

Lab 3

Addressing Modes

Objectives

- Students will learn basic addressing modes
- Students will come to know about little-endian and big-endian notations
- Students will be able to access any physical location of memory for data

When we run a program, the operating system locates the complete program on disk and loads (copies) it to the RAM. It also initializes the value of the CS and SS registers with the starting addresses of the code and stack segments. However, it does not initialize the DS segment. The DS value (and ES if used) must be initialized by the program to access memory for data. This is done as follows:

```
MOV AX, @DATA  
MOV DS, AX
```

Here, @DATA refers to the start of the data segment and is replaced by a number. Since we cannot assign a number directly to segment registers, therefore, we must first assign it to a general-purpose register and then from that general-purpose register to a segment register.

Addressing Modes

Immediate addressing mode

In the immediate addressing mode, the source operand is a constant. In immediate addressing mode, as the name implies, when the instruction is assembled, the operand comes immediately after the opcode. For this reason, this addressing mode executes quickly. However, in programming, it has limited use. Immediate addressing mode can be used to load information into general-purpose registers.

Examples:

```
MOV AX,2550H  
MOV CX,625  
MOV BL,40H
```

Register addressing mode

The register addressing mode involves the use of registers to hold the data to be manipulated. Memory is not accessed when this addressing mode is executed; therefore, it is relatively fast.

Examples:

```
ADD BX, DX  
MOV ES, AX  
MOV AL, BH
```

Direct Addressing mode

In the direct addressing mode, the data is in some memory location(s) and the address of the data in memory comes immediately after the instruction. Note that in immediate addressing, the operand itself is provided with the instruction, whereas in direct addressing mode, the address of the operand is provided with the instruction. This address is an offset.

MOV DL, [2400h] ; move contents of DS:2400H into DL

The physical address is calculated by $DS \times 10h + 2400h$.

Register indirect addressing mode

In the register indirect addressing mode, the address of the memory location where the operand resides is held by a register. The registers used for this purpose are SI, DI, and BX.

For example:

MOV AL, [BX] ;moves the contents of the memory location into AL, the memory location pointed to DS:BX.

- The physical address is calculated by $DS \times 10h + BX$. The same rules apply when using register SI or DI.
- BP register can also be used as a pointer register. However, the physical address will be calculated by $SS \times 10h + BP$.
- AX, CX, DX, and SP cannot be used as pointer registers/ offset.

We can override the segment register while using BX, SI, DI, and BP registers as pointers by writing the preferred segment register name with it as shown below.

MOV AL, ES:[BX]
MOV AL, DS:[BP]
MOV AL, CS:[SI]
MOV AL, SS:[DI]

The size of a register in an instruction specifies how many bytes will be read or written from or to memory. If a register is not there, one byte is accessed from memory by default in Emu8086. However, to avoid confusion, one must specify the number of bytes to read or write to memory using the following:

- **byte ptr** - for byte.
- **word ptr** - for word (two bytes).

For example:

```
MOV byte ptr [2400h], 1
```

```
ADD word ptr [2400h], 2
```

Storing multi-byte data in RAM

Little-endian and big-endian are the two ways of storing data in memory. In little-endian, the lower byte is stored at a lower address and the higher byte is stored at a higher address. However, in big-endian, the lower byte is stored at the higher address and the higher byte is stored at the lower address.

Let the AX register contain '0xABCD' and the DS register contain "0000h". The following instruction will write the contents of the AX register into the memory and start writing from physical location 0x00100 onwards. Since the Intel 8086 processor is based on little-endian, it will store the lower byte 'CD' at 0x00100 and higher byte 'AB' at 0x00101 as shown in the following figure. However, the processor that is based on big-endian will store it the other way around. Memory will be read in the same way.

Instruction:

```
MOV [0100h], AX
```

AX	
AH	AL
AB	CD

0x00100	CD
0x00101	AB
0x00102	
0x00103	
0x00104	
0x00105	
0x00106	

Little Endian
Intel 8086

0x00100	AB
0x00101	CD
0x00102	
0x00103	
0x00104	
0x00105	
0x00106	

Big Endian

Accessing a specific physical location of RAM

Using the concepts above, one can access any physical location of memory. The physical address is a 20-bit value that has to be converted into Segment: Offset in a way such that when the processor combines it to form a physical address, it should be in the same location.

To write a byte “0x12” to the physical address 0xABCDE of RAM, one needs to break the physical address into Segment and Offset parts each of 16-bit. One combination is:

Segment : Offset
0ABCDh : 000Eh.

.code

```
mov ax,0abcdh
mov ds,ax
mov bx,000eh
mov byte ptr [bx],012h
```

Similarly, to write a word with the physical address “0xABCDE”, the prefix: byte ptr will be replaced with word ptr.

```
mov ax,0abcdh
mov ds,ax
mov bx,000eh
mov word ptr [bx],012h
```

Program to write 0x1234 physical memory “0xABCDE”.

```
.model small  
  
.data  
  
.code  
  
mov bx,0abcdh  
mov es,bx  
mov word ptr es:[000eh],01234h  
  
.exit
```

The description of the above program is shown in the following table.

Instruction	Description
mov bx,0abcdh	Moving segment address to bx before moving it to ES
mov es,bx	Moving segment address to ES
mov word ptr es:[000eh]	Writing word “0x1234” to physical memory 0xABCDE

Emu8086 Tutorial Step by Step

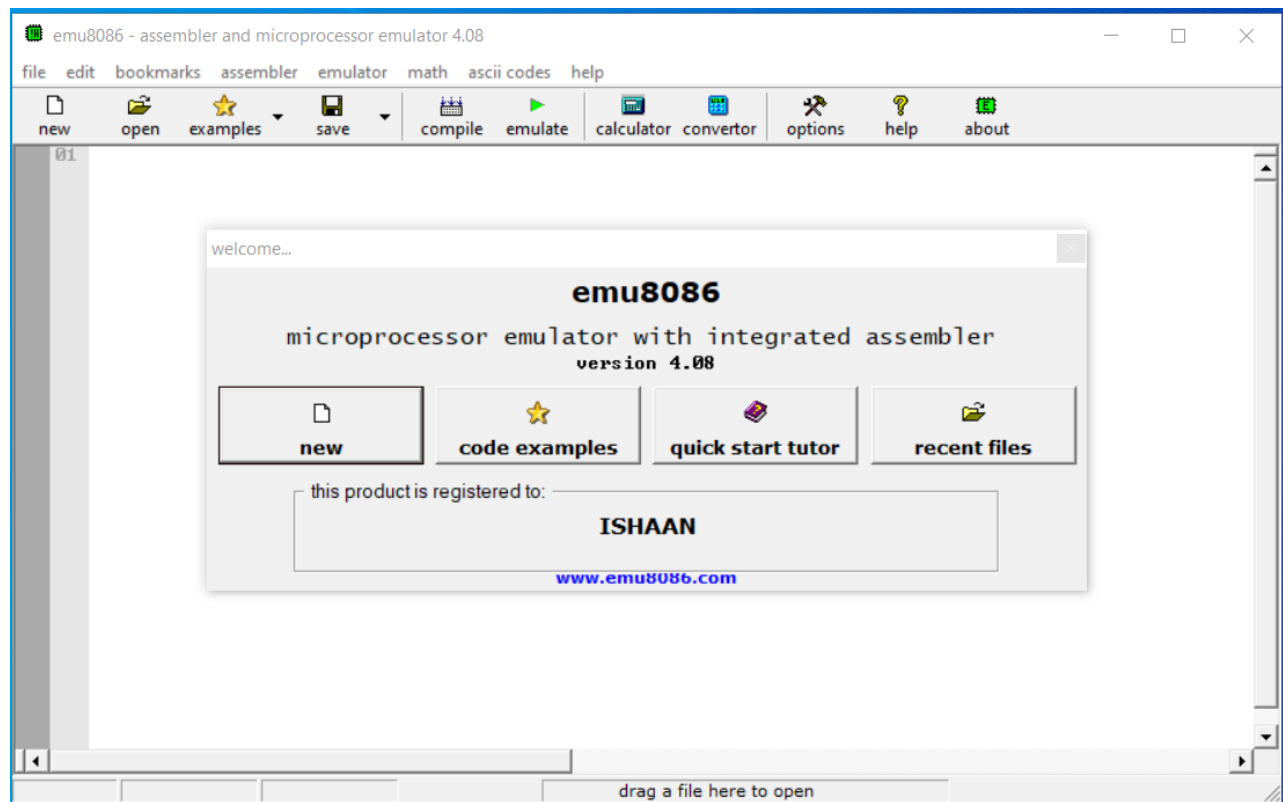
Step-1



Double-click on the icon on the desktop

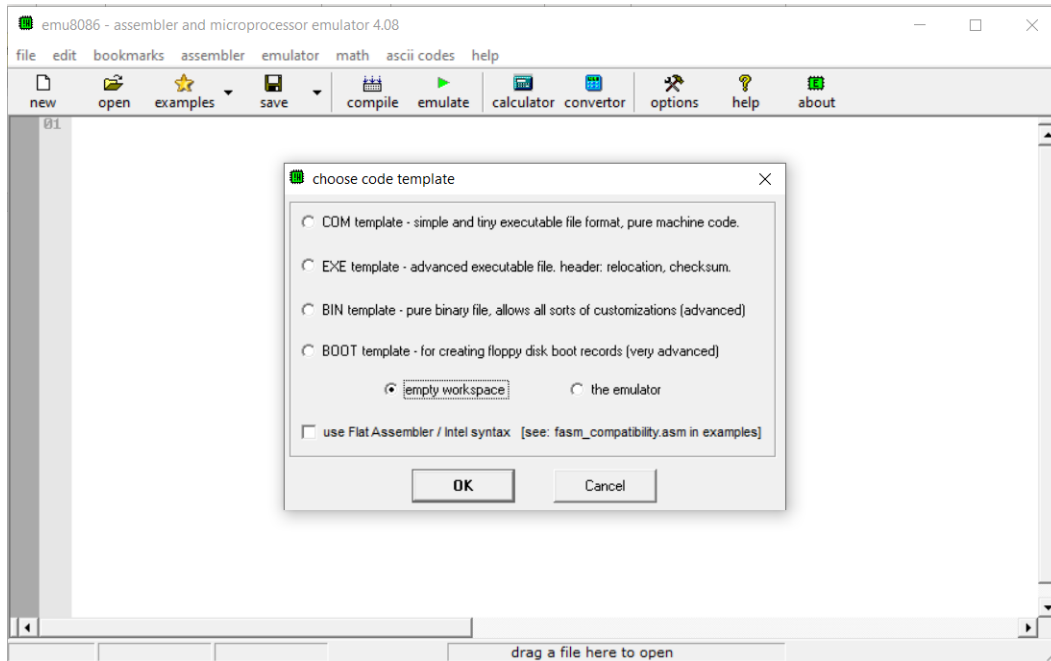
Step-2

The following window will appear. Click on new.



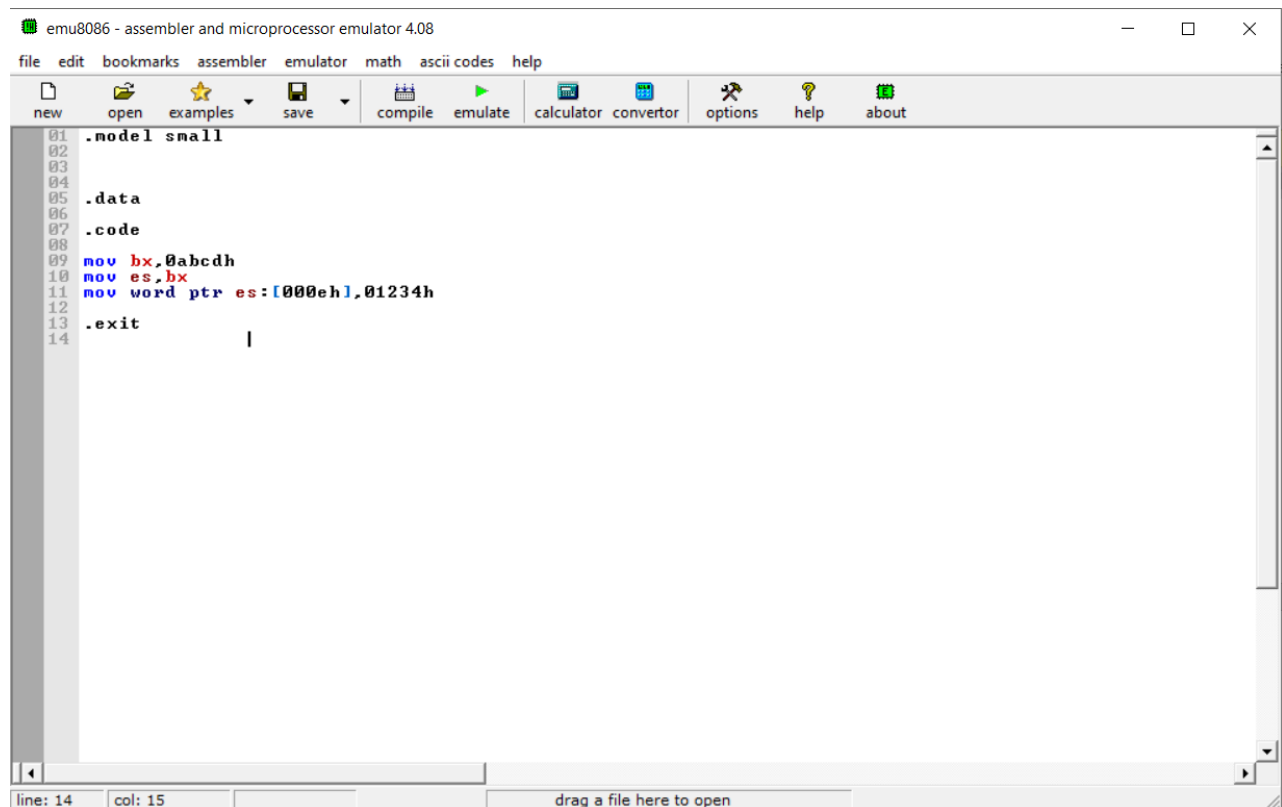
Step-3

Click on the empty workspace and press OK.



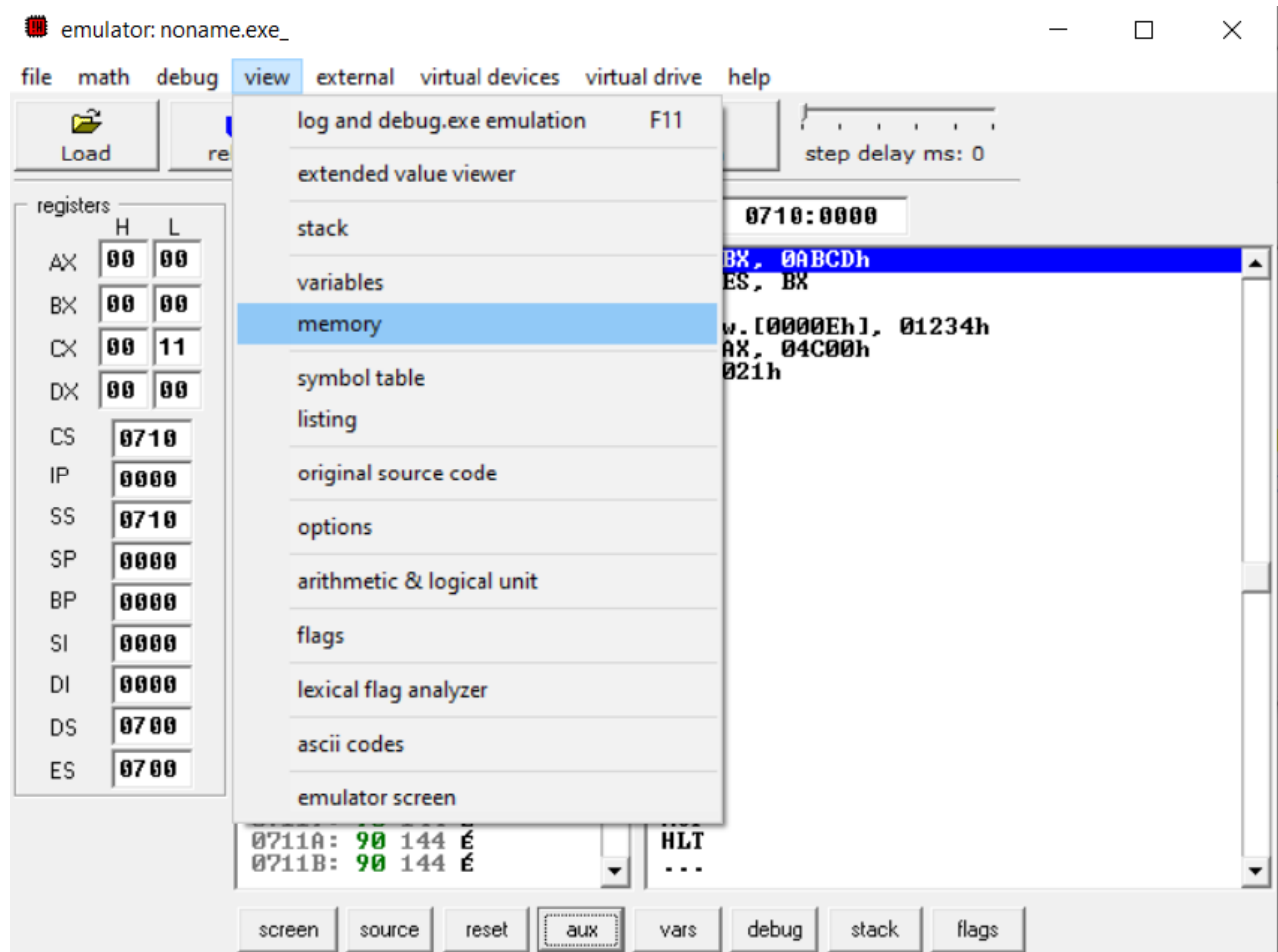
Step-4

Type the code given above and click on emulate.



Step-5

Click on “Memory” from the view menu!



Step-6

Write the logical address of the desired part of the segment that you want to view.

emulator: noname.exe_

file math debug view external virtual devices virtual drive help

Load reload step back single step run step delay ms: 0

registers

	H	L
AX	00	00
BX	00	00
CX	00	11
DX	00	00
CS	0710	
IP	0000	

0710:0000 0710:0000

07100: BB 187 7	MOV BX, 0ABCDh
07101: CD 205 =	MOV ES, BX
07102: AB 171 2	ES:
07103: 8E 142 a	MOV w.[0000Eh], 01234h
07104: C3 195	MOV AX, 04C00h
07105: 26 038 &	INT 021h
07106: C7 199	NOP
07107: 06 006 4	NOP
07108: 0E 014 J	NOP
07109: 00 000 NULL	NOP
0710A: 34 052 4	NOP

Random Access Memory

abcd:000e update table list

Update Map

ABCD:000E	00 00 00 00 00 00 00 00 00-00 00 00 00 00 00 00
ABCD:001F	00 00 00 00 00 00 00 00 00-00 00 00 00 00 00 00
ABCD:0027	00 00 00 00 00 00 00 00 00-00 00 00 00 00 00 00
ABCD:003E	00 00 00 00 00 00 00 00 00-00 00 00 00 00 00 00
ABCD:004E	00 00 00 00 00 00 00 00 00-00 00 00 00 00 00 00
ABCD:005E	00 00 00 00 00 00 00 00 00-00 00 00 00 00 00 00
ABCD:006E	00 00 00 00 00 00 00 00 00-00 00 00 00 00 00 00
ABCD:007E	00 00 00 00 00 00 00 00 00-00 00 00 00 00 00 00

Segment: Offset

0711A: 90 144 E HLT

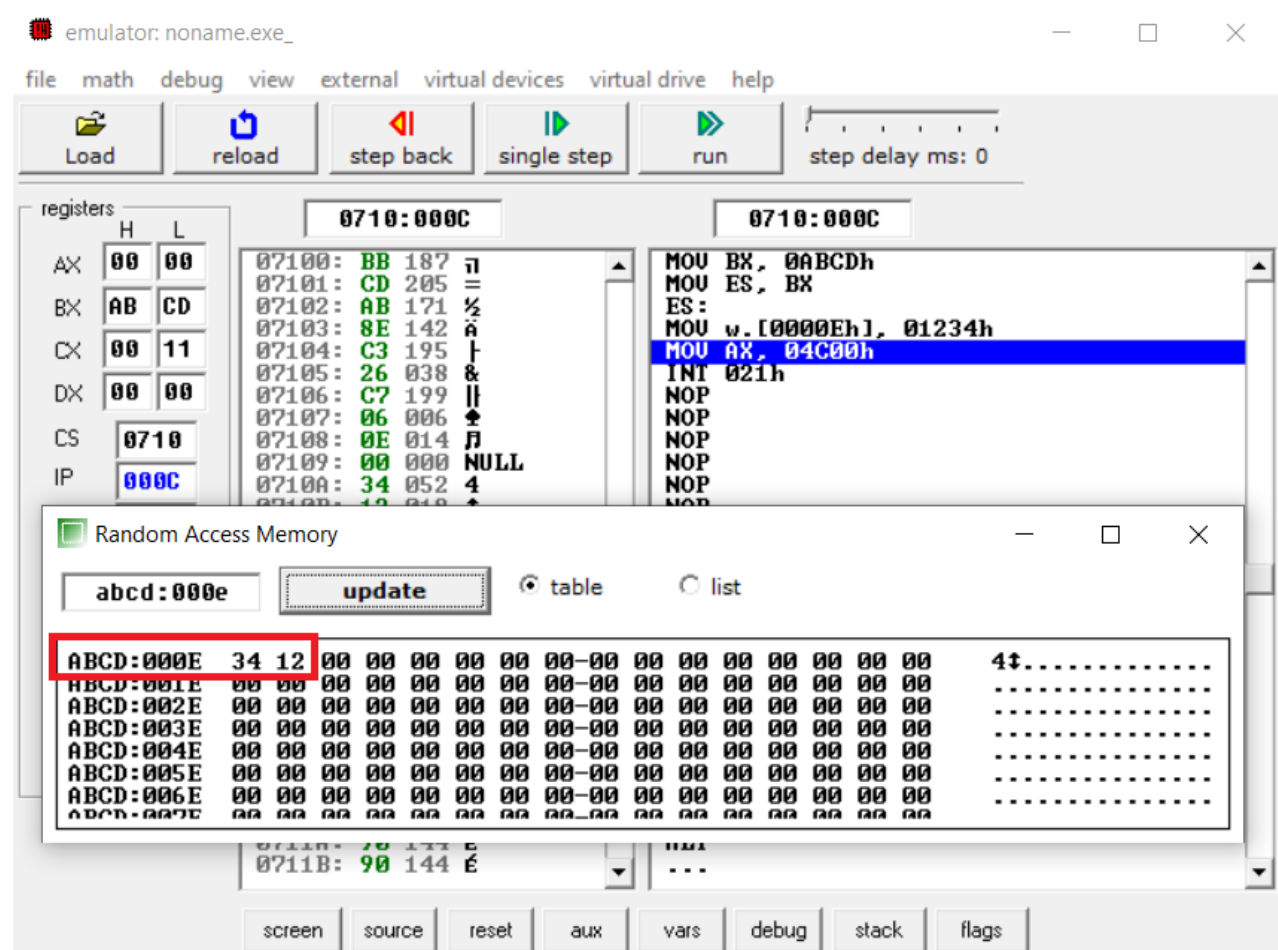
0711B: 90 144 E ...

screen source reset aux vars debug stack flags

Step-7

Keep clicking on “Single step” to execute program instructions one by one.

(The **mov word ptr es:[000eh],01234h** instruction will write two bytes to memory.)

**Observation:**

- Which byte is written to what address?

Practice Exercise

Task-1

Write a program that stores the following numbers into the current data segment at offset: 0x1000 onwards. The program then calculates the sum of these numbers and stores it at the physical location: 0xCD1F3.

Numbers: 0x1F00, 0xA0B1, 0x1254, 0x34EF, 0x8700

Task-2

The "HLT" has a machine code of "0xF4". Write a code to replace the first instruction of your program with this instruction.