

PDA's

Ex.: Find PDA's that accept the following languages:

1. $L_1 = \{ww^R \mid w \in \{a, b\}^+\}$
2. $L_2 = \{a^n b^{2n} \mid n \in N^*\}$
3. $L_3 = \{a^n b^{2n} \mid n \in N\}$
4. $L_4 = \{a^{2n} b^n \mid n \in N^*\}$

Sol.:

2,3 - see the test for Seminar 13

1. $M = (\{q_0, q_1, q_2\}, \{a, b\}, \{A, B, Z\}, d, q_0, Z, \{q_2\})$ // Ariana Hategan

$d(q_0, a, Z) = \{(q_0, AZ)\}$

$d(q_0, b, Z) = \{(q_0, BZ)\}$

$d(q_0, a, A) = \{(q_0, AA), (q_1, \text{epsilon})\}$

$d(q_0, a, B) = \{(q_0, AB)\}$

$d(q_0, b, A) = \{(q_0, BA)\}$

$d(q_0, b, B) = \{(q_0, BB), (q_1, \text{epsilon})\}$

$d(q_1, a, A) = \{(q_1, \text{epsilon})\}$

$d(q_1, b, B) = \{(q_1, \text{epsilon})\}$

$d(q_1, \text{epsilon}, Z) = \{(q_2, Z)\}$

$(q_0, abba, Z) \vdash (q_0, bba, AZ) \vdash (q_0, ba, BAZ) \vdash (q_1, a, AZ) \vdash (q_1, \text{epsilon}, Z) \vdash (q_2, \text{epsilon}, Z)$
=> accepted

$(q_0, abaa, Z) \vdash (q_0, baa, AZ) \vdash (q_0, aa, BAZ) \vdash (q_0, a, ABAZ) \vdash (q_0, \text{epsilon}, AABAZ)$
| $\vdash (q_1, \text{epsilon}, BAZ)$

=> not accepted

- $d(q_0, \text{epsilon}, Z) = \{(q_2, Z)\}$ - transition to add in order to also accept the empty sequence ($L'_1 = \{ww^R \mid w \in \{a, b\}^*\}$)

4. $M = (\{q_0, q_1, q_2, q_3\}, \{a, b\}, \{A, Z\}, d, q_0, Z, \{q_3\})$ // Ariana Hategan

$d(q_0, a, Z) = \{(q_1, AZ)\}$

$d(q_1, a, A) = \{(q_0, AZ)\}$

$d(q_0, a, A) = \{(q_1, AA)\}$

$d(q_1, a, A) = \{(q_0, AA)\}$

$d(q_0, b, A) = \{(q_2, \epsilon)\}$

$d(q_2, b, A) = \{(q_2, \epsilon)\}$

$d(q_2, \epsilon, Z) = \{(q_3, Z)\}$

$(q_0, aaaabb, Z) \vdash (q_1, aaabb, AZ) \vdash (q_0, aabb, AZ) \vdash (q_1, abb, AAZ) \vdash (q_0, bb, AAZ) \vdash (q_2, b, AZ) \vdash (q_2, \epsilon, Z) \vdash (q_3, \epsilon, Z) \Rightarrow aaaabb$ is accepted

Attribute Grammars

Ex.: Give an attribute grammar for evaluating simple arithmetic expressions with $id, (,), +, *$

Sol.:

S - attributed grammar

$E \rightarrow E + T \quad \{E1.val = E2.val + T.val\}$

$E \rightarrow T \quad \{E.val = T.val\}$

$T \rightarrow T * F \quad \{T1.val = T2.val * F.val\}$

$T \rightarrow F \quad \{T.val = F.val\}$

$F \rightarrow (E) \quad \{F.val = E.val\}$

$F \rightarrow id \quad \{F.val = id.val\}$

Obs.: The green arrows from the syntax tree below indicate how evaluation is performed (bottom-up - value of attribute in parent is computed based on values of attributes in descendants, attributes are synthesized)

$$w = 2 + 3 * 4$$

