  $p \rightarrow q$
- (If not  $q$  then not  $p$ )  $\equiv (p \rightarrow q)$  (why?)



## Example

“Bob pays taxes **only if** his income  $\geq \$1000$ ”

$\triangleq$  “if Bob’s income  $< \$1000$  then he does not pay taxes”

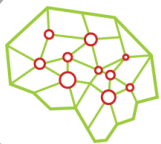
$\equiv$  “if Bob pays tax then his income  $\geq \$1000$ ”



Bob



# Conditional Operator: Sufficient and Necessary Conditions



When  $p \rightarrow q$ ,  $p$  is called a **sufficient condition** for  $q$ ,  
 $q$  is a **necessary condition** for  $p$ .

- Being an apple is a **sufficient condition** for being a fruit.  
≡ If it is an apple then it must be a fruit.
- Being a fruit is a **necessary condition** for being an apple.  
≡ If it is not a fruit then it cannot be an apple.



# Conditional Operator: Example

Let  $f$ : “you fix my ceiling”,  $p$ : “I will pay my rent”

“You fix my ceiling **or** I won’t pay my rent!”

$$f \vee \neg p \equiv p \rightarrow f$$



Landlord

“**If** you do not fix my ceiling, **then** I won’t pay my rent.”

$$\neg f \rightarrow \neg p \equiv p \rightarrow f$$



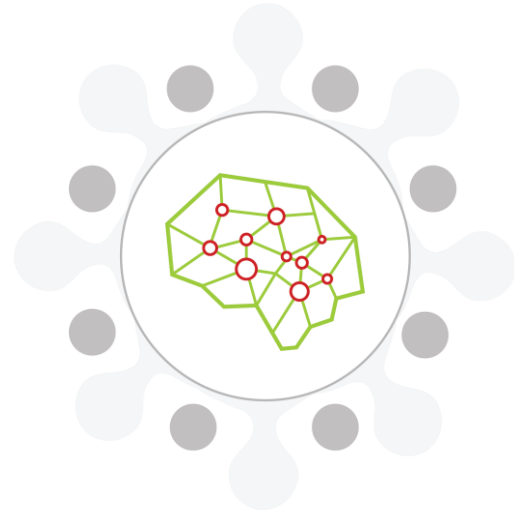
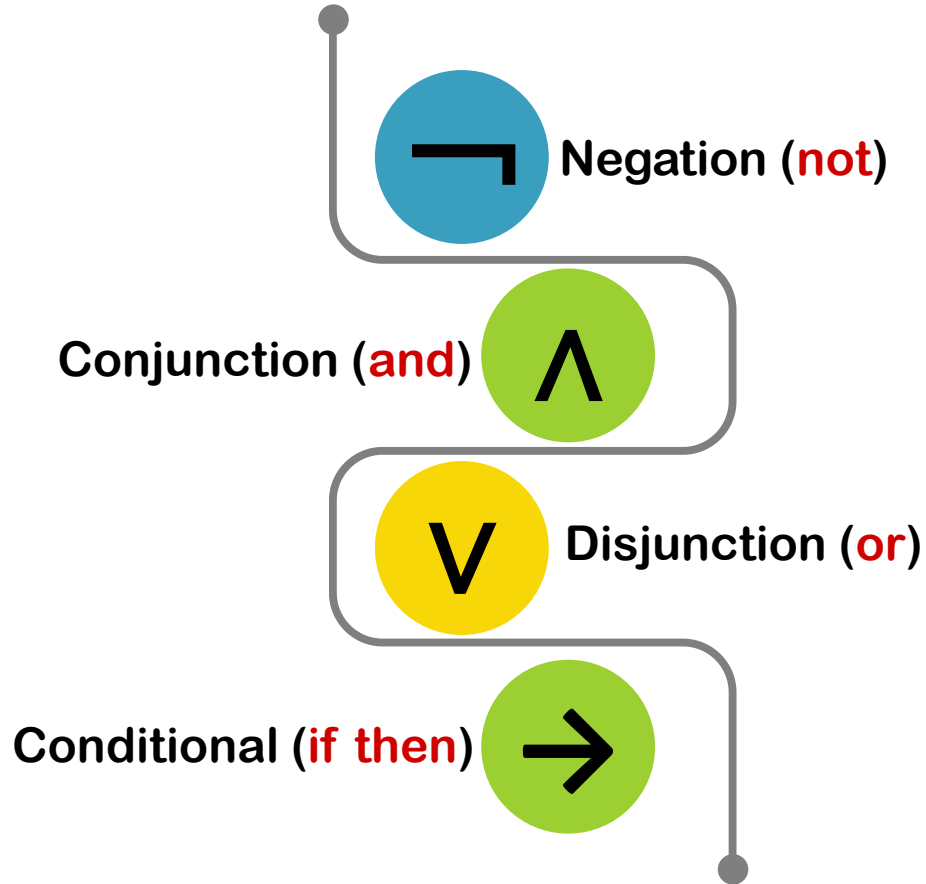
Tenant

“I will pay my rent **only if** you fix my ceiling.”

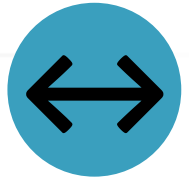
$$\neg f \rightarrow \neg p \equiv p \rightarrow f$$

# Biconditional Operator

# Biconditional Operator: Known Operators

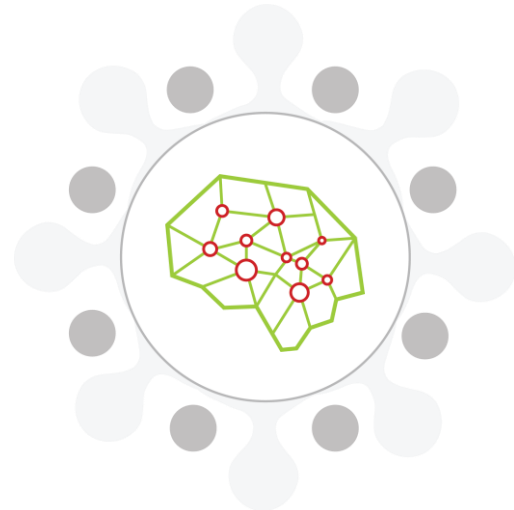


# Biconditional Operator: If and Only If



- The **biconditional** of  $p$  and  $q$ :  $p \leftrightarrow q \triangleq (p \rightarrow q) \wedge (q \rightarrow p)$ 
  - True only when  $p$  and  $q$  have identical truth value
- If and only if (**iff**)

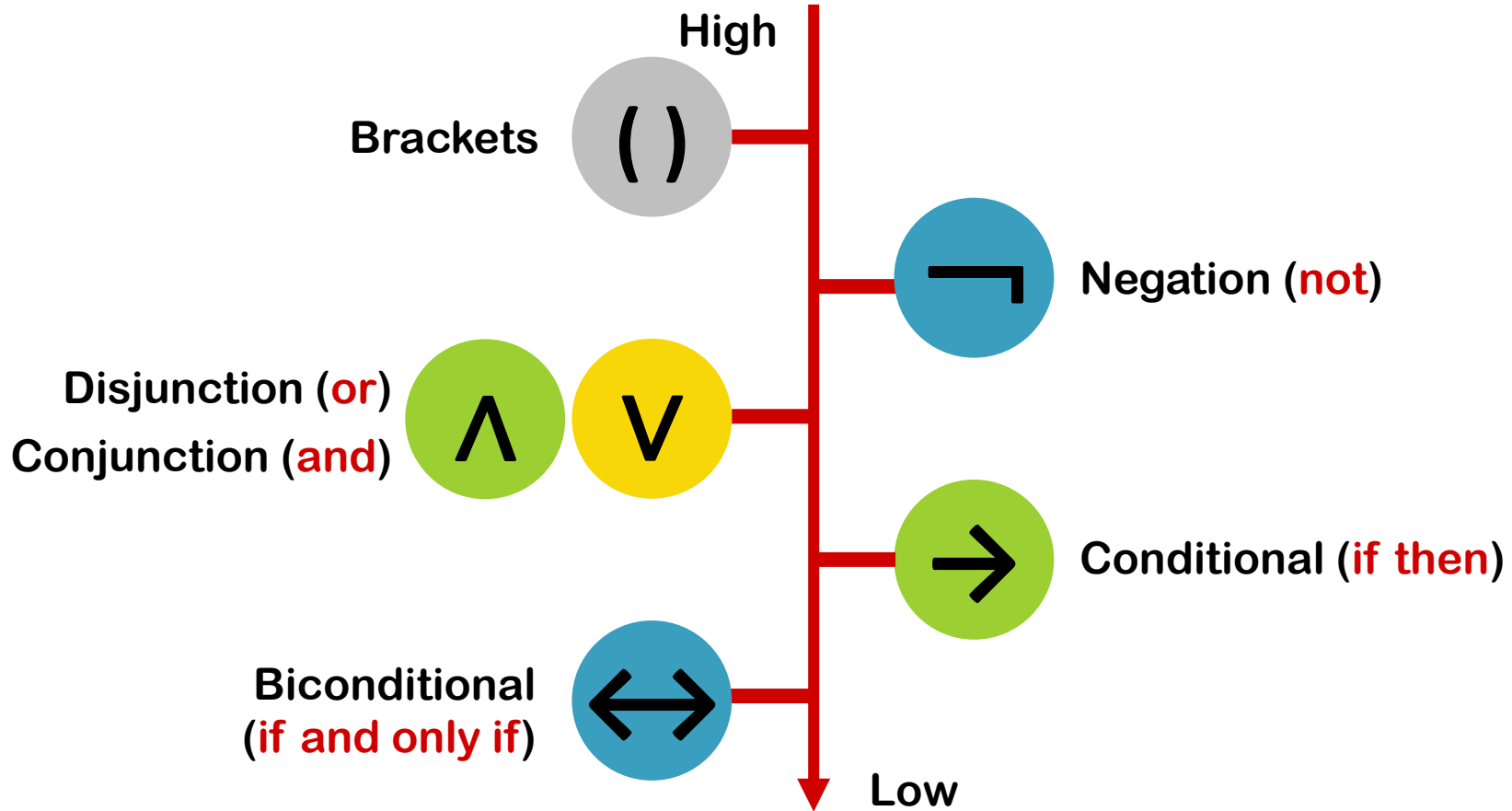
$p \ q$	$p \rightarrow q$	$q \rightarrow p$	$p \leftrightarrow q$
T T	T	T	T
T F	F	T	F
F T	T	F	F
F F	T	T	T



# Operator Precedence



# Operator Precedence: High to Low



# Operator Precedence: Leftmost and Rightmost

## Leftmost Precedence

When equal priority instances of **binary connectives** are not separated by  $()$ , the **leftmost** one has precedence.

E.g.,  $p \rightarrow q \rightarrow r \equiv (p \rightarrow q) \rightarrow r$



## Rightmost Precedence

When instances of  $\neg$  are not separated by  $()$ , the **rightmost** one has precedence.

E.g.,  $\neg\neg\neg p \equiv \neg(\neg(\neg p))$



# Operator Precedence: Example



Show that  $p \vee q \rightarrow r \equiv (p \rightarrow r) \wedge (q \rightarrow r)$

$$p \vee q \rightarrow r$$

$$\equiv (p \vee q) \rightarrow r$$

Operator precedence

$$\equiv \neg(p \vee q) \vee r$$

Why?

$$\equiv (\neg p \wedge \neg q) \vee r$$

De Morgan's

$$\equiv (\neg p \vee r) \wedge (\neg q \vee r)$$

Why?

$$\equiv (p \rightarrow r) \wedge (q \rightarrow r)$$

Why?

# Topic Summary

# Let's recap...

- Useful logical equivalence laws:
  - Proving equivalence using these laws
- Conditional and biconditional operators:
  - Sufficient and necessary conditions
- Operator precedence

