For the eqn.

$$(\lambda + \mu)P_n = \lambda P_{n-1} + \mu P_{n+1}$$

Substitute by the following in the RHS

$$P_n = \rho P_0$$

$$\therefore \lambda \rho^{n-1} P_0 + \mu \rho^{n+1} P_0 = (\lambda \rho^{n-1} + \mu \rho^{n+1}) P_0$$

$$= P_0 \rho^n (\lambda \rho^{-1} + \mu \rho) = P_0 \rho^n (\mu + \lambda)$$

= L.H.S.

 $\therefore P_n = \rho P_0$ satisfies the equation.

For the eqn.

$$(\lambda + \mu)P_n = \lambda P_{n-1} + \mu P_{n+1}$$

substitute by

$$P_{n+1} = \rho P_n$$
 i.e. $(P_n = P_{n+1} / \rho)$ in the R.H.S. of the equation

$$\therefore \lambda P_n / \rho + \mu \rho P_n = \mu P_n + \lambda P_n = P_n (\lambda + \mu) = \text{L.H.S.}$$

$$\therefore P_{n+1} = \rho P_n$$
 satisfies the equation.

For the M/M/1 queue

$$P_{0} = 1 - \sum_{n=1}^{N} P_{n}$$

$$= 1 - \frac{P_{1}(1 - (\frac{P_{n}}{P_{n-1}})^{N})}{1 - (\frac{P_{n}}{P_{n-1}})}$$

$$= 1 - \frac{P_{0}\rho(1 - \rho^{N})}{1 - \rho}$$

$$\therefore P_{0}[1 + \frac{(\rho - \rho^{N+1})}{1 - \rho}] = 1$$

$$\therefore P_{0} = \frac{1 - \rho}{1 - \rho^{N+1}}$$
and
$$\therefore P_{n} = P_{0}\rho^{n} = \frac{(1 - \rho)\rho^{n}}{1 - \rho^{N+1}}$$

Problem 2 (Schwartz 2.9)

For finite m/m/1 queue

$$\gamma = \lambda(1 - P_{\scriptscriptstyle B}) = \mu(1 - P_{\scriptscriptstyle 0})$$

$$\therefore 1 - P_B = \frac{\mu}{\lambda} (1 - P_0)$$

so
$$P_B = 1 - \frac{\mu}{\lambda} (1 - P_0)$$

$$\therefore P_0 = \frac{1 - \rho}{1 - \rho^{N+1}}$$

$$\therefore P_{B} = 1 - \frac{1}{\rho} + \frac{1}{\rho} \frac{1 - \rho}{(1 - \rho^{N+1})} = (1 - \rho) \rho^{N} \left[\frac{-1}{\rho^{N+1}} + \frac{1}{1 - \rho^{N+1}} \right]$$

$$= (1 - \rho) \rho^{N} \left[\frac{\rho^{N+1} - 1 + 1}{\rho^{N+1} (1 - \rho^{N+1})} \right] = \frac{(1 - \rho) \rho^{N}}{1 - \rho^{N+1}}$$

$$= P_{N}$$

Problem 3 (Schwartz 2.10)

$$P_{B} = \frac{(1-\rho)\rho^{N}}{1-\rho^{N+1}}$$

- 1. For $\rho = 0.5$, $P_B = 10^{-3}$
 - a. We can get the length of the buffer by trial and error N= 9 customers
 - b. For $\rho = 0.5$, $P_B = 10^{-6}$ N=19 customers
- 2. For $\rho = 0.8$,
 - a. $P_B = 10^{-3}$ N=23 customers
 - b. For $\rho = 0.8$, $P_B = 10^{-6}$ N= 29 customers

From the results we conclude that, in order to decrease the blocking probability of the system or to increase the utilization of the system we have to use a longer buffer.

Problem 4 (Shwartz 2-11)

$$E[n] = \sum_{n=0}^{\infty} nP_n = \sum_{n=0}^{\infty} nP_0 \rho^n = \sum_{n=0}^{\infty} n(1-\rho)\rho^n = (1-\rho)\sum_{n=0}^{\infty} n\rho^n = (1-\rho)\rho\sum_{n=1}^{\infty} n\rho^{n-1}$$

$$\therefore \sum_{n=0}^{\infty} \rho^n = \frac{1}{1-\rho}$$

$$\therefore \frac{d}{dn} \sum_{n=1}^{\infty} \rho^n = \frac{1}{(1-\rho)^2}$$
and so, $E[n] = \frac{\rho}{1-\rho}$

Problem 5

1. For
$$M / M / \infty$$

$$\lambda_n = \lambda , \quad \mu = n\mu$$

$$\therefore P_n = P_0 \frac{\rho^n}{n!}$$

From the property of probability, we get $\sum_{i=0}^{\infty} P_i = 1$

$$\therefore \sum_{n=0}^{\infty} P_0 \frac{\rho^n}{n!} = 1$$

$$P_0 = e^{-\rho}$$

 $\therefore P_n = e^{-\rho} \frac{\rho^n}{n!}$ which is the Poisson's distribution

$$E[n] = \rho;$$
 $E[T] = \frac{\rho}{\lambda} = \frac{1}{\mu}$

$$E[w] = E[T] - \frac{1}{u} = 0; \quad E[q] = 0$$

2. For queue with discouragement

$$\lambda_n = \frac{\lambda}{n+1}, \quad \mu_n = \mu$$

$$\therefore P_n = \frac{\lambda^n P_0}{n! \mu^n} = \frac{P_0 \rho^n}{n!}$$

Thus we get the value of $P_0 = e^{-\rho}$ which is the same as the M/M/ \propto case

Similarly $P_n = \frac{e^{-\rho} \rho^n}{n!}$ is the Poisson's distribution

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Problem 6 (Galleger 3.8)

- a. Given λ =10packets/sec τ (Transmission Time)=20msec=.02sec
 - : Probability [Inter-arrival time < Transmission Time] = 1-Prob.[no arrival of Packets] = $1 - e^{-\lambda \tau}$
 - ∴ Probability of no collision with the predecessor = $1 (1 e^{-\lambda \tau})$

Also Probability of no collision with the predecessor and successor

$$= e^{-\lambda\tau} * e^{-\lambda\tau}$$

= $e^{-(0.2+0.2)} = e^{-0.4} = 0.67$

Also Prob. that a packet does-not collide with another Packet will be the same = 0.67

Problem 7 (Galleger 3.9)

$$\lambda=150~packets/min/session$$
 $\mu=50*10^3~bits/sec$ / $1000~bits=50~packets$ / $sec=3000~packets$ / $min=300~packets/min/session$

$$\rho = \lambda/\mu = 150/300 = 0.5$$

a) For TDM with 10 channels (m=10)

$$E[n] = m\rho/(1-\rho) = 10*0.5/(1-0.5) = 10$$
 packets

$$E[T] = m / (\mu - \lambda) = 10 / (300-150) = 6.67*10^{-3} min = 0.4 sec$$

$$E[q] = E[n]-m\rho=10-10*0.5=5$$
 packets

For Statistical Multiplexing

0.166)=1.1packets

$$E[n] = \rho/(1\text{-}\rho) = 0.5 \ / \ (1\text{-}0.5) = 1 \ packet$$

$$E[T] = 1 / (\mu - \lambda) = 1/(3000-1500) = 6.67*10^{-4} \text{ min} = 0.04 \text{ sec}$$

$$E[q] = E[n]-\rho=1-0.5=0.5$$
 packet

b) In case of five sessions have rates 250 packets/min and other five have rates of 50 packets/min

$$\begin{array}{lll} & \text{For the } 1^{\text{st}} \text{ five sessions} & \text{For the } 2^{\text{nd}} \text{ five sessions} \\ \lambda = & 250 \cdot 5 = 1250 \text{ packets/min} & \lambda = & 50 \cdot 5 = 250 \text{ packets/min} \\ \mu = & 3000 \cdot 5 / 10 = & 1500 \text{ packets/min} & \mu \\ = & 3000 \cdot 5 / 10 = & 1500 \text{ packets/min} & \mu \\ \rho = & \lambda / \mu = & 1250 / 1500 = & 0.833 & \rho = \lambda / \mu = & 1250 / 1500 = & 0.166 \\ E[n] = & 5 \cdot 0.83 / (1 - 0.83) = & 24.9 \text{ packets} & E[n] = & 5 \cdot 0.166 / (1 - 0.83) = & 1.00 \cdot 10^{-1} \cdot$$

$$E[n] total = 24.9 + 1.1 = 26 packets$$

$$E[T]=m/(\mu-\lambda)=5/(1500-1250)=0.02min$$
 $E[T]=5/(1500-250)=4*10^{-3} min$

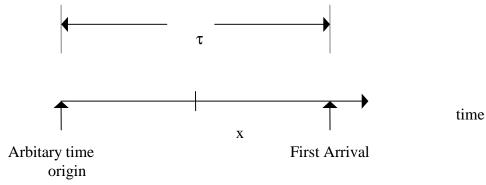
$$E[T] total = (0.02+0.004)*60=1.44 sec$$

 $E[q]=E[n]-m\rho=24.9-5*0.83=20.735$ packets E[q]=1.1-5*0.16=0.27 packets

$$E[q] total = 20.735 + 0.27 = 21 packets$$

Problem 8 (Galleger 3.10)

Consider the time diagram shown



a. Let τ representing the time the first arrival after some arbitary time origin.

Take any value x.

No arrival occur in the interval (0,x) if and only if $\tau > x$.

The probability that no arrivals occur in (0,x); i.e.

 $P(\tau > x) = \text{prob.}(\text{no arrivals in }(o,x))$

For Poisson's distribution

$$p(k) = (\lambda t)^k e^{-\lambda t} / k!$$

here k=0

$$\therefore p(\tau > x) = e^{-\lambda x}$$

Then the prob. that $\tau \le x = 1 - e^{-\lambda x}$

On the basis, the probability of packet arrivals in the small time interval (t , t+ δ t) is just $\lambda\delta$ t+o(δ t)

Let $N^{(i)}(t, t + \varphi \delta t)$ be the number of events in Poisson process I, I=1,2,....m in the interval $(t, t + \delta t)$.

Let $N(t, t + \delta t)$ be the total number of events from the whole stream.

Then
$$\Pr{ob.[N(t,t+\delta t)=0]} = \prod_{i=1}^{n} prob.[N^{(i)}(t,t+\delta t)=0]$$

Since the probability of no packet arrival in the time interval $(t.t + \delta t)$ is $(1 - \lambda_i \delta t + o(\delta t))$

$$\therefore prob.[N(t, t + \delta t) = 0] = \prod_{i=1}^{n} (1 - \lambda_i \delta t + o(\delta t))$$
$$= 1 - \lambda \delta t + o(\delta t) + \text{higher powers of } (\delta t) \text{ that goes to zero}$$

Where $\lambda = \sum_{i=1}^{m} \lambda_i$, since the individual process are independent

Now
$$prob.[N(t, t + \delta t) = 1] = \prod_{i=1}^{n} \lambda_i \delta t + o(\delta t)$$

= $\lambda \delta t + o(\delta t) + \text{higher powers of } (\delta t) \text{ which goes to}$

zero

$$prob.[N(t,t+\delta t) > 2]$$

$$= 1 - prob.[N(t,t+\delta t) = 0] - prob.[N(t,t+\delta t) = 1]$$

$$= o(\delta t)$$

b. For a R .V. z with Poisson distribution $P_z[(N_1 + N_2) = n]$ is given by

$$P_{z}(n) = \sum_{k=0}^{n} P[N_{1} = k] P[N_{2} = n - k]$$

$$= \sum_{k=0}^{n} P_{x}(k) P_{y}(n - k)$$

$$= \sum_{k=0}^{n} \frac{1}{k!} \frac{1}{(n - k)!} e^{-(\lambda_{1} + \lambda_{2})} \lambda_{1}^{k} \lambda_{2}^{n - k}$$

Recall from the Binomial theorem

$$\sum_{k=0}^{n} \binom{n}{k} \lambda_1^{k} \lambda_2^{n-k} = (\lambda_1 + \lambda_2)^n$$

Then
$$P_z(n) = \frac{1}{n!} e^{-(\lambda_1 + \lambda_2)} \sum_{k=0}^n \binom{n}{k} \lambda_1^k \lambda_2^{n-k}$$
$$= \frac{(\lambda_1 + \lambda_2)^n}{n!} e^{-(\lambda_1 + \lambda_2)} \text{ for } n \ge 0$$

Then the number of arrivals in the union of the intervals is Poisson distributed with parameter $\lambda_1 + \lambda_2$

c. $P[1 \text{ arrival from } A_1 \text{ prior to t } | 1 \text{ occurred}]$

$$= \frac{p[1arrivalfromA_1 priortot, 0 fromA_2]}{P[1arrivaloccured]}$$
$$= \frac{\lambda_1 t e^{-\lambda_1 t} e^{-\lambda_2 t}}{\lambda_1 t e^{-\lambda_1 t}} = \frac{\lambda_1}{\lambda_1}$$

d. For all $s \in [t_1, t_2]$

$$= \frac{P[1arrivaloccuredin[t_1, s), 0arrivalsoccuredin[s, t_2]]}{P[1arrivaloccured]}$$

$$= \frac{\lambda(s - t_1)e^{-\lambda(s - t_1)}e^{-\lambda(t_2 - s)}}{\lambda(t_2 - t_1)e^{-\lambda(t_2 - t_1)}} = \frac{s - t_1}{t_2 - t_1}$$

Hence conditional on the knowledge that only one arrival occurred, the time of this arrival is uniformly distributed in $[t_1,t_2]$.