

## Homework #6. Due Saturday, March 19th, by 11:59pm in filedrop

All reading assignments and references to exercises, definitions etc. are from our main book ‘Coding Theory: A First Course’ by Ling and Xing

### Reading and plan for the next week:

1. For this homework assignment read 5.1-5.4 and 5.7.
2. Plan for next week: Tue (Mar 15) – Plotkin bound (5.6) and Gilbert-Varshamov bound (5.2). Thu (Mar 17) – Polynomial rings and basic theory of finite fields (3.2 and parts of 3.3; see also online Lecture 15 from Spring 2020).

### Problems:

1. Recall that if  $C$  is a code of length  $n$  and  $I$  is a proper subset of  $\{1, \dots, n\}$ , the punctured code  $C_I$  is a code of length  $n - |I|$  obtained from  $C$  by puncturing the  $i^{\text{th}}$  coordinate for every  $i \in I$  (from every  $c \in C$ ). Prove part (b) of Lemma 13.1 from class:

$$d(C_I) \geq d(C) - |I|.$$

2. In parts (a) and (b) of this problem assume that  $d \geq 2$ .
  - (a) Prove that  $A_q(n, d) \leq A_q(n, d - 1)$ . In other words, fix an alphabet  $A$  with  $|A| = q$  and prove the following: if there exists an  $(n, M, d)$ -code  $C$  over  $A$ , there also exists an  $(n, M, d - 1)$ -code  $C'$  over  $A$ . You should give a precise argument; do not try to say this is obvious or something like that.
  - (b) Now prove that  $A_q(n, d) \leq A_q(n - 1, d - 1)$ . **Hint:** use punctured codes.
  - (c) Now use (b) and induction to give another proof of the Singleton bound:  $A_q(n, d) \leq q^{n-d+1}$ .

3. Problem 5.7.

4.

- (a) Show that the all-one vector  $(1, 1, \dots, 1)$  of length 24 lies in the extended binary Golay code  $G_{24}$ .
- (b) Assume without proof that  $G_{24}$  contains (exactly) 759 words of weight 8. Use this fact and (a) to prove that the distribution of weights in  $G_{24}$  is given by Table 5.5 on page 109.

5. This problem deals with the Golay code  $G_{23}$ .

- (a) Use Problem 5(b) to prove that possible weights of elements of  $G_{23}$  are 0, 7, 8, 11, 12, 15, 16 and 23. Make sure to prove that each of those numbers actually arises as the weight of some element of  $G_{23}$ .
- (b) Let  $w \in \mathbb{F}_2^{23}$  with  $wt(w) = 4$ . Let  $c_w \in G_{23}$  be the result of applying NND decoding (with respect to  $G_{23}$ ) to  $w$ . Use (a) and the fact that  $G_{23}$  is perfect (as shown in class) to prove that  $d(w, c_w) = 3$  and  $wt(c_w) = 7$ .

6. Problem 5.19. **Note:** Simplex codes  $S(r, q)$  are defined at the end of 5.3.2, page 88.

7. Use the result of Problem 5.19 to show that the simplex codes  $S(r, q)$  attain the Griesmer bound.