A realization about RSA ep system

Based on the Euler theorem, RSA use some mathematical transformation to ep some data message. It can also be used for digital signature and identity authentication. This algorithm is presented by three young professor from MIT in 1977, and named by their last name, Rivest, Shamir and Adlernan. A fact in number theory field is that though compute the product of two large number is simple, it is extremely hard to factor it back into the primes. Compare with *Diffie-hellman* algorithm, RSA have the obvious superiority because the pair don’t have to involve in the ep process at the same time. Others, it is suit for electronic mail system.

Before 1976, all the ep algorithm is in a same mode, 1),A choose a mode of ep regulation

2)B choose the same regulation to DP it. Because of the same rule, we call this way to *symmetric-key algorithm.* The biggest weakness of this way is that A must tell B the ep rule to B, otherwise, he can’t DP it. But after 1976, *Diffie-hellman* algorithm had been invented, and based on that, RSA appeared. Now, the longest security key which had been cracked length is 768, that means a security key which have a length more than 1024 is almost safe to most average person, if you use a key which have 2048 digit, it is too safe to us.

The principle of this algorithm is based on a little number theory.

1. co-prime relationship

If a pair of positive integer have no common factor except 1, we call this pair of number is co-prime, for example, 13 and 61.

1. Euler function

Given any positive integer n, how many positive integer smaller than n can constitute a relatively prime relationship? The way to calculate this value is Euler function. the calculate method is not difficult, but in order to get the final formula, we must discuss the situation,

1. n=1

f(n)=1.

1. n is a prime number

f(n)=n-1.

1. If n=p^k (p is a prime number, k>1 and is a integer)

f(p^k)=p^k-p^(k-1)

this formula can also written in another form   
 f(p^k)=p^k-p^(k-1)=p^k(1-1/p)

1. If n can decomposed into two mutually integer integers

n=p1\*p2

then f(n)=f(p1p2)=f(p1)f(p2)

for example f(56)=f(8)f(7)=24.

1. By the fact that any positive integer greater than 1 can be written as a series of prime numbers, according to the corollary of article 4,

f(n)=f(Pk1)f(Pk2)……f(Pkr)

and according to the corollary of article 3

f(n)=n(1-1/P1)(1-1/P2)…(1-1/Pr).

This is the calculation process, if we want to calculate 1323’s Euler function.

f(1323)=f(3^3\*7^2)=1323(1-1/3)(1-1/7)=756

The reason we use the Euler formula is that it can greatly simplify some operations, for example 7 and 10 are qualities, according to Euler theorem, 7^(f(10))=1(mod 10), we know that f(10)=4, so we can immediately know 7^4k=1(mode 10)

The latest concept is *Modal elements*. If the two positive integer a and n are prime, we must find a integer b which can let ab-1 divisible by n. then we call b as modal elements to a. it is worth mentioning that modal element are not unique.

That is all the mathematical tools we would use in this paper. Then we would look at the generate steps of key. If Tony and Friday want to perform encrypted communication, how could he generate a public key?

1. Select two prime numbers p and q randomly

In our example, he choose 61 and 53. But we must know that the bigger these prime number is, more difficult to crack to other one.

1. Calculate the product of this pair of number n, the result is 3233.

The length of n is the length of key 3233 written in binary is 110010100001, the length of this number is 12, so the length of the key is 12, in practical application, the length is always 1024, even 2048 in important occasion.

1. Calculate f(n)

By the formula

f(n)=(p-1)(q-1)

f(3233)=60\*52=3120.

1. Choose a integer e randomly, it must between 1 and f\*n) and e must quality with f(n).