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# System and Devices (SYS 3)

## Lecture 3: Computer Organization Basics Review

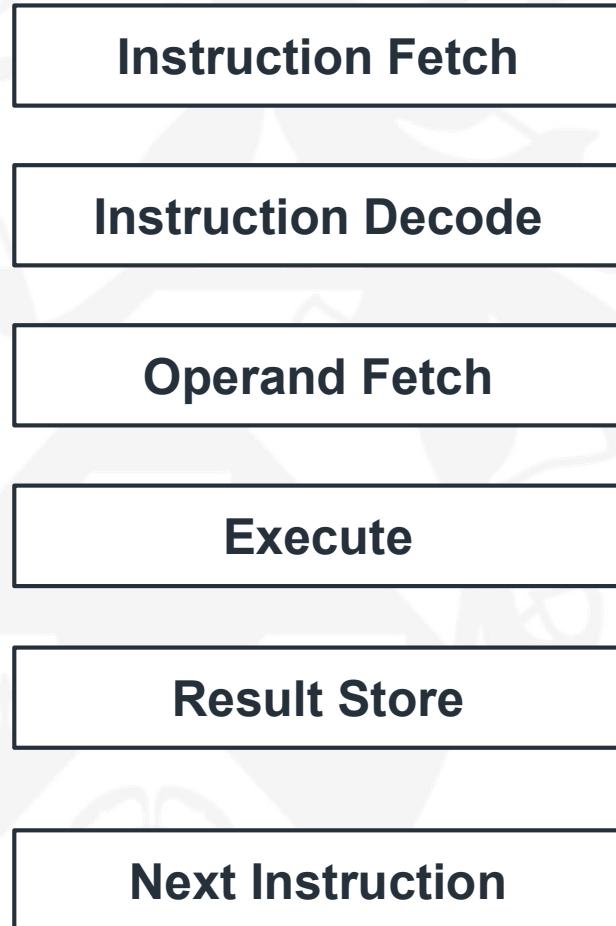
Dr Pengcheng Liu

COM00018I

# Instruction Execution Cycle



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Obtain instruction from program storage

Determine required actions and instruction size

Locate and obtain operand data

Compute result value or status

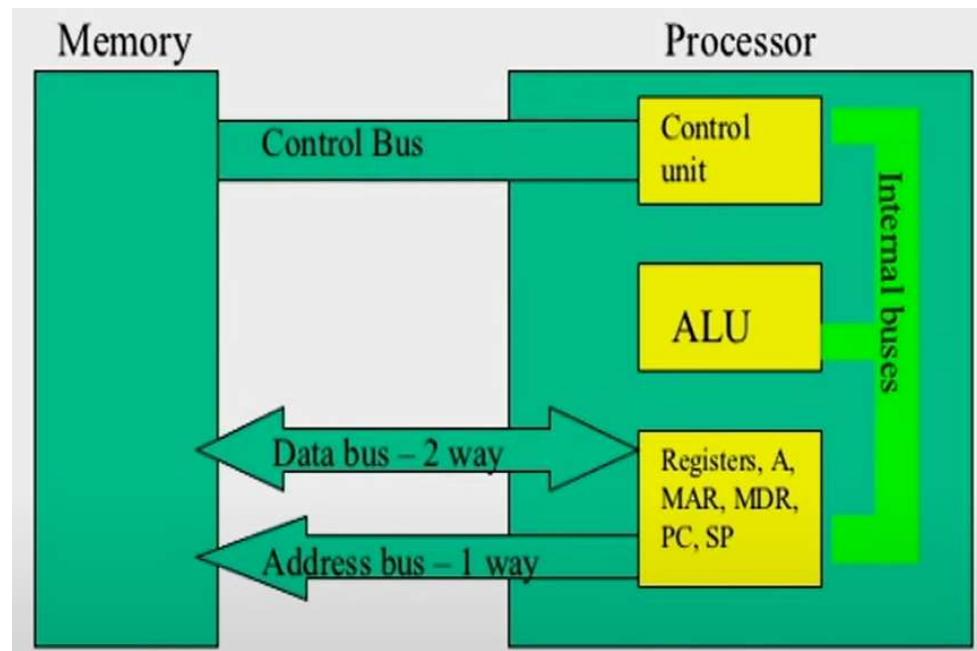
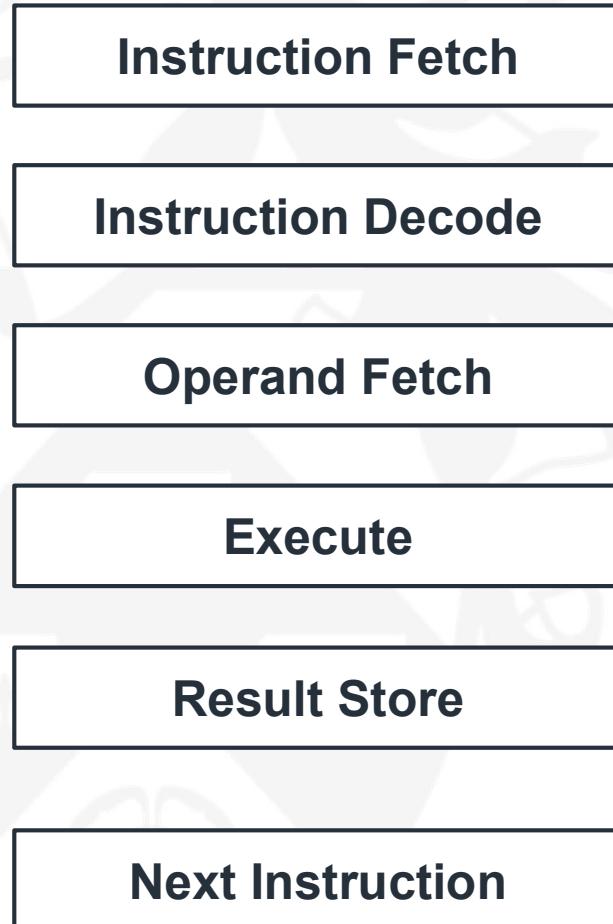
Deposit results in storage for later use

Determine successor instruction

# Processor Memory interaction



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# Processor Memory interaction



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Instruction Fetch

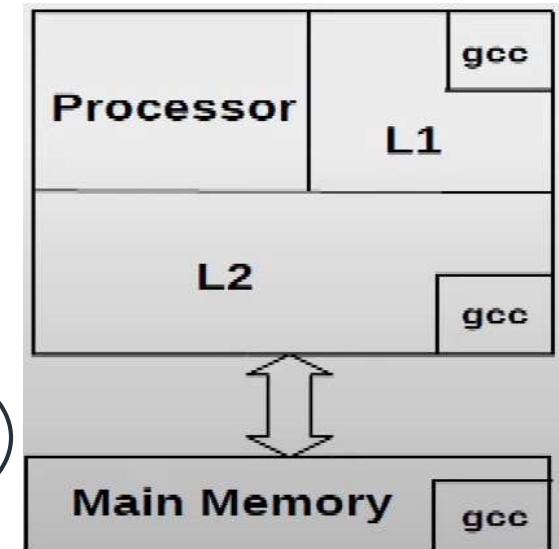
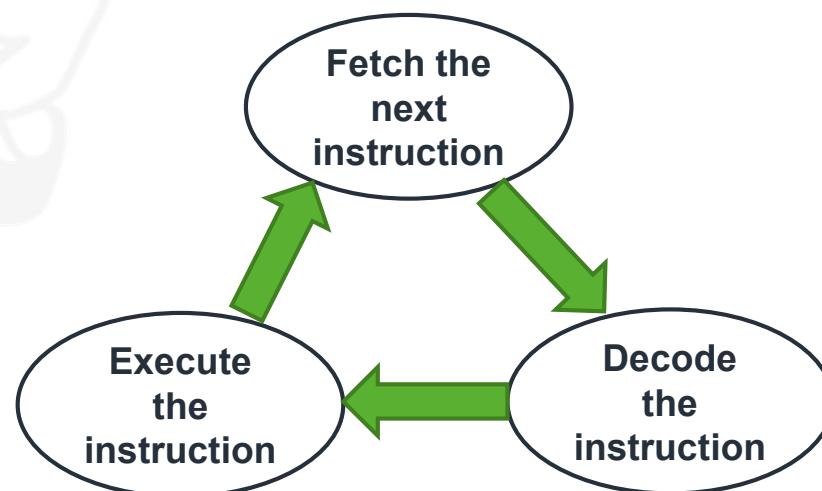
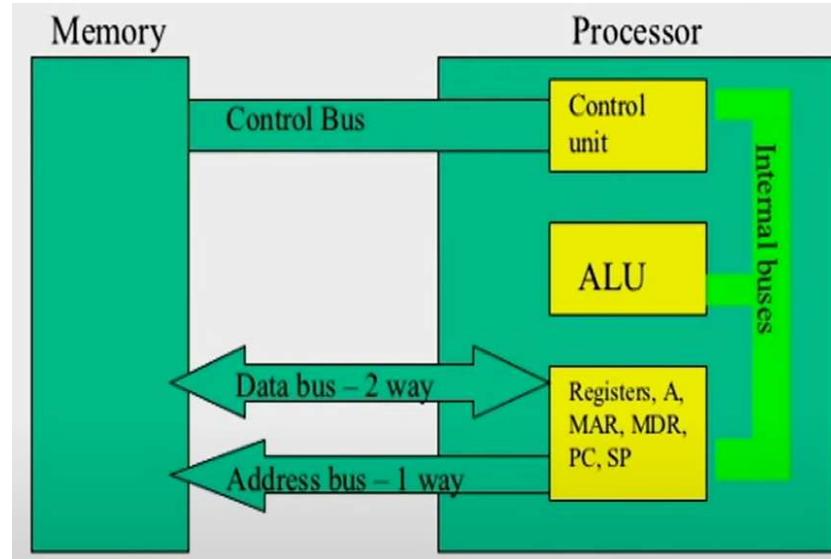
Instruction Decode

Operand Fetch

Execute

Result Store

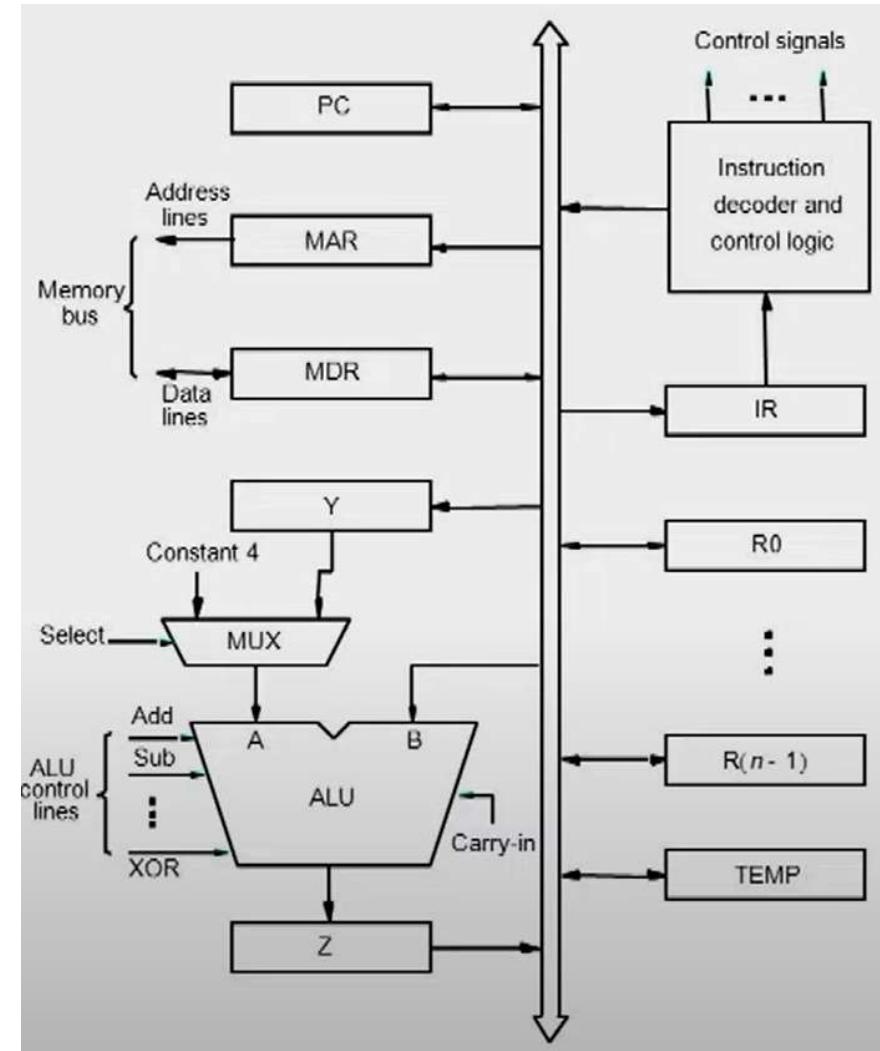
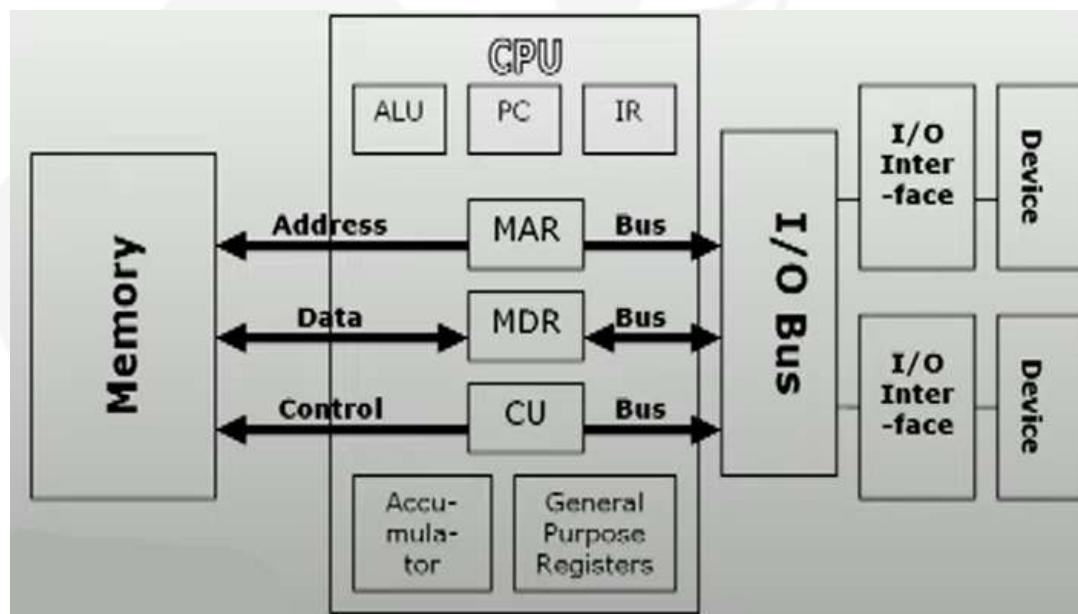
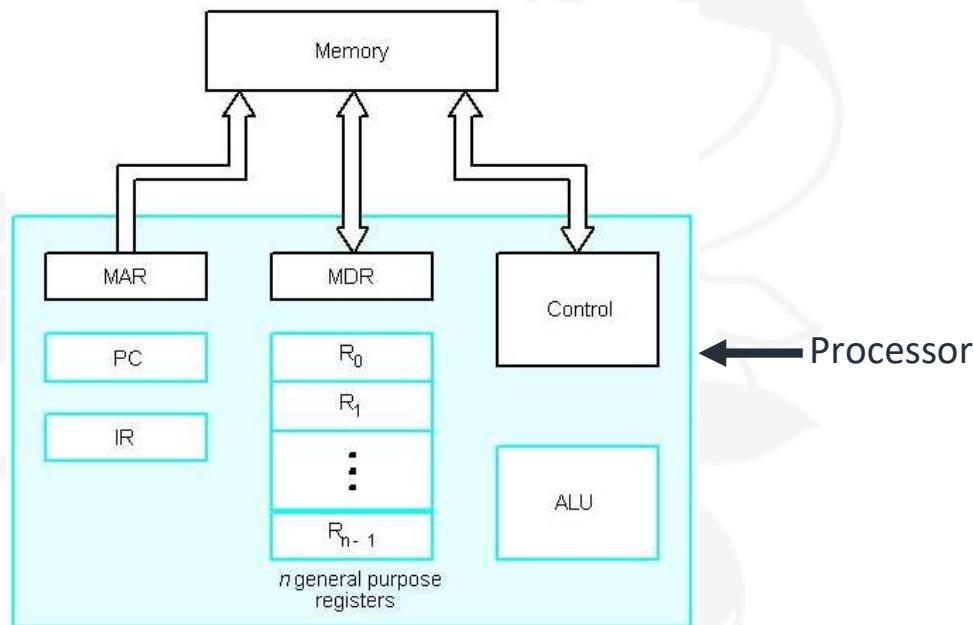
Next Instruction



# Processor Memory interaction



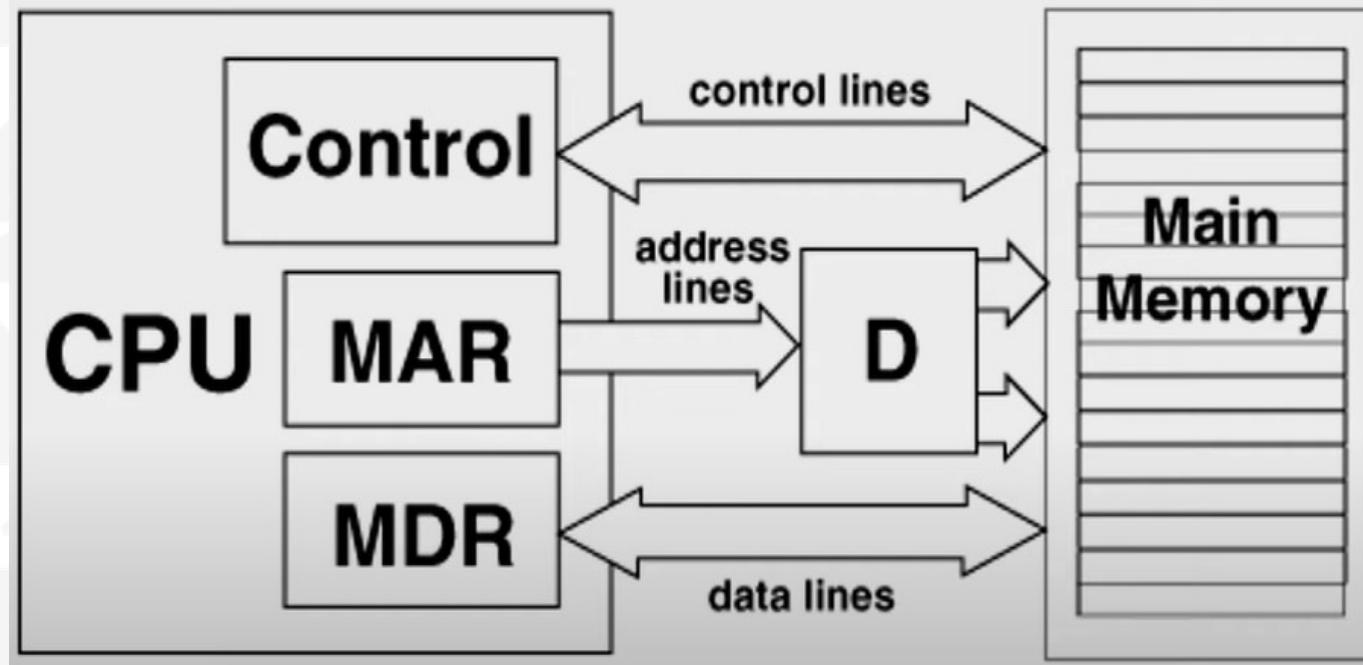
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# Processor Memory interaction



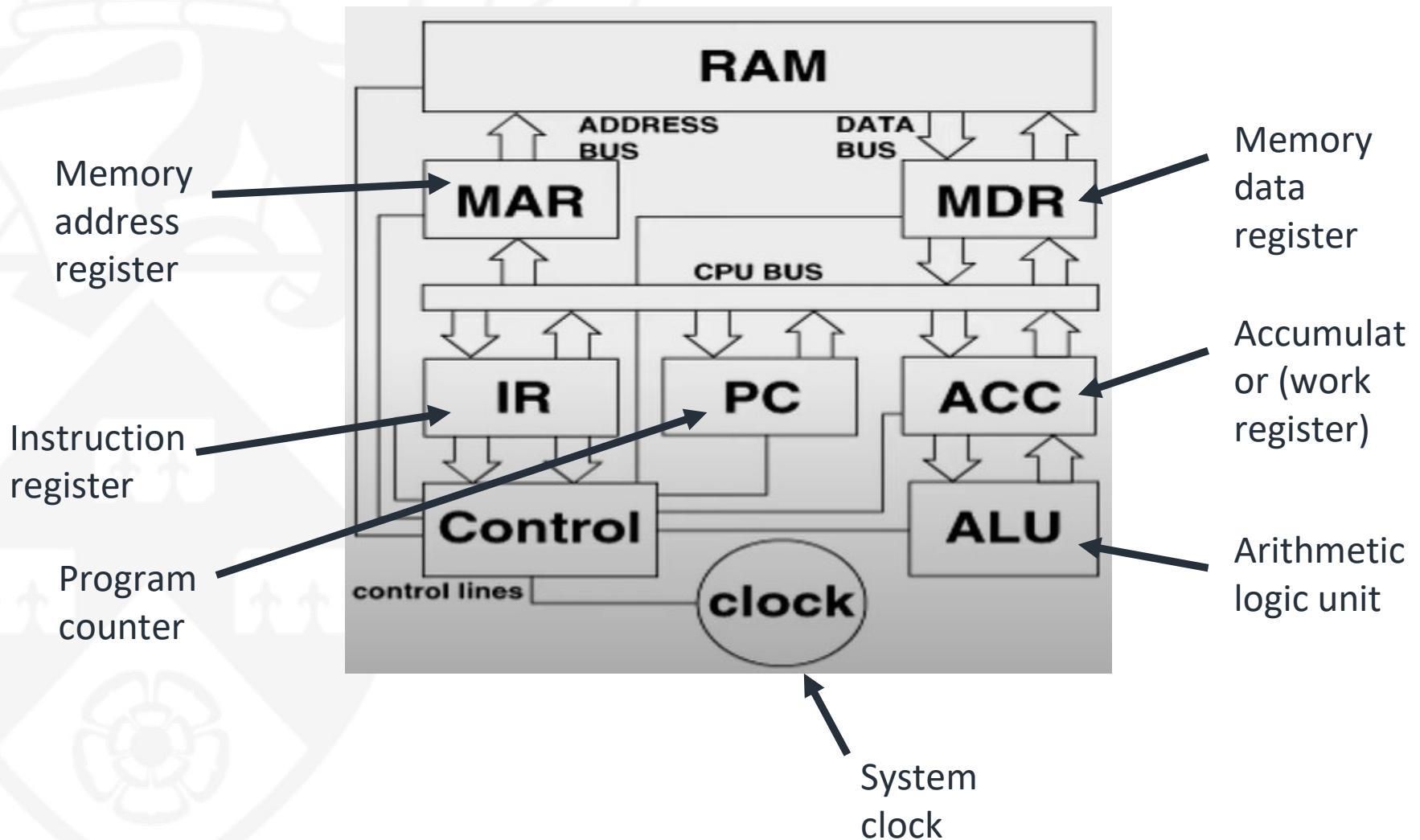
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# Inside the CPU



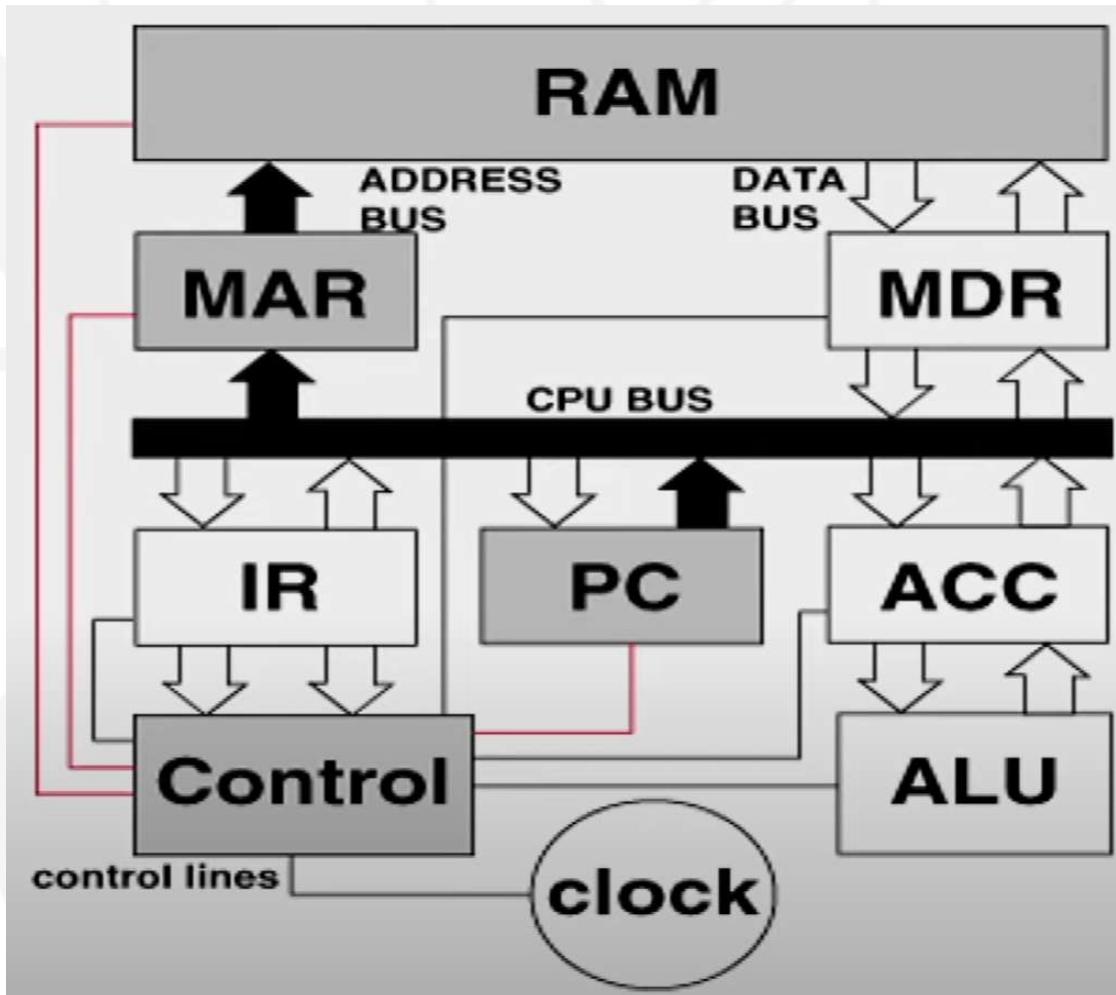
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# Fetch the instruction



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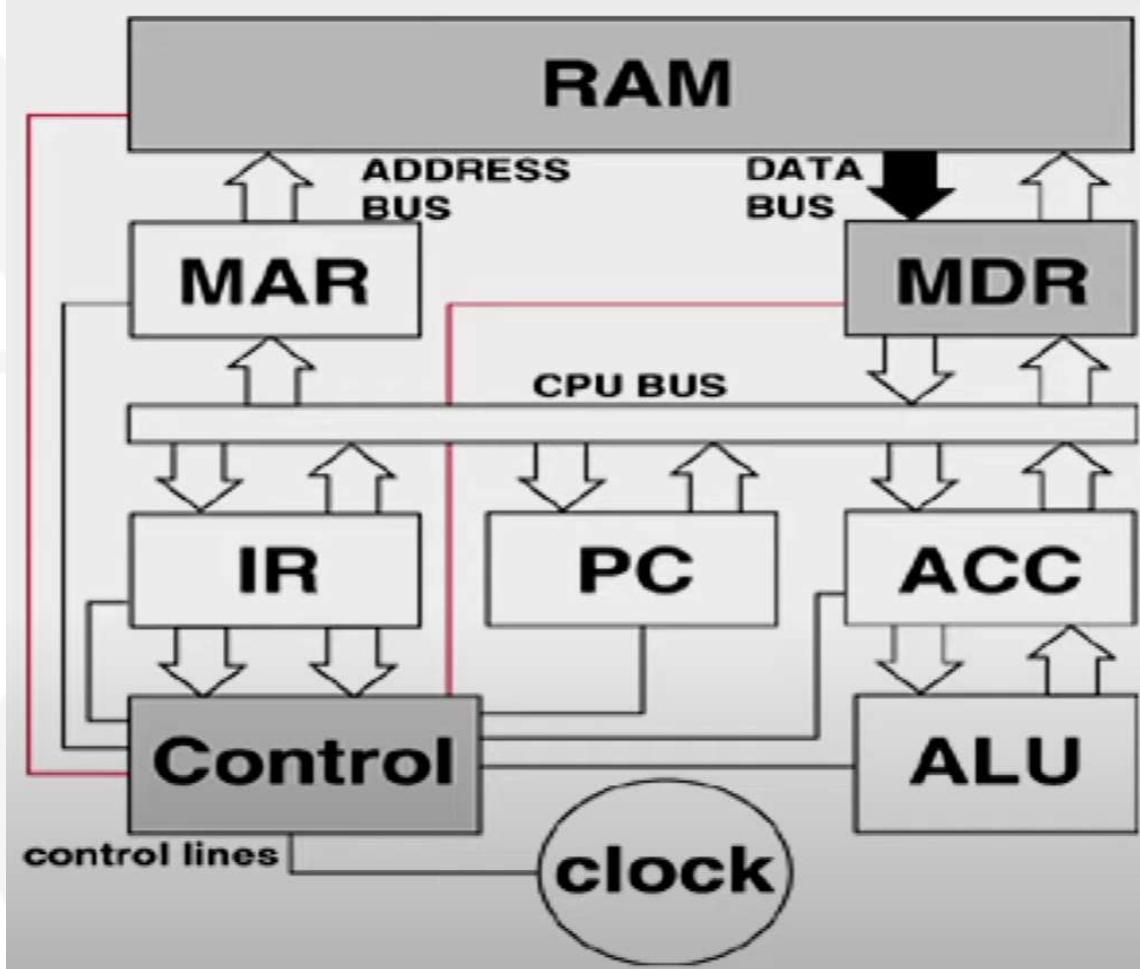


1. Address of the next instruction is transferred from PC to MAR.
2. The instruction is located in memory.

# Fetch the instruction



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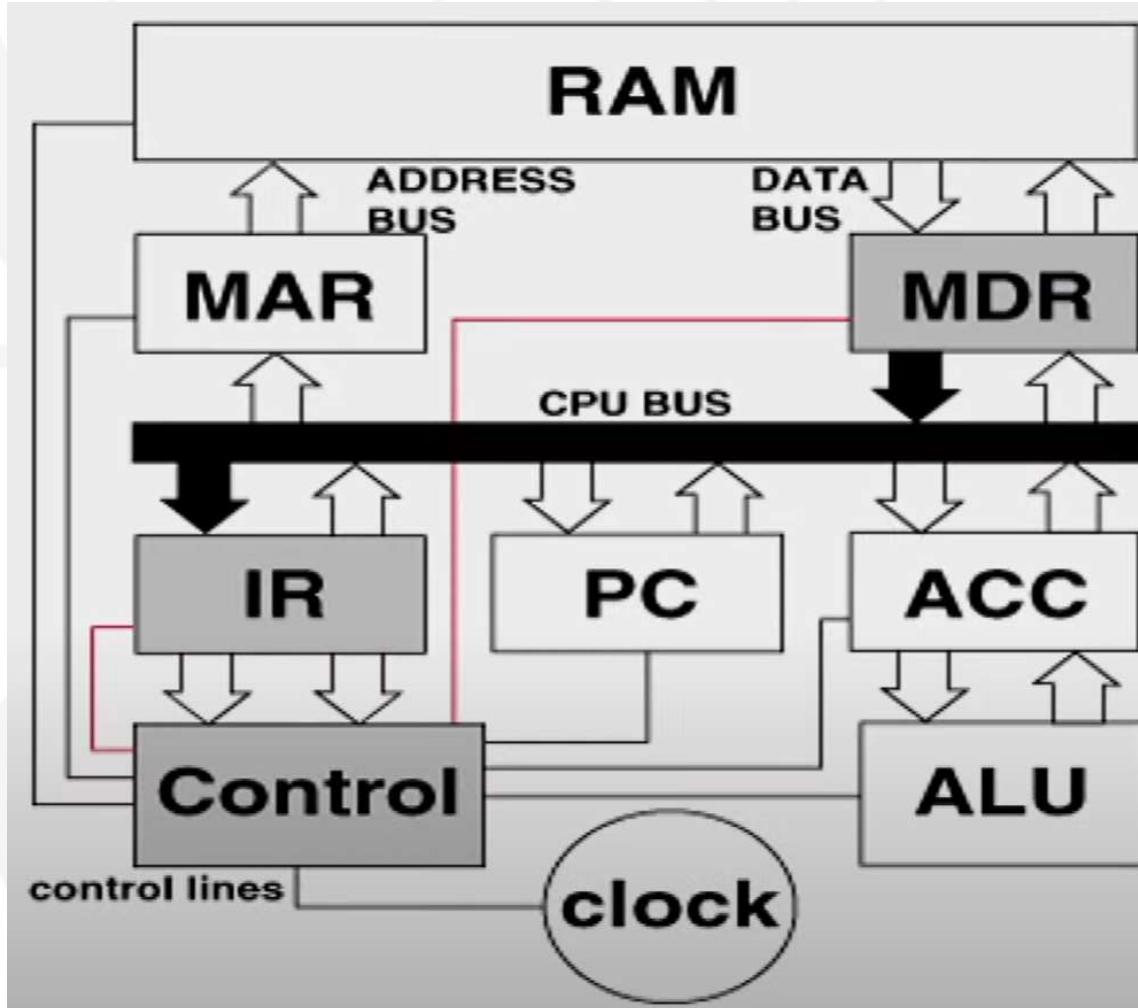


3. Instruction is copied from memory to MDR.

# Decode the instruction

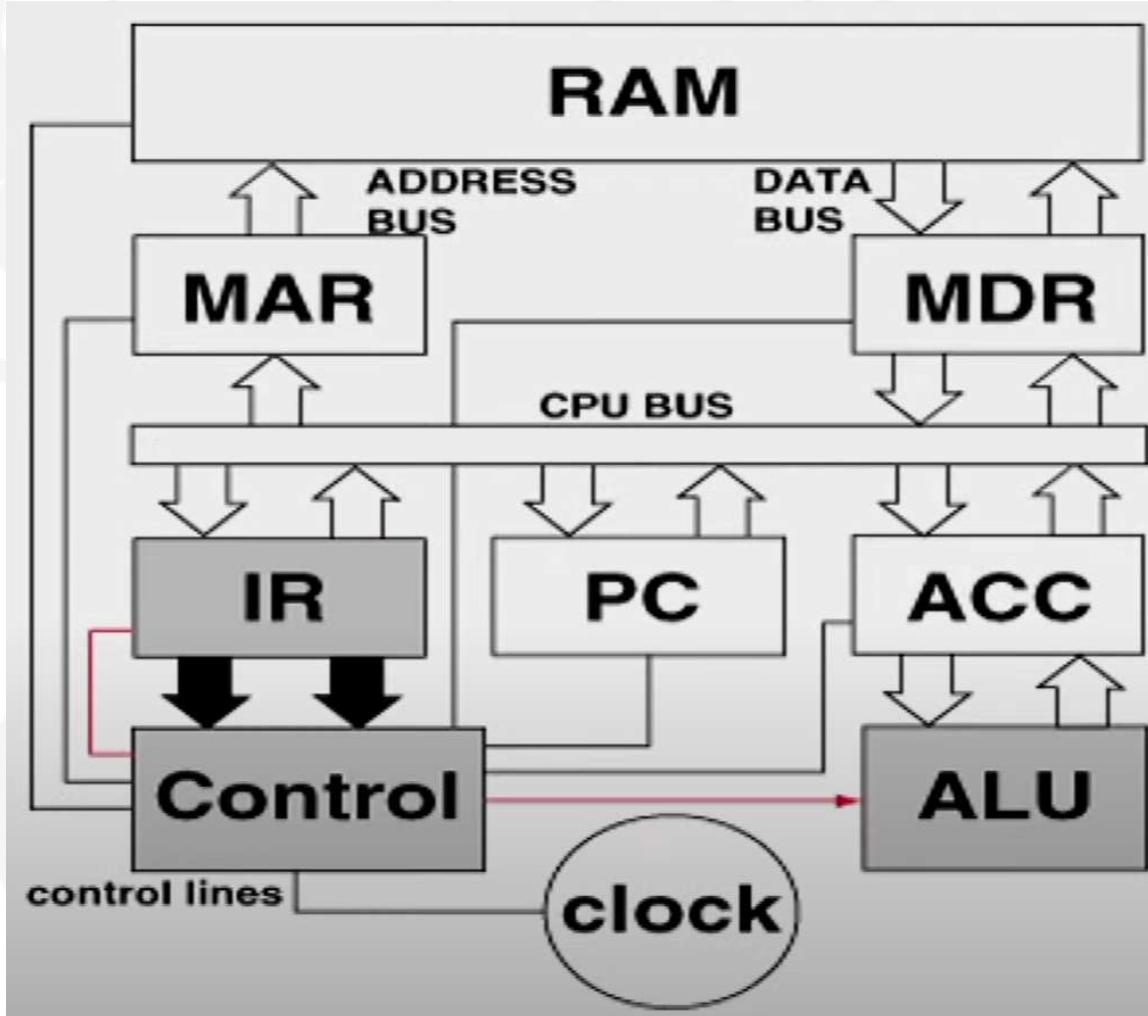


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4. Instruction is transferred to and decoded in the IR.

# Execute the instruction

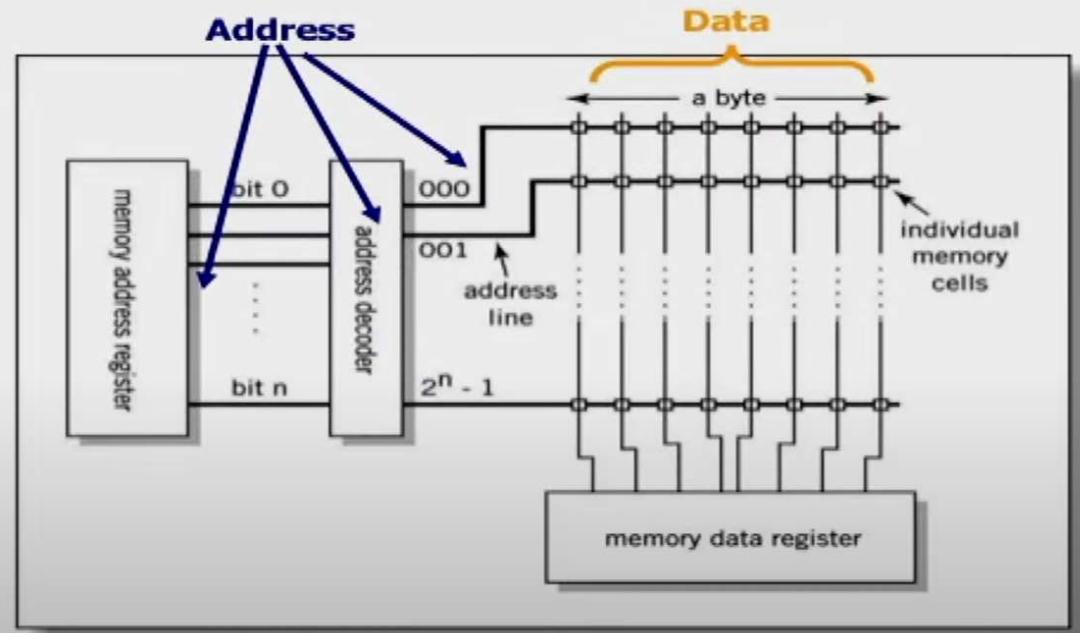
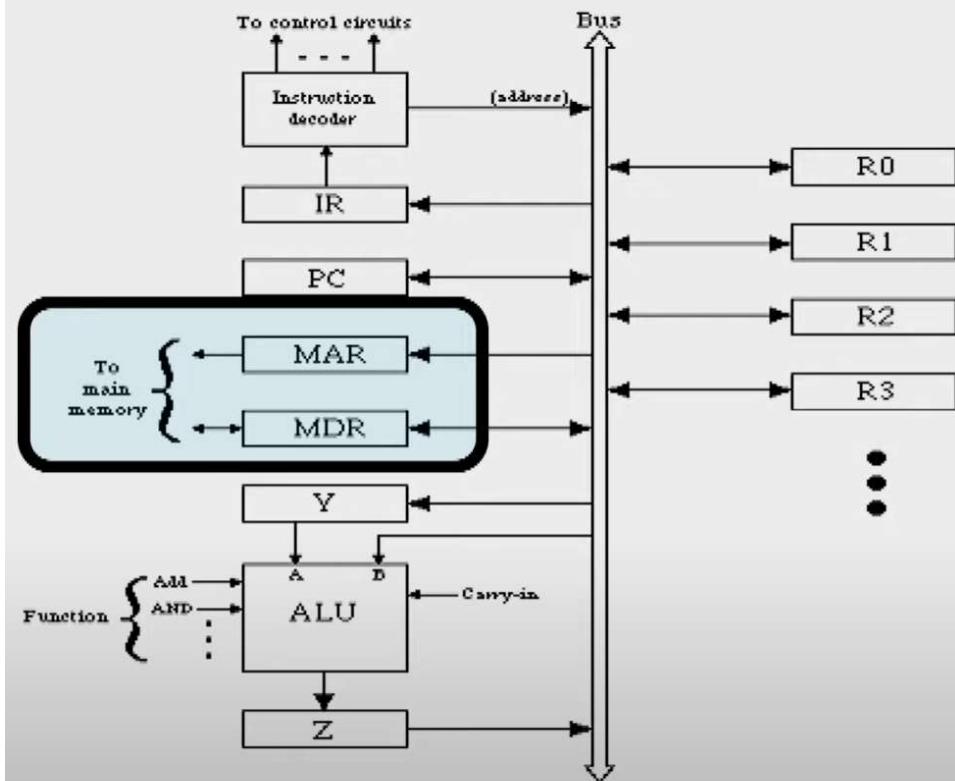


5. Control unit sends signals to appropriate devices to cause execution of the instruction.

# Processor Memory Interaction



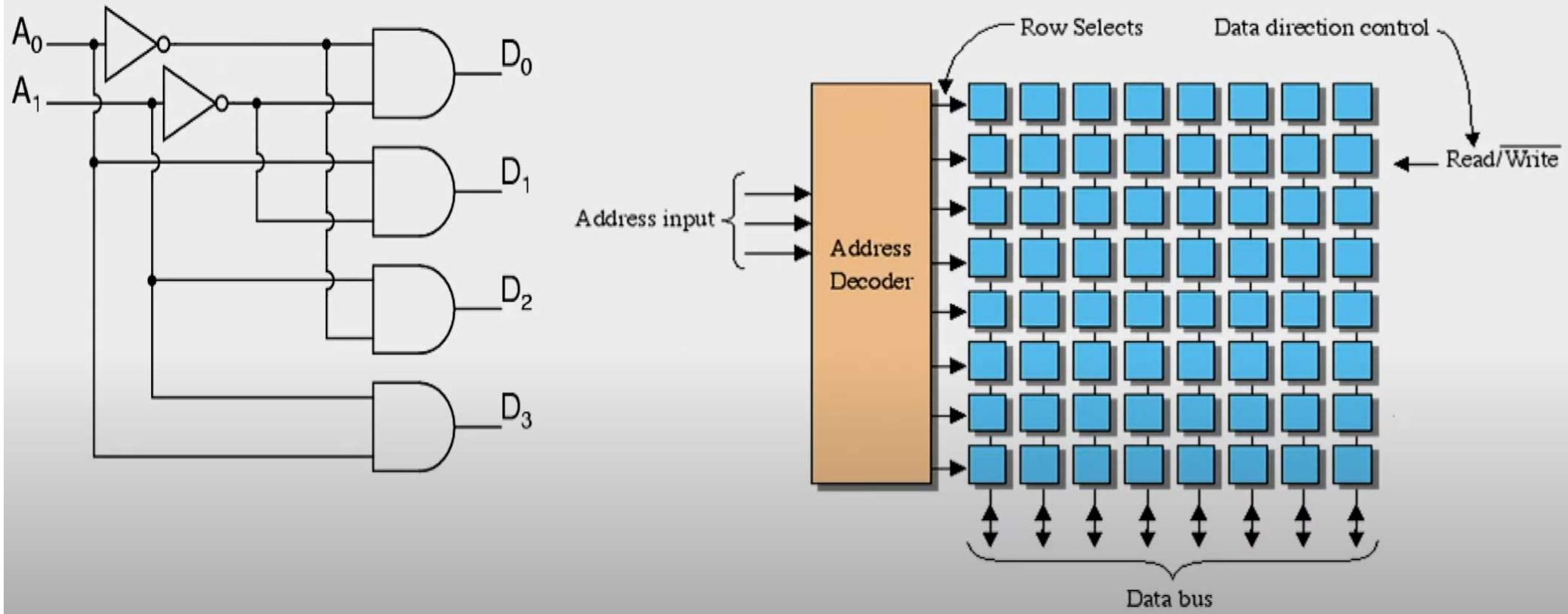
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# How memory works?



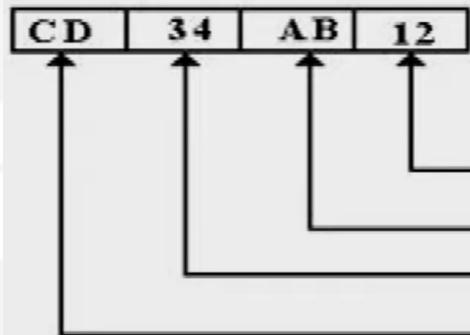
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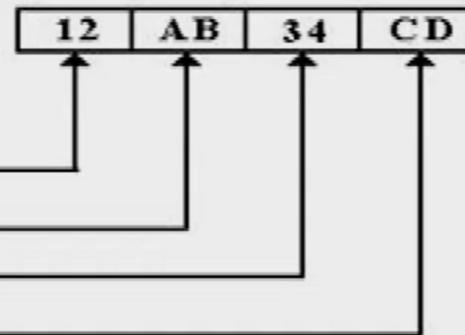
# Byte Ordering

## Big Endian VS Little Endian

Little Endian

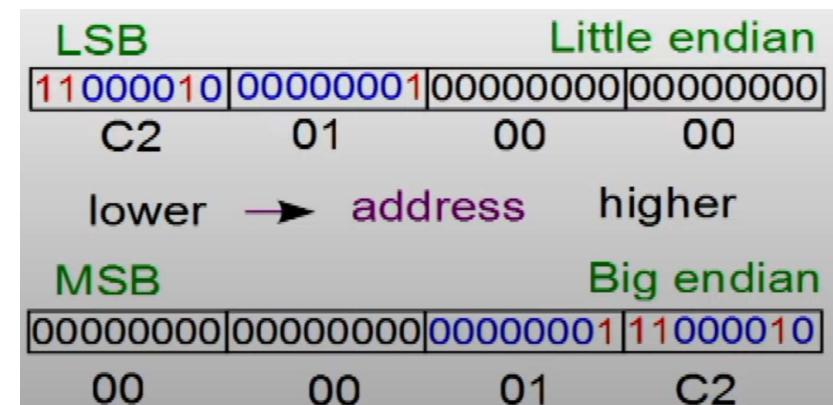
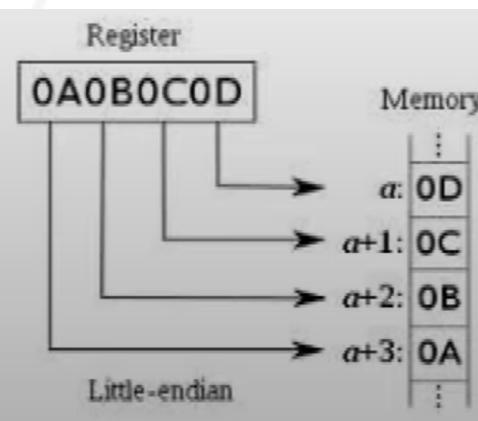
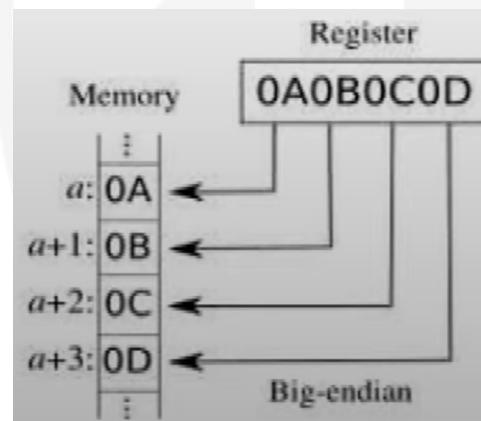


Big Endian



32-bit number:  
0XCD34AB12

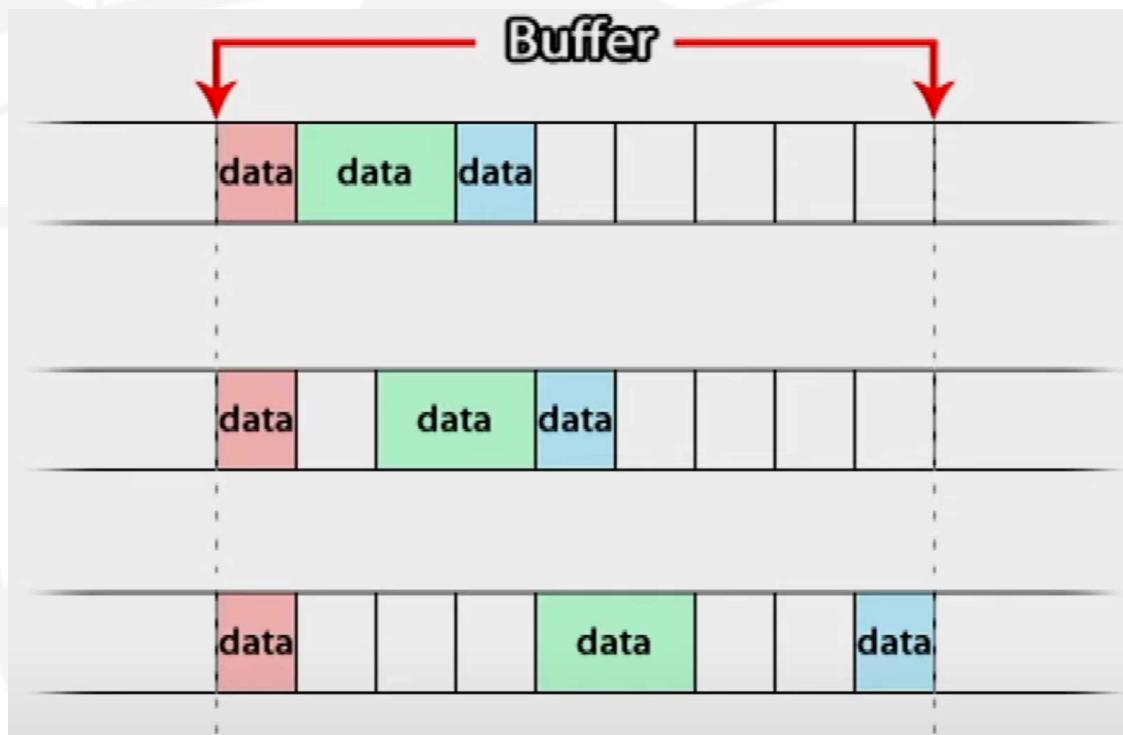
$\text{Int } i = 450 = 2^8 + 2^7 + 2^6 + 2 = X000001C2$



# Byte / Word Alignment



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Aligned to 1 byte

Aligned to 2 bytes

Aligned to 4 bytes



# Byte / Word Alignment

Value of 3 low-order bits of byte address

Width of object	0	1	2	3	4	5	6	7
1 byte (byte)	Aligned	Aligned	Aligned	Aligned	Aligned	Aligned	Aligned	Aligned
2 bytes (half word)	Aligned		Aligned		Aligned		Aligned	
2 bytes (half word)		Misaligned		Misaligned		Misaligned		Misaligned
4 bytes (word)		Aligned				Aligned		
4 bytes (word)			Misaligned				Misaligned	
4 bytes (word)				Misaligned				Misaligned
4 bytes (word)					Misaligned			
8 bytes (double word)			Aligned					
8 bytes (double word)				Misaligned				
8 bytes (double word)					Misaligned			
8 bytes (double word)						Misaligned		
8 bytes (double word)							Misaligned	
8 bytes (double word)								Misaligned
8 bytes (double word)								Misaligned
8 bytes (double word)								Misaligned



# Types of Processor Operations

## Arithmetic and Logical Operations

- Integer arithmetic
- Comparing two quantities
- Shifting, rotating bits in a quantity
- Testing, comparing, and converting bits



# Types of Processor Operations

## Data Movement Operations

- Moving data from memory to the CPU
- Moving data from memory to memory
- Input and output



# Types of Processor Operations

## Program Control Operations

- Starting a program
- Halting a program
- Skipping to other instructions
- Testing data to decide whether to skip over some instructions

# Instruction Set Architecture (ISA)



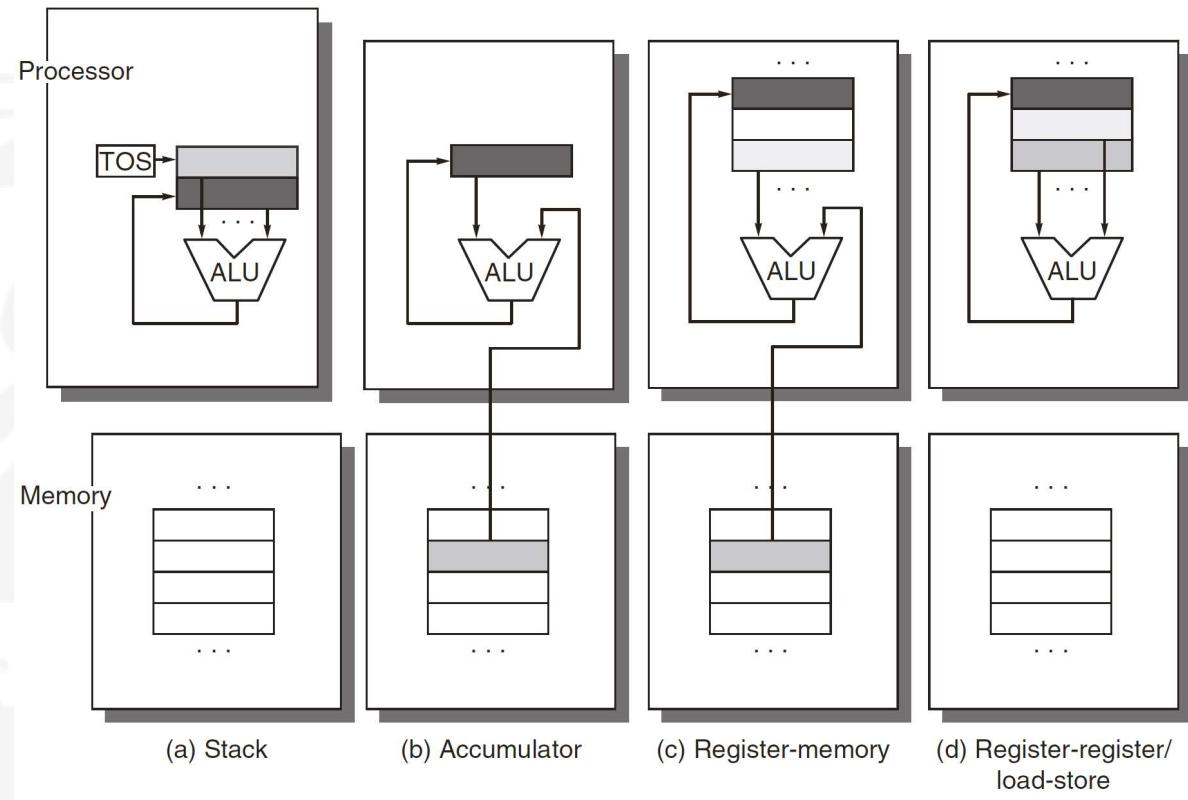
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- **Instruction vs Program vs Software**
- **Opcode, Operand**
- **Classification of ISA**
  - Stack architecture
  - Accumulator architecture
  - Register-Memory architecture
  - Register-Register/Load-Store architecture

# Instruction Set Architecture (ISA)



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Operand locations for  
four instruction set  
architecture classes.

Stack	Accumulator	Register (register-memory)	Register (load-store)
Push A	Load A	Load R1,A	Load R1,A
Push B	Add B	Add R3,R1,B	Load R2,B
Add	Store C	Store R3,C	Add R3,R1,R2
Pop C			Store R3,C

The code sequence for  $C = A + B$  for four classes of instruction sets.

# Instruction Set Architecture (ISA)



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$$A = D * (B + C) - E$$

## Stack Machine

- Push D
- Push B
- Push C
- Add
- Mul
- Push E
- Sub
- Pop A

## Accumulator Machine

- Load B
- Add C
- Mul D
- Sub E
- Store A

## Load-Store Machine

- Load R1, D
- Load R2, B
- Load R3, C
- Add R4, R2, R3
- Mul R5, R1, R4
- Load R6, E
- Sub R7, R5, R6
- Store R7, A

# Addressing Modes



## The way by which an operand is specified in an instruction

Addressing mode	Example instruction	Meaning	When used
Register	Add R4, R3	$\text{Regs}[\text{R4}] \leftarrow \text{Regs}[\text{R4}] + \text{Regs}[\text{R3}]$	When a value is in a register.
Immediate	Add R4, #3	$\text{Regs}[\text{R4}] \leftarrow \text{Regs}[\text{R4}] + 3$	For constants.
Displacement	Add R4, 100(R1)	$\text{Regs}[\text{R4}] \leftarrow \text{Regs}[\text{R4}] + \text{Mem}[100 + \text{Regs}[\text{R1}]]$	Accessing local variables (+ simulates register indirect, direct addressing modes).
Register indirect	Add R4, (R1)	$\text{Regs}[\text{R4}] \leftarrow \text{Regs}[\text{R4}] + \text{Mem}[\text{Regs}[\text{R1}]]$	Accessing using a pointer or a computed address.
Indexed	Add R3, (R1 + R2)	$\text{Regs}[\text{R3}] \leftarrow \text{Regs}[\text{R3}] + \text{Mem}[\text{Regs}[\text{R1}] + \text{Regs}[\text{R2}]]$	Sometimes useful in array addressing: R1 = base of array; R2 = index amount.
Direct or absolute	Add R1, (1001)	$\text{Regs}[\text{R1}] \leftarrow \text{Regs}[\text{R1}] + \text{Mem}[1001]$	Sometimes useful for accessing static data; address constant may need to be large.
Memory indirect	Add R1, @(R3)	$\text{Regs}[\text{R1}] \leftarrow \text{Regs}[\text{R1}] + \text{Mem}[\text{Mem}[\text{Regs}[\text{R3}]]]$	If R3 is the address of a pointer $p$ , then mode yields $*p$ .
Autoincrement	Add R1, (R2)+	$\text{Regs}[\text{R1}] \leftarrow \text{Regs}[\text{R1}] + \text{Mem}[\text{Regs}[\text{R2}]]$ $\text{Regs}[\text{R2}] \leftarrow \text{Regs}[\text{R2}] + d$	Useful for stepping through arrays within a loop. R2 points to start of array; each reference increments R2 by size of an element, $d$ .
Autodecrement	Add R1, -(R2)	$\text{Regs}[\text{R2}] \leftarrow \text{Regs}[\text{R2}] - d$ $\text{Regs}[\text{R1}] \leftarrow \text{Regs}[\text{R1}] + \text{Mem}[\text{Regs}[\text{R2}]]$	Same use as autoincrement. Autodecrement/-increment can also act as push/pop to implement a stack.
Scaled	Add R1, 100(R2) [R3]	$\text{Regs}[\text{R1}] \leftarrow \text{Regs}[\text{R1}] + \text{Mem}[100 + \text{Regs}[\text{R2}] + \text{Regs}[\text{R3}] * d]$	Used to index arrays. May be applied to any indexed addressing mode in some computers.

Selection of addressing modes with examples, meaning, and usage.

# Summary



# Suggested Readings



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Computer Architecture: A Quantitative Approach (5<sup>th</sup> Edition) by  
John Hennessy and David A Patterson:

- pp. 2 – 16: Sections 1.1 – 1.3
- pp. A2 – A16: Sections A.1 – A.5 Instruction Set Principles