

ECS140A
 Assignment 1
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1. (a) We want to generate aabccd. Since the second letter is a and the third is a b, so we have:

$\langle S \rangle \Rightarrow a\langle S \rangle c\langle B \rangle \Rightarrow aa\langle S \rangle c\langle B \rangle c\langle B \rangle \Rightarrow aabc\langle B \rangle c\langle B \rangle \Rightarrow aabc\langle A \rangle cd \Rightarrow aabcccd.$ which is not aabccd, so we can't generate the string from the grammar..

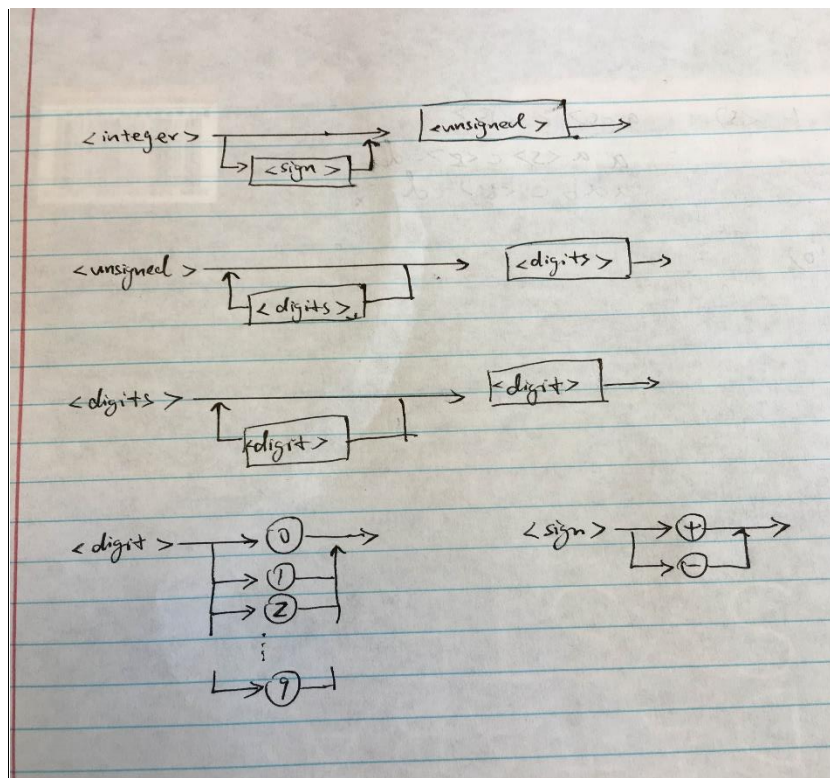
- (b) Since the second letter is c, so $\langle S \rangle \Rightarrow a\langle A \rangle c\langle B \rangle \Rightarrow acc\langle B \rangle.$ However there is no way to generate a b through the grammar. So we can't derive the string from the given grammar.

- (c) $\langle S \rangle \Rightarrow a\langle S \rangle c\langle B \rangle \Rightarrow a\langle A \rangle c\langle A \rangle \Rightarrow ac\langle A \rangle cc\langle A \rangle \Rightarrow accccc$
 Yes, we can generate the form from the grammar.

2. (a)

$\langle \text{integer} \rangle ::= [\langle \text{sign} \rangle] \langle \text{unsigned} \rangle$
 $\langle \text{unsigned} \rangle ::= \{ \langle \text{digits} \rangle \} \langle \text{digits} \rangle$
 $\langle \text{digits} \rangle ::= \{ \langle \text{digit} \rangle \} \langle \text{digit} \rangle$
 $\langle \text{digit} \rangle ::= 0 \mid 1 \mid \dots \mid 9$
 $\langle \text{sign} \rangle ::= + \mid -$

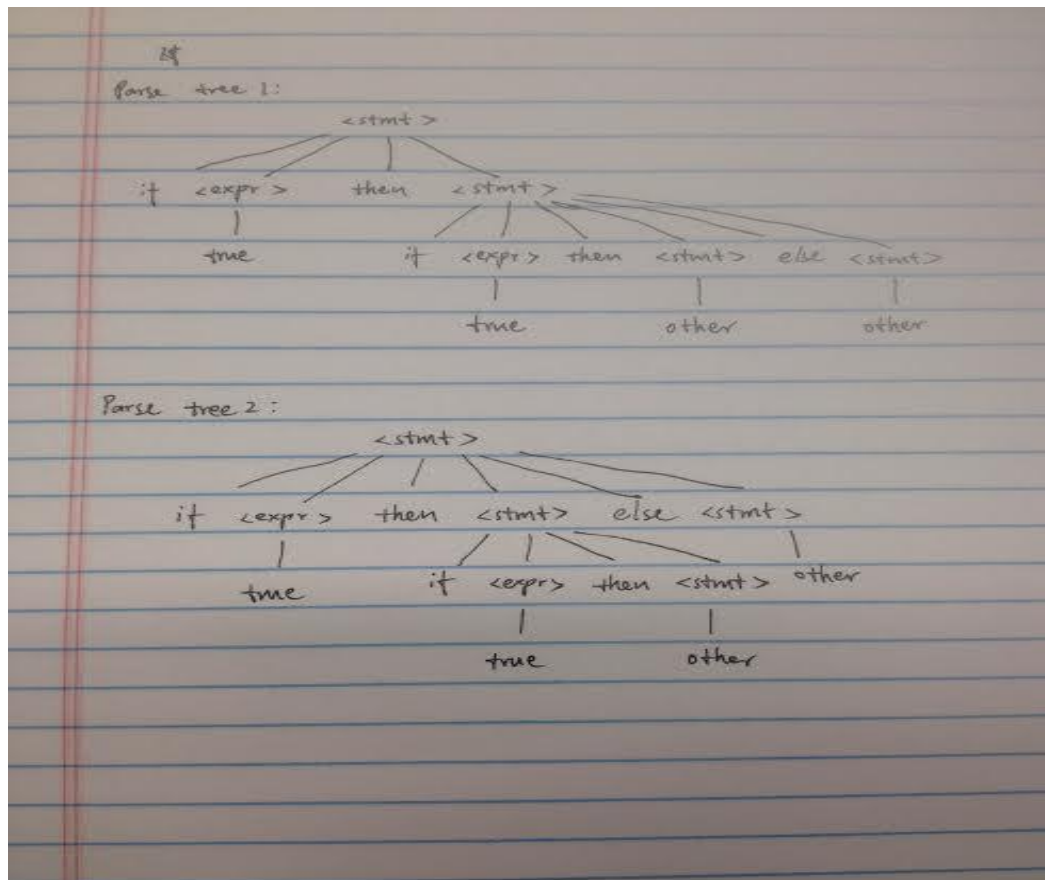
- (b)



3. (a) Language = $\{a^n b^n \mid n \geq 1\}$
 (b) Language = $\{a^n b^m c^k \mid n \geq 1, m \geq 1, k \geq 1\}$
 (c) Language = $\{0^n 1^m \mid n \geq m+1 \text{ or } 0^n 1^m \mid m \geq 2n+1\}$

4.(a)

String: if true then if true then other else other.



(b)

$\langle \text{stmt} \rangle ::= \langle \text{matchingStmt} \rangle \mid \langle \text{unmStmt} \rangle$

$\langle \text{matchingStmt} \rangle ::= \text{if } \langle \text{expr} \rangle \text{ then } \langle \text{matchingStmt} \rangle \text{ else } \langle \text{matchingStmt} \rangle \mid \text{other}$

$\langle \text{unmStmt} \rangle ::= \text{if } \langle \text{expr} \rangle \text{ then } \langle \text{stmt} \rangle \mid \text{if } \langle \text{expr} \rangle \text{ then } \langle \text{matchingStmt} \rangle \text{ else } \langle \text{unmStmt} \rangle$

$\langle \text{expr} \rangle ::= \text{true} \mid \text{false}$

5. (a) BNF:

$\langle \text{exp} \rangle ::= \langle A \rangle \langle \text{exp} \rangle \mid \langle A \rangle$

$\langle A \rangle ::= a \mid \epsilon$

EBNF:

$\langle \text{exp} \rangle ::= \{ \langle A \rangle \}$

$\langle A \rangle ::= a \mid \epsilon$

(b) BNF:

$$\langle \text{exp} \rangle ::= \langle A \rangle \langle \text{exp} \rangle \mid a$$
$$\langle A \rangle ::= a, \mid \epsilon$$

EBNF:

$$\langle \text{exp} \rangle ::= \{ \langle A \rangle \} a$$
$$\langle A \rangle ::= a, \mid \epsilon$$