# ECE 642 Lecture 11 - Excerpts

### Part I Open queueing networks

For an M queue VC, the state probability vector is as follows:

$$p(\underline{n}) = \prod_{i=1}^{M} (1 - \rho_i) \rho_i^{n_i}$$

Given the service time at a the ith queue in the M-hop VC is  $1/\mu_i$ , the average end-to-end time delay is calculated as follows:

$$E(T) = \sum_{i=1}^{M} \frac{1}{\mu_i - \lambda_i} = \sum_{i=1}^{M} \frac{1/\mu_i}{1 - \rho_i}$$

for  $\rho$  defined as:

$$\rho = \frac{\lambda_i}{\mu_i}$$

We obtain the network-wide delay by invoking Little's formula in the network:

$$E(n_i) = \lambda_i T_i = \lambda_i \frac{1}{\mu_i - \lambda_i}$$

$$E(n) = \sum_{i=1}^M E(n_i) = \sum_{i=1}^M \frac{\lambda_i}{\mu_i - \lambda_i}$$

$$E(T) = \frac{E(n)}{\gamma} = \frac{1}{\gamma} \sum_{i=1}^M \frac{\lambda_i}{\mu_i - \lambda_i}$$

## Part II Closed queueing networks

The state probability vector is defined as:

$$p(\underline{n}) = p(n_1, n_2 \dots n_M)$$
$$= \prod_{i=1}^{M} (\frac{\lambda_i}{\mu_i})^{n_i} p(\underline{0})$$

- $n_i$ : number of packets in queue i
- $1/\mu_i$ : average service time in queue i
- $\lambda_i$ : flows over link (queue) i

Define  $g(N, M) = \frac{1}{p(0)}$ , the state probability vector is given by:

$$p(\underline{n}) = \left[\prod_{i=1}^{M} \left(\frac{\lambda_i}{\mu_i}\right)^{n_i}\right] \frac{1}{g(N, M)}$$

### Part III Buzen's algorithm

- 1.  $g(n,m) = g(n,m-1) + \rho_m g(n-1,m)$  with  $\rho_m = \frac{\lambda_m}{\mu_m}$  where m = 1, 2, ... M.
- 2. The initial condition is define as follows:

$$g(n,1) = \rho_1^n$$

where n = 0, 1, 2, ... N, and

$$g(0,m) = 1$$

where  $m = 1, 2 \dots M$ .

Let's consider a 3-hop VC as an example, assume:

$$\rho_1 = \rho_2 = \dots = \rho_M = \frac{\lambda_1}{\mu} = 1$$

$$\rho_{M+1} = \frac{1}{\rho}$$

The g(n,m) matrix is shown as follows:

When  $\rho = 1$ , redo the matrix above:

After obtaining this matrix by using Buzen's algorithm, the throughput and delay is calculated as follows:

$$P(n_{i} \ge k) = \rho_{i}^{k} \frac{g(N - k, M)}{g(N, M)}$$

$$P(n_{i} = k) = P(n_{i} \ge k) - P(n_{i} \ge k + 1)$$

$$= \rho_{i}^{k} \left[ \frac{g(N - k, M)}{g(N, M)} \right] - \rho_{i}^{k+1} \left[ \frac{g(N - k - 1, M)}{g(N, M)} \right]$$

$$= \frac{\rho_{i}^{k}}{g(N, M)} [g(N - k, M) - \rho_{i}g(N - k - 1, M)]$$

For each VC, the expected number of packets is calculated as:

$$E(n_i) = \sum_{k=0}^{N} kP(n_i = k)$$

$$= \sum_{k=0}^{N} \frac{k\rho_i^k}{g(N, M)} [g(N - k, M) - \rho_i g(N - k - 1, M)]$$

$$= \sum_{k=1}^{N} \rho_i^k \left[ \frac{g(N - k, M)}{g(N, M)} \right]$$

The throughput is then calculated:

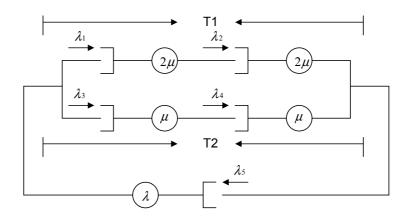
$$\gamma_i = \mu_i P(n_i \ge 1)$$

$$= \mu_i \rho_i \left[ \frac{g(N-1, M)}{g(N, M)} \right]$$

The normalized throughput is:

$$\frac{\gamma_i}{\mu_i} = \rho_i \left[ \frac{g(N-1,M)}{g(N,M)} \right]$$

## Part IV Example



Given  $\lambda_1 = \lambda_2$ ,  $\lambda_3 = \lambda_4$ , and  $q_{51} = q_{53} = 1/2$ , we know that:

$$\rho_1 = \rho_2 = \frac{\lambda_1}{2\mu}$$

$$\rho_3 = \rho_4 = \frac{\lambda_3}{\mu} = \frac{\lambda_1}{\mu}$$

If  $\rho_1 = \rho_2 = \frac{\lambda_1}{2\mu} = 1$ , then:

$$\rho_1 = \rho_2 = \frac{\lambda_1}{\mu} = 2$$

$$\rho_5 = \frac{\lambda_5}{\lambda} = \frac{2\lambda_1}{\lambda}$$

$$= \frac{2\lambda_1}{\mu} \frac{\mu}{\lambda} = \frac{4}{\rho}$$

By using Buzen's algorithm, we can get the matrix like this:

n/m 1 2 3 4 5  
0 1 1 1 1 1  
1 1 2 4 6 
$$6+4/\rho$$
  
2 1 3 11 23  $23+(6+4/\rho)$   
3 1 4 26 72  $72+(23+(6+4/\rho))$   
4 1 5 57 201  $201+(72+(23+(6+4/\rho)))$ 

The throughput is calculated like this:

$$\frac{\gamma_i}{\mu_i} = \rho_i \left[ \frac{g(N-1, M)}{g(N, M)} \right]$$

The expected number of packets in system is calculated like this:

$$E(n_i) = \sum_{k=1}^{N} \rho_i^k \left[ \frac{g(N-k, M)}{g(N, M)} \right]$$

The expected time delay is calculated like this:

$$E(T_1) = \frac{E(n_1) + E(n_2)}{\gamma_1}$$

$$E(T_2) = \frac{E(n_3) + E(n_4)}{\gamma_3}$$