CNL
Communications & Networking
Laboratory

ECE 462 – Data and Computer Communications

Lecture 19-20: LAN Protocols

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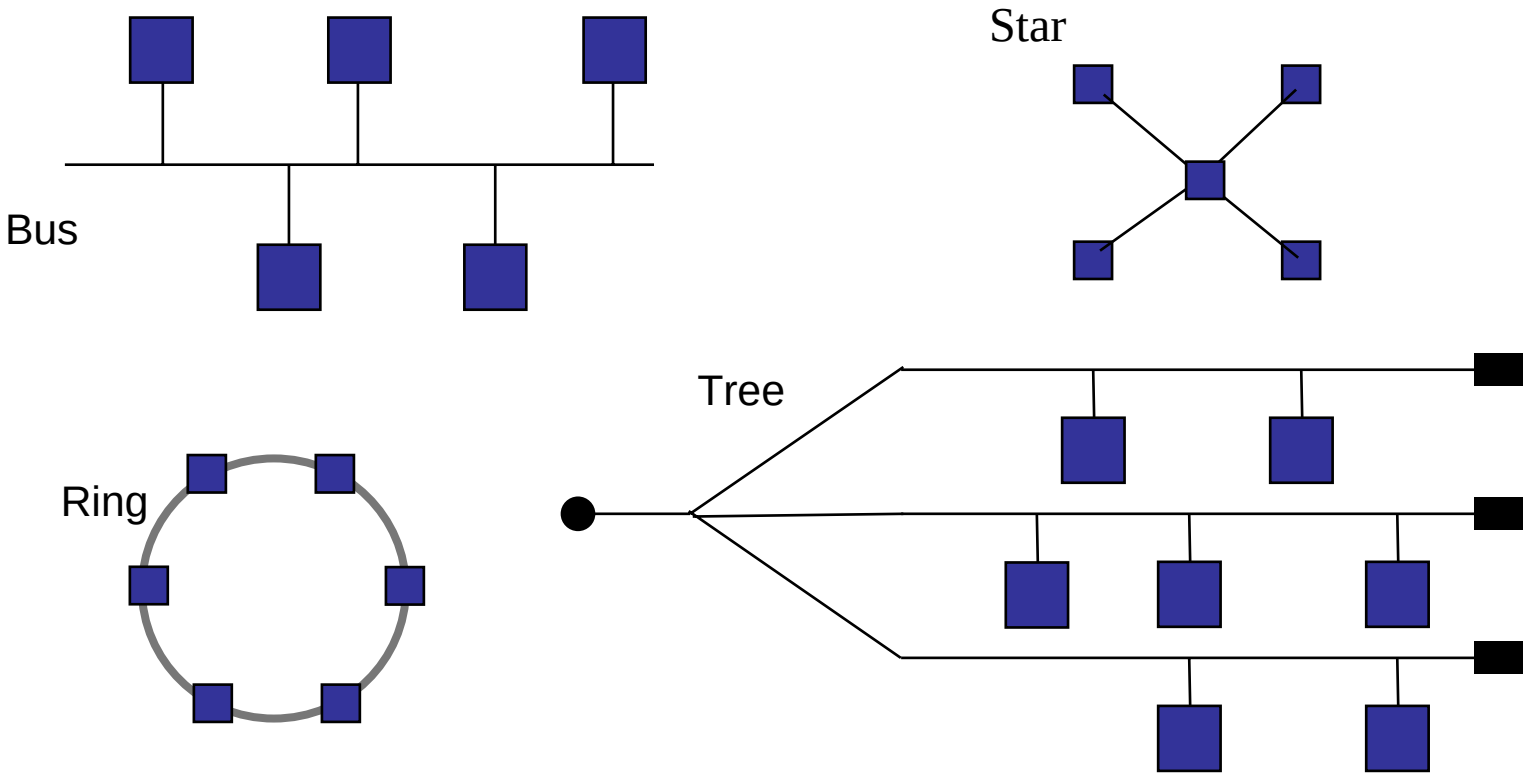
LAN Logical Topologies

- the basic LAN function is to reach all other stations without complex routing



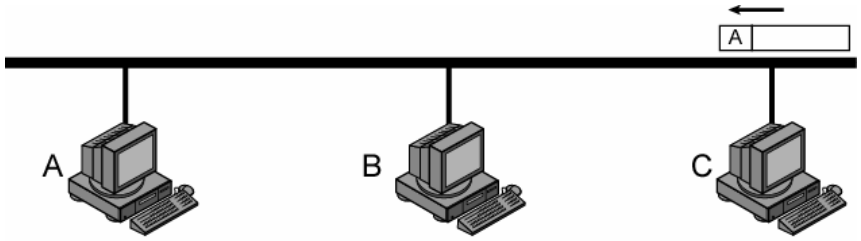
LAN Architecture

- Topologies
  - Tree, Ring, Star
  - Bus
    - Special case of tree (one trunk, no branches)
- Transmission medium
- Layout
- Medium access control

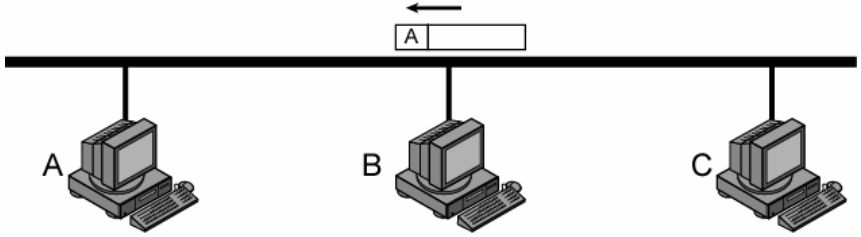
Bus and Tree

- Multipoint medium
- Transmission propagates throughout medium
- Heard by all stations
  - Need to identify target station
    - Each station has unique address
- Full duplex connection between station and tap
  - Allows for transmission and reception
- Need to regulate transmission
  - To avoid collisions
  - To avoid hogging
    - Data in small blocks - frames
- Terminator absorbs frames at end of medium

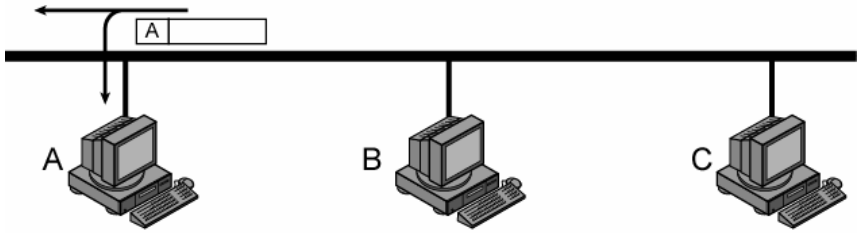
Frame Transmission on Bus LAN



C transmits frame addressed to A



Frame is not addressed to B; B ignores it

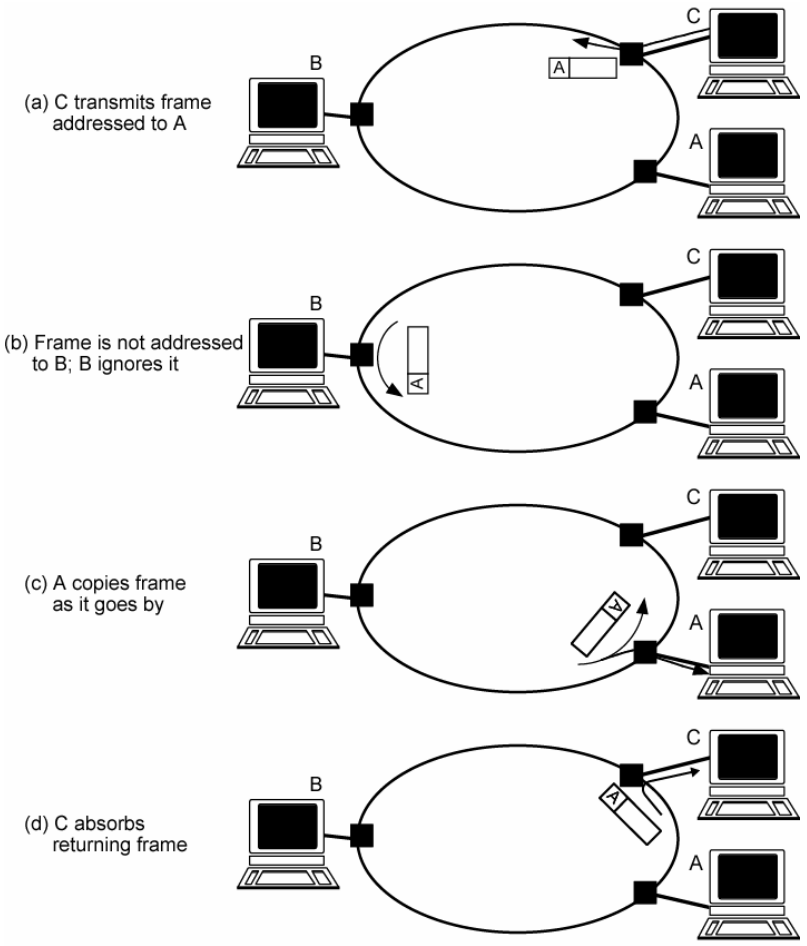


A copies frame as it goes by

Ring Topology

- Repeaters joined by point to point links in closed loop
  - Receive data on one link and retransmit on another
  - Links unidirectional
  - Stations attach to repeaters
- Data in frames
  - Circulate past all stations
  - Destination recognizes address and copies frame
  - Frame circulates back to source where it is removed
- Media access control determines when station can insert frame

Frame Transmission Ring LAN



Star Topology

- Each station connected directly to central node
  - Usually via two point to point links
- Central node can broadcast
  - Physical star, logical bus
  - Only one station can transmit at a time
- Central node can act as frame switch

Choice of Topology

- Reliability
- Expandability
- Performance
- Needs considering in context of:
  - Medium
  - Wiring layout
  - Access control

Bus LAN Transmission Media

- Twisted pair
  - Early LANs used voice grade cable
  - Didn't scale for fast LANs
  - Not used in bus LANs now
- Baseband coaxial cable
  - Uses digital signalling
  - Original Ethernet
- Broadband coaxial cable
  - As in cable TV systems
  - Analog signals at radio frequencies
  - Expensive, hard to install and maintain
  - No longer used in LANs
- Optical fiber
  - Expensive taps
  - Better alternatives available
  - Not used in bus LANs
- All hard to work with compared with star topology twisted pair
- Coaxial baseband still used but not often in new installations

Ring and Star Usage

- Ring
  - Very high speed links over long distances
  - Single link or repeater failure disables network
- Star
  - Uses natural layout of wiring in building
  - Best for short distances
  - High data rates for small number of devices

Choice of Medium

- Constrained by LAN topology
- Capacity
- Reliability
- Types of data supported
- Environmental scope

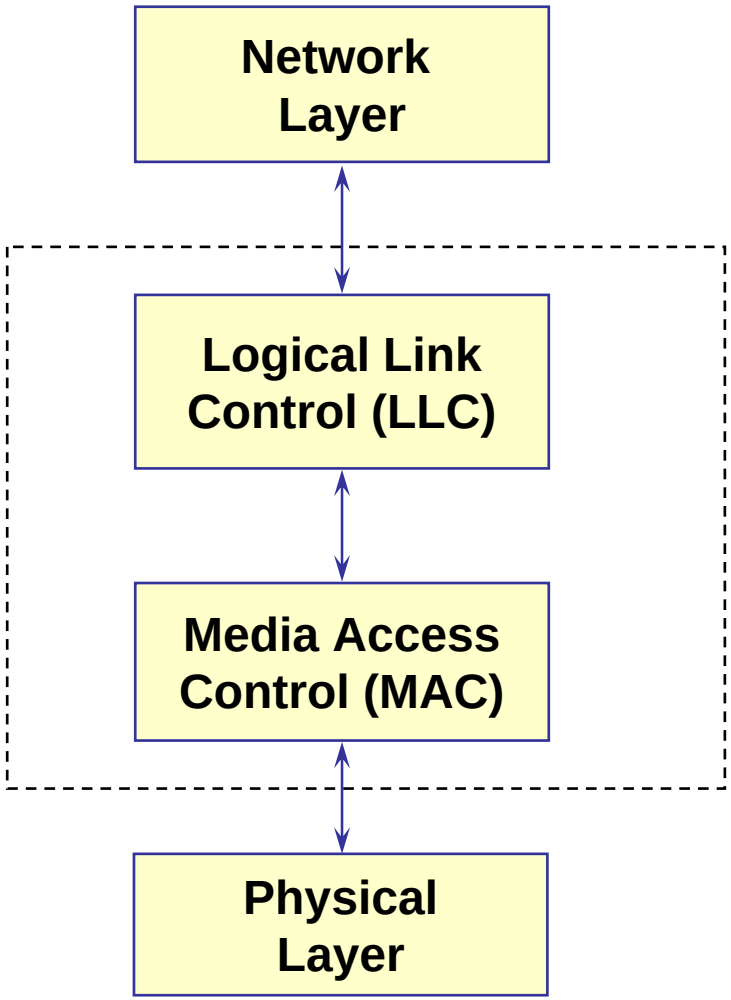
Media Available

- Voice grade unshielded twisted pair (UTP)
  - Cat 3
  - Cheap
  - Well understood
  - Use existing telephone wiring in office building
  - Low data rates
- Shielded twisted pair and baseband coaxial
  - More expensive than UTP but higher data rates
- Broadband cable
  - Still more expensive and higher data rate
- High performance UTP
  - Cat 5 and above
  - High data rate for small number of devices
  - Switched star topology for large installations
- Optical fiber
  - Electromagnetic isolation
  - High capacity
  - Small size
  - High cost of components
  - High skill needed to install and maintain
    - Prices are coming down as demand and product range increases

Protocol Architecture

- Lower layers of OSI model
- IEEE 802 reference model
- Physical
- Logical link control (LLC)
- Media access control (MAC)

DLC Refinement for Local Area Networks



Data Link Control Layer

Scope of IEEE 802

802 Layers - Physical

- Encoding/decoding
- Preamble generation/removal
- Bit transmission/reception
- Transmission medium and topology

802 Layers – Logical Link Control

- Interface to higher layers
- Flow and error control

Logical Link Control

- Transmission of link level PDUs between two stations
- Must support multiaccess, shared medium
- Relieved of some link access details by MAC layer
- Addressing involves specifying source and destination LLC users
  - Referred to as service access points (SAP)
  - Typically higher level protocol

LLC Services

- Based on HDLC
- Unacknowledged connectionless service
- Connection mode service
- Acknowledged connectionless service

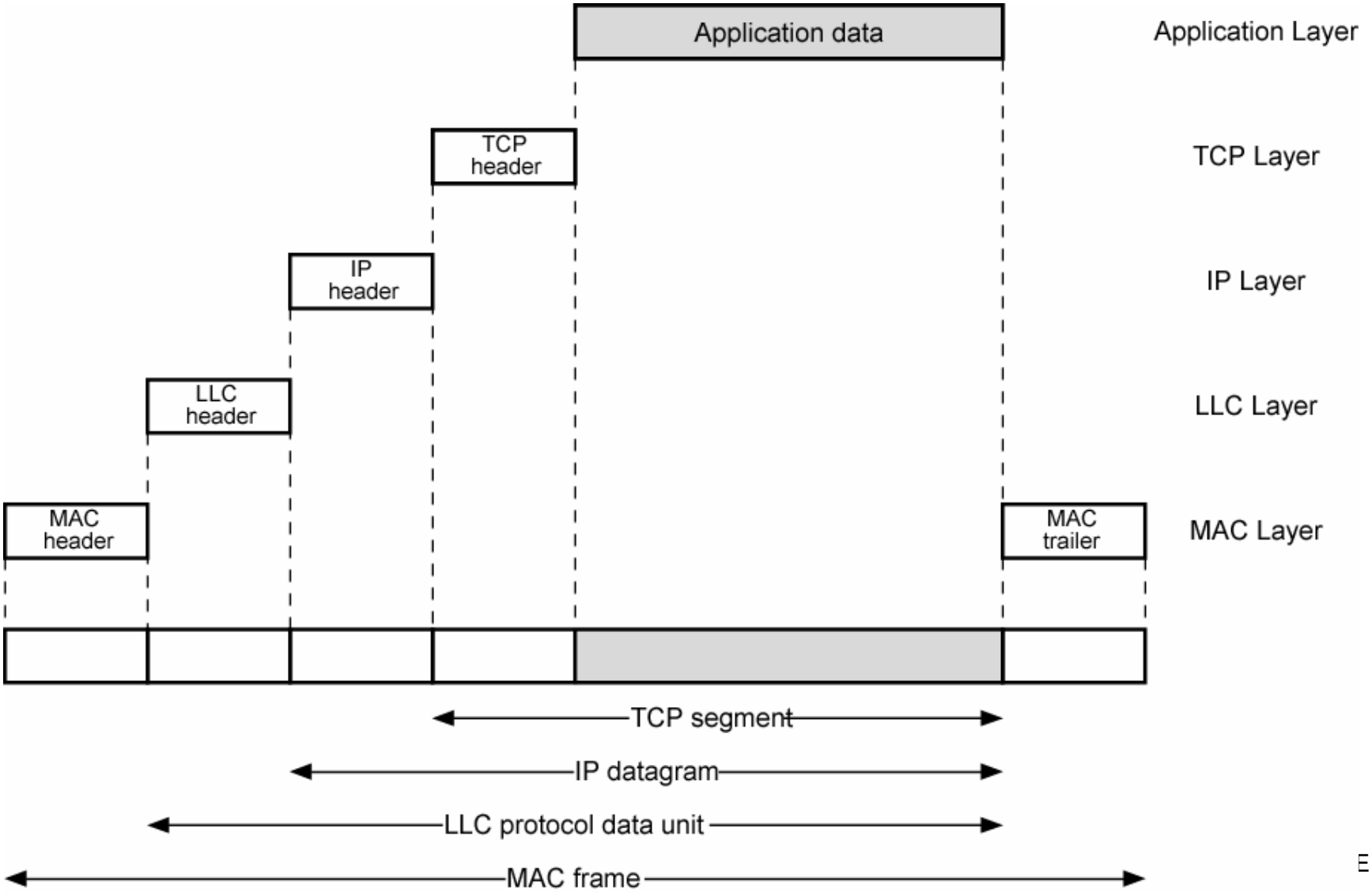
LLC Protocol

- Modeled after HDLC
- Asynchronous balanced mode to support connection mode LLC service (type 2 operation)
- Unnumbered information PDUs to support Acknowledged connectionless service (type 1)
- Multiplexing using LSAPs

Media Access Control

- Assembly of data into frame with address and error detection fields
- Disassembly of frame
  - Address recognition
  - Error detection
- Govern access to transmission medium
  - Not found in traditional layer 2 data link control
- For the same LLC, several MAC options may be available

LAN Protocols in Context



Media Access Control

- Where
  - Central
    - Greater control
    - Simple access logic at station
    - Avoids problems of co-ordination
    - Single point of failure
    - Potential bottleneck
  - Distributed
- How
  - Synchronous
    - Specific capacity dedicated to connection
  - Asynchronous
    - In response to demand

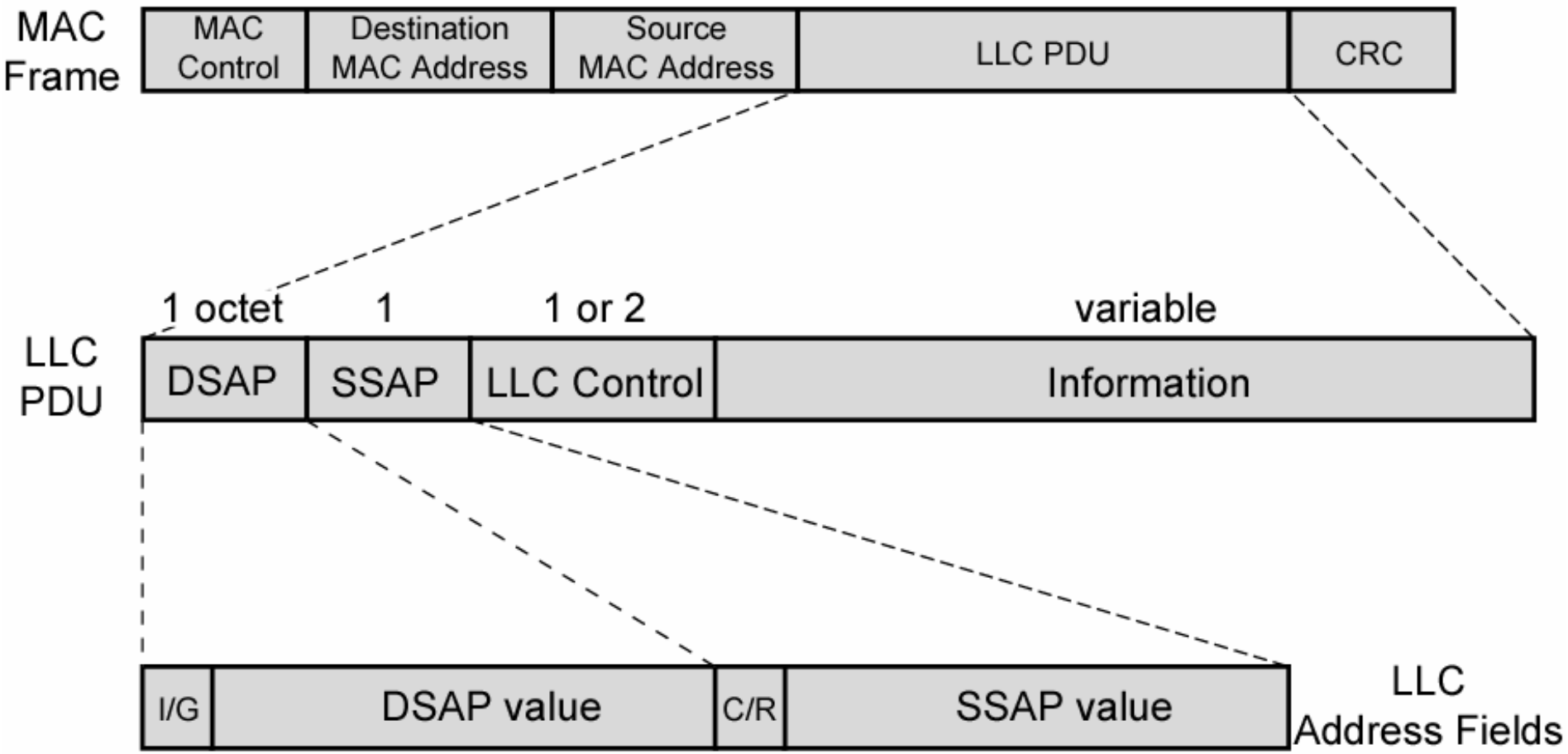
Asynchronous Systems

- Round robin
  - Good if many stations have data to transmit over extended period
- Reservation
  - Good for stream traffic
- Contention
  - Good for bursty traffic
  - All stations contend for time
  - Distributed
  - Simple to implement
  - Efficient under moderate load
  - Tend to collapse under heavy load

MAC Frame Format

- MAC layer receives data from LLC layer
- MAC control
- Destination MAC address
- Source MAC address
- LLS
- CRC
- MAC layer detects errors and discards frames
- LLC optionally retransmits unsuccessful frames

Generic MAC Frame Format



I/G = Individual/Group
C/R = Command/Response

Ethernet Frame Structure

64 bits	48 bits	48 bits	16 bits	368 to 12,000 bits	32 bits
Preamble	Destination	Source	Type	Frame Data	CRC

- 48-bit address is installed at the factory for each interface
- Destination must be on the same LAN as the Source
- Frame Type describes the payload; thus each frame is self-identifying (example: TCP/IP packet)
- Minimum frame size = slot time = 512 bits
- Interframe gap = 96 bits
- Jamming signal size = 32 - 48 bits

Bridges

- Ability to expand beyond single LAN
- Provide interconnection to other LANs/WANs
- Use Bridge or router
- Bridge is simpler
  - Connects similar LANs
  - Identical protocols for physical and link layers
  - Minimal processing
- Router more general purpose
  - Interconnect various LANs and WANs
  - see later

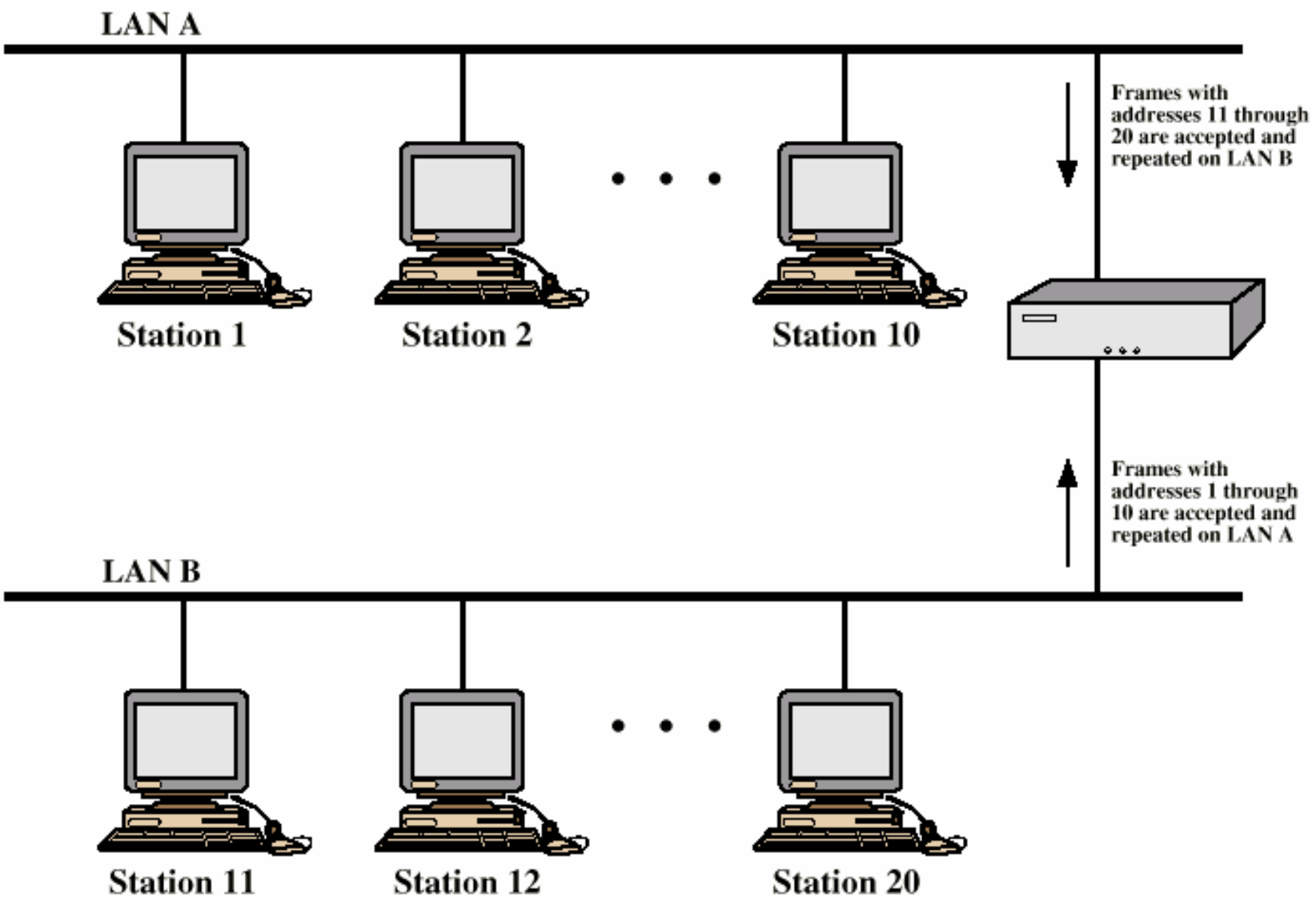
Why Bridge?

- Reliability
- Performance
- Security
- Geography

Functions of a Bridge

- Read all frames transmitted on one LAN and accept those address to any station on the other LAN
- Using MAC protocol for second LAN, retransmit each frame
- Do the same the other way round

Bridge Operation



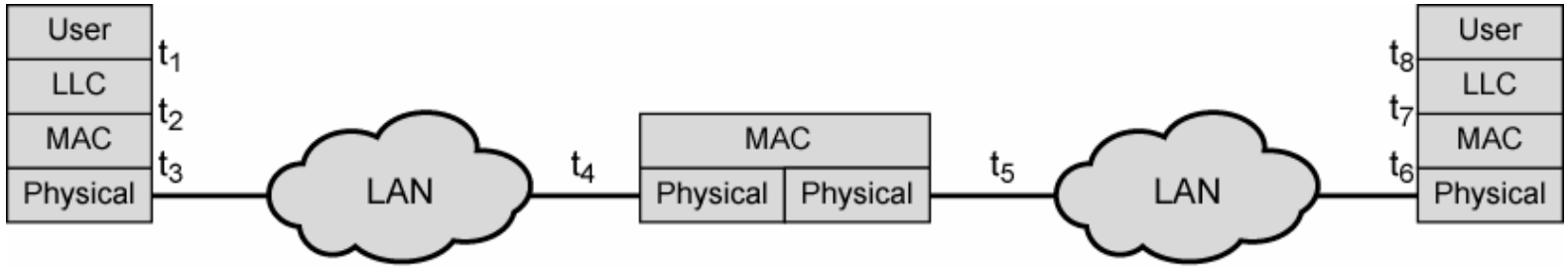
Bridge Design Aspects

- No modification to content or format of frame
- No encapsulation
- Exact bitwise copy of frame
- Minimal buffering to meet peak demand
- Contains routing and address intelligence
  - Must be able to tell which frames to pass
  - May be more than one bridge to cross
- May connect more than two LANs
- Bridging is transparent to stations
  - Appears to all stations on multiple LANs as if they are on one single LAN

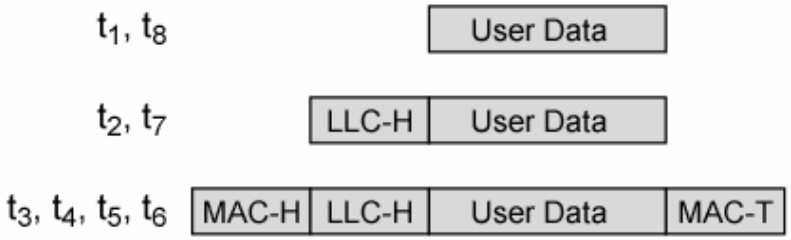
Bridge Protocol Architecture

- IEEE 802.1D
- MAC level
  - Station address is at this level
- Bridge does not need LLC layer
  - It is relaying MAC frames
- Can pass frame over external comms system
  - e.g. WAN link
  - Capture frame
  - Encapsulate it
  - Forward it across link
  - Remove encapsulation and forward over LAN link

Connection of Two LANs



(a) Architecture

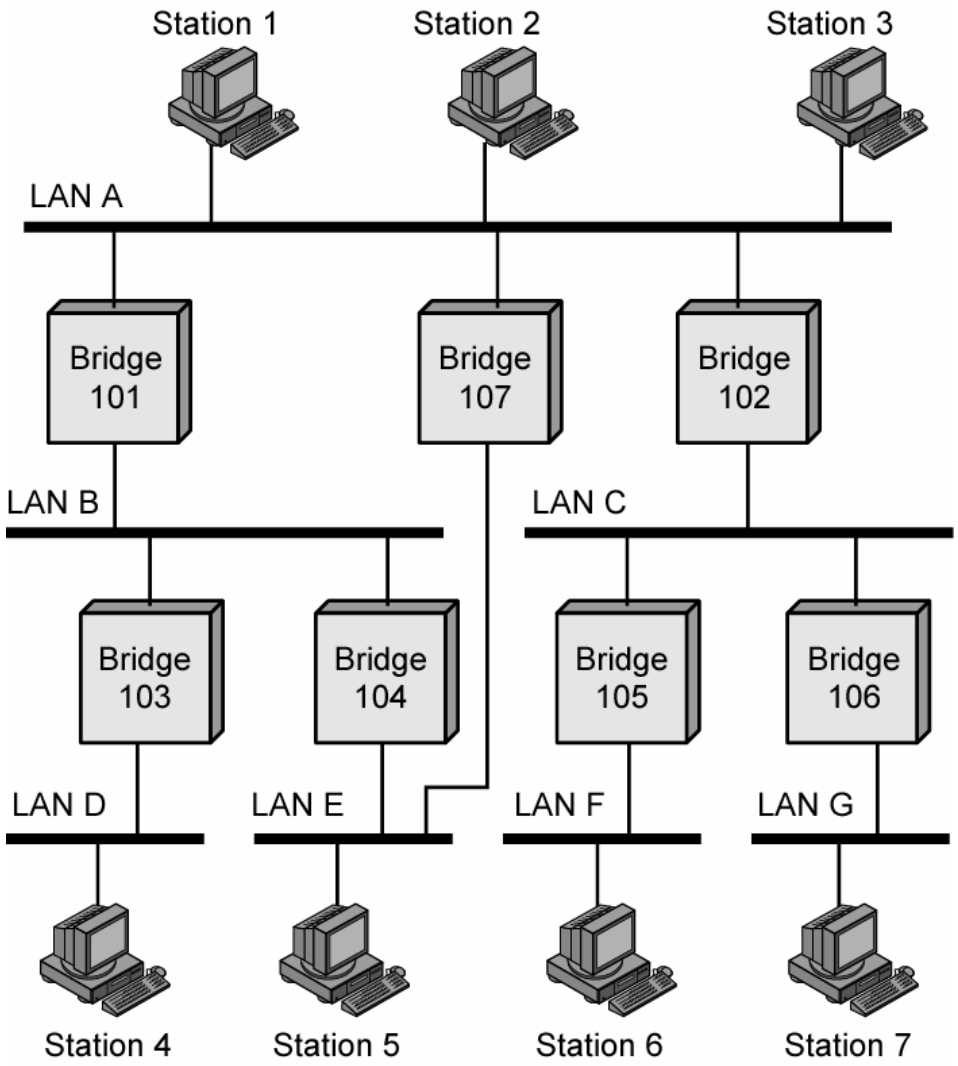


(b) Operation

Fixed Routing

- Complex large LANs need alternative routes
  - Load balancing
  - Fault tolerance
- Bridge must decide whether to forward frame
- Bridge must decide which LAN to forward frame on
- Routing selected for each source-destination pair of LANs
  - Done in configuration
  - Usually least hop route
  - Only changed when topology changes

Bridges and LANs with Alternative Routes



Spanning Tree

- Bridge automatically develops routing table
- Automatically update in response to changes
- Frame forwarding
- Address learning
- Loop resolution

Frame forwarding

- Maintain forwarding database for each port
  - List station addresses reached through each port
- For a frame arriving on port X:
  - Search forwarding database to see if MAC address is listed for any port except X
  - If address not found, forward to all ports except X
  - If address listed for port Y, check port Y for blocking or forwarding state
    - Blocking prevents port from receiving or transmitting
  - If not blocked, transmit frame through port Y

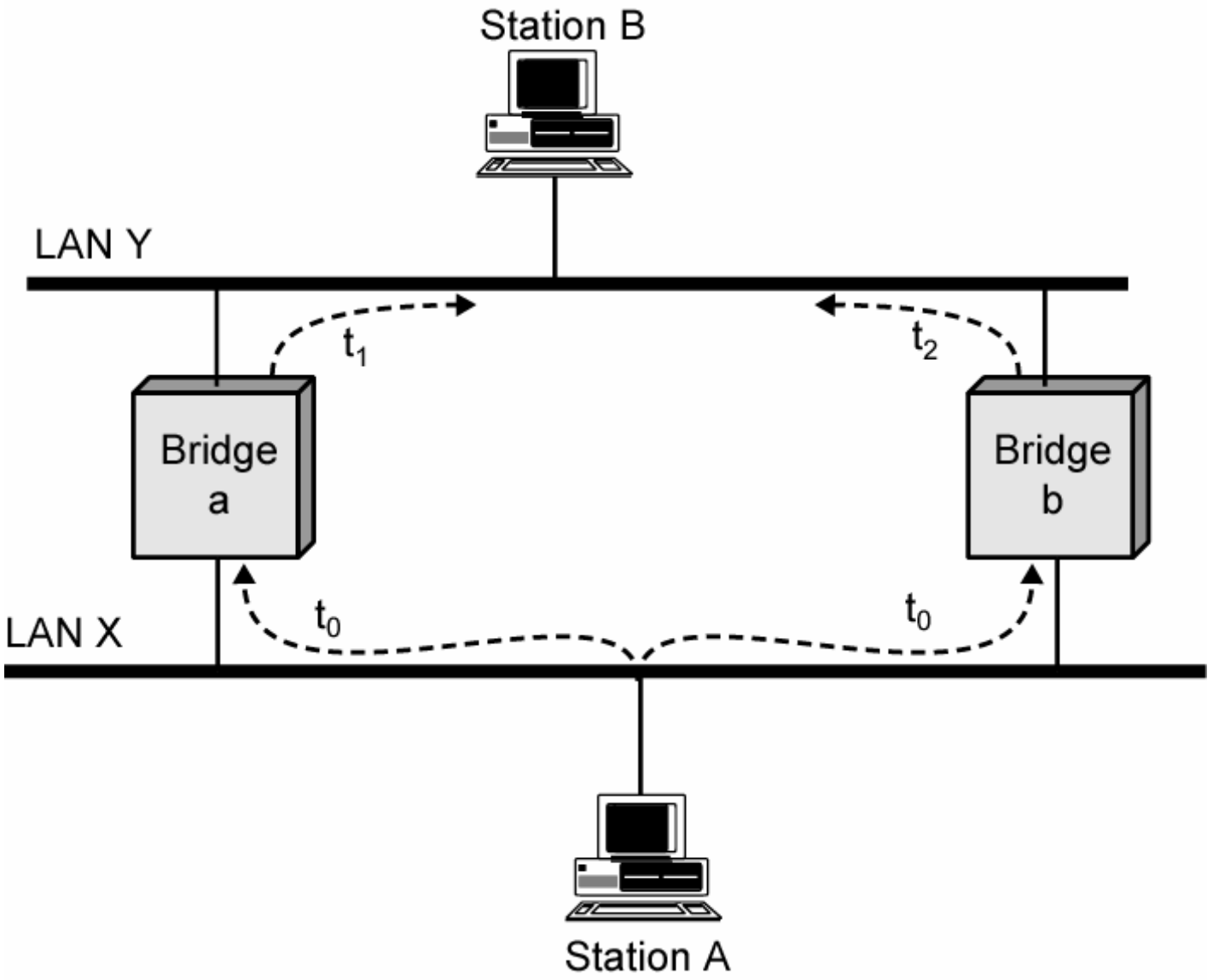
Address Learning

- Can preload forwarding database
- Can be learned
- When frame arrives at port X, it has come form the LAN attached to port X
- Use the source address to update forwarding database for port X to include that address
- Timer on each entry in database
- Each time frame arrives, source address checked against forwarding database

Spanning Tree Algorithm

- Address learning works for tree layout
  - i.e. no closed loops
- For any connected graph there is a spanning tree that maintains connectivity but contains no closed loops
- Each bridge assigned unique identifier
- Exchange between bridges to establish spanning tree

Loop of Bridges



Layer 2 and Layer 3 Switches

- Now many types of devices for interconnecting LANs
- Beyond bridges and routers
- Layer 2 switches
- Layer 3 switches

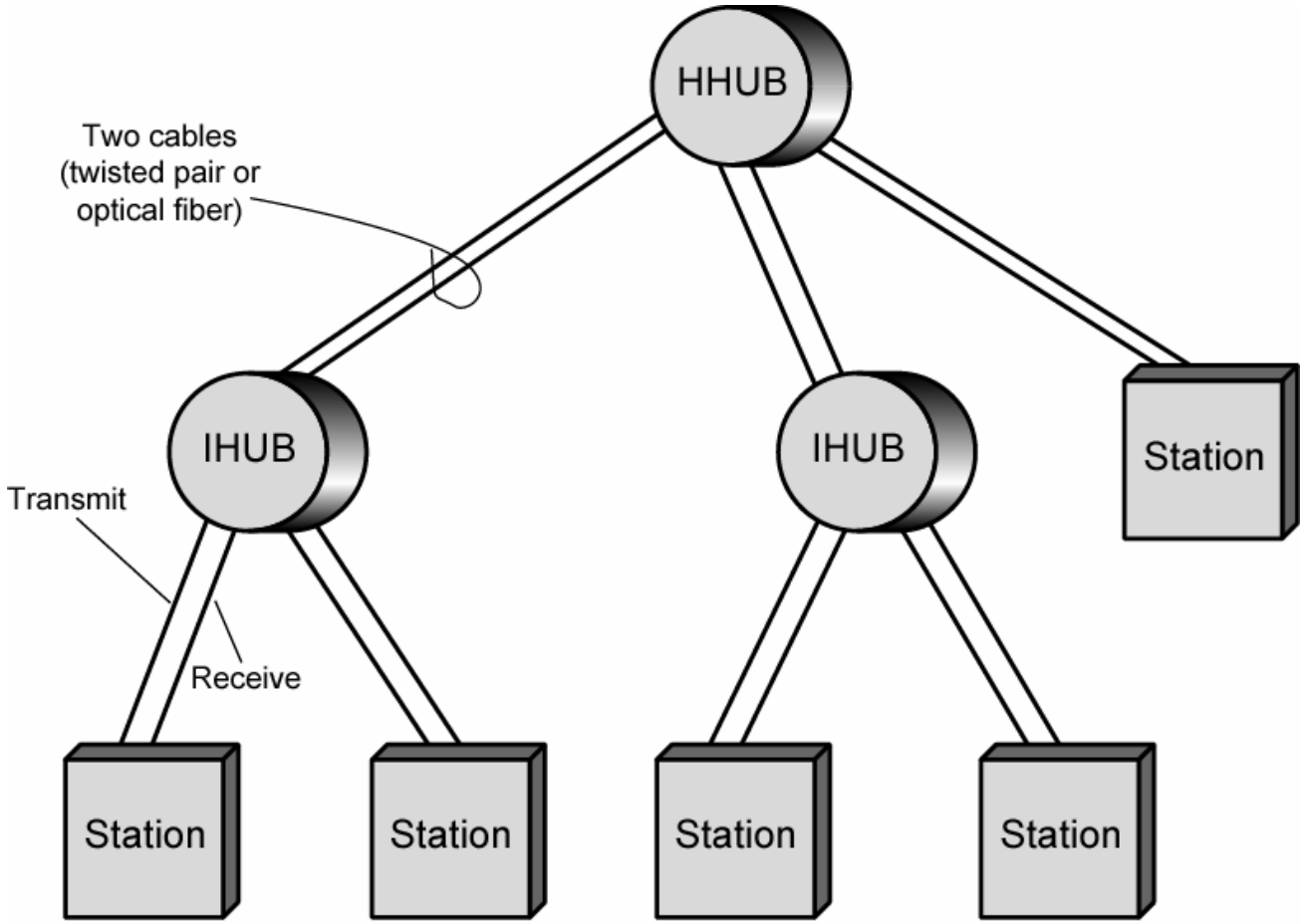
Hubs

- Active central element of star layout
- Each station connected to hub by two lines
  - Transmit and receive
- Hub acts as a repeater
- When single station transmits, hub repeats signal on outgoing line to each station
- Line consists of two unshielded twisted pairs
- Limited to about 100 m
  - High data rate and poor transmission qualities of UTP
- Optical fiber may be used
  - Max about 500 m
- Physically star, logically bus
- Transmission from any station received by all other stations
- If two stations transmit at the same time, collision

Hub Layouts

- Multiple levels of hubs cascaded
- Each hub may have a mixture of stations and other hubs attached to from below
- Fits well with building wiring practices
  - Wiring closet on each floor
  - Hub can be placed in each one
  - Each hub services stations on its floor

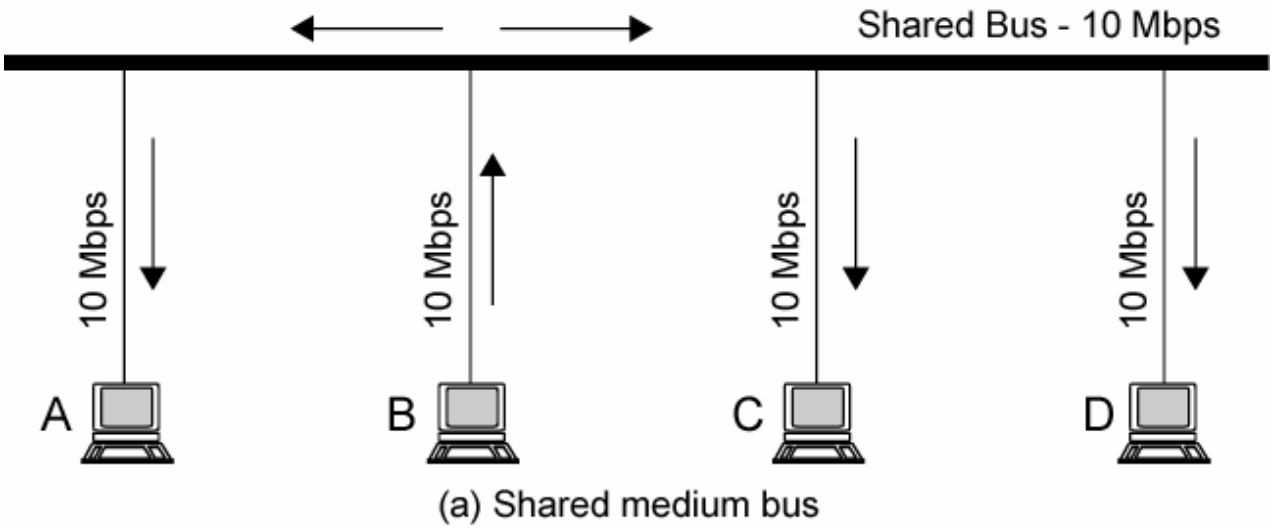
Two Level Star Topology



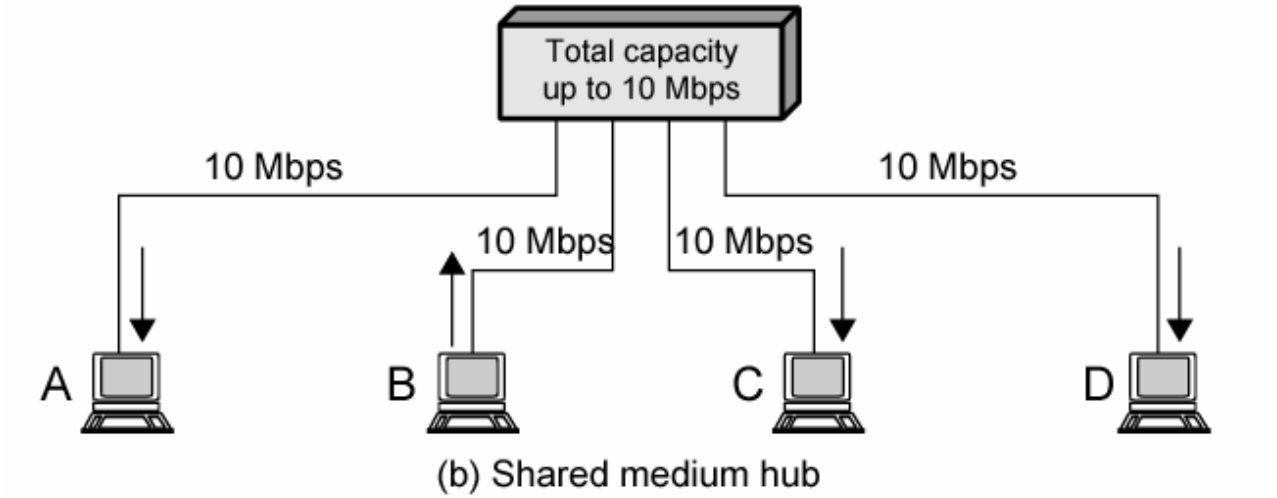
Buses and Hubs

- Bus configuration
  - All stations share capacity of bus (e.g. 10Mbps)
  - Only one station transmitting at a time
- Hub uses star wiring to attach stations to hub
  - Transmission from any station received by hub and retransmitted on all outgoing lines
  - Only one station can transmit at a time
  - Total capacity of LAN is 10 Mbps
- Improve performance with layer 2 switch

Shared Medium Bus and Hub

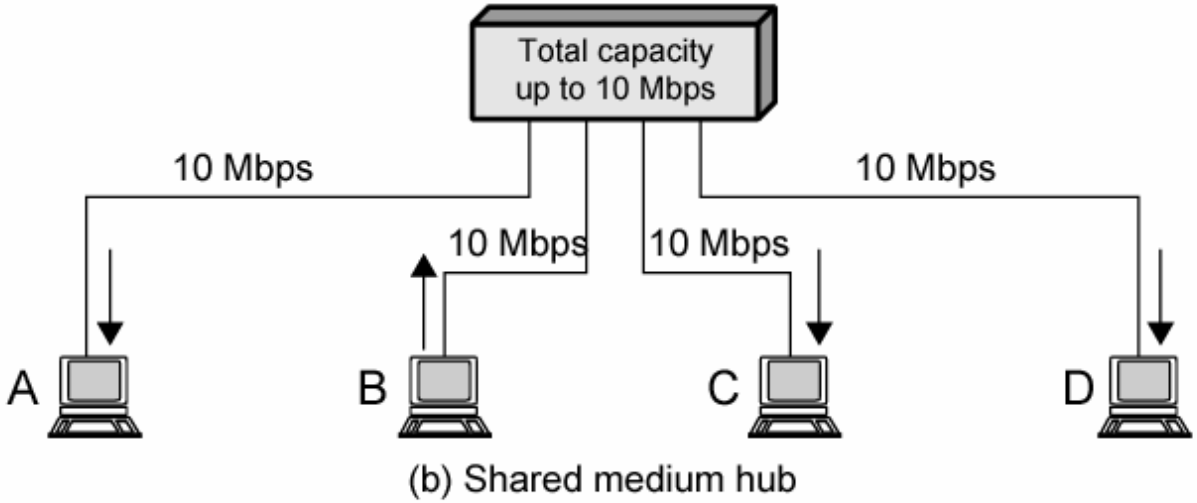


(a) Shared medium bus

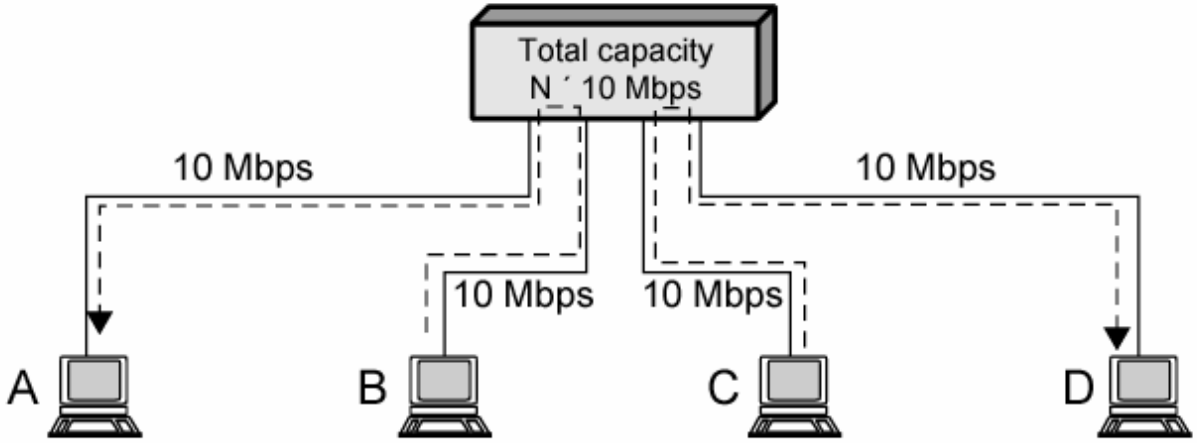


(b) Shared medium hub

Shared Medium Hub and Layer 2 Switch



(b) Shared medium hub



(c) Layer 2 switch

Layer 2 Switches

- Central hub acts as switch
- Incoming frame from particular station switched to appropriate output line
- Unused lines can switch other traffic
- More than one station transmitting at a time
- Multiplying capacity of LAN

Layer 2 Switch Benefits

- No change to attached devices to convert bus LAN or hub LAN to switched LAN
- For Ethernet LAN, each device uses Ethernet MAC protocol
- Device has dedicated capacity equal to original LAN
  - Assuming switch has sufficient capacity to keep up with all devices
  - For example if switch can sustain throughput of 20 Mbps, each device appears to have dedicated capacity for either input or output of 10 Mbps
- Layer 2 switch scales easily
  - Additional devices attached to switch by increasing capacity of layer 2

Types of Layer 2 Switch

- Store-and-forward switch
  - Accepts frame on input line
  - Buffers it briefly,
  - Then routes it to appropriate output line
  - Delay between sender and receiver
  - Boosts integrity of network
- Cut-through switch
  - Takes advantage of destination address appearing at beginning of frame
  - Switch begins repeating frame onto output line as soon as it recognizes destination address
  - Highest possible throughput
  - Risk of propagating bad frames
    - Switch unable to check CRC prior to retransmission

Layer 2 Switch v Bridge

- Layer 2 switch can be viewed as full-duplex hub
- Can incorporate logic to function as multiport bridge
- Bridge frame handling done in software
- Switch performs address recognition and frame forwarding in hardware
- Bridge only analyzes and forwards one frame at a time
- Switch has multiple parallel data paths
  - Can handle multiple frames at a time
- Bridge uses store-and-forward operation
- Switch can have cut-through operation
- Bridge suffered commercially
  - New installations typically include layer 2 switches with bridge functionality rather than bridges

Problems with Layer 2 Switches (1)

- As number of devices in building grows, layer 2 switches reveal some inadequacies
- Broadcast overload
- Lack of multiple links
- Set of devices and LANs connected by layer 2 switches have flat address space
  - Allusers share common MAC broadcast address
  - If any device issues broadcast frame, that frame is delivered to all devices attached to network connected by layer 2 switches and/or bridges
  - In large network, broadcast frames can create big overhead
  - Malfunctioning device can create broadcast storm
    - Numerous broadcast frames clog network

Problems with Layer 2 Switches (2)

- Current standards for bridge protocols dictate no closed loops
  - Only one path between any two devices
  - Impossible in standards-based implementation to provide multiple paths through multiple switches between devices
    - Limits both performance and reliability.
- Solution: break up network into subnetworks connected by routers
- MAC broadcast frame limited to devices and switches contained in single subnetwork
- IP-based routers employ sophisticated routing algorithms
  - Allow use of multiple paths between subnetworks going through different routers

Problems with Routers

- Routers do all IP-level processing in software
  - High-speed LANs and high-performance layer 2 switches pump millions of packets per second
  - Software-based router only able to handle well under a million packets per second
- Solution: layer 3 switches
  - Implement packet-forwarding logic of router in hardware
- Two categories
  - Packet by packet
  - Flow based

Packet by Packet or Flow Based

- Operates insame way as traditional router
- Order of magnitude increase in performance compared to software-based router
- Flow-based switch tries to enhance performance by identifying flows of IP packets
  - Same source and destination
  - Done by observing ongoing traffic or using a special flow label in packet header (IPv6)
  - Once flow is identified, predefined route can be established

Typical Large LAN Organization

- Thousands to tens of thousands of devices
- Desktop systems links 10 Mbps to 100 Mbps
  - Into layer 2 switch
- Wireless LAN connectivity available for mobile users
- Layer 3 switches at local network's core
  - Form local backbone
  - Interconnected at 1 Gbps
  - Connect to layer 2 switches at 100 Mbps to 1 Gbps
- Servers connect directly to layer 2 or layer 3 switches at 1 Gbps
- Lower-cost software-based router provides WAN connection
- Circles in diagram identify separate LAN subnetworks
- MAC broadcast frame limited to own subnetwork