## **Monitoring locks**

SQL Server Profiler

- tracing tool that can be used to monitor events on the server;
- e.g., the user can create a *trace* for *Lock:Acquired* and *Lock:Released* events (indicating a lock has been acquired / released) and run it, then analyze its output;

sp\_lock (feature in maintenance mode)

- returns info about all the locks held by sessions that are currently active / one or two sessions specified through input parameters (i.e., 0, 1 or 2 parameters);
- example:

sys.dm tran locks

- 1 row per currently active request issued to the lock manager, i.e., the corresponding lock has been granted or is waiting to be granted;
- example:

```
-- resource_type: KEY resource_database_id: 23 request_mode: X request_type: LOCK request_status: GRANT request_session_id: 53 request_owner_type: TRANSACTION request_owner_id: 106940
```

sys.dm\_tran\_active\_transactions

- info about transactions;

Ex. join with sys.dm\_tran\_locks on transaction\_id

#### **Query Governor Cost Limit**

example – AdventureWorks

```
SELECT *
```

```
FROM Sales.SalesOrderDetail d INNER JOIN Sales.SalesOrderHeader h
   ON d.SalesOrderID = h.SalesOrderID
--query allowed to run; result set of about 120.000 rows
```

#### SET QUERY\_GOVERNOR\_COST\_LIMIT 1

--query is not allowed to run, since its estimated cost exceeds the specified value: The query has been canceled because the estimated cost of this query (2) exceeds the configured threshold of 1.

## Row Level Versioning (RLV) in SQL Server

Read Committed Snapshot & Full Snapshot – illustrated on examples:

- example dirty read:
- \* row with MovieID = 1, Nominations = 28
- \* dirty read scenario:

T1	T2
UPDATE Movie	
SET Nominations = Nominations + 2	
WHERE MovieID = 1	
	SELECT * FROM Movie

- \* if T2 runs under READ COMMITTED, it's suspended until T1 releases its X lock;
- \* if T2 runs under READ UNCOMMITTED, it reads dirty data (value 30 for Nominations);
- \* if T2 runs under READ COMMITTED SNAPSHOT, it doesn't block, it doesn't read dirty data; it instead reads a value that has been previously committed: the most recent committed data as of the beginning of the SELECT statement, which is 28 in this particular case;

- \* if T2 runs under FULL SNAPSHOT, it doesn't block, it doesn't read dirty data; it instead reads a value that has been previously committed: the most recent committed data as of the beginning of the transaction, which is again 28;
- \* set the isolation level to READ COMMITTED SNAPSHOT:

```
ALTER DATABASE MyImdb
SET READ_COMMITTED_SNAPSHOT ON
```

- make sure T2 runs under READ COMMITTED;
- \* set the isolation level to FULL SNAPSHOT:

```
ALTER DATABASE MyImdb
SET ALLOW_SNAPSHOT_ISOLATION ON
SET TRANSACTION ISOLATION LEVEL SNAPSHOT -- for T2
```

\* check if RCS / FS has been enabled on the DB:

```
SELECT database_id, is_read_committed_snapshot_on, snapshot_isolation_state
FROM sys.databases
WHERE database_id = DB_ID('MyImdb')
```

- example non-repeatable read:
- \* row with MovieID = 1, Nominations = 28
- \* non-repeatable read scenario:

T1	T2	Read Nominations value under RCS	Read Nominations value under FS
	SELECT * FROM Movie WHERE MovieID = 1	28	28
UPDATE Movie  SET Nominations = Nominations + 2  WHERE MovieID = 1			
COMMIT TRAN	SELECT * FROM Movie WHERE MovieID = 1	30 - since T2 doesn't acquire an S lock for its read operation, T1 is able to change the row and commit;	28 - T2 doesn't acquire an S lock for its read operation, so T1 is able to change the row and commit; - when the 2 <sup>nd</sup> SELECT is executed, the most recent

- when the 2 <sup>nd</sup> SELECT is	committed value as of the
executed, the most recent	beginning of the
committed value as of the	transaction is 28, i.e., the
beginning of the operation	read is repeatable under
(i.e., the SELECT statement)	FS!
is 30, i.e., the read is not	
repeatable under RCS!	

- example update conflict:
- \* row with MovieID = 1, Nominations = 28
- \* update conflict scenario:

T1	T2 – under Full Snapshot	
	BEGIN TRAN	
	SELECT Nominations	
	FROM Movie	
	WHERE MovieID = 1	
	returned value: 28	
BEGIN TRAN		
UPDATE Movie		
SET Nominations = Nominations + 2		
WHERE MovieID = 1		
nominations is now 30		
	LIDDATE Marrie	
	UPDATE Movie	
	SET Nominations = Nominations + 5 WHERE MovieID = 1	
	T2 is suspended, since it's trying to acquire an X lock that conflicts	
	with the X lock held by T1	
COMMIT TRAN	With the Albokineta by 12	
	=> error: Snapshot isolation transaction aborted due to update conflict	
	- T2 cannot proceed with its update;	
	- T2 started when the value for Nominations was 28;	
	- in the meantime, T1 changed that value to 30: another row version	
	was created for the row and added to tempdb;	
	- if T2 is allowed to proceed, it adds 5 to the most recent committed	
	value as of the beginning of the transaction, which is 28: we would	
	have a lost update (the update of T1 would be lost in this case);	

# **PIVOT**

- example – table ProductSales:

ProductType	DOW	Sales
Chocolate	Monday	1000
Chocolate	Monday	500
Chocolate	Tuesday	1000
IceCream	Monday	2000
IceCream	Tuesday	600
IceCream	Wednesday	300

## DBMSs – seminar 4 - extra doc

2000

IceCream

```
IceCream
             Wednesday
                          600
             Wednesday
                           300
IceCream
SELECT *
FROM ProductSales
PIVOT (SUM(Sales)
      FOR DOW in ([Monday], [Tuesday], [Wednesday])
         ) t
ProductType
             Monday
                          Tuesday
                                        Wednesday
Chocolate
             1500
                                        NULL
                          1000
```

600

1200