

American Computer Science League

2021 Finals • Solutions to Short Problems • Senior Division

1. Boolean Algebra

$$\begin{aligned}\overline{ABC(A + B)(\overline{B} + \overline{C})} &= \overline{ABC(\overline{AB} + \overline{AC} + \overline{BB} + \overline{BC})} \\ &= \overline{ABC + AB\overline{C}\overline{C} + \overline{A}B\overline{B}C + \overline{A}B\overline{B}C\overline{C}} \\ &= \overline{ABC} \\ &= \overline{A} + B + \overline{C}\end{aligned}$$

1. $\overline{A} + B + \overline{C}$ (D)

2. Boolean Algebra

$$\begin{aligned}(\overline{AB + C}) + (\overline{A}\overline{B} + \overline{C}) &= \overline{A}\overline{B}\overline{C} + (\overline{A}\overline{B} + \overline{C}) \\ &= (\overline{A} + \overline{B})\overline{C} + \overline{A}\overline{B} + \overline{C} \\ &= \overline{A}\overline{C} + \overline{B}\overline{C} + \overline{A}\overline{B} + \overline{C} \\ &= \overline{A}\overline{B} + \overline{C}(\overline{A} + \overline{B} + 1) \\ &= \overline{A}\overline{B} + \overline{C}\end{aligned}$$

If $\overline{C} = 1$, then $C = 0$. 4 ordered triples make it TRUE. (*, *, 0)

If $\overline{C} = 0$, then $C = 1$. $\Rightarrow \overline{A} = 1, \overline{B} = 1$. (0, 0, 1)

2. 5 (C)

3. Bit-String Flicking

$$\begin{aligned}&((\text{LCIRC-25 } 011001) \text{ OR } (\text{RCIRC-16 } 101101) \text{ AND } (\text{NOT } 011100)) \\ &= ((\text{LCIRC-1 } 011001) \text{ OR } (\text{RCIRC-4 } 101101) \text{ AND } 100011) \\ &= (110010 \text{ OR } (110110 \text{ AND } 100011)) \\ &= (110010 \text{ OR } 100010) \\ &= 110010\end{aligned}$$

3. 110010 (A)

4. Bit-String Flicking

Let $X = abcde$.

$$((\text{LSHIFT-1 } X) \text{ OR } (\text{NOT } (\text{RSHIFT-1 } 10111))) = X$$

$$((\text{LSHIFT-1 } abcde) \text{ OR } (\text{NOT } (\text{RSHIFT-1 } 10111))) = abcde$$

$$\text{LHS} = ((\text{LSHIFT-1 } abcde) \text{ OR } (\text{NOT } (\text{RSHIFT-1 } 10111)))$$

$$= (bcde0 \text{ OR } (\text{NOT } 01011))$$

$$= (bcde0 \text{ OR } 10100)$$

$$= 1c1e0$$

$$\text{LHS} = \text{RHS} \Rightarrow 1c1e0 = abcde$$

$$\Rightarrow a = 1, b = c, c = 1, d = e, e = 0$$

$$\Rightarrow a = b = c = 1, d = e = 0$$

$$\Rightarrow 11100$$

4. 11100 (B)

5. Recursive Functions

$$f(20) = f(f(20 - 3)) + 3 = f(f(17)) + 3 = f(5) + 3 = 0$$

$$f(17) = f(f(17 - 3)) + 3 = f(f(14)) + 3 = f(3) + 3 = 2$$

$$f(14) = f(f(14 - 3)) + 3 = f(f(11)) + 3 = f(1) + 3 = 0$$

$$f(11) = f(f(11 - 3)) + 3 = f(f(8)) + 3 = f(-1) + 3 =$$

$$f(8) = f(8 - 2) - 2 = f(6) - 2 = 1 - 2 = -1$$

$$f(6) = f(6 - 2) + 2 = f(4) - 2 = 3 - 2 = 1$$

$$f(4) = 4 - 1 = 3$$

$$f(-1) = -1 - 1 = -2$$

$$f(1) = 1 - 1 = 0$$

$$f(3) = 3 - 1 = 2$$

$$f(5) = f(5 - 2) - 2 = f(3) - 2 = 2 - 2 = 0$$

5. 3 (B)

6. Recursive Functions

Each zig zag curve has twice the length of the segment on which it is constructed. The perimeter will double the length at each stage.

Stage	1	2	3	4	5	6	7
Perimeter	64	128	256	512	1024	2048	4096

6. 7 (C)

7. Digital Electronics

The circuit can be represented by the following boolean expression:

$$A \oplus \overline{A} \oplus (B \oplus C)$$

A	B	C	$B \oplus C$	$A \oplus (B \oplus C)$	$\overline{A \oplus (B \oplus C)}$	$A \oplus \overline{A \oplus (B \oplus C)}$
0	0	0	0	0	1	1
0	0	1	1	1	0	0
0	1	0	1	1	0	0
0	1	1	0	0	1	1
1	0	0	0	1	0	1
1	0	1	1	0	1	0
1	1	0	1	0	1	0
1	1	1	0	1	0	1

7. 4 (C)

8. Digital Electronics

The boolean expression for the circuit is: $\overline{A(\overline{A + B + BC}) + C}$

$$\overline{A(\overline{A + B + BC}) + C} = \overline{A(\overline{A} \overline{B} \overline{B} \overline{C}) + C} = \overline{C}$$

8. \overline{C} (B)

9. Prefix-Infix-Postfix

$$\begin{aligned} & 3 \ 4 \ 5 + 7 \ 2 \ 1 \ \& \ 6 \ ! \ * \ * \ 7 \ 8 - 4 \ 3 + 2 \ \& \ 3 \ 2 \ ^ \ ! \ / \ \& \\ & = 3 \ (4 \ 5 +) \ (7 \ 2 \ 1 \ \&) \ (6 \ ! \ * \ * \ (7 \ 8 -) \ (4 \ 3 +) \ 2 \ \& \ (3 \ 2 \ ^) \ ! \ / \ \& \\ & = 3 \ \{ \ (4 \ 5 +) \ [\ (7 \ 2 \ 1 \ \&) \ (6 \ ! \ * \ * \ { \ [\ (7 \ 8 -) \ (4 \ 3 +) \ 2 \ \& \ [\ (3 \ 2 \ ^) \ ! \] \ / \ } \ \& \\ & = 3 \ \{ \ (+4 \ 5) \ [\ (\&7 \ 2 \ 1) \ (! \ 6) \ * \ * \ { \ [\ (-7 \ 8) \ (+4 \ 3) \ 2 \ \& \ [\ (^3 \ 2) \ ! \] \ / \ } \ \& \\ & = 3 \ \{ \ (+4 \ 5) \ [\ * \ \&7 \ 2 \ 1 \ ! \ 6 \) \] \ * \ { \ [\ \& \ -7 \ 8 + 4 \ 3 \ 2 \] \ [\ ! \ ^3 \ 2 \] \ / \ } \ \& \\ & = 3 \ \{ \ * \ +4 \ 5 \ * \ \&7 \ 2 \ 1 \ ! \ 6 \ } \ { \ / \ \& \ -7 \ 8 + 4 \ 3 \ 2 \ ! \ ^3 \ 2 \ } \ \& \\ & = \ \&3 \ * \ +4 \ 5 \ * \ \&7 \ 2 \ 1 \ ! \ 6 \ / \ \& \ -7 \ 8 + 4 \ 3 \ 2 \ ! \ ^3 \ 2 \end{aligned}$$

9. $\&3*+4\ 5*\&7\ 2\ 1!\ 6/\&\ -7\ 8 +4\ 2!\wedge3\ 2$ (D)

10. Prefix-Infix-Postfix

$$\begin{aligned} & * - + 4 7 - 1 8 / + ^ 5 2 * 5 - 6 2 ^ 3 2 \\ & = * - (+ 4 7) (- 1 8) / + (^ 5 2) * 5 (- 6 2) (^ 3 2) \\ & = (* (- 11 (-7)) / + 25 (* 5 4) 9 \\ & = * 18 / (+ 25 20) 9 \\ & = * 18 (/ 45 9) \\ & = * 18 5 \\ & = 90 \end{aligned}$$

10. 90 (C)

11. Computer Number Systems

$$\begin{aligned} A74F2C_{16} &= 1010\ 0111\ 0100\ 1111\ 0010\ 1100_2 \\ 4_{16} &= 100_2 \\ \text{Dividing } 1010\ 0111\ 0100\ 1111\ 0010\ 1100_2 &\text{ by } 100_2 \\ \text{shifts the bits 2 spaces to the right. Therefore the answer is:} \\ 1010\ 0111\ 0100\ 1111\ 0010\ 11_2 &= 10\ 1001\ 1101\ 0011\ 1100\ 1011_2 \\ &= 2\ 9\ D\ 3\ C\ B_{16} \end{aligned}$$

11. 29D3CB₁₆ (B)**12. Computer Number Systems**

$$2001 = 3721_8 \qquad 2021 = 3745_8$$

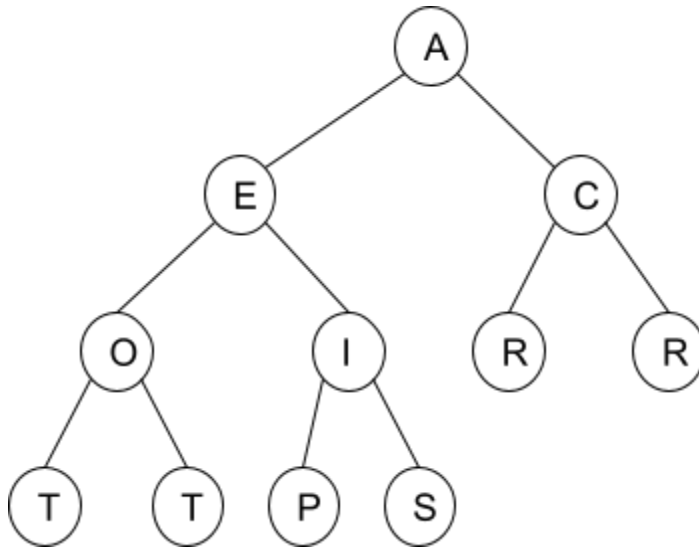
YEAR ₈	3721	3722	3723	3724	3725	3726	3727
DIGITS	1	1	1	1	2	2	2
YEAR ₈	3730	3731	3732	3733	3734	3735	3736
DIGITS	1	1	1	1	1	2	2
YEAR ₈	3737	3740	3741	3742	3743	3744	3745
DIGITS	2	1	1	1	1	1	2

There are 28 digits greater than 4.

12. 28 (D)

13. Data Structures

The min-heap for TRICERATOPS is:



13. O I R R (C)

14. Data Structures

FIFO: July: \$20/share (sold)
March: \$10/share (bought)
Profit: \$10/share
Total profit: \$1000 for 100 shares
Tax: 20% of \$1000 = \$200

LIFO: July: \$20/share (sold)
May: \$5/share (bought)
Profit: \$15/share
Total profit: \$1500 for 100 shares
Tax: 20% of \$1500 = \$300

14. FIFO, saving \$100 (A)

<p>15. Graph Theory</p> <p>Squaring a matrix gives the number of paths of length 2.</p> $ \begin{vmatrix} 0 & 1 & 1 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 \end{vmatrix}^2 = \begin{vmatrix} 1 & 1 & 1 & 0 & 2 & 1 \\ 2 & 2 & 0 & 0 & 0 & 2 \\ 0 & 1 & 2 & 1 & 2 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 2 & 1 & 2 & 1 \\ 1 & 1 & 0 & 0 & 0 & 1 \end{vmatrix} $ <p>There are 15 paths that do not have a length of 2, but only 4 unique pairs of vertices do not have a path of length 2 in either direction. DD, AD/DA, BD/DB, DF/FD</p>	<p>15. 4 (B)</p>
<p>16. Graph Theory</p> <p>There are 12 cycles: ACA, ABCA, ABEA, ADEA, ACBEA, ACDEA, ABCDEA, ADEBCA, BEB, BCB, BCDEB, and DED</p>	<p>16. None of the above (E)</p>
<p>17. What Does This Program Do?</p> <p>Pascal's triangle rows start and end with 1. A new entry, pt(i), in a row is found by adding the entry above to the left, pt(i-1), and the one to the right, pt(i). This is choice B.</p>	<p>17. for i = X to 1 step -1 pt(i) = pt(i) + pt(i-1) next i (B)</p>

18. LISP Programming

```
G = (CADDADADDR S)
  = (CADDADADDR '(H ((A C) K) (E (R R (A))) (N (K))))
  = (CADDADADR '(((A C) K) (E (R R (A))) (N (K))))
  = (CADDADAR '(E (R R (A))) (N (K)))
  = (CADDADR '(E (R R (A))))
  = (CADDAR '((R R (A))))
  = (CADDR '(R R (A)))
  = (CADR '(R (A)))
  = (CAR '((A)))
  = (A)
```

```
H = (CADADADR Q))
  = (CADADADR '((P L) (A (T (F O) R)) M))
  = (CADADAR '(A (T (F O) R)) M))
  = (CADADR '(A (T (F O) R)))
  = (CADAR '((T (F O) R)))
  = (CADR '(T (F O) R))
  = (CAR '((F O) R))
  = (F O)
```

```
(CONS G H) = (CONS (A) (F O))
            = ((A) F O)
```

18. ((A) F O) (A)

19. FSAs and Regular Expressions

Given the pattern `[^aeiou][aeiou][^a-j][p-t]?.(er|s)*` the following strings fail at the given character:

crayons: r - must be a vowel

erasers: e - cannot start with a vowel

staples: t - must be a vowel

staplers: t - must be a vowel

notebooks: e - must be p-t

computers: t - there is no symbol to use

tablets: b - cannot be a-j

boards: a - cannot be a-j

rulers: e - must have another character before the er or the s

There are 3 left which satisfy the regular expression which are:
markers, desks, and compass.

19. None of the above (E)

20. Assembly Language

The assembly programs can be converted to ACSL WDTPD code as follows:

```
input x, y
z = x - y
c = 1: b = 1: a = 1
while z > 0
    c = c * z
    z = z - 1
end while
while y > 0
    b = b * y
    y = y - 1
end while
while x > 0
    a = a * x
    x = x - 1
end while
d = a / b / c
output d
```

Each while loop calculates a factorial.

$d = a!/b!/c!$ which is the formula to calculate 8 things taken 3 at a time where order doesn't matter.

20. 56 (C)