PAC Checker

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Abstract

Abstract—Generating and checking proof certificates is important to increase the trust in automated reasoning tools. In recent years formal verification using computer algebra became more important and is heavily used in automated circuit verification. An existing proof format which covers algebraic reasoning and allows efficient proof checking is the practical algebraic calculus. In this development, we present the verified checker Pastèque that is obtained by synthesis via the Refinement Framework.

This is the formalization going with our FMCAD'20 tool presentation [1].

Contents

theory *Duplicate-Free-Multiset*
imports *Nested-Multisets-Ordinals.Multiset-More*
begin

1 Duplicate Free Multisets

Duplicate free multisets are isomorphic to finite sets, but it can be useful to reason about duplication to speak about intermediate execution steps in the refinements.

lemma *distinct-mset-remdups-mset-id*: ‹$\mathit{distinct\text{-}mset}\ C \Longrightarrow \mathit{remdups\text{-}mset}\ C = C$›
⟨*proof*⟩

lemma *notin-add-mset-remdups-mset*:
‹$a \notin\# A \Longrightarrow \mathit{add\text{-}mset}\ a\ (\mathit{remdups\text{-}mset}\ A) = \mathit{remdups\text{-}mset}\ (\mathit{add\text{-}mset}\ a\ A)$›
⟨*proof*⟩

lemma *distinct-mset-image-mset*:
‹$\mathit{distinct\text{-}mset}\ (\mathit{image\text{-}mset}\ f\ (\mathit{mset}\ xs)) \longleftrightarrow \mathit{distinct}\ (\mathit{map}\ f\ xs)$›
⟨*proof*⟩

lemma *distinct-image-mset-not-equal*:
assumes

LL': ‹$L \neq L'$› **and**
$dist$: ‹$distinct\text{-}mset\ (image\text{-}mset\ f\ M)$› **and**
L: ‹$L \in\#\ M$› **and**
L': ‹$L' \in\#\ M$› **and**
$fLL'[simp]$: ‹$f\ L = f\ L'$›
shows ‹$False$›
⟨*proof*⟩
lemma *distinct-mset-mono*: ‹$D' \subseteq\#\ D \Longrightarrow distinct\text{-}mset\ D \Longrightarrow distinct\text{-}mset\ D'$›
⟨*proof*⟩

lemma *distinct-mset-mono-strict*: ‹$D' \subset\#\ D \Longrightarrow distinct\text{-}mset\ D \Longrightarrow distinct\text{-}mset\ D'$›
⟨*proof*⟩

lemma *distinct-set-mset-eq-iff*:
assumes ‹$distinct\text{-}mset\ M$› ‹$distinct\text{-}mset\ N$›
shows ‹$set\text{-}mset\ M = set\text{-}mset\ N \longleftrightarrow M = N$›
⟨*proof*⟩

lemma *distinct-mset-union2*:
‹$distinct\text{-}mset\ (A + B) \Longrightarrow distinct\text{-}mset\ B$›
⟨*proof*⟩

lemma *distinct-mset-mset-set*: ‹$distinct\text{-}mset\ (mset\text{-}set\ A)$›
⟨*proof*⟩

lemma *distinct-mset-inter-remdups-mset*:
assumes $dist$: ‹$distinct\text{-}mset\ A$›
shows ‹$A \cap\#\ remdups\text{-}mset\ B = A \cap\#\ B$›
⟨*proof*⟩

lemma *finite-mset-set-inter*:
‹$finite\ A \Longrightarrow finite\ B \Longrightarrow mset\text{-}set\ (A \cap B) = mset\text{-}set\ A \cap\#\ mset\text{-}set\ B$›
⟨*proof*⟩

lemma *removeAll-notin*: ‹$a \notin\#\ A \Longrightarrow removeAll\text{-}mset\ a\ A = A$›
⟨*proof*⟩

lemma *same-mset-distinct-iff*:
‹$mset\ M = mset\ M' \Longrightarrow distinct\ M \longleftrightarrow distinct\ M'$›
⟨*proof*⟩

1.1 More Lists

lemma *in-set-conv-iff*:
‹$x \in set\ (take\ n\ xs) \longleftrightarrow (\exists\, i < n.\ i < length\ xs \wedge xs\ !\ i = x)$›
⟨*proof*⟩

lemma *in-set-take-conv-nth*:
‹$x \in set\ (take\ n\ xs) \longleftrightarrow (\exists\, m{<}min\ n\ (length\ xs).\ xs\ !\ m = x)$›
⟨*proof*⟩

lemma *in-set-remove1D*:
‹$a \in set\ (remove1\ x\ xs) \Longrightarrow a \in set\ xs$›
⟨*proof*⟩

1.2 Generic Multiset

lemma *mset-drop-upto*: ‹$mset\ (drop\ a\ N) = \{\#N!i.\ i \in\#\ mset\text{-}set\ \{a..<length\ N\}\#\}$›
⟨*proof*⟩

1.3 Other

I believe this should be activated by default, as the set becomes much easier...

lemma *Collect-eq-comp′*: ‹ $\{(x, y).\ P\ x\ y\}\ O\ \{(c, a).\ c = f\ a\} = \{(x, a).\ P\ x\ (f\ a)\}$›
⟨*proof*⟩

end

theory *WB-Sort*
imports *Refine-Imperative-HOL.IICF HOL−Library.Rewrite Duplicate-Free-Multiset*
begin

This a complete copy-paste of the IsaFoL version because sharing is too hard.

Every element between *lo* and *hi* can be chosen as pivot element.

definition *choose-pivot* :: ‹$('b \Rightarrow 'b \Rightarrow bool) \Rightarrow ('a \Rightarrow 'b) \Rightarrow 'a\ list \Rightarrow nat \Rightarrow nat \Rightarrow nat\ nres$› **where**
‹$choose\text{-}pivot\ \text{-}\ \text{-}\ \text{-}\ lo\ hi = SPEC(\lambda k.\ k \geq lo \wedge k \leq hi)$›

The element at index *p* partitions the subarray *lo..hi*. This means that every element

definition *isPartition-wrt* :: ‹$('b \Rightarrow 'b \Rightarrow bool) \Rightarrow 'b\ list \Rightarrow nat \Rightarrow nat \Rightarrow nat \Rightarrow bool$› **where**
‹$isPartition\text{-}wrt\ R\ xs\ lo\ hi\ p \equiv (\forall\ i.\ i \geq lo \wedge i < p \longrightarrow R\ (xs!i)\ (xs!p)) \wedge (\forall\ j.\ j > p \wedge j \leq hi \longrightarrow R\ (xs!p)\ (xs!j))$›

lemma *isPartition-wrtI*:
‹$(\bigwedge\ i.\ [\![i \geq lo;\ i < p]\!] \Longrightarrow R\ (xs!i)\ (xs!p)) \Longrightarrow (\bigwedge\ j.\ [\![j > p;\ j \leq hi]\!] \Longrightarrow R\ (xs!p)\ (xs!j)) \Longrightarrow isPartition\text{-}wrt\ R\ xs\ lo\ hi\ p$›
⟨*proof*⟩

definition *isPartition* :: ‹$'a :: order\ list \Rightarrow nat \Rightarrow nat \Rightarrow nat \Rightarrow bool$› **where**
‹$isPartition\ xs\ lo\ hi\ p \equiv isPartition\text{-}wrt\ (\leq)\ xs\ lo\ hi\ p$›

abbreviation *isPartition-map* :: ‹$('b \Rightarrow 'b \Rightarrow bool) \Rightarrow ('a \Rightarrow 'b) \Rightarrow 'a\ list \Rightarrow nat \Rightarrow nat \Rightarrow nat \Rightarrow bool$›
where
‹$isPartition\text{-}map\ R\ h\ xs\ i\ j\ k \equiv isPartition\text{-}wrt\ (\lambda a\ b.\ R\ (h\ a)\ (h\ b))\ xs\ i\ j\ k$›

lemma *isPartition-map-def′*:
‹$lo \leq p \Longrightarrow p \leq hi \Longrightarrow hi < length\ xs \Longrightarrow isPartition\text{-}map\ R\ h\ xs\ lo\ hi\ p = isPartition\text{-}wrt\ R\ (map\ h\ xs)\ lo\ hi\ p$›
⟨*proof*⟩

Example: 6 is the pivot element (with index 4); $7{::}'a$ is equal to the $length\ xs - 1$.

lemma ‹$isPartition\ [0,5,3,4,6,9,8,10{::}nat]\ 0\ 7\ 4$›
⟨*proof*⟩

definition *sublist* :: ‹$'a\ list \Rightarrow nat \Rightarrow nat \Rightarrow 'a\ list$› **where**
‹$sublist\ xs\ i\ j \equiv take\ (Suc\ j - i)\ (drop\ i\ xs)$›

lemma *take-Suc0*:
$l \neq [] \Longrightarrow take\ (Suc\ 0)\ l = [l!0]$
$0 < length\ l \Longrightarrow take\ (Suc\ 0)\ l = [l!0]$
$Suc\ n \leq length\ l \Longrightarrow take\ (Suc\ 0)\ l = [l!0]$
⟨*proof*⟩

lemma *sublist-single*: ‹$i < length\ xs \Longrightarrow sublist\ xs\ i\ i = [xs!i]$›
⟨*proof*⟩

lemma *insert-eq*: ‹$insert\ a\ b = b \cup \{a\}$›
⟨*proof*⟩

lemma *sublist-nth*: ‹$[\![lo \leq hi;\ hi < length\ xs;\ k{+}lo \leq hi]\!] \Longrightarrow (sublist\ xs\ lo\ hi)!k = xs!(lo{+}k)$›
⟨*proof*⟩

lemma *sublist-length*: ‹$[\![i \leq j;\ j < length\ xs]\!] \Longrightarrow length\ (sublist\ xs\ i\ j) = 1 + j - i$›
⟨*proof*⟩

lemma *sublist-not-empty*: ‹$[\![i \leq j;\ j < length\ xs;\ xs \neq []]\!] \Longrightarrow sublist\ xs\ i\ j \neq []$›
⟨*proof*⟩

lemma *sublist-app*: ‹$[\![i1 \leq i2;\ i2 \leq i3]\!] \Longrightarrow sublist\ xs\ i1\ i2\ @\ sublist\ xs\ (Suc\ i2)\ i3 = sublist\ xs\ i1\ i3$›
⟨*proof*⟩

definition *sorted-sublist-wrt* :: ‹$('b \Rightarrow 'b \Rightarrow bool) \Rightarrow 'b\ list \Rightarrow nat \Rightarrow nat \Rightarrow bool$› **where**
‹$sorted\text{-}sublist\text{-}wrt\ R\ xs\ lo\ hi = sorted\text{-}wrt\ R\ (sublist\ xs\ lo\ hi)$›

definition *sorted-sublist* :: ‹$'a :: linorder\ list \Rightarrow nat \Rightarrow nat \Rightarrow bool$› **where**
‹$sorted\text{-}sublist\ xs\ lo\ hi = sorted\text{-}sublist\text{-}wrt\ (\leq)\ xs\ lo\ hi$›

abbreviation *sorted-sublist-map* :: ‹$('b \Rightarrow 'b \Rightarrow bool) \Rightarrow ('a \Rightarrow 'b) \Rightarrow 'a\ list \Rightarrow nat \Rightarrow nat \Rightarrow bool$›
where
‹$sorted\text{-}sublist\text{-}map\ R\ h\ xs\ lo\ hi \equiv sorted\text{-}sublist\text{-}wrt\ (\lambda a\ b.\ R\ (h\ a)\ (h\ b))\ xs\ lo\ hi$›

lemma *sorted-sublist-map-def'*:
‹$lo < length\ xs \Longrightarrow sorted\text{-}sublist\text{-}map\ R\ h\ xs\ lo\ hi \equiv sorted\text{-}sublist\text{-}wrt\ R\ (map\ h\ xs)\ lo\ hi$›
⟨*proof*⟩

lemma *sorted-sublist-wrt-refl*: ‹$i < length\ xs \Longrightarrow sorted\text{-}sublist\text{-}wrt\ R\ xs\ i\ i$›
⟨*proof*⟩

lemma *sorted-sublist-refl*: ‹$i < length\ xs \Longrightarrow sorted\text{-}sublist\ xs\ i\ i$›
⟨*proof*⟩

lemma *sublist-map*: ‹$sublist\ (map\ f\ xs)\ i\ j = map\ f\ (sublist\ xs\ i\ j)$›
⟨*proof*⟩

lemma *take-set*: ‹$j \leq length\ xs \Longrightarrow x \in set\ (take\ j\ xs) \equiv (\exists\ k.\ k < j \wedge xs!k = x)$›
⟨*proof*⟩

lemma *drop-set*: ‹$j \leq length\ xs \Longrightarrow x \in set\ (drop\ j\ xs) \equiv (\exists k.\ j{\leq}k{\wedge}k{<}length\ xs \wedge xs!k{=}x)$›
⟨*proof*⟩

lemma *sublist-el*: ‹$i \leq j \Longrightarrow j < length\ xs \Longrightarrow x \in set\ (sublist\ xs\ i\ j) \equiv (\exists\ k.\ k < Suc\ j-i \wedge xs!(i+k)=x)$›
⟨*proof*⟩

lemma *sublist-el′*: ‹$i \leq j \Longrightarrow j < length\ xs \Longrightarrow x \in set\ (sublist\ xs\ i\ j) \equiv (\exists\ k.\ i\leq k\wedge k\leq j \wedge xs!k=x)$›
⟨*proof*⟩

lemma *sublist-lt*: ‹$hi < lo \Longrightarrow sublist\ xs\ lo\ hi = []$›
⟨*proof*⟩

lemma *nat-le-eq-or-lt*: ‹$(a :: nat) \leq b = (a = b \vee a < b)$›
⟨*proof*⟩

lemma *sorted-sublist-wrt-le*: ‹$hi \leq lo \Longrightarrow hi < length\ xs \Longrightarrow sorted\text{-}sublist\text{-}wrt\ R\ xs\ lo\ hi$›
⟨*proof*⟩

Elements in a sorted sublists are actually sorted

lemma *sorted-sublist-wrt-nth-le*:
assumes ‹$sorted\text{-}sublist\text{-}wrt\ R\ xs\ lo\ hi$› **and** ‹$lo \leq hi$› **and** ‹$hi < length\ xs$› **and**
‹$lo \leq i$› **and** ‹$i < j$› **and** ‹$j \leq hi$›
shows ‹$R\ (xs!i)\ (xs!j)$›
⟨*proof*⟩

We can make the assumption $i < j$ weaker if we have a reflexivie relation.

lemma *sorted-sublist-wrt-nth-le′*:
assumes *ref*: ‹$\bigwedge\ x.\ R\ x\ x$›
and ‹$sorted\text{-}sublist\text{-}wrt\ R\ xs\ lo\ hi$› **and** ‹$lo \leq hi$› **and** ‹$hi < length\ xs$›
and ‹$lo \leq i$› **and** ‹$i \leq j$› **and** ‹$j \leq hi$›
shows ‹$R\ (xs!i)\ (xs!j)$›
⟨*proof*⟩

lemma *sorted-sublist-le*: ‹$hi \leq lo \Longrightarrow hi < length\ xs \Longrightarrow sorted\text{-}sublist\ xs\ lo\ hi$›
⟨*proof*⟩

lemma *sorted-sublist-map-le*: ‹$hi \leq lo \Longrightarrow hi < length\ xs \Longrightarrow sorted\text{-}sublist\text{-}map\ R\ h\ xs\ lo\ hi$›
⟨*proof*⟩

lemma *sublist-cons*: ‹$lo < hi \Longrightarrow hi < length\ xs \Longrightarrow sublist\ xs\ lo\ hi = xs!lo\ \#\ sublist\ xs\ (Suc\ lo)\ hi$›
⟨*proof*⟩

lemma *sorted-sublist-wrt-cons′*:
‹$sorted\text{-}sublist\text{-}wrt\ R\ xs\ (lo+1)\ hi \Longrightarrow lo \leq hi \Longrightarrow hi < length\ xs \Longrightarrow (\forall j.\ lo<j\wedge j\leq hi \longrightarrow R\ (xs!lo)\ (xs!j)) \Longrightarrow sorted\text{-}sublist\text{-}wrt\ R\ xs\ lo\ hi$›
⟨*proof*⟩

lemma *sorted-sublist-wrt-cons*:
assumes *trans*: ‹$(\bigwedge\ x\ y\ z.\ [\![R\ x\ y;\ R\ y\ z]\!] \Longrightarrow R\ x\ z)$› **and**

$\langle$*sorted-sublist-wrt R xs* $(lo+1)$ *hi*$\rangle$ **and**
$\langle lo \le hi\rangle$ **and** $\langle hi <$ *length xs* $\rangle$ **and** $\langle R\ (xs!lo)\ (xs!(lo+1))\rangle$
shows $\langle$*sorted-sublist-wrt R xs lo hi*$\rangle$
$\langle proof\rangle$

lemma *sorted-sublist-map-cons*:
$\langle(\bigwedge x\ y\ z.\ [\![R\ (h\ x)\ (h\ y);\ R\ (h\ y)\ (h\ z)]\!] \Longrightarrow R\ (h\ x)\ (h\ z)) \Longrightarrow$
sorted-sublist-map R h xs $(lo+1)$ *hi* $\Longrightarrow lo \le hi \Longrightarrow hi <$ *length xs* $\Longrightarrow R\ (h\ (xs!lo))\ (h\ (xs!(lo+1)))$
$\Longrightarrow$ *sorted-sublist-map R h xs lo hi*$\rangle$
$\langle proof\rangle$

lemma *sublist-snoc*: $\langle lo < hi \Longrightarrow hi <$ *length xs* $\Longrightarrow$ *sublist xs lo hi* $=$ *sublist xs lo* $(hi-1)$ @ $[xs!hi]\rangle$
$\langle proof\rangle$

lemma *sorted-sublist-wrt-snoc′*:
$\langle$*sorted-sublist-wrt R xs lo* $(hi-1) \Longrightarrow lo \le hi \Longrightarrow hi <$ *length xs* $\Longrightarrow (\forall j.\ lo{\le}j{\wedge}j{<}hi \longrightarrow R\ (xs!j)$
$(xs!hi)) \Longrightarrow$ *sorted-sublist-wrt R xs lo hi*$\rangle$
$\langle proof\rangle$

lemma *sorted-sublist-wrt-snoc*:
assumes *trans*: $\langle(\bigwedge x\ y\ z.\ [\![R\ x\ y;\ R\ y\ z]\!] \Longrightarrow R\ x\ z)\rangle$ **and**
$\langle$*sorted-sublist-wrt R xs lo* $(hi-1)\rangle$ **and**
$\langle lo \le hi\rangle$ **and** $\langle hi <$ *length xs*$\rangle$ **and** $\langle(R\ (xs!(hi-1))\ (xs!hi))\rangle$
shows $\langle$*sorted-sublist-wrt R xs lo hi*$\rangle$
$\langle proof\rangle$

lemma *sublist-split*: $\langle lo \le hi \Longrightarrow lo < p \Longrightarrow p < hi \Longrightarrow hi <$ *length xs* $\Longrightarrow$ *sublist xs lo p* @ *sublist xs*
$(p+1)$ *hi* $=$ *sublist xs lo hi*$\rangle$
$\langle proof\rangle$

lemma *sublist-split-part*: $\langle lo \le hi \Longrightarrow lo < p \Longrightarrow p < hi \Longrightarrow hi <$ *length xs* $\Longrightarrow$ *sublist xs lo* $(p-1)$ @
$xs!p$ # *sublist xs* $(p+1)$ *hi* $=$ *sublist xs lo hi*$\rangle$
$\langle proof\rangle$

A property for partitions (we always assume that R is transitive.

lemma *isPartition-wrt-trans*:
$\langle(\bigwedge x\ y\ z.\ [\![R\ x\ y;\ R\ y\ z]\!] \Longrightarrow R\ x\ z) \Longrightarrow$
isPartition-wrt R xs lo hi p $\Longrightarrow$
$(\forall i\ j.\ lo \le i \wedge i < p \wedge p < j \wedge j \le hi \longrightarrow R\ (xs!i)\ (xs!j))\rangle$
$\langle proof\rangle$

lemma *isPartition-map-trans*:
$\langle(\bigwedge x\ y\ z.\ [\![R\ (h\ x)\ (h\ y);\ R\ (h\ y)\ (h\ z)]\!] \Longrightarrow R\ (h\ x)\ (h\ z)) \Longrightarrow$
$hi <$ *length xs* $\Longrightarrow$
isPartition-map R h xs lo hi p $\Longrightarrow$
$(\forall i\ j.\ lo \le i \wedge i < p \wedge p < j \wedge j \le hi \longrightarrow R\ (h\ (xs!i))\ (h\ (xs!j)))\rangle$
$\langle proof\rangle$

lemma *merge-sorted-wrt-partitions-between′*:
$\langle lo \le hi \Longrightarrow lo < p \Longrightarrow p < hi \Longrightarrow hi <$ *length xs* $\Longrightarrow$
isPartition-wrt R xs lo hi p $\Longrightarrow$
sorted-sublist-wrt R xs lo $(p-1) \Longrightarrow$ *sorted-sublist-wrt R xs* $(p+1)$ *hi* $\Longrightarrow$

$(\forall\, i\ j.\ lo \le i \wedge i < p \wedge p < j \wedge j \le hi \longrightarrow R\ (xs!i)\ (xs!j)) \Longrightarrow$
sorted-sublist-wrt R xs lo hi›
⟨*proof*⟩

lemma *merge-sorted-wrt-partitions-between*:
‹$(\bigwedge x\ y\ z.\ [\![R\ x\ y;\ R\ y\ z]\!] \Longrightarrow R\ x\ z) \Longrightarrow$
isPartition-wrt R xs lo hi p $\Longrightarrow$
sorted-sublist-wrt R xs lo $(p-1) \Longrightarrow$ *sorted-sublist-wrt R xs* $(p+1)$ *hi* $\Longrightarrow$
$lo \le hi \Longrightarrow hi < length\ xs \Longrightarrow lo < p \Longrightarrow p < hi \Longrightarrow hi < length\ xs \Longrightarrow$
sorted-sublist-wrt R xs lo hi›
⟨*proof*⟩

The main theorem to merge sorted lists

lemma *merge-sorted-wrt-partitions*:
‹*isPartition-wrt R xs lo hi p* $\Longrightarrow$
sorted-sublist-wrt R xs lo $(p - Suc\ 0) \Longrightarrow$ *sorted-sublist-wrt R xs* $(Suc\ p)$ *hi* $\Longrightarrow$
$lo \le hi \Longrightarrow lo \le p \Longrightarrow p \le hi \Longrightarrow hi < length\ xs \Longrightarrow$
$(\forall\, i\ j.\ lo \le i \wedge i < p \wedge p < j \wedge j \le hi \longrightarrow R\ (xs!i)\ (xs!j)) \Longrightarrow$
sorted-sublist-wrt R xs lo hi›
⟨*proof*⟩

theorem *merge-sorted-map-partitions*:
‹$(\bigwedge x\ y\ z.\ [\![R\ (h\ x)\ (h\ y);\ R\ (h\ y)\ (h\ z)]\!] \Longrightarrow R\ (h\ x)\ (h\ z)) \Longrightarrow$
isPartition-map R h xs lo hi p $\Longrightarrow$
sorted-sublist-map R h xs lo $(p - Suc\ 0) \Longrightarrow$ *sorted-sublist-map R h xs* $(Suc\ p)$ *hi* $\Longrightarrow$
$lo \le hi \Longrightarrow lo \le p \Longrightarrow p \le hi \Longrightarrow hi < length\ xs \Longrightarrow$
sorted-sublist-map R h xs lo hi›
⟨*proof*⟩

lemma *partition-wrt-extend*:
‹*isPartition-wrt R xs lo′ hi′ p* $\Longrightarrow$
$hi < length\ xs \Longrightarrow$
$lo \le lo' \Longrightarrow lo' \le hi \Longrightarrow hi' \le hi \Longrightarrow$
$lo' \le p \Longrightarrow p \le hi' \Longrightarrow$
$(\bigwedge i.\ lo{\le}i \Longrightarrow i < lo' \Longrightarrow R\ (xs!i)\ (xs!p)) \Longrightarrow$
$(\bigwedge j.\ hi'{<}j \Longrightarrow j{\le}hi \Longrightarrow R\ (xs!p)\ (xs!j)) \Longrightarrow$
isPartition-wrt R xs lo hi p›
⟨*proof*⟩

lemma *partition-map-extend*:
‹*isPartition-map R h xs lo′ hi′ p* $\Longrightarrow$
$hi < length\ xs \Longrightarrow$
$lo \le lo' \Longrightarrow lo' \le hi \Longrightarrow hi' \le hi \Longrightarrow$
$lo' \le p \Longrightarrow p \le hi' \Longrightarrow$
$(\bigwedge i.\ lo{\le}i \Longrightarrow i < lo' \Longrightarrow R\ (h\ (xs!i))\ (h\ (xs!p))) \Longrightarrow$
$(\bigwedge j.\ hi'{<}j \Longrightarrow j{\le}hi \Longrightarrow R\ (h\ (xs!p))\ (h\ (xs!j))) \Longrightarrow$
isPartition-map R h xs lo hi p›
⟨*proof*⟩

lemma *isPartition-empty*:
‹$(\bigwedge j.\ [\![lo < j;\ j \le hi]\!] \Longrightarrow R\ (xs\ !\ lo)\ (xs\ !\ j)) \Longrightarrow$
isPartition-wrt R xs lo hi lo›
⟨*proof*⟩

lemma *take-ext*:
‹($\forall$ i<k. $xs'!i$=$xs!i$) $\Longrightarrow$
k < *length xs* $\Longrightarrow$ k < *length xs'* $\Longrightarrow$
take k xs' = *take k xs*›
⟨*proof*⟩

lemma *drop-ext'*:
‹($\forall$ i. $i \geq k$ $\wedge$ i<*length xs* $\longrightarrow$ $xs'!i$=$xs!i$) $\Longrightarrow$
0<k $\Longrightarrow$ $xs \neq []$ $\Longrightarrow$ — These corner cases will be dealt with in the next lemma
length xs'=*length xs* $\Longrightarrow$
drop k xs' = *drop k xs*›
⟨*proof*⟩

lemma *drop-ext*:
‹($\forall$ i. $i \geq k$ $\wedge$ i<*length xs* $\longrightarrow$ $xs'!i$=$xs!i$) $\Longrightarrow$
length xs'=*length xs* $\Longrightarrow$
drop k xs' = *drop k xs*›
⟨*proof*⟩

lemma *sublist-ext'*:
‹($\forall$ i. $lo \leq i \wedge i \leq hi$ $\longrightarrow$ $xs'!i$=$xs!i$) $\Longrightarrow$
length xs' = *length xs* $\Longrightarrow$
$lo \leq hi$ $\Longrightarrow$ *Suc hi* < *length xs* $\Longrightarrow$
sublist xs' lo hi = *sublist xs lo hi*›
⟨*proof*⟩

lemma *lt-Suc*: ‹(a < b) = (*Suc a* = b $\vee$ *Suc a* < b)›
⟨*proof*⟩

lemma *sublist-until-end-eq-drop*: ‹*Suc hi* = *length xs* $\Longrightarrow$ *sublist xs lo hi* = *drop lo xs*›
⟨*proof*⟩

lemma *sublist-ext*:
‹($\forall$ i. $lo \leq i \wedge i \leq hi$ $\longrightarrow$ $xs'!i$=$xs!i$) $\Longrightarrow$
length xs' = *length xs* $\Longrightarrow$
$lo \leq hi$ $\Longrightarrow$ hi < *length xs* $\Longrightarrow$
sublist xs' lo hi = *sublist xs lo hi*›
⟨*proof*⟩

lemma *sorted-wrt-lower-sublist-still-sorted*:
assumes ‹*sorted-sublist-wrt R xs lo* (lo' − *Suc 0*)› **and**
‹$lo \leq lo'$› **and** ‹lo' < *length xs*› **and**
‹($\forall$ i. $lo \leq i \wedge i < lo'$ $\longrightarrow$ $xs'!i$=$xs!i$)› **and** ‹*length xs'* = *length xs*›
shows ‹*sorted-sublist-wrt R xs' lo* (lo' − *Suc 0*)›
⟨*proof*⟩

lemma *sorted-map-lower-sublist-still-sorted*:
assumes ‹*sorted-sublist-map R h xs lo* (lo' − *Suc 0*)› **and**
‹$lo \leq lo'$› **and** ‹lo' < *length xs*› **and**
‹($\forall$ i. $lo \leq i \wedge i < lo'$ $\longrightarrow$ $xs'!i$=$xs!i$)› **and** ‹*length xs'* = *length xs*›

shows ‹*sorted-sublist-map R h xs' lo* (*lo'* − *Suc 0*)›
⟨*proof*⟩

lemma *sorted-wrt-upper-sublist-still-sorted*:
assumes ‹*sorted-sublist-wrt R xs* (*hi'*+*1*) *hi*› **and**
‹*lo* ≤ *lo'*› **and** ‹*hi* < *length xs*› **and**
‹∀ *j*. *hi'*<*j*∧*j*≤*hi* ⟶ *xs'*!*j*=*xs*!*j*› **and** ‹*length xs'* = *length xs*›
shows ‹*sorted-sublist-wrt R xs'* (*hi'*+*1*) *hi*›
⟨*proof*⟩

lemma *sorted-map-upper-sublist-still-sorted*:
assumes ‹*sorted-sublist-map R h xs* (*hi'*+*1*) *hi*› **and**
‹*lo* ≤ *lo'*› **and** ‹*hi* < *length xs*› **and**
‹∀ *j*. *hi'*<*j*∧*j*≤*hi* ⟶ *xs'*!*j*=*xs*!*j*› **and** ‹*length xs'* = *length xs*›
shows ‹*sorted-sublist-map R h xs'* (*hi'*+*1*) *hi*›
⟨*proof*⟩

The specification of the partition function

definition *partition-spec* :: ‹('*b* ⇒ '*b* ⇒ *bool*) ⇒ ('*a* ⇒ '*b*) ⇒ '*a list* ⇒ *nat* ⇒ *nat* ⇒ '*a list* ⇒ *nat* ⇒ *bool*› **where**
‹*partition-spec R h xs lo hi xs' p* ≡
mset xs' = *mset xs* ∧ — The list is a permutation
isPartition-map R h xs' lo hi p ∧ — We have a valid partition on the resulting list
lo ≤ *p* ∧ *p* ≤ *hi* ∧ — The partition index is in bounds
(∀ *i*. *i*<*lo* ⟶ *xs'*!*i*=*xs*!*i*) ∧ (∀ *i*. *hi*<*i*∧*i*<*length xs'* ⟶ *xs'*!*i*=*xs*!*i*)› — Everything else is unchanged.

lemma *in-set-take-conv-nth*:
‹*x* ∈ *set* (*take n xs*) ⟷ (∃ *m*<*min n* (*length xs*). *xs* ! *m* = *x*)›
⟨*proof*⟩

lemma *mset-drop-upto*: ‹*mset* (*drop a N*) = {#*N*!*i*. *i* ∈# *mset-set* {*a*..<*length N*}#}›
⟨*proof*⟩

lemma *mathias*:
assumes
Perm: ‹*mset xs'* = *mset xs*›
and *I*: ‹*lo*≤*i*› ‹*i*≤*hi*› ‹*xs'*!*i*=*x*›
and *Bounds*: ‹*hi* < *length xs*›
and *Fix*: ‹⋀ *i*. *i*<*lo* ⟹ *xs'*!*i* = *xs*!*i*› ‹⋀ *j*. ⟦*hi*<*j*; *j*<*length xs*⟧ ⟹ *xs'*!*j* = *xs*!*j*›
shows ‹∃ *j*. *lo*≤*j*∧*j*≤*hi* ∧ *xs*!*j* = *x*›
⟨*proof*⟩

If we fix the left and right rest of two permutated lists, then the sublists are also permutations.

But we only need that the sets are equal.

lemma *mset-sublist-incl*:
assumes *Perm*: ‹*mset xs'* = *mset xs*›
and *Fix*: ‹⋀ *i*. *i*<*lo* ⟹ *xs'*!*i* = *xs*!*i*› ‹⋀ *j*. ⟦*hi*<*j*; *j*<*length xs*⟧ ⟹ *xs'*!*j* = *xs*!*j*›
and *bounds*: ‹*lo* ≤ *hi*› ‹*hi* < *length xs*›
shows ‹*set* (*sublist xs' lo hi*) ⊆ *set* (*sublist xs lo hi*)›
⟨*proof*⟩

lemma *mset-sublist-eq*:

assumes ‹$mset\ xs' = mset\ xs$›
and ‹$\bigwedge i.\ i{<}lo \Longrightarrow xs'!i = xs!i$›
and ‹$\bigwedge j.\ [\![hi{<}j;\ j{<}length\ xs]\!] \Longrightarrow xs'!j = xs!j$›
and *bounds*: ‹$lo \le hi$› ‹$hi < length\ xs$›
shows ‹$set\ (sublist\ xs'\ lo\ hi) = set\ (sublist\ xs\ lo\ hi)$›
⟨*proof*⟩

Our abstract recursive quicksort procedure. We abstract over a partition procedure.

definition *quicksort* :: ‹$('b \Rightarrow 'b \Rightarrow bool) \Rightarrow ('a \Rightarrow 'b) \Rightarrow nat \times nat \times 'a\ list \Rightarrow 'a\ list\ nres$› **where**
‹$quicksort\ R\ h = (\lambda(lo,hi,xs0).\ do\ \{$
$RECT\ (\lambda f\ (lo,hi,xs).\ do\ \{$
$ASSERT(lo \le hi \wedge hi < length\ xs \wedge mset\ xs = mset\ xs0)$; — Premise for a partition function
$(xs,\ p) \leftarrow SPEC(uncurry\ (partition\text{-}spec\ R\ h\ xs\ lo\ hi))$; — Abstract partition function
$ASSERT(mset\ xs = mset\ xs0)$;
$xs \leftarrow (if\ p{-}1{\le}lo\ then\ RETURN\ xs\ else\ f\ (lo,\ p{-}1,\ xs))$;
$ASSERT(mset\ xs = mset\ xs0)$;
$if\ hi{\le}p{+}1\ then\ RETURN\ xs\ else\ f\ (p{+}1,\ hi,\ xs)$
$\})\ (lo,hi,xs0)$
$\})$›

As premise for quicksor, we only need that the indices are ok.

definition *quicksort-pre* :: ‹$('b \Rightarrow 'b \Rightarrow bool) \Rightarrow ('a \Rightarrow 'b) \Rightarrow 'a\ list \Rightarrow nat \Rightarrow nat \Rightarrow 'a\ list \Rightarrow bool$›
where
‹$quicksort\text{-}pre\ R\ h\ xs0\ lo\ hi\ xs \equiv lo \le hi \wedge hi < length\ xs \wedge mset\ xs = mset\ xs0$›

definition *quicksort-post* :: ‹$('b \Rightarrow 'b \Rightarrow bool) \Rightarrow ('a \Rightarrow 'b) \Rightarrow nat \Rightarrow nat \Rightarrow 'a\ list \Rightarrow 'a\ list \Rightarrow bool$›
where
‹$quicksort\text{-}post\ R\ h\ lo\ hi\ xs\ xs' \equiv$
$mset\ xs' = mset\ xs\ \wedge$
$sorted\text{-}sublist\text{-}map\ R\ h\ xs'\ lo\ hi\ \wedge$
$(\forall\ i.\ i{<}lo \longrightarrow xs'!i = xs!i)\ \wedge$
$(\forall\ j.\ hi{<}j \wedge j{<}length\ xs \longrightarrow xs'!j = xs!j)$›

Convert Pure to HOL

lemma *quicksort-postI*:
‹$[\![mset\ xs' = mset\ xs;\ sorted\text{-}sublist\text{-}map\ R\ h\ xs'\ lo\ hi;\ (\bigwedge i.\ [\![i{<}lo]\!] \Longrightarrow xs'!i = xs!i);\ (\bigwedge j.\ [\![hi{<}j;\ j{<}length\ xs]\!] \Longrightarrow xs'!j = xs!j)]\!] \Longrightarrow quicksort\text{-}post\ R\ h\ lo\ hi\ xs\ xs'$›
⟨*proof*⟩

The first case for the correctness proof of (abstract) quicksort: We assume that we called the partition function, and we have $p - (1{::}'a) \le lo$ and $hi \le p + (1{::}'a)$.

lemma *quicksort-correct-case1*:
assumes *trans*: ‹$\bigwedge x\ y\ z.\ [\![R\ (h\ x)\ (h\ y);\ R\ (h\ y)\ (h\ z)]\!] \Longrightarrow R\ (h\ x)\ (h\ z)$› **and** *lin*: ‹$\bigwedge x\ y.\ x \ne y \Longrightarrow R\ (h\ x)\ (h\ y) \vee R\ (h\ y)\ (h\ x)$›
and *pre*: ‹$quicksort\text{-}pre\ R\ h\ xs0\ lo\ hi\ xs$›
and *part*: ‹$partition\text{-}spec\ R\ h\ xs\ lo\ hi\ xs'\ p$›
and *ifs*: ‹$p{-}1 \le lo$› ‹$hi \le p{+}1$›
shows ‹$quicksort\text{-}post\ R\ h\ lo\ hi\ xs\ xs'$›
⟨*proof*⟩

In the second case, we have to show that the precondition still holds for (p+1, hi, x’) after the partition.

lemma *quicksort-correct-case2*:
assumes

pre: ‹*quicksort-pre R h xs0 lo hi xs*›
and *part*: ‹*partition-spec R h xs lo hi xs' p*›
and *ifs*: ‹$\neg$ *hi* $\leq$ *p* + *1*›
shows ‹*quicksort-pre R h xs0* (*Suc p*) *hi xs'*›
⟨*proof*⟩

lemma *quicksort-post-set*:
assumes ‹*quicksort-post R h lo hi xs xs'*›
and *bounds*: ‹*lo* $\leq$ *hi*› ‹*hi* < *length xs*›
shows ‹*set* (*sublist xs' lo hi*) = *set* (*sublist xs lo hi*)›
⟨*proof*⟩

In the third case, we have run quicksort recursively on (p+1, hi, xs') after the partition, with hi<=p+1 and p-1<=lo.

lemma *quicksort-correct-case3*:
assumes *trans*: ‹$\bigwedge$ *x y z*. ⟦*R* (*h x*) (*h y*); *R* (*h y*) (*h z*)⟧ $\Longrightarrow$ *R* (*h x*) (*h z*)› **and** *lin*: ‹$\bigwedge$*x y*. *x* $\neq$ *y* $\Longrightarrow$ *R* (*h x*) (*h y*) $\vee$ *R* (*h y*) (*h x*)›
and *pre*: ‹*quicksort-pre R h xs0 lo hi xs*›
and *part*: ‹*partition-spec R h xs lo hi xs' p*›
and *ifs*: ‹*p* − *Suc 0* $\leq$ *lo*› ‹$\neg$ *hi* $\leq$ *Suc p*›
and *IH1'*: ‹*quicksort-post R h* (*Suc p*) *hi xs' xs''*›
shows ‹*quicksort-post R h lo hi xs xs''*›
⟨*proof*⟩

In the 4th case, we have to show that the premise holds for (*lo*, *p* − (*1*::*'b*), *xs'*), in case $\neg$ *p* − (*1*::*'a*) $\leq$ *lo*

Analogous to case 2.

lemma *quicksort-correct-case4*:
assumes
pre: ‹*quicksort-pre R h xs0 lo hi xs*›
and *part*: ‹*partition-spec R h xs lo hi xs' p*›
and *ifs*: ‹$\neg$ *p* − *Suc 0* $\leq$ *lo* ›
shows ‹*quicksort-pre R h xs0 lo* (*p*−*Suc 0*) *xs'*›
⟨*proof*⟩

In the 5th case, we have run quicksort recursively on (lo, p-1, xs').

lemma *quicksort-correct-case5*:
assumes *trans*: ‹$\bigwedge$ *x y z*. ⟦*R* (*h x*) (*h y*); *R* (*h y*) (*h z*)⟧ $\Longrightarrow$ *R* (*h x*) (*h z*)› **and** *lin*: ‹$\bigwedge$*x y*. *x* $\neq$ *y* $\Longrightarrow$ *R* (*h x*) (*h y*) $\vee$ *R* (*h y*) (*h x*)›
and *pre*: ‹*quicksort-pre R h xs0 lo hi xs*›
and *part*: ‹*partition-spec R h xs lo hi xs' p*›
and *ifs*: ‹$\neg$ *p* − *Suc 0* $\leq$ *lo*› ‹*hi* $\leq$ *Suc p*›
and *IH1'*: ‹*quicksort-post R h lo* (*p* − *Suc 0*) *xs' xs''*›
shows ‹*quicksort-post R h lo hi xs xs''*›
⟨*proof*⟩

In the 6th case, we have run quicksort recursively on (lo, p-1, xs'). We show the precondition on the second call on (p+1, hi, xs")

lemma *quicksort-correct-case6*:
assumes
pre: ‹*quicksort-pre R h xs0 lo hi xs*›

and *part*: ‹$partition\text{-}spec\ R\ h\ xs\ lo\ hi\ xs'\ p$›
and *ifs*: ‹$\neg\ p - Suc\ 0 \leq lo$› ‹$\neg\ hi \leq Suc\ p$›
and *IH1*: ‹$quicksort\text{-}post\ R\ h\ lo\ (p - Suc\ 0)\ xs'\ xs''$›
shows ‹$quicksort\text{-}pre\ R\ h\ xs0\ (Suc\ p)\ hi\ xs''$›
⟨*proof*⟩

In the 7th (and last) case, we have run quicksort recursively on (lo, p-1, xs'). We show the postcondition on the second call on (p+1, hi, xs")

lemma *quicksort-correct-case7*:
assumes *trans*: ‹$\bigwedge\ x\ y\ z.\ [\![R\ (h\ x)\ (h\ y);\ R\ (h\ y)\ (h\ z)]\!] \Longrightarrow R\ (h\ x)\ (h\ z)$› **and** *lin*: ‹$\bigwedge x\ y.\ x \neq y \Longrightarrow R\ (h\ x)\ (h\ y) \vee R\ (h\ y)\ (h\ x)$›
and *pre*: ‹$quicksort\text{-}pre\ R\ h\ xs0\ lo\ hi\ xs$›
and *part*: ‹$partition\text{-}spec\ R\ h\ xs\ lo\ hi\ xs'\ p$›
and *ifs*: ‹$\neg\ p - Suc\ 0 \leq lo$› ‹$\neg\ hi \leq Suc\ p$›
and *IH1'*: ‹$quicksort\text{-}post\ R\ h\ lo\ (p - Suc\ 0)\ xs'\ xs''$›
and *IH2'*: ‹$quicksort\text{-}post\ R\ h\ (Suc\ p)\ hi\ xs''\ xs'''$›
shows ‹$quicksort\text{-}post\ R\ h\ lo\ hi\ xs\ xs'''$›
⟨*proof*⟩

We can now show the correctness of the abstract quicksort procedure, using the refinement framework and the above case lemmas.

lemma *quicksort-correct*:
assumes *trans*: ‹$\bigwedge\ x\ y\ z.\ [\![R\ (h\ x)\ (h\ y);\ R\ (h\ y)\ (h\ z)]\!] \Longrightarrow R\ (h\ x)\ (h\ z)$› **and** *lin*: ‹$\bigwedge x\ y.\ x \neq y \Longrightarrow R\ (h\ x)\ (h\ y) \vee R\ (h\ y)\ (h\ x)$›
and *Pre*: ‹$lo0 \leq hi0$› ‹$hi0 < length\ xs0$›
shows ‹$quicksort\ R\ h\ (lo0,hi0,xs0) \leq \Downarrow Id\ (SPEC(\lambda xs.\ quicksort\text{-}post\ R\ h\ lo0\ hi0\ xs0\ xs))$›
⟨*proof*⟩

definition *partition-main-inv* :: ‹$('b \Rightarrow 'b \Rightarrow bool) \Rightarrow ('a \Rightarrow 'b) \Rightarrow nat \Rightarrow nat \Rightarrow 'a\ list \Rightarrow (nat \times nat \times 'a\ list) \Rightarrow bool$› **where**
‹$partition\text{-}main\text{-}inv\ R\ h\ lo\ hi\ xs0\ p \equiv$
$case\ p\ of\ (i,j,xs) \Rightarrow$
$j < length\ xs \wedge j \leq hi \wedge i < length\ xs \wedge lo \leq i \wedge i \leq j \wedge mset\ xs = mset\ xs0 \wedge$
$(\forall k.\ k \geq lo \wedge k < i \longrightarrow R\ (h\ (xs!k))\ (h\ (xs!hi))) \wedge$ — All elements from *lo* to $i - (1::'c)$ are smaller than the pivot
$(\forall k.\ k \geq i \wedge k < j \longrightarrow R\ (h\ (xs!hi))\ (h\ (xs!k))) \wedge$ — All elements from i to $j - (1::'c)$ are greater than the pivot
$(\forall k.\ k < lo \longrightarrow xs!k = xs0!k) \wedge$ — Everything below *lo* is unchanged
$(\forall k.\ k \geq j \wedge k < length\ xs \longrightarrow xs!k = xs0!k)$ — All elements from j are unchanged (including everyting above *hi*)
›

The main part of the partition function. The pivot is assumed to be the last element. This is exactly the "Lomuto partition scheme" partition function from Wikipedia.

definition *partition-main* :: ‹$('b \Rightarrow 'b \Rightarrow bool) \Rightarrow ('a \Rightarrow 'b) \Rightarrow nat \Rightarrow nat \Rightarrow 'a\ list \Rightarrow ('a\ list \times nat)\ nres$› **where**
‹$partition\text{-}main\ R\ h\ lo\ hi\ xs0 = do\ \{$
$ASSERT(hi < length\ xs0);$
$pivot \leftarrow RETURN\ (h\ (xs0\ !\ hi));$
$(i,j,xs) \leftarrow WHILE_T{}^{partition\text{-}main\text{-}inv\ R\ h\ lo\ hi\ xs0}$ — We loop from $j = lo$ to $j = hi - (1::'c)$.

$(\lambda(i,j,xs).\ j < hi)$
$(\lambda(i,j,xs).\ do\ \{$
$ASSERT(i < length\ xs \wedge j < length\ xs);$
$if\ R\ (h\ (xs!j))\ pivot$
$then\ RETURN\ (i+1,\ j+1,\ swap\ xs\ i\ j)$
$else\ RETURN\ (i,\ \ j+1,\ xs)$
$\})$
$(lo,\ lo,\ xs0);$ — i and j are both initialized to lo
$ASSERT(i < length\ xs \wedge j = hi \wedge lo \leq i \wedge hi < length\ xs \wedge mset\ xs = mset\ xs0);$
$RETURN\ (swap\ xs\ i\ hi,\ i)$
$\}\rangle$

lemma *partition-main-correct*:
assumes *bounds*: $\langle hi < length\ xs\rangle\ \langle lo \leq hi\rangle$ **and**
trans: $\langle \bigwedge x\ y\ z.\ [\![R\ (h\ x)\ (h\ y);\ R\ (h\ y)\ (h\ z)]\!] \Longrightarrow R\ (h\ x)\ (h\ z)\rangle$ **and** *lin*: $\langle \bigwedge x\ y.\ R\ (h\ x)\ (h\ y) \vee R\ (h\ y)\ (h\ x)\rangle$
shows $\langle partition\text{-}main\ R\ h\ lo\ hi\ xs \leq SPEC(\lambda(xs',\ p).\ mset\ xs = mset\ xs' \wedge$
$lo \leq p \wedge p \leq hi \wedge isPartition\text{-}map\ R\ h\ xs'\ lo\ hi\ p \wedge (\forall\ i.\ i{<}lo \longrightarrow xs'!i{=}xs!i) \wedge (\forall\ i.\ hi{<}i{\wedge}i{<}length\ xs' \longrightarrow xs'!i{=}xs!i))\rangle$
$\langle proof\rangle$

definition *partition-between* :: $\langle('b \Rightarrow 'b \Rightarrow bool) \Rightarrow ('a \Rightarrow 'b) \Rightarrow nat \Rightarrow nat \Rightarrow 'a\ list \Rightarrow ('a\ list \times nat)\ nres\rangle$ **where**
$\langle partition\text{-}between\ R\ h\ lo\ hi\ xs0 = do\ \{$
$ASSERT(hi < length\ xs0 \wedge lo \leq hi);$
$k \leftarrow choose\text{-}pivot\ R\ h\ xs0\ lo\ hi;$ — choice of pivot
$ASSERT(k < length\ xs0);$
$xs \leftarrow RETURN\ (swap\ xs0\ k\ hi);$ — move the pivot to the last position, before we start the actual loop
$ASSERT(length\ xs = length\ xs0);$
$partition\text{-}main\ R\ h\ lo\ hi\ xs$
$\}\rangle$

lemma *partition-between-correct*:
assumes $\langle hi < length\ xs\rangle$ **and** $\langle lo \leq hi\rangle$ **and**
$\langle \bigwedge x\ y\ z.\ [\![R\ (h\ x)\ (h\ y);\ R\ (h\ y)\ (h\ z)]\!] \Longrightarrow R\ (h\ x)\ (h\ z)\rangle$ **and** $\langle \bigwedge x\ y.\ R\ (h\ x)\ (h\ y) \vee R\ (h\ y)\ (h\ x)\rangle$
shows $\langle partition\text{-}between\ R\ h\ lo\ hi\ xs \leq SPEC(uncurry\ (partition\text{-}spec\ R\ h\ xs\ lo\ hi))\rangle$
$\langle proof\rangle$

We use the median of the first, the middle, and the last element.

definition *choose-pivot3* **where**
$\langle choose\text{-}pivot3\ R\ h\ xs\ lo\ (hi{::}nat) = do\ \{$
$ASSERT(lo < length\ xs);$
$ASSERT(hi < length\ xs);$
$let\ k' = (hi - lo)\ div\ 2;$
$let\ k = lo + k';$
$ASSERT(k < length\ xs);$
$let\ start = h\ (xs\ !\ lo);$
$let\ mid = h\ (xs\ !\ k);$
$let\ end = h\ (xs\ !\ hi);$
$if\ (R\ start\ mid \wedge R\ mid\ end) \vee (R\ end\ mid \wedge R\ mid\ start)\ then\ RETURN\ k$

else if $(R\ start\ end \wedge R\ end\ mid) \vee (R\ mid\ end \wedge R\ end\ start)$ *then RETURN hi*
else RETURN lo
}›

— We only have to show that this procedure yields a valid index between *lo* and *hi*.
lemma *choose-pivot3-choose-pivot*:
assumes ‹$lo < length\ xs$› ‹$hi < length\ xs$› ‹$hi \geq lo$›
shows ‹$choose\text{-}pivot3\ R\ h\ xs\ lo\ hi \leq \Downarrow Id\ (choose\text{-}pivot\ R\ h\ xs\ lo\ hi)$›
⟨*proof*⟩

The refined partion function: We use the above pivot function and fold instead of non-deterministic iteration.

definition *partition-between-ref*
:: ‹$('b \Rightarrow 'b \Rightarrow bool) \Rightarrow ('a \Rightarrow 'b) \Rightarrow nat \Rightarrow nat \Rightarrow 'a\ list \Rightarrow ('a\ list \times nat)\ nres$›
where
‹*partition-between-ref R h lo hi xs0* = *do* {
$ASSERT(hi < length\ xs0 \wedge hi < length\ xs0 \wedge lo \leq hi)$;
$k \leftarrow choose\text{-}pivot3\ R\ h\ xs0\ lo\ hi$; — choice of pivot
$ASSERT(k < length\ xs0)$;
$xs \leftarrow RETURN\ (swap\ xs0\ k\ hi)$; — move the pivot to the last position, before we start the actual loop
$ASSERT(length\ xs = length\ xs0)$;
partition-main R h lo hi xs
}›

lemma *partition-main-ref'*:
‹*partition-main R h lo hi xs*
$\leq \Downarrow ((\lambda\ a\ b\ c\ d.\ Id)\ a\ b\ c\ d)$ (*partition-main R h lo hi xs*)›
⟨*proof*⟩

lemma *Down-id-eq*:
‹$\Downarrow Id\ x = x$›
⟨*proof*⟩

lemma *partition-between-ref-partition-between*:
‹*partition-between-ref R h lo hi xs* ≤ (*partition-between R h lo hi xs*)›
⟨*proof*⟩

Technical lemma for sepref

lemma *partition-between-ref-partition-between'*:
‹(*uncurry2* (*partition-between-ref R h*), *uncurry2* (*partition-between R h*)) ∈
$(nat\text{-}rel \times_r nat\text{-}rel) \times_r \langle Id\rangle list\text{-}rel \rightarrow_f \langle\langle Id\rangle list\text{-}rel \times_r nat\text{-}rel\rangle nres\text{-}rel$›
⟨*proof*⟩

Example instantiation for pivot

definition *choose-pivot3-impl* **where**
‹*choose-pivot3-impl* = *choose-pivot3* (≤) *id*›

lemma *partition-between-ref-correct*:
assumes *trans*: ‹$\bigwedge\ x\ y\ z.\ [\![R\ (h\ x)\ (h\ y);\ R\ (h\ y)\ (h\ z)]\!] \Longrightarrow R\ (h\ x)\ (h\ z)$› **and** *lin*: ‹$\bigwedge x\ y.\ R\ (h\ x)\ (h\ y) \vee R\ (h\ y)\ (h\ x)$›

and *bounds*: ‹$hi < length\ xs$› ‹$lo \leq hi$›
shows ‹$partition\text{-}between\text{-}ref\ R\ h\ lo\ hi\ xs \leq SPEC\ (uncurry\ (partition\text{-}spec\ R\ h\ xs\ lo\ hi))$›
⟨*proof*⟩

Refined quicksort algorithm: We use the refined partition function.

definition *quicksort-ref* :: ‹$- \Rightarrow - \Rightarrow nat \times nat \times 'a\ list \Rightarrow 'a\ list\ nres$› **where**
‹$quicksort\text{-}ref\ R\ h = (\lambda(lo,hi,xs0).$
$do\ \{$
$RECT\ (\lambda f\ (lo,hi,xs).\ do\ \{$
$ASSERT(lo \leq hi \wedge hi < length\ xs0 \wedge mset\ xs = mset\ xs0);$
$(xs,\ p) \leftarrow partition\text{-}between\text{-}ref\ R\ h\ lo\ hi\ xs;$ — This is the refined partition function. Note that we need the premises (trans,lin,bounds) here.
$ASSERT(mset\ xs = mset\ xs0 \wedge p \geq lo \wedge p < length\ xs0);$
$xs \leftarrow (if\ p{-}1{\leq}lo\ then\ RETURN\ xs\ else\ f\ (lo,\ p{-}1,\ xs));$
$ASSERT(mset\ xs = mset\ xs0);$
$if\ hi{\leq}p{+}1\ then\ RETURN\ xs\ else\ f\ (p{+}1,\ hi,\ xs)$
$\})\ (lo,hi,xs0)$
$\})$›

lemma *fref-to-Down-curry2*:
‹$(uncurry2\ f,\ uncurry2\ g) \in [P]_f\ A \rightarrow \langle B\rangle nres\text{-}rel \Longrightarrow$
$(\bigwedge x\ x'\ y\ y'\ z\ z'.\ P\ ((x',\ y'),\ z') \Longrightarrow (((x,\ y),\ z),\ ((x',\ y'),\ z')) \in A \Longrightarrow$
$f\ x\ y\ z \leq \Downarrow B\ (g\ x'\ y'\ z'))$›
⟨*proof*⟩

lemma *fref-to-Down-curry*:
‹$(f,\ g) \in [P]_f\ A \rightarrow \langle B\rangle nres\text{-}rel \Longrightarrow$
$(\bigwedge x\ x'.\ P\ x' \Longrightarrow (x,\ x') \in A \Longrightarrow$
$f\ x\ \leq \Downarrow B\ (g\ x'))$›
⟨*proof*⟩

lemma *quicksort-ref-quicksort*:
assumes *bounds*: ‹$hi < length\ xs$› ‹$lo \leq hi$› **and**
trans: ‹$\bigwedge\ x\ y\ z.\ [\![R\ (h\ x)\ (h\ y);\ R\ (h\ y)\ (h\ z)]\!] \Longrightarrow R\ (h\ x)\ (h\ z)$› **and** *lin*: ‹$\bigwedge x\ y.\ R\ (h\ x)\ (h\ y) \vee R\ (h\ y)\ (h\ x)$›
shows ‹$quicksort\text{-}ref\ R\ h\ x0 \leq \Downarrow Id\ (quicksort\ R\ h\ x0)$›
⟨*proof*⟩
definition *full-quicksort* **where**
‹$full\text{-}quicksort\ R\ h\ xs \equiv if\ xs = []\ then\ RETURN\ xs\ else\ quicksort\ R\ h\ (0,\ length\ xs - 1,\ xs)$›

definition *full-quicksort-ref* **where**
‹$full\text{-}quicksort\text{-}ref\ R\ h\ xs \equiv$
$if\ List.null\ xs\ then\ RETURN\ xs$
$else\ quicksort\text{-}ref\ R\ h\ (0,\ length\ xs - 1,\ xs)$›

definition *full-quicksort-impl* :: ‹$nat\ list \Rightarrow nat\ list\ nres$› **where**
‹$full\text{-}quicksort\text{-}impl\ xs = full\text{-}quicksort\text{-}ref\ (\leq)\ id\ xs$›

lemma *full-quicksort-ref-full-quicksort*:
assumes *trans*: ‹$\bigwedge\ x\ y\ z.\ [\![R\ (h\ x)\ (h\ y);\ R\ (h\ y)\ (h\ z)]\!] \Longrightarrow R\ (h\ x)\ (h\ z)$› **and** *lin*: ‹$\bigwedge x\ y.\ R\ (h\ x)\ (h\ y) \vee R\ (h\ y)\ (h\ x)$›

shows ‹(*full-quicksort-ref R h*, *full-quicksort R h*) ∈
⟨*Id*⟩*list-rel* $\rightarrow_f$ ⟨ ⟨*Id*⟩*list-rel*⟩*nres-rel*›
⟨*proof*⟩

lemma *sublist-entire*:
‹*sublist xs 0* (*length xs* − *1*) = *xs*›
⟨*proof*⟩

lemma *sorted-sublist-wrt-entire*:
assumes ‹*sorted-sublist-wrt R xs 0* (*length xs* − *1*)›
shows ‹*sorted-wrt R xs*›
⟨*proof*⟩

lemma *sorted-sublist-map-entire*:
assumes ‹*sorted-sublist-map R h xs 0* (*length xs* − *1*)›
shows ‹*sorted-wrt* (λ *x y*. *R* (*h x*) (*h y*)) *xs*›
⟨*proof*⟩

Final correctness lemma

theorem *full-quicksort-correct-sorted*:
assumes
trans: ‹$\bigwedge$*x y z*. ⟦*R* (*h x*) (*h y*); *R* (*h y*) (*h z*)⟧ ⟹ *R* (*h x*) (*h z*)› **and** *lin*: ‹$\bigwedge$*x y*. *x* ≠ *y* ⟹ *R* (*h x*) (*h y*) ∨ *R* (*h y*) (*h x*)›
shows ‹*full-quicksort R h xs* ≤ ⇓ *Id* (*SPEC*(λ*xs*'. *mset xs*' = *mset xs* ∧ *sorted-wrt* (λ *x y*. *R* (*h x*) (*h y*)) *xs*'))›
⟨*proof*⟩

lemma *full-quicksort-correct*:
assumes
trans: ‹$\bigwedge$*x y z*. ⟦*R* (*h x*) (*h y*); *R* (*h y*) (*h z*)⟧ ⟹ *R* (*h x*) (*h z*)› **and**
lin: ‹$\bigwedge$*x y*. *R* (*h x*) (*h y*) ∨ *R* (*h y*) (*h x*)›
shows ‹*full-quicksort R h xs* ≤ ⇓ *Id* (*SPEC*(λ*xs*'. *mset xs*' = *mset xs*))›
⟨*proof*⟩

end
theory *More-Loops*
imports
Refine-Monadic.Refine-While
Refine-Monadic.Refine-Foreach
HOL−Library.Rewrite
begin

1.4 More Theorem about Loops

Most theorem below have a counterpart in the Refinement Framework that is weaker (by missing assertions for example that are critical for code generation).

lemma *Down-id-eq*:
‹⇓*Id x* = *x*›
⟨*proof*⟩

lemma *while-upt-while-direct1*:
b ≥ *a* ⟹
do {

$(\text{-},\sigma) \leftarrow WHILE_T$ $(FOREACH\text{-}cond\ c)$ $(\lambda x.\ do\ \{ASSERT\ (FOREACH\text{-}cond\ c\ x);\ FOREACH\text{-}body\ f\ x\})$
$([a..{<}b],\sigma)$;
$RETURN\ \sigma$
$\} \leq do\ \{$
$(\text{-},\sigma) \leftarrow WHILE_T\ (\lambda(i, x).\ i < b \wedge c\ x)\ (\lambda(i, x).\ do\ \{ASSERT\ (i < b);\ \sigma'{\leftarrow}f\ i\ x;\ RETURN\ (i{+}1,\sigma')\})\ (a,\sigma)$;
$RETURN\ \sigma$
$\}$
$\langle proof\rangle$

lemma *while-upt-while-direct2*:
$b \geq a \Longrightarrow$
$do\ \{$
$(\text{-},\sigma) \leftarrow WHILE_T$ $(FOREACH\text{-}cond\ c)$ $(\lambda x.\ do\ \{ASSERT\ (FOREACH\text{-}cond\ c\ x);\ FOREACH\text{-}body\ f\ x\})$
$([a..{<}b],\sigma)$;
$RETURN\ \sigma$
$\} \geq do\ \{$
$(\text{-},\sigma) \leftarrow WHILE_T\ (\lambda(i, x).\ i < b \wedge c\ x)\ (\lambda(i, x).\ do\ \{ASSERT\ (i < b);\ \sigma'{\leftarrow}f\ i\ x;\ RETURN\ (i{+}1,\sigma')\})\ (a,\sigma)$;
$RETURN\ \sigma$
$\}$
$\langle proof\rangle$

lemma *while-upt-while-direct*:
$b \geq a \Longrightarrow$
$do\ \{$
$(\text{-},\sigma) \leftarrow WHILE_T$ $(FOREACH\text{-}cond\ c)$ $(\lambda x.\ do\ \{ASSERT\ (FOREACH\text{-}cond\ c\ x);\ FOREACH\text{-}body\ f\ x\})$
$([a..{<}b],\sigma)$;
$RETURN\ \sigma$
$\} = do\ \{$
$(\text{-},\sigma) \leftarrow WHILE_T\ (\lambda(i, x).\ i < b \wedge c\ x)\ (\lambda(i, x).\ do\ \{ASSERT\ (i < b);\ \sigma'{\leftarrow}f\ i\ x;\ RETURN\ (i{+}1,\sigma')\})\ (a,\sigma)$;
$RETURN\ \sigma$
$\}$
$\langle proof\rangle$

lemma *while-nfoldli*:
$do\ \{$
$(\text{-},\sigma) \leftarrow WHILE_T$ $(FOREACH\text{-}cond\ c)$ $(\lambda x.\ do\ \{ASSERT\ (FOREACH\text{-}cond\ c\ x);\ FOREACH\text{-}body\ f\ x\})\ (l,\sigma)$;
$RETURN\ \sigma$
$\} \leq nfoldli\ l\ c\ f\ \sigma$
$\langle proof\rangle$
lemma *nfoldli-while*: $nfoldli\ l\ c\ f\ \sigma$
$\leq$
$(WHILE_T{}^I$
$(FOREACH\text{-}cond\ c)\ (\lambda x.\ do\ \{ASSERT\ (FOREACH\text{-}cond\ c\ x);\ FOREACH\text{-}body\ f\ x\})\ (l, \sigma)$
$\gg\!=$
$(\lambda(\text{-}, \sigma).\ RETURN\ \sigma))$
$\langle proof\rangle$

lemma *while-eq-nfoldli*: $do\ \{$

(-,σ) $\leftarrow$ *WHILE*$_T$ (*FOREACH-cond c*) ($\lambda x.$ *do* {*ASSERT* (*FOREACH-cond c x*); *FOREACH-body f x*}) (*l*,σ);
RETURN σ
} = *nfoldli l c f* σ
⟨*proof*⟩

end
theory *PAC-More-Poly*
imports *HOL−Library.Poly-Mapping HOL−Algebra.Polynomials Polynomials.MPoly-Type-Class*
HOL−Algebra.Module
HOL−Library.Countable-Set
begin

2 Libraries

2.1 More Polynoms

Here are more theorems on polynomials. Most of these facts are extremely trivial and should probably be generalised and moved to the Isabelle distribution.

lemma *Const*$_0$*-add*:
‹*Const*$_0$ (*a* + *b*) = *Const*$_0$ *a* + *Const*$_0$ *b*›
⟨*proof*⟩

lemma *Const-mult*:
‹*Const* (*a* ∗ *b*) = *Const a* ∗ *Const b*›
⟨*proof*⟩

lemma *Const*$_0$*-mult*:
‹*Const*$_0$ (*a* ∗ *b*) = *Const*$_0$ *a* ∗ *Const*$_0$ *b*›
⟨*proof*⟩

lemma *Const0*[*simp*]:
‹*Const 0* = *0*›
⟨*proof*⟩

lemma (**in** −) *Const-uminus*[*simp*]:
‹*Const* (−*n*) = − *Const n*›
⟨*proof*⟩

lemma [*simp*]: ‹*Const*$_0$ *0* = *0*›
‹*MPoly 0* = *0*›
⟨*proof*⟩

lemma *Const-add*:
‹*Const* (*a* + *b*) = *Const a* + *Const b*›
⟨*proof*⟩

instance *mpoly* :: (*comm-semiring-1*) *comm-semiring-1*
⟨*proof*⟩

lemma *degree-uminus*[*simp*]:
‹*degree* (−*A*) x' = *degree A* x'›
⟨*proof*⟩

lemma *degree-sum-notin*:
‹$x' \notin vars\ B \Longrightarrow degree\ (A + B)\ x' = degree\ A\ x'$›
⟨*proof*⟩

lemma *degree-notin-vars*:
‹$x \notin (vars\ B) \Longrightarrow degree\ (B :: 'a :: \{monoid\text{-}add\}\ mpoly)\ x = 0$›
⟨*proof*⟩

lemma *not-in-vars-coeff0*:
‹$x \notin vars\ p \Longrightarrow MPoly\text{-}Type.coeff\ p\ (monomial\ (Suc\ 0)\ x) = 0$›
⟨*proof*⟩

lemma *keys-mapping-sum-add*:
‹$finite\ A \Longrightarrow keys\ (mapping\text{-}of\ (\sum v \in A.\ f\ v)) \subseteq \bigcup(keys\ `\ mapping\text{-}of\ `\ f\ `\ UNIV)$›
⟨*proof*⟩

lemma *vars-sum-vars-union*:
fixes f :: ‹$int\ mpoly \Rightarrow int\ mpoly$›
assumes ‹$finite\ \{v.\ f\ v \neq 0\}$›
shows ‹$vars\ (\sum v \mid f\ v \neq 0.\ f\ v * v) \subseteq \bigcup(vars\ `\ \{v.\ f\ v \neq 0\}) \cup \bigcup(vars\ `\ f\ `\ \{v.\ f\ v \neq 0\})$›
(**is** ‹$?A \subseteq ?B$›)
⟨*proof*⟩

lemma *vars-in-right-only*:
$x \in vars\ q \Longrightarrow x \notin vars\ p \Longrightarrow x \in vars\ (p{+}q)$
⟨*proof*⟩

lemma [*simp*]:
‹$vars\ 0 = \{\}$›
⟨*proof*⟩

lemma *vars-Un-nointer*:
‹$keys\ (mapping\text{-}of\ p) \cap\ keys\ (mapping\text{-}of\ q) = \{\} \Longrightarrow vars\ (p + q) = vars\ p \cup vars\ q$›
⟨*proof*⟩

lemmas [*simp*] = *zero-mpoly.rep-eq*

lemma *polynom-sum-monoms*:
fixes p :: ‹$'a :: \{comm\text{-}monoid\text{-}add,cancel\text{-}comm\text{-}monoid\text{-}add\}\ mpoly$›
shows
‹$p = (\sum x{\in}keys\ (mapping\text{-}of\ p).\ MPoly\text{-}Type.monom\ x\ (MPoly\text{-}Type.coeff\ p\ x))$›
‹$keys\ (mapping\text{-}of\ p) \subseteq I \Longrightarrow finite\ I \Longrightarrow p = (\sum x{\in}I.\ MPoly\text{-}Type.monom\ x\ (MPoly\text{-}Type.coeff\ p\ x))$›
⟨*proof*⟩

lemma *vars-mult-monom*:
fixes p :: ‹$int\ mpoly$›
shows ‹$vars\ (p * (monom\ (monomial\ (Suc\ 0)\ x')\ 1)) = (if\ p = 0\ then\ \{\}\ else\ insert\ x'\ (vars\ p))$›
⟨*proof*⟩

lemma *in-mapping-mult-single*:
‹$x \in (\lambda x.\ lookup\ x\ x')\ `\ keys\ (A * (Var_0\ x' :: (nat \Rightarrow_0 nat) \Rightarrow_0 'b :: \{monoid\text{-}mult,zero\text{-}neq\text{-}one,semiring\text{-}0\}))$

$\longleftrightarrow$
$x > 0 \wedge x - 1 \in (\lambda x.\ lookup\ x\ x')$ ‘ $keys\ (A)$›
⟨proof⟩

lemma *Max-Suc-Suc-Max*:
‹$finite\ A \Longrightarrow A \neq \{\} \Longrightarrow Max\ (insert\ 0\ (Suc$ ‘ $A)) =$
$Suc\ (Max\ (insert\ 0\ A))$›
⟨proof⟩

lemma [*simp*]:
‹$keys\ (Var_0\ x' :: ('a \Rightarrow_0 nat) \Rightarrow_0 'b :: \{zero\text{-}neq\text{-}one\}) = \{Poly\text{-}Mapping.single\ x'\ 1\}$›
⟨proof⟩

lemma *degree-mult-Var*:
‹$degree\ (A * Var\ x')\ x' = (if\ A = 0\ then\ 0\ else\ Suc\ (degree\ A\ x'))$› **for** A :: ‹*int mpoly*›
⟨proof⟩

lemma *degree-mult-Var'*:
‹$degree\ (Var\ x' * A)\ x' = (if\ A = 0\ then\ 0\ else\ Suc\ (degree\ A\ x'))$› **for** A :: ‹*int mpoly*›
⟨proof⟩

lemma *degree-add-max*:
‹$degree\ (A + B)\ x \leq max\ (degree\ A\ x)\ (degree\ B\ x)$›
⟨proof⟩

lemma *degree-times-le*:
‹$degree\ (A * B)\ x \leq degree\ A\ x + degree\ B\ x$›
⟨proof⟩

lemma *monomial-inj*:
$monomial\ c\ s = monomial\ (d{::}'b{::}zero\text{-}neq\text{-}one)\ t \longleftrightarrow (c = 0 \wedge d = 0) \vee (c = d \wedge s = t)$
⟨proof⟩

lemma *MPoly-monomial-power'*:
‹$MPoly\ (monomial\ 1\ x')\ \hat{}\ (n{+}1) =\ MPoly\ (monomial\ (1)\ (((\lambda x.\ x + x')\ \hat{}\hat{}\ n)\ x'))$›
⟨proof⟩

lemma *MPoly-monomial-power*:
‹$n > 0 \Longrightarrow MPoly\ (monomial\ 1\ x')\ \hat{}\ (n) =\ MPoly\ (monomial\ (1)\ (((\lambda x.\ x + x')\ \hat{}\hat{}\ (n - 1))\ x'))$›
⟨proof⟩

lemma *vars-uminus*[*simp*]:
‹$vars\ (-p) = vars\ p$›
⟨proof⟩

lemma *coeff-uminus*[*simp*]:
‹$MPoly\text{-}Type.coeff\ (-p)\ x = -MPoly\text{-}Type.coeff\ p\ x$›
⟨proof⟩

definition *decrease-key*::$'a \Rightarrow ('a \Rightarrow_0 'b{::}\{monoid\text{-}add,\ minus,one\}) \Rightarrow ('a \Rightarrow_0 'b)$ **where**
$decrease\text{-}key\ k0\ f = Abs\text{-}poly\text{-}mapping\ (\lambda k.\ if\ k = k0 \wedge lookup\ f\ k \neq 0\ then\ lookup\ f\ k - 1\ else\ lookup\ f\ k)$

lemma *remove-key-lookup*:
  $lookup\ (decrease\text{-}key\ k0\ f)\ k = (if\ k = k0 \wedge lookup\ f\ k \neq 0\ then\ lookup\ f\ k - 1\ else\ lookup\ f\ k)$
  ⟨*proof*⟩

lemma *polynom-split-on-var*:
  fixes p :: ‹$'a$ :: {*comm-monoid-add*,*cancel-comm-monoid-add*,*semiring-0*,*comm-semiring-1*} *mpoly*›
  obtains $q\ r$ **where**
    ‹$p = monom\ (monomial\ (Suc\ 0)\ x')\ 1 * q + r$› **and**
    ‹$x' \notin vars\ r$›
⟨*proof*⟩

lemma *polynom-split-on-var2*:
  fixes p :: ‹*int mpoly*›
  assumes ‹$x' \notin vars\ s$›
  obtains $q\ r$ **where**
    ‹$p = (monom\ (monomial\ (Suc\ 0)\ x')\ 1 - s) * q + r$› **and**
    ‹$x' \notin vars\ r$›
⟨*proof*⟩

lemma *polynom-split-on-var-diff-sq2*:
 fixes p :: ‹*int mpoly*›
  obtains $q\ r\ s$ **where**
    ‹$p = monom\ (monomial\ (Suc\ 0)\ x')\ 1 * q + r + s * (monom\ (monomial\ (Suc\ 0)\ x')\ 1\hat{}2 - monom\ (monomial\ (Suc\ 0)\ x')\ 1)$› **and**
    ‹$x' \notin vars\ r$› **and**
    ‹$x' \notin vars\ q$›
⟨*proof*⟩

lemma *polynom-decomp-alien-var*:
  fixes $q\ A\ b$ :: ‹*int mpoly*›
  assumes
    q: ‹$q = A * (monom\ (monomial\ (Suc\ 0)\ x')\ 1) + b$› **and**
    x: ‹$x' \notin vars\ q$› ‹$x' \notin vars\ b$›
  shows
    ‹$A = 0$› **and**
    ‹$q = b$›
⟨*proof*⟩

lemma *polynom-decomp-alien-var2*:
  fixes $q\ A\ b$ :: ‹*int mpoly*›
  assumes
    q: ‹$q = A * (monom\ (monomial\ (Suc\ 0)\ x')\ 1 + p) + b$› **and**
    x: ‹$x' \notin vars\ q$› ‹$x' \notin vars\ b$› ‹$x' \notin vars\ p$›
  shows
    ‹$A = 0$› **and**
    ‹$q = b$›
⟨*proof*⟩

lemma *vars-unE*: ‹$x \in vars\ (a * b) \Longrightarrow (x \in vars\ a \Longrightarrow thesis) \Longrightarrow (x \in vars\ b \Longrightarrow thesis) \Longrightarrow thesis$›
   ⟨*proof*⟩

lemma *in-keys-minusI1*:

assumes $t \in keys\ p$ **and** $t \notin keys\ q$
shows $t \in keys\ (p - q)$
⟨*proof*⟩

lemma *in-keys-minusI2*:
fixes t :: ‹$'a$› **and** q :: ‹$'a \Rightarrow_0 {'b}$ :: {*cancel-comm-monoid-add*,*group-add*}›
assumes $t \in keys\ q$ **and** $t \notin keys\ p$
shows $t \in keys\ (p - q)$
⟨*proof*⟩

lemma *in-vars-addE*:
‹$x \in vars\ (p + q) \Longrightarrow (x \in vars\ p \Longrightarrow thesis) \Longrightarrow (x \in vars\ q \Longrightarrow thesis) \Longrightarrow thesis$›
⟨*proof*⟩

lemma *lookup-monomial-If*:
‹$lookup\ (monomial\ v\ k) = (\lambda k'.\ if\ k = k'\ then\ v\ else\ 0)$›
⟨*proof*⟩

lemma *vars-mult-Var*:
‹$vars\ (Var\ x * p) = (if\ p = 0\ then\ \{\}\ else\ insert\ x\ (vars\ p))$› **for** p :: ‹*int mpoly*›
⟨*proof*⟩

lemma *keys-mult-monomial*:
‹$keys\ (monomial\ (n :: int)\ k * mapping\text{-}of\ a) = (if\ n = 0\ then\ \{\}\ else\ ((+)\ k)\ `\ keys\ (mapping\text{-}of\ a))$›
⟨*proof*⟩

lemma *vars-mult-Const*:
‹$vars\ (Const\ n * a) = (if\ n = 0\ then\ \{\}\ else\ vars\ a)$› **for** a :: ‹*int mpoly*›
⟨*proof*⟩

lemma *coeff-minus*: $coeff\ p\ m - coeff\ q\ m = coeff\ (p-q)\ m$
⟨*proof*⟩

lemma *Const-1-eq-1*: ‹$Const\ (1 :: int) = (1 :: int\ mpoly)$›
⟨*proof*⟩

lemma [*simp*]:
‹$vars\ (1 :: int\ mpoly) = \{\}$›
⟨*proof*⟩

2.2 More Ideals

lemma
fixes A :: ‹$(('x \Rightarrow_0 nat) \Rightarrow_0 {'a}::comm\text{-}ring\text{-}1)\ set$›
assumes ‹$p \in ideal\ A$›
shows ‹$p * q \in ideal\ A$›
⟨*proof*⟩

The following theorem is very close to $More\text{-}Modules.ideal\ (insert\ ?a\ ?S) = \{x.\ \exists k.\ x - k * ?a \in More\text{-}Modules.ideal\ ?S\}$, except that it is more useful if we need to take an element of $More\text{-}Modules.ideal\ (insert\ a\ S)$.

lemma *ideal-insert'*:
‹$More\text{-}Modules.ideal\ (insert\ a\ S) = \{y.\ \exists x\ k.\ y = x + k * a \land x \in More\text{-}Modules.ideal\ S\}$›

⟨*proof*⟩

lemma *ideal-mult-right-in*:
‹$a \in ideal\ A \Longrightarrow a * b \in More\text{-}Modules.ideal\ A$›
⟨*proof*⟩

lemma *ideal-mult-right-in2*:
‹$a \in ideal\ A \Longrightarrow b * a \in More\text{-}Modules.ideal\ A$›
⟨*proof*⟩

lemma [*simp*]: ‹$vars\ (Var\ x :: 'a :: \{zero\text{-}neq\text{-}one\}\ mpoly) = \{x\}$›
⟨*proof*⟩

lemma *vars-minus-Var-subset*:
‹$vars\ (p' - Var\ x :: 'a :: \{ab\text{-}group\text{-}add,one,zero\text{-}neq\text{-}one\}\ mpoly) \subseteq \mathcal{V} \Longrightarrow vars\ p' \subseteq insert\ x\ \mathcal{V}$›
⟨*proof*⟩

lemma *vars-add-Var-subset*:
‹$vars\ (p' + Var\ x :: 'a :: \{ab\text{-}group\text{-}add,one,zero\text{-}neq\text{-}one\}\ mpoly) \subseteq \mathcal{V} \Longrightarrow vars\ p' \subseteq insert\ x\ \mathcal{V}$›
⟨*proof*⟩

lemma *coeff-monomila-in-varsD*:
‹$coeff\ p\ (monomial\ (Suc\ 0)\ x) \neq 0 \Longrightarrow x \in vars\ (p :: int\ mpoly)$›
⟨*proof*⟩

lemma (**in** −)*coeff-MPoly-monomila*[*simp*]:
‹$Const\ (MPoly\text{-}Type.coeff\ (MPoly\ (monomial\ a\ m))\ m) = Const\ a$›
⟨*proof*⟩

end
theory *PAC-Specification*
imports *PAC-More-Poly*
begin

3 Specification of the PAC checker

3.1 Ideals

type-synonym *int-poly* = ‹*int mpoly*›
definition *polynom-bool* :: ‹*int-poly set*› **where**
‹$polynom\text{-}bool = (\lambda c.\ Var\ c\ \hat{}\ 2 - Var\ c)\ `\ UNIV$›

definition *pac-ideal* **where**
‹$pac\text{-}ideal\ A \equiv ideal\ (A \cup polynom\text{-}bool)$›

lemma *X2-X-in-pac-ideal*:
‹$Var\ c\ \hat{}\ 2 - Var\ c \in pac\text{-}ideal\ A$›
⟨*proof*⟩

lemma *pac-idealI1*[*intro*]:
‹$p \in A \Longrightarrow p \in pac\text{-}ideal\ A$›
⟨*proof*⟩

lemma *pac-idealI2*[*intro*]:

$\langle p \in ideal\ A \Longrightarrow p \in pac\text{-}ideal\ A\rangle$
$\langle proof \rangle$

lemma *pac-idealI3* [*intro*]:
$\langle p \in ideal\ A \Longrightarrow p{*}q \in pac\text{-}ideal\ A\rangle$
$\langle proof \rangle$

lemma *pac-ideal-Xsq2-iff*:
$\langle Var\ c\ \widehat{}\ 2 \in pac\text{-}ideal\ A \longleftrightarrow Var\ c \in pac\text{-}ideal\ A\rangle$
$\langle proof \rangle$

lemma *diff-in-polynom-bool-pac-idealI*:
assumes *a1*: $p \in pac\text{-}ideal\ A$
assumes *a2*: $p - p' \in More\text{-}Modules.ideal\ polynom\text{-}bool$
shows $\langle p' \in pac\text{-}ideal\ A\rangle$
$\langle proof \rangle$

lemma *diff-in-polynom-bool-pac-idealI2*:
assumes *a1*: $p \in A$
assumes *a2*: $p - p' \in More\text{-}Modules.ideal\ polynom\text{-}bool$
shows $\langle p' \in pac\text{-}ideal\ A\rangle$
$\langle proof \rangle$

lemma *pac-ideal-alt-def*:
$\langle pac\text{-}ideal\ A = ideal\ (A \cup ideal\ polynom\text{-}bool)\rangle$
$\langle proof \rangle$

The equality on ideals is restricted to polynomials whose variable appear in the set of ideals. The function restrict sets:

definition *restricted-ideal-to* **where**
$\langle restricted\text{-}ideal\text{-}to\ B\ A = \{p \in A.\ vars\ p\ \subseteq B\}\rangle$

abbreviation $restricted\text{-}ideal\text{-}to_I$ **where**
$\langle restricted\text{-}ideal\text{-}to_I\ B\ A \equiv restricted\text{-}ideal\text{-}to\ B\ (pac\text{-}ideal\ (set\text{-}mset\ A))\rangle$

abbreviation $restricted\text{-}ideal\text{-}to_V$ **where**
$\langle restricted\text{-}ideal\text{-}to_V\ B \equiv restricted\text{-}ideal\text{-}to\ (\bigcup(vars\ `\ set\text{-}mset\ B))\rangle$

abbreviation $restricted\text{-}ideal\text{-}to_{V\,I}$ **where**
$\langle restricted\text{-}ideal\text{-}to_{V\,I}\ B\ A \equiv restricted\text{-}ideal\text{-}to\ (\bigcup(vars\ `\ set\text{-}mset\ B))\ (pac\text{-}ideal\ (set\text{-}mset\ A))\rangle$

lemma *restricted-idealI*:
$\langle p \in pac\text{-}ideal\ (set\text{-}mset\ A) \Longrightarrow vars\ p \subseteq C \Longrightarrow p \in restricted\text{-}ideal\text{-}to_I\ C\ A\rangle$
$\langle proof \rangle$

lemma *pac-ideal-insert-already-in*:
$\langle pq \in pac\text{-}ideal\ (set\text{-}mset\ A) \Longrightarrow pac\text{-}ideal\ (insert\ pq\ (set\text{-}mset\ A)) = pac\text{-}ideal\ (set\text{-}mset\ A)\rangle$
$\langle proof \rangle$

lemma *pac-ideal-add*:
$\langle p \in\#\ A \Longrightarrow q \in\#\ A \Longrightarrow p + q \in pac\text{-}ideal\ (set\text{-}mset\ A)\rangle$
$\langle proof \rangle$
lemma *pac-ideal-mult*:
$\langle p \in\#\ A \Longrightarrow p * q \in pac\text{-}ideal\ (set\text{-}mset\ A)\rangle$

⟨*proof*⟩

lemma *pac-ideal-mono*:
‹$A \subseteq B \Longrightarrow$ *pac-ideal* $A \subseteq$ *pac-ideal* B›
⟨*proof*⟩

3.2 PAC Format

The PAC format contains three kind of steps:

- add that adds up two polynomials that are known.
- mult that multiply a known polynomial with another one.
- del that removes a polynomial that cannot be reused anymore.

To model the simplification that happens, we add the $p - p' \in$ *polynom-bool* stating that p and p' are equivalent.

type-synonym *pac-st* = ‹(*nat set* × *int-poly multiset*)›

inductive *PAC-Format* :: ‹*pac-st* ⇒ *pac-st* ⇒ *bool*› **where**
add:
‹*PAC-Format* $(\mathcal{V}, A)$ $(\mathcal{V},$ *add-mset* p' $A)$›
if
‹$p \in\# A$› ‹$q \in\# A$›
‹$p+q - p' \in$ *ideal polynom-bool*›
‹*vars* $p' \subseteq \mathcal{V}$› |
mult:
‹*PAC-Format* $(\mathcal{V}, A)$ $(\mathcal{V},$ *add-mset* p' $A)$›
if
‹$p \in\# A$›
‹$p*q - p' \in$ *ideal polynom-bool*›
‹*vars* $p' \subseteq \mathcal{V}$›
‹*vars* $q \subseteq \mathcal{V}$› |
del:
‹$p \in\# A \Longrightarrow$ *PAC-Format* $(\mathcal{V}, A)$ $(\mathcal{V}, A - \{\#p\#\})$› |
extend-pos:
‹*PAC-Format* $(\mathcal{V}, A)$ $(\mathcal{V} \cup \{x' \in$ *vars* $(-$*Var* $x + p')$. $x' \notin \mathcal{V}\}$, *add-mset* $(-$*Var* $x + p')$ $A)$›
if
‹$(p')^2 - p' \in$ *ideal polynom-bool*›
‹*vars* $p' \subseteq \mathcal{V}$›
‹$x \notin \mathcal{V}$›

In the PAC format above, we have a technical condition on the normalisation: *vars* $p' \subseteq$ *vars* $(p + q)$ is here to ensure that we don't normalise *0* to $(\mathit{Var}\ x)^2 - \mathit{Var}\ x$ for a new variable x. This is completely obvious for the normalisation processe we have in mind when we write the specification, but we must add it explicitely because we are too general.

lemmas *PAC-Format-induct-split* =
PAC-Format.induct[*split-format*(*complete*), *of* V A V' A' **for** V A V' A']

lemma *PAC-Format-induct*[*consumes 1*, *case-names add mult del ext*]:
assumes
‹*PAC-Format* $(\mathcal{V}, A)$ $(\mathcal{V}', A')$› **and**
cases:

⟨$\bigwedge p\ q\ p'\ \ A\ \mathcal{V}.\ p \in\#\ A \Longrightarrow q \in\#\ A \Longrightarrow p{+}q\ -\ p' \in$ *ideal polynom-bool* $\Longrightarrow$ *vars* $p' \subseteq \mathcal{V} \Longrightarrow P\ \mathcal{V}$ $A\ \mathcal{V}$ (*add-mset* p' A)⟩
⟨$\bigwedge p\ q\ p'\ A\ \mathcal{V}.\ p \in\#\ A \Longrightarrow p{*}q\ -\ p' \in$ *ideal polynom-bool* $\Longrightarrow$ *vars* $p' \subseteq \mathcal{V} \Longrightarrow$ *vars* $q \subseteq \mathcal{V} \Longrightarrow$ $P\ \mathcal{V}\ A\ \mathcal{V}$ (*add-mset* p' A)⟩
⟨$\bigwedge p\ A\ \mathcal{V}.\ p \in\#\ A \Longrightarrow P\ \mathcal{V}\ A\ \mathcal{V}\ (A - \{\#p\#\})$⟩
⟨$\bigwedge p'\ x\ r.$
$(p')\hat{}2\ -\ (p') \in$ *ideal polynom-bool* $\Longrightarrow$ *vars* $p' \subseteq \mathcal{V} \Longrightarrow$
$x \notin \mathcal{V} \Longrightarrow P\ \mathcal{V}\ A\ (\mathcal{V} \cup \{x' \in$ *vars* $(p' -\ Var\ x).\ x' \notin \mathcal{V}\})$ (*add-mset* $(p' - Var\ x)$ A)⟩
shows
⟨$P\ \mathcal{V}\ A\ \mathcal{V}'\ A'$⟩
⟨*proof*⟩

The theorem below (based on the proof ideal by Manuel Kauers) is the correctness theorem of extensions. Remark that the assumption *vars* $q \subseteq \mathcal{V}$ is only used to show that $x' \notin$ *vars* q.

lemma *extensions-are-safe*:
assumes ⟨$x' \in$ *vars* p⟩ **and**
x': ⟨$x' \notin \mathcal{V}$⟩ **and**
⟨$\bigcup$ (*vars* ‘ *set-mset* A) $\subseteq \mathcal{V}$⟩ **and**
p-x-coeff: ⟨*coeff* p (*monomial* (*Suc 0*) x') = *1*⟩ **and**
vars-q: ⟨*vars* $q \subseteq \mathcal{V}$⟩ **and**
q: ⟨$q \in$ *More-Modules.ideal* (*insert* p (*set-mset* $A \cup$ *polynom-bool*))⟩ **and**
leading: ⟨$x' \notin$ *vars* $(p -\ Var\ x')$⟩ **and**
diff: ⟨$(Var\ x' - p)^2 - (Var\ x' - p) \in$ *More-Modules.ideal polynom-bool*⟩
shows
⟨$q \in$ *More-Modules.ideal* (*set-mset* $A \cup$ *polynom-bool*)⟩
⟨*proof*⟩

lemma *extensions-are-safe-uminus*:
assumes ⟨$x' \in$ *vars* p⟩ **and**
x': ⟨$x' \notin \mathcal{V}$⟩ **and**
⟨$\bigcup$ (*vars* ‘ *set-mset* A) $\subseteq \mathcal{V}$⟩ **and**
p-x-coeff: ⟨*coeff* p (*monomial* (*Suc 0*) x') = -1⟩ **and**
vars-q: ⟨*vars* $q \subseteq \mathcal{V}$⟩ **and**
q: ⟨$q \in$ *More-Modules.ideal* (*insert* p (*set-mset* $A \cup$ *polynom-bool*))⟩ **and**
leading: ⟨$x' \notin$ *vars* $(p +\ Var\ x')$⟩ **and**
diff: ⟨$(Var\ x' + p)\hat{}2 - (Var\ x' + p) \in$ *More-Modules.ideal polynom-bool*⟩
shows
⟨$q \in$ *More-Modules.ideal* (*set-mset* $A \cup$ *polynom-bool*)⟩
⟨*proof*⟩

This is the correctness theorem of a PAC step: no polynomials are added to the ideal.

lemma *vars-subst-in-left-only*:
⟨$x \notin$ *vars* $p \Longrightarrow x \in$ *vars* $(p -\ Var\ x)$⟩ **for** p :: ⟨*int mpoly*⟩
⟨*proof*⟩

lemma *vars-subst-in-left-only-diff-iff*:
⟨$x \notin$ *vars* $p \Longrightarrow$ *vars* $(p -\ Var\ x)$ = *insert* x (*vars* p)⟩ **for** p :: ⟨*int mpoly*⟩
⟨*proof*⟩

lemma *vars-subst-in-left-only-iff*:
⟨$x \notin$ *vars* $p \Longrightarrow$ *vars* $(p +\ Var\ x)$ = *insert* x (*vars* p)⟩ **for** p :: ⟨*int mpoly*⟩
⟨*proof*⟩

lemma *coeff-add-right-notin*:
⟨$x \notin$ *vars* $p \Longrightarrow$ *MPoly-Type.coeff* $(Var\ x - p)$ (*monomial* (*Suc 0*) x) = *1*⟩

$\langle proof \rangle$

lemma *coeff-add-left-notin*:
‹$x \notin vars\ p \Longrightarrow MPoly\text{-}Type.coeff\ (p - Var\ x)\ (monomial\ (Suc\ 0)\ x) = -1$› **for** p :: ‹*int mpoly*›
$\langle proof \rangle$

lemma *ideal-insert-polynom-bool-swap*: ‹$r - s \in ideal\ polynom\text{-}bool \Longrightarrow$
$More\text{-}Modules.ideal\ (insert\ r\ (A \cup polynom\text{-}bool)) = More\text{-}Modules.ideal\ (insert\ s\ (A \cup polynom\text{-}bool))$›
$\langle proof \rangle$

lemma *PAC-Format-subset-ideal*:
‹$PAC\text{-}Format\ (\mathcal{V}, A)\ (\mathcal{V}', B) \Longrightarrow \bigcup(vars\ `\ set\text{-}mset\ A) \subseteq \mathcal{V} \Longrightarrow$
$restricted\text{-}ideal\text{-}to_I\ \mathcal{V}\ B \subseteq restricted\text{-}ideal\text{-}to_I\ \mathcal{V}\ A \wedge \mathcal{V} \subseteq \mathcal{V}' \wedge \bigcup(vars\ `\ set\text{-}mset\ B) \subseteq \mathcal{V}'$›
$\langle proof \rangle$

In general, if deletions are disallowed, then the stronger $B = pac\text{-}ideal\ A$ holds.

lemma *restricted-ideal-to-restricted-ideal-to*$_I$*D*:
‹$restricted\text{-}ideal\text{-}to\ \mathcal{V}\ (set\text{-}mset\ A) \subseteq restricted\text{-}ideal\text{-}to_I\ \mathcal{V}\ A$›
$\langle proof \rangle$

lemma *rtranclp-PAC-Format-subset-ideal*:
‹$rtranclp\ PAC\text{-}Format\ (\mathcal{V}, A)\ (\mathcal{V}', B) \Longrightarrow \bigcup(vars\ `\ set\text{-}mset\ A) \subseteq \mathcal{V} \Longrightarrow$
$restricted\text{-}ideal\text{-}to_I\ \mathcal{V}\ B \subseteq restricted\text{-}ideal\text{-}to_I\ \mathcal{V}\ A \wedge \mathcal{V} \subseteq \mathcal{V}' \wedge \bigcup(vars\ `\ set\text{-}mset\ B) \subseteq \mathcal{V}'$›
$\langle proof \rangle$

end
theory *Finite-Map-Multiset*
imports *HOL−Library.Finite-Map Duplicate-Free-Multiset*
begin

notation *image-mset* (**infixr** *‘# 90*)

4 Finite maps and multisets

4.1 Finite sets and multisets

abbreviation *mset-fset* :: ‹$'a\ fset \Rightarrow 'a\ multiset$› **where**
‹$mset\text{-}fset\ N \equiv mset\text{-}set\ (fset\ N)$›

definition *fset-mset* :: ‹$'a\ multiset \Rightarrow 'a\ fset$› **where**
‹$fset\text{-}mset\ N \equiv Abs\text{-}fset\ (set\text{-}mset\ N)$›

lemma *fset-mset-mset-fset*: ‹$fset\text{-}mset\ (mset\text{-}fset\ N) = N$›
$\langle proof \rangle$

lemma *mset-fset-fset-mset*[*simp*]:
‹$mset\text{-}fset\ (fset\text{-}mset\ N) = remdups\text{-}mset\ N$›
$\langle proof \rangle$

lemma *in-mset-fset-fmember*[*simp*]: ‹$x \in\#\ mset\text{-}fset\ N \longleftrightarrow x\ |\in|\ N$›
$\langle proof \rangle$

lemma *in-fset-mset-mset*[*simp*]: ‹$x\ |\in|\ fset\text{-}mset\ N \longleftrightarrow x \in\#\ N$›

⟨*proof*⟩

4.2 Finite map and multisets

Roughly the same as *ran* and *dom*, but with duplication in the content (unlike their finite sets counterpart) while still working on finite domains (unlike a function mapping). Remark that *dom-m* (the keys) does not contain duplicates, but we keep for symmetry (and for easier use of multiset operators as in the definition of *ran-m*).

definition *dom-m* **where**
‹$dom\text{-}m\ N = mset\text{-}fset\ (fmdom\ N)$›

definition *ran-m* **where**
‹$ran\text{-}m\ N = the\ \text{`}\#\ fmlookup\ N\ \text{`}\#\ dom\text{-}m\ N$›

lemma *dom-m-fmdrop*[*simp*]: ‹$dom\text{-}m\ (fmdrop\ C\ N) = remove1\text{-}mset\ C\ (dom\text{-}m\ N)$›
⟨*proof*⟩

lemma *dom-m-fmdrop-All*: ‹$dom\text{-}m\ (fmdrop\ C\ N) = removeAll\text{-}mset\ C\ (dom\text{-}m\ N)$›
⟨*proof*⟩

lemma *dom-m-fmupd*[*simp*]: ‹$dom\text{-}m\ (fmupd\ k\ C\ N) = add\text{-}mset\ k\ (remove1\text{-}mset\ k\ (dom\text{-}m\ N))$›
⟨*proof*⟩

lemma *distinct-mset-dom*: ‹$distinct\text{-}mset\ (dom\text{-}m\ N)$›
⟨*proof*⟩

lemma *in-dom-m-lookup-iff*: ‹$C \in\#\ dom\text{-}m\ N' \longleftrightarrow fmlookup\ N'\ C \neq None$›
⟨*proof*⟩

lemma *in-dom-in-ran-m*[*simp*]: ‹$i \in\#\ dom\text{-}m\ N \Longrightarrow the\ (fmlookup\ N\ i) \in\#\ ran\text{-}m\ N$›
⟨*proof*⟩

lemma *fmupd-same*[*simp*]:
‹$x1 \in\#\ dom\text{-}m\ x1aa \Longrightarrow fmupd\ x1\ (the\ (fmlookup\ x1aa\ x1))\ x1aa = x1aa$›
⟨*proof*⟩

lemma *ran-m-fmempty*[*simp*]: ‹$ran\text{-}m\ fmempty = \{\#\}$› **and**
dom-m-fmempty[*simp*]: ‹$dom\text{-}m\ fmempty = \{\#\}$›
⟨*proof*⟩

lemma *fmrestrict-set-fmupd*:
‹$a \in xs \Longrightarrow fmrestrict\text{-}set\ xs\ (fmupd\ a\ C\ N) = fmupd\ a\ C\ (fmrestrict\text{-}set\ xs\ N)$›
‹$a \notin xs \Longrightarrow fmrestrict\text{-}set\ xs\ (fmupd\ a\ C\ N) = fmrestrict\text{-}set\ xs\ N$›
⟨*proof*⟩

lemma *fset-fmdom-fmrestrict-set*:
‹$fset\ (fmdom\ (fmrestrict\text{-}set\ xs\ N)) = fset\ (fmdom\ N) \cap xs$›
⟨*proof*⟩

lemma *dom-m-fmrestrict-set*: ‹$dom\text{-}m\ (fmrestrict\text{-}set\ (set\ xs)\ N) = mset\ xs \cap\#\ dom\text{-}m\ N$›
⟨*proof*⟩

lemma *dom-m-fmrestrict-set'*: ‹$dom\text{-}m\ (fmrestrict\text{-}set\ xs\ N) = mset\text{-}set\ (xs \cap set\text{-}mset\ (dom\text{-}m\ N))$›
⟨*proof*⟩

lemma *indom-mI*: ‹$fmlookup\ m\ x = Some\ y \Longrightarrow x \in\#\ dom\text{-}m\ m$›
⟨*proof*⟩

lemma *fmupd-fmdrop-id*:
assumes ‹$k\ |\in|\ fmdom\ N'$›
shows ‹$fmupd\ k\ (the\ (fmlookup\ N'\ k))\ (fmdrop\ k\ N') = N'$›
⟨*proof*⟩

lemma *fm-member-split*: ‹$k\ |\in|\ fmdom\ N' \Longrightarrow \exists N''\ v.\ N' = fmupd\ k\ v\ N'' \wedge the\ (fmlookup\ N'\ k) = v$
$\wedge$
$k\ |\notin|\ fmdom\ N''$›
⟨*proof*⟩

lemma ‹$fmdrop\ k\ (fmupd\ k\ va\ N'') = fmdrop\ k\ N''$›
⟨*proof*⟩

lemma *fmap-ext-fmdom*:
‹$(fmdom\ N = fmdom\ N') \Longrightarrow (\bigwedge x.\ x\ |\in|\ fmdom\ N \Longrightarrow fmlookup\ N\ x = fmlookup\ N'\ x) \Longrightarrow$
$N = N'$›
⟨*proof*⟩

lemma *fmrestrict-set-insert-in*:
‹$xa\ \in fset\ (fmdom\ N) \Longrightarrow$
$fmrestrict\text{-}set\ (insert\ xa\ l1)\ N = fmupd\ xa\ (the\ (fmlookup\ N\ xa))\ (fmrestrict\text{-}set\ l1\ N)$›
⟨*proof*⟩

lemma *fmrestrict-set-insert-notin*:
‹$xa\ \notin fset\ (fmdom\ N) \Longrightarrow$
$fmrestrict\text{-}set\ (insert\ xa\ l1)\ N = fmrestrict\text{-}set\ l1\ N$›
⟨*proof*⟩

lemma *fmrestrict-set-insert-in-dom-m*[*simp*]:
‹$xa\ \in\#\ dom\text{-}m\ N \Longrightarrow$
$fmrestrict\text{-}set\ (insert\ xa\ l1)\ N = fmupd\ xa\ (the\ (fmlookup\ N\ xa))\ (fmrestrict\text{-}set\ l1\ N)$›
⟨*proof*⟩

lemma *fmrestrict-set-insert-notin-dom-m*[*simp*]:
‹$xa\ \notin\#\ dom\text{-}m\ N \Longrightarrow$
$fmrestrict\text{-}set\ (insert\ xa\ l1)\ N = fmrestrict\text{-}set\ l1\ N$›
⟨*proof*⟩

lemma *fmlookup-restrict-set-id*: ‹$fset\ (fmdom\ N) \subseteq A \Longrightarrow fmrestrict\text{-}set\ A\ N = N$›
⟨*proof*⟩

lemma *fmlookup-restrict-set-id'*: ‹$set\text{-}mset\ (dom\text{-}m\ N) \subseteq A \Longrightarrow fmrestrict\text{-}set\ A\ N = N$›
⟨*proof*⟩

lemma *ran-m-mapsto-upd*:
assumes
NC: ‹$C \in\#\ dom\text{-}m\ N$›
shows ‹$ran\text{-}m\ (fmupd\ C\ C'\ N) =$
$add\text{-}mset\ C'\ (remove1\text{-}mset\ (the\ (fmlookup\ N\ C))\ (ran\text{-}m\ N))$›
⟨*proof*⟩

lemma *ran-m-mapsto-upd-notin*:

assumes *NC*: ‹$C \notin\# \ dom\text{-}m\ N$›
shows ‹$ran\text{-}m\ (fmupd\ C\ C'\ N) = add\text{-}mset\ C'\ (ran\text{-}m\ N)$›
⟨*proof*⟩

lemma *image-mset-If-eq-notin*:
‹$C \notin\# \ A \Longrightarrow \{\#f\ (if\ x = C\ then\ a\ x\ else\ b\ x).\ x \in\#\ A\#\} = \{\#\ f(b\ x).\ x \in\#\ A\ \#\}$›
⟨*proof*⟩

lemma *filter-mset-cong2*:
$(\bigwedge x.\ x \in\#\ M \Longrightarrow f\ x = g\ x) \Longrightarrow M = N \Longrightarrow filter\text{-}mset\ f\ M = filter\text{-}mset\ g\ N$
⟨*proof*⟩

lemma *ran-m-fmdrop*:
‹$C \in\#\ dom\text{-}m\ N \Longrightarrow\ ran\text{-}m\ (fmdrop\ C\ N) = remove1\text{-}mset\ (the\ (fmlookup\ N\ C))\ (ran\text{-}m\ N)$›
⟨*proof*⟩

lemma *ran-m-fmdrop-notin*:
‹$C \notin\#\ dom\text{-}m\ N \Longrightarrow ran\text{-}m\ (fmdrop\ C\ N) = ran\text{-}m\ N$›
⟨*proof*⟩

lemma *ran-m-fmdrop-If*:
‹$ran\text{-}m\ (fmdrop\ C\ N) = (if\ C \in\#\ dom\text{-}m\ N\ then\ remove1\text{-}mset\ (the\ (fmlookup\ N\ C))\ (ran\text{-}m\ N)\ else\ ran\text{-}m\ N)$›
⟨*proof*⟩

lemma *dom-m-empty-iff* [*iff*]:
‹$dom\text{-}m\ NU = \{\#\} \longleftrightarrow NU = fmempty$›
⟨*proof*⟩

end
theory *PAC-Map-Rel*
imports
Refine-Imperative-HOL.IICF Finite-Map-Multiset
begin

5 Hash-Map for finite mappings

This function declares hash-maps for $('a,\ 'b)\ fmap$, that are nicer to use especially here where everything is finite.

definition *fmap-rel* **where**
[*to-relAPP*]:
$fmap\text{-}rel\ K\ V \equiv \{(m1,\ m2).$
$(\forall\, i\ j.\ i\ |{\in}|\ fmdom\ m2 \longrightarrow (j,\ i) \in K \longrightarrow (the\ (fmlookup\ m1\ j),\ the\ (fmlookup\ m2\ i)) \in V) \wedge$
$fset\ (fmdom\ m1) \subseteq Domain\ K \wedge fset\ (fmdom\ m2) \subseteq Range\ K \wedge$
$(\forall\, i\ j.\ (i,\ j) \in K \longrightarrow j\ |{\in}|\ fmdom\ m2 \longleftrightarrow i\ |{\in}|\ fmdom\ m1)\}$

lemma *fmap-rel-alt-def*:
‹$\langle K,\ V\rangle fmap\text{-}rel \equiv$
$\{(m1,\ m2).$
$(\forall\, i\ j.\ i \in\#\ dom\text{-}m\ m2 \longrightarrow$
$(j,\ i) \in K \longrightarrow (the\ (fmlookup\ m1\ j),\ the\ (fmlookup\ m2\ i)) \in V) \wedge$
$fset\ (fmdom\ m1) \subseteq Domain\ K \wedge$
$fset\ (fmdom\ m2) \subseteq Range\ K \wedge$

$(\forall\, i\ j.\ (i,\ j) \in K \longrightarrow (j \in\#\ \textit{dom-m m2}) = (i \in\#\ \textit{dom-m m1}))\}$
›
⟨*proof*⟩

lemma *fmap-rel-empty1-simp*[*simp*]:
(*fmempty*,*m*)∈⟨*K*,*V*⟩*fmap-rel* ⟷ *m*=*fmempty*
⟨*proof*⟩

lemma *fmap-rel-empty2-simp*[*simp*]:
(*m*,*fmempty*)∈⟨*K*,*V*⟩*fmap-rel* ⟷ *m*=*fmempty*
⟨*proof*⟩

sepref-decl-intf ('*k*,'*v*) *f-map* **is** ('*k*, '*v*) *fmap*

lemma [*synth-rules*]: ⟦*INTF-OF-REL K TYPE*('*k*); *INTF-OF-REL V TYPE*('*v*)⟧
⟹ *INTF-OF-REL* (⟨*K*,*V*⟩*fmap-rel*) *TYPE*(('*k*,'*v*) *f-map*) ⟨*proof*⟩

5.1 Operations

sepref-decl-op *fmap-empty*: *fmempty* :: ⟨*K*,*V*⟩*fmap-rel* ⟨*proof*⟩

sepref-decl-op *fmap-is-empty*: (=) *fmempty* :: ⟨*K*,*V*⟩*fmap-rel* → *bool-rel*
⟨*proof*⟩

lemma *fmap-rel-fmupd-fmap-rel*:
‹$(A,\ B) \in \langle K,\ R\rangle$*fmap-rel* ⟹ $(p,\ p') \in K \Longrightarrow (q,\ q') \in R \Longrightarrow$
(*fmupd* $p\ q\ A$, *fmupd* $p'\ q'\ B$) ∈ ⟨*K*, *R*⟩*fmap-rel*›
if *single-valued K single-valued* (K^{-1})
⟨*proof*⟩

sepref-decl-op *fmap-update*: *fmupd* :: *K* → *V* → ⟨*K*,*V*⟩*fmap-rel* → ⟨*K*,*V*⟩*fmap-rel*
where *single-valued K single-valued* (K^{-1})
⟨*proof*⟩

lemma *fmap-rel-fmdrop-fmap-rel*:
‹$(A,\ B) \in \langle K,\ R\rangle$*fmap-rel* ⟹ $(p,\ p') \in K \Longrightarrow$
(*fmdrop* $p\ A$, *fmdrop* $p'\ B$) ∈ ⟨*K*, *R*⟩*fmap-rel*›
if *single-valued K single-valued* (K^{-1})
⟨*proof*⟩

sepref-decl-op *fmap-delete*: *fmdrop* :: *K* → ⟨*K*,*V*⟩*fmap-rel* → ⟨*K*,*V*⟩*fmap-rel*
where *single-valued K single-valued* (K^{-1})
⟨*proof*⟩

lemma *fmap-rel-nat-the-fmlookup*[*intro*]:
‹$(A,\ B) \in \langle S,\ R\rangle$*fmap-rel* ⟹ $(p,\ p') \in S \Longrightarrow p' \in\#$ *dom-m B* ⟹
(*the* (*fmlookup* $A\ p$), *the* (*fmlookup* $B\ p'$)) ∈ *R*›
⟨*proof*⟩

lemma *fmap-rel-in-dom-iff*:
‹$(aa,\ a'a) \in \langle K,\ V\rangle$*fmap-rel* ⟹
$(a,\ a') \in K \Longrightarrow$
$a' \in\#$ *dom-m* $a'a$ ⟷

$a \in\# \ dom\text{-}m \ aa$›
⟨*proof*⟩

lemma *fmap-rel-fmlookup-rel*:
‹$(a, a') \in K \Longrightarrow (aa, a'a) \in \langle K, V\rangle fmap\text{-}rel \Longrightarrow$
$(fmlookup \ aa \ a, fmlookup \ a'a \ a') \in \langle V\rangle option\text{-}rel$›
⟨*proof*⟩

sepref-decl-op *fmap-lookup*: $fmlookup :: \langle K,V\rangle fmap\text{-}rel \to K \to \langle V\rangle option\text{-}rel$
⟨*proof*⟩

lemma *in-fdom-alt*: $k{\in}\#dom\text{-}m \ m \longleftrightarrow \neg is\text{-}None \ (fmlookup \ m \ k)$
⟨*proof*⟩

sepref-decl-op *fmap-contains-key*: $\lambda k \ m. \ k{\in}\#dom\text{-}m \ m :: K \to \langle K,V\rangle fmap\text{-}rel \to bool\text{-}rel$
⟨*proof*⟩

5.2 Patterns

lemma *pat-fmap-empty*[*pat-rules*]: $fmempty \equiv op\text{-}fmap\text{-}empty$ ⟨*proof*⟩

lemma *pat-map-is-empty*[*pat-rules*]:
$(=) \ \$m\$fmempty \equiv op\text{-}fmap\text{-}is\text{-}empty\$m$
$(=) \ \$fmempty\$m \equiv op\text{-}fmap\text{-}is\text{-}empty\$m$
$(=) \ \$(dom\text{-}m\$m)\$\{\#\} \equiv op\text{-}fmap\text{-}is\text{-}empty\$m$
$(=) \ \$\{\#\}\$(dom\text{-}m\$m) \equiv op\text{-}fmap\text{-}is\text{-}empty\$m$
⟨*proof*⟩

lemma *op-map-contains-key*[*pat-rules*]:
$(\in\#) \ \$ \ k \ \$ \ (dom\text{-}m\$m) \equiv op\text{-}fmap\text{-}contains\text{-}key\$'k\$'m$
⟨*proof*⟩

5.3 Mapping to Normal Hashmaps

abbreviation *map-of-fmap* :: ‹$('k \Rightarrow 'v \ option) \Rightarrow ('k, 'v) \ fmap$› **where**
‹$map\text{-}of\text{-}fmap \ h \equiv Abs\text{-}fmap \ h$›

definition *map-fmap-rel* **where**
‹$map\text{-}fmap\text{-}rel = br \ map\text{-}of\text{-}fmap \ (\lambda a. \ finite \ (dom \ a))$›

lemma *fmdrop-set-None*:
‹$(op\text{-}map\text{-}delete, fmdrop) \in Id \to map\text{-}fmap\text{-}rel \to map\text{-}fmap\text{-}rel$›
⟨*proof*⟩

lemma *map-upd-fmupd*:
‹$(op\text{-}map\text{-}update, fmupd) \in Id \to Id \to map\text{-}fmap\text{-}rel \to map\text{-}fmap\text{-}rel$›
⟨*proof*⟩

Technically *op-map-lookup* has the arguments in the wrong direction.

definition *fmlookup'* **where**
[*simp*]: ‹$fmlookup' \ A \ k = fmlookup \ k \ A$›

lemma [*def-pat-rules*]:
‹$((\in\#)\$k\$(dom\text{-}m\$A)) \equiv Not\$(is\text{-}None\$(fmlookup'\$k\$A))$›

⟨*proof*⟩

lemma *op-map-lookup-fmlookup*:
‹$(op\text{-}map\text{-}lookup,\ fmlookup') \in Id \rightarrow map\text{-}fmap\text{-}rel \rightarrow \langle Id\rangle option\text{-}rel$›
⟨*proof*⟩

abbreviation *hm-fmap-assn* **where**
‹$hm\text{-}fmap\text{-}assn\ K\ V \equiv hr\text{-}comp\ (hm.assn\ K\ V)\ map\text{-}fmap\text{-}rel$›

lemmas *fmap-delete-hnr* [*sepref-fr-rules*] =
hm.delete-hnr[*FCOMP fmdrop-set-None*]

lemmas *fmap-update-hnr* [*sepref-fr-rules*] =
hm.update-hnr[*FCOMP map-upd-fmupd*]

lemmas *fmap-lookup-hnr* [*sepref-fr-rules*] =
hm.lookup-hnr[*FCOMP op-map-lookup-fmlookup*]

lemma *fmempty-empty*:
‹$(uncurry0\ (RETURN\ op\text{-}map\text{-}empty),\ uncurry0\ (RETURN\ fmempty)) \in unit\text{-}rel \rightarrow_f \langle map\text{-}fmap\text{-}rel\rangle nres\text{-}rel$›
⟨*proof*⟩

lemmas [*sepref-fr-rules*] =
hm.empty-hnr[*FCOMP fmempty-empty*, *unfolded op-fmap-empty-def*[*symmetric*]]

abbreviation *iam-fmap-assn* **where**
‹$iam\text{-}fmap\text{-}assn\ K\ V \equiv hr\text{-}comp\ (iam.assn\ K\ V)\ map\text{-}fmap\text{-}rel$›

lemmas *iam-fmap-delete-hnr* [*sepref-fr-rules*] =
iam.delete-hnr[*FCOMP fmdrop-set-None*]

lemmas *iam-ffmap-update-hnr* [*sepref-fr-rules*] =
iam.update-hnr[*FCOMP map-upd-fmupd*]

lemmas *iam-ffmap-lookup-hnr* [*sepref-fr-rules*] =
iam.lookup-hnr[*FCOMP op-map-lookup-fmlookup*]

definition *op-iam-fmap-empty* **where**
‹$op\text{-}iam\text{-}fmap\text{-}empty = fmempty$›

lemma *iam-fmempty-empty*:
‹$(uncurry0\ (RETURN\ op\text{-}map\text{-}empty),\ uncurry0\ (RETURN\ op\text{-}iam\text{-}fmap\text{-}empty)) \in unit\text{-}rel \rightarrow_f \langle map\text{-}fmap\text{-}rel\rangle nres\text{-}rel$›
⟨*proof*⟩

lemmas [*sepref-fr-rules*] =
iam.empty-hnr[*FCOMP fmempty-empty*, *unfolded op-iam-fmap-empty-def*[*symmetric*]]

definition *upper-bound-on-dom* **where**
‹$upper\text{-}bound\text{-}on\text{-}dom\ A = SPEC(\lambda n.\ \forall\, i \in\#(dom\text{-}m\ A).\ i < n)$›

lemma [*sepref-fr-rules*]:
‹((*Array.len*), *upper-bound-on-dom*) $\in$ (*iam-fmap-assn nat-assn V*)^k $\rightarrow_a$ *nat-assn*›
⟨*proof*⟩

lemma *fmap-rel-nat-rel-dom-m*[*simp*]:
‹(A, B) $\in$ ⟨*nat-rel*, R⟩*fmap-rel* $\Longrightarrow$ *dom-m* A = *dom-m* B›
⟨*proof*⟩

lemma *ref-two-step'*:
‹$A \leq B \Longrightarrow \Downarrow R\ A \leq \Downarrow R\ B$›
⟨*proof*⟩

end
theory *PAC-Checker-Specification*
imports *PAC-Specification*
Refine-Imperative-HOL.IICF
Finite-Map-Multiset
begin

6 Checker Algorithm

In this level of refinement, we define the first level of the implementation of the checker, both with the specification as on ideals and the first version of the loop.

6.1 Specification

datatype *status* =
is-failed: *FAILED* |
is-success: *SUCCESS* |
is-found: *FOUND*

lemma *is-success-alt-def*:
‹*is-success a* $\longleftrightarrow$ *a* = *SUCCESS*›
⟨*proof*⟩

datatype (*'a*, *'b*, *'lbls*) *pac-step* =
Add (*pac-src1*: *'lbls*) (*pac-src2*: *'lbls*) (*new-id*: *'lbls*) (*pac-res*: *'a*) |
Mult (*pac-src1*: *'lbls*) (*pac-mult*: *'a*) (*new-id*: *'lbls*) (*pac-res*: *'a*) |
Extension (*new-id*: *'lbls*) (*new-var*: *'b*) (*pac-res*: *'a*) |
Del (*pac-src1*: *'lbls*)

type-synonym *pac-state* = ‹(*nat set* $\times$ *int-poly multiset*)›

definition *PAC-checker-specification*
:: ‹*int-poly* $\Rightarrow$ *int-poly multiset* $\Rightarrow$ (*status* $\times$ *nat set* $\times$ *int-poly multiset*) *nres*›
where
‹*PAC-checker-specification spec A* = *SPEC*(λ(*b*, $\mathcal{V}$, *B*).
($\neg$*is-failed b* $\longrightarrow$ *restricted-ideal-to*$_I$ ($\bigcup$(*vars* ‘ *set-mset A*) $\cup$ *vars spec*) *B* $\subseteq$ *restricted-ideal-to*$_I$ ($\bigcup$(*vars* ‘ *set-mset A*) $\cup$ *vars spec*) *A*) $\wedge$
(*is-found b* $\longrightarrow$ *spec* $\in$ *pac-ideal* (*set-mset A*)))›

definition *PAC-checker-specification-spec*

:: ‹*int-poly* ⇒ *pac-state* ⇒ (*status* × *pac-state*) ⇒ *bool*›
where
‹*PAC-checker-specification-spec spec* = (λ($\mathcal{V}$, *A*) (*b*, *B*). ($\neg$*is-failed b* ⟶ $\bigcup$(*vars* ‘ *set-mset A*) ⊆ $\mathcal{V}$) ∧
(*is-success b* ⟶ *PAC-Format*** ($\mathcal{V}$, *A*) *B*) ∧
(*is-found b* ⟶ *PAC-Format*** ($\mathcal{V}$, *A*) *B* ∧ *spec* ∈ *pac-ideal* (*set-mset A*)))›

abbreviation *PAC-checker-specification2*
:: ‹*int-poly* ⇒ (*nat set* × *int-poly multiset*) ⇒ (*status* × (*nat set* × *int-poly multiset*)) *nres*›
where
‹*PAC-checker-specification2 spec A* ≡ *SPEC*(*PAC-checker-specification-spec spec A*)›

definition *PAC-checker-specification-step-spec*
:: ‹*pac-state* ⇒ *int-poly* ⇒ *pac-state* ⇒ (*status* × *pac-state*) ⇒ *bool*›
where
‹*PAC-checker-specification-step-spec* = (λ($\mathcal{V}_0$, A_0) *spec* ($\mathcal{V}$, *A*) (*b*, *B*).
(*is-success b* ⟶
$\bigcup$(*vars* ‘ *set-mset* A_0) ⊆ $\mathcal{V}_0$ ∧
$\bigcup$(*vars* ‘ *set-mset A*) ⊆ $\mathcal{V}$ ∧ *PAC-Format*** ($\mathcal{V}_0$, A_0) ($\mathcal{V}$, *A*) ∧ *PAC-Format*** ($\mathcal{V}$, *A*) *B*) ∧
(*is-found b* ⟶
$\bigcup$(*vars* ‘ *set-mset* A_0) ⊆ $\mathcal{V}_0$ ∧
$\bigcup$(*vars* ‘ *set-mset A*) ⊆ $\mathcal{V}$ ∧ *PAC-Format*** ($\mathcal{V}_0$, A_0) ($\mathcal{V}$, *A*) ∧ *PAC-Format*** ($\mathcal{V}$, *A*) *B* ∧
spec ∈ *pac-ideal* (*set-mset* A_0)))›

abbreviation *PAC-checker-specification-step2*
:: ‹*pac-state* ⇒ *int-poly* ⇒ *pac-state* ⇒ (*status* × *pac-state*) *nres*›
where
‹*PAC-checker-specification-step2* A_0 *spec A* ≡ *SPEC*(*PAC-checker-specification-step-spec* A_0 *spec A*)›

definition *normalize-poly-spec* :: ‹-› **where**
‹*normalize-poly-spec p* = *SPEC* (λ*r*. *p* − *r* ∈ *ideal polynom-bool* ∧ *vars r* ⊆ *vars p*)›

lemma *normalize-poly-spec-alt-def*:
‹*normalize-poly-spec p* = *SPEC* (λ*r*. *r* − *p* ∈ *ideal polynom-bool* ∧ *vars r* ⊆ *vars p*)›
⟨*proof*⟩

definition *mult-poly-spec* :: ‹*int mpoly* ⇒ *int mpoly* ⇒ *int mpoly nres*› **where**
‹*mult-poly-spec p q* = *SPEC* (λ*r*. *p* ∗ *q* − *r* ∈ *ideal polynom-bool*)›

definition *check-add* :: ‹(*nat*, *int mpoly*) *fmap* ⇒ *nat set* ⇒ *nat* ⇒ *nat* ⇒ *nat* ⇒ *int mpoly* ⇒ *bool nres*› **where**
‹*check-add A* $\mathcal{V}$ *p q i r* =
SPEC(λ*b*. *b* ⟶ *p* ∈# *dom-m A* ∧ *q* ∈# *dom-m A* ∧ *i* ∉# *dom-m A* ∧ *vars r* ⊆ $\mathcal{V}$ ∧
the (*fmlookup A p*) + *the* (*fmlookup A q*) − *r* ∈ *ideal polynom-bool*)›

definition *check-mult* :: ‹(*nat*, *int mpoly*) *fmap* ⇒ *nat set* ⇒ *nat* ⇒ *int mpoly* ⇒ *nat* ⇒ *int mpoly* ⇒ *bool nres*› **where**
‹*check-mult A* $\mathcal{V}$ *p q i r* =
SPEC(λ*b*. *b* ⟶ *p* ∈# *dom-m A* ∧*i* ∉# *dom-m A* ∧ *vars q* ⊆ $\mathcal{V}$ ∧ *vars r* ⊆ $\mathcal{V}$ ∧
the (*fmlookup A p*) ∗ *q* − *r* ∈ *ideal polynom-bool*)›

definition *check-extension* :: ‹(*nat*, *int mpoly*) *fmap* ⇒ *nat set* ⇒ *nat* ⇒ *nat* ⇒ *int mpoly* ⇒ (*bool*) *nres*› **where**
‹*check-extension A* $\mathcal{V}$ *i v p* =

$SPEC(\lambda b.\ b \longrightarrow (i \notin\#\ dom\text{-}m\ A \wedge$
$(v \notin \mathcal{V} \wedge$
$(p+ Var\ v)^2 - (p+ Var\ v) \in ideal\ polynom\text{-}bool \wedge$
$vars\ (p+ Var\ v) \subseteq \mathcal{V})))$›

fun *merge-status* **where**
‹*merge-status* (*FAILED*) - = *FAILED*› |
‹*merge-status* - (*FAILED*) = *FAILED*› |
‹*merge-status FOUND* - = *FOUND*› |
‹*merge-status* - *FOUND* = *FOUND*› |
‹*merge-status* - - = *SUCCESS*›

type-synonym *fpac-step* = ‹*nat set* × (*nat*, *int-poly*) *fmap*›

definition *check-del* :: ‹(*nat*, *int mpoly*) *fmap* ⇒ *nat* ⇒ *bool nres*› **where**
‹*check-del A p* =
$SPEC(\lambda b.\ b \longrightarrow True)$›

6.2 Algorithm

definition *PAC-checker-step*
:: ‹*int-poly* ⇒ (*status* × *fpac-step*) ⇒ (*int-poly*, *nat*, *nat*) *pac-step* ⇒
(*status* × *fpac-step*) *nres*›
where
‹*PAC-checker-step* = (λ*spec* (*stat*, ($\mathcal{V}$, *A*)) *st*. *case st of*
Add - - - - ⇒
do {
$r \leftarrow$ *normalize-poly-spec* (*pac-res st*);
$eq \leftarrow$ *check-add A* $\mathcal{V}$ (*pac-src1 st*) (*pac-src2 st*) (*new-id st*) *r*;
$st' \leftarrow SPEC(\lambda st'.\ (\neg is\text{-}failed\ st' \wedge is\text{-}found\ st' \longrightarrow r - spec \in ideal\ polynom\text{-}bool))$;
if eq
then RETURN (*merge-status stat st'*,
$\mathcal{V}$, *fmupd* (*new-id st*) *r A*)
else RETURN (*FAILED*, ($\mathcal{V}$, *A*))
}
| *Del* - ⇒
do {
$eq \leftarrow$ *check-del A* (*pac-src1 st*);
if eq
then RETURN (*stat*, ($\mathcal{V}$, *fmdrop* (*pac-src1 st*) *A*))
else RETURN (*FAILED*, ($\mathcal{V}$, *A*))
}
| *Mult* - - - - ⇒
do {
$r \leftarrow$ *normalize-poly-spec* (*pac-res st*);
$q \leftarrow$ *normalize-poly-spec* (*pac-mult st*);
$eq \leftarrow$ *check-mult A* $\mathcal{V}$ (*pac-src1 st*) *q* (*new-id st*) *r*;
$st' \leftarrow SPEC(\lambda st'.\ (\neg is\text{-}failed\ st' \wedge is\text{-}found\ st' \longrightarrow r - spec \in ideal\ polynom\text{-}bool))$;
if eq
then RETURN (*merge-status stat st'*,
$\mathcal{V}$, *fmupd* (*new-id st*) *r A*)
else RETURN (*FAILED*, ($\mathcal{V}$, *A*))
}
| *Extension* - - - ⇒
do {
$r \leftarrow$ *normalize-poly-spec* (*pac-res st* − *Var* (*new-var st*));

(*eq*) ← *check-extension A* 𝒱 (*new-id st*) (*new-var st*) *r*;
if eq
then do {
RETURN (*stat*,
insert (*new-var st*) 𝒱, *fmupd* (*new-id st*) (*r*) *A*)}
else RETURN (*FAILED*, (𝒱, *A*))
}
)›

definition *polys-rel* :: ‹((*nat*, *int mpoly*)*fmap* × -) *set*› **where**
‹*polys-rel* = {(*A*, *B*). *B* = (*ran-m A*)}›

definition *polys-rel-full* :: ‹((*nat set* × (*nat*, *int mpoly*)*fmap*) × -) *set*› **where**
‹*polys-rel-full* = {((𝒱, *A*), (𝒱′ , *B*)). (*A*, *B*) ∈ *polys-rel* ∧ 𝒱 = 𝒱′}›

lemma *polys-rel-update-remove*:
‹*x13* ∉#*dom-m A* ⟹ *x11* ∈# *dom-m A* ⟹ *x12* ∈# *dom-m A* ⟹ *x11* ≠ *x12* ⟹ (*A*,*B*) ∈ *polys-rel* ⟹
(*fmupd x13 r* (*fmdrop x11* (*fmdrop x12 A*)),
add-mset r B − {#*the* (*fmlookup A x11*), *the* (*fmlookup A x12*)#})
∈ *polys-rel*›
‹*x13* ∉#*dom-m A* ⟹ *x11* ∈# *dom-m A* ⟹ (*A*,*B*) ∈ *polys-rel* ⟹
(*fmupd x13 r* (*fmdrop x11 A*),*add-mset r B* − {#*the* (*fmlookup A x11*)#})
∈ *polys-rel*›
‹*x13* ∉#*dom-m A* ⟹ (*A*,*B*) ∈ *polys-rel* ⟹
(*fmupd x13 r A*, *add-mset r B*) ∈ *polys-rel*›
‹*x13* ∈#*dom-m A* ⟹ (*A*,*B*) ∈ *polys-rel* ⟹
(*fmdrop x13 A*, *remove1-mset* (*the* (*fmlookup A x13*)) *B*) ∈ *polys-rel*›
⟨*proof*⟩

lemma *polys-rel-in-dom-inD*:
‹(*A*, *B*) ∈ *polys-rel* ⟹
x12 ∈# *dom-m A* ⟹
the (*fmlookup A x12*) ∈# *B*›
⟨*proof*⟩

lemma *PAC-Format-add-and-remove*:
‹*r* − *x14* ∈ *More-Modules.ideal polynom-bool* ⟹
(*A*, *B*) ∈ *polys-rel* ⟹
x12 ∈# *dom-m A* ⟹
x13 ∉# *dom-m A* ⟹
vars r ⊆ 𝒱 ⟹
2 ∗ *the* (*fmlookup A x12*) − *r* ∈ *More-Modules.ideal polynom-bool* ⟹
PAC-Format** (𝒱, *B*) (𝒱, *remove1-mset* (*the* (*fmlookup A x12*)) (*add-mset r B*))›
‹*r* − *x14* ∈ *More-Modules.ideal polynom-bool* ⟹
(*A*, *B*) ∈ *polys-rel* ⟹
the (*fmlookup A x11*) + *the* (*fmlookup A x12*) − *r* ∈ *More-Modules.ideal polynom-bool* ⟹
x11 ∈# *dom-m A* ⟹
x12 ∈# *dom-m A* ⟹
vars r ⊆ 𝒱 ⟹
PAC-Format** (𝒱, *B*) (𝒱, *add-mset r B*)›
‹*r* − *x14* ∈ *More-Modules.ideal polynom-bool* ⟹
(*A*, *B*) ∈ *polys-rel* ⟹
x11 ∈# *dom-m A* ⟹
x12 ∈# *dom-m A* ⟹

the (*fmlookup A x11*) + *the* (*fmlookup A x12*) − *r* ∈ *More-Modules.ideal polynom-bool* ⟹
vars r ⊆ 𝒱 ⟹
x11 ≠ *x12* ⟹
*PAC-Format*** (𝒱, *B*)
(𝒱, *add-mset r B* − {#*the* (*fmlookup A x11*), *the* (*fmlookup A x12*)#})›
‹(*A*, *B*) ∈ *polys-rel* ⟹
r − *x34* ∈ *More-Modules.ideal polynom-bool* ⟹
x11 ∈# *dom-m A* ⟹
the (*fmlookup A x11*) ∗ *x32* − *r* ∈ *More-Modules.ideal polynom-bool* ⟹
vars x32 ⊆ 𝒱 ⟹
vars r ⊆ 𝒱 ⟹
*PAC-Format*** (𝒱, *B*) (𝒱, *add-mset r B*)›
‹(*A*, *B*) ∈ *polys-rel* ⟹
r − *x34* ∈ *More-Modules.ideal polynom-bool* ⟹
x11 ∈# *dom-m A* ⟹
the (*fmlookup A x11*) ∗ *x32* − *r* ∈ *More-Modules.ideal polynom-bool* ⟹
vars x32 ⊆ 𝒱 ⟹
vars r ⊆ 𝒱 ⟹
*PAC-Format*** (𝒱, *B*) (𝒱, *remove1-mset* (*the* (*fmlookup A x11*)) (*add-mset r B*))›
‹(*A*, *B*) ∈ *polys-rel* ⟹
x12 ∈# *dom-m A* ⟹
*PAC-Format*** (𝒱, *B*) (𝒱, *remove1-mset* (*the* (*fmlookup A x12*)) *B*)›
‹(*A*, *B*) ∈ *polys-rel* ⟹
$(p' + Var\ x)^2$ − (*p'* + *Var x*) ∈ *ideal polynom-bool* ⟹
x ∉ 𝒱 ⟹
x ∉ *vars*(*p'* + *Var x*) ⟹
vars(*p'* + *Var x*) ⊆ 𝒱 ⟹
*PAC-Format*** (𝒱, *B*)
(*insert x* 𝒱, *add-mset p' B*)›
⟨*proof*⟩

abbreviation *status-rel* :: ‹(*status* × *status*) *set*› **where**
‹*status-rel* ≡ *Id*›

lemma *is-merge-status*[*simp*]:
‹*is-failed* (*merge-status a st'*) ⟷ *is-failed a* ∨ *is-failed st'*›
‹*is-found* (*merge-status a st'*) ⟷ ¬*is-failed a* ∧ ¬*is-failed st'* ∧ (*is-found a* ∨ *is-found st'*)›
‹*is-success* (*merge-status a st'*) ⟷ (*is-success a* ∧ *is-success st'*)›
⟨*proof*⟩

lemma *status-rel-merge-status*:
‹(*merge-status a b*, *SUCCESS*) ∉ *status-rel* ⟷
(*a* = *FAILED*) ∨ (*b* = *FAILED*) ∨
a = *FOUND* ∨ (*b* = *FOUND*)›
⟨*proof*⟩

lemma *Ex-status-iff*:
‹(∃ *a*. *P a*) ⟷ *P SUCCESS* ∨ *P FOUND* ∨ (*P* (*FAILED*))›
⟨*proof*⟩

lemma *is-failed-alt-def*:
‹*is-failed st'* ⟷ ¬*is-success st'* ∧ ¬*is-found st'*›
⟨*proof*⟩

lemma *merge-status-eq-iff* [*simp*]:
⟨*merge-status a SUCCESS = SUCCESS* $\longleftrightarrow$ *a = SUCCESS*⟩
⟨*merge-status a SUCCESS = FOUND* $\longleftrightarrow$ *a = FOUND*⟩
⟨*merge-status SUCCESS a = SUCCESS* $\longleftrightarrow$ *a = SUCCESS*⟩
⟨*merge-status SUCCESS a = FOUND* $\longleftrightarrow$ *a = FOUND*⟩
⟨*merge-status SUCCESS a = FAILED* $\longleftrightarrow$ *a = FAILED*⟩
⟨*merge-status a SUCCESS = FAILED* $\longleftrightarrow$ *a = FAILED*⟩
⟨*merge-status FOUND a = FAILED* $\longleftrightarrow$ *a = FAILED*⟩
⟨*merge-status a FOUND = FAILED* $\longleftrightarrow$ *a = FAILED*⟩
⟨*merge-status a FOUND = SUCCESS* $\longleftrightarrow$ *False*⟩
⟨*merge-status a b = FOUND* $\longleftrightarrow$ (*a = FOUND* $\vee$ *b = FOUND*) $\wedge$ (*a* $\neq$ *FAILED* $\wedge$ *b* $\neq$ *FAILED*)⟩
⟨*proof*⟩

lemma *fmdrop-irrelevant*: ⟨*x11* $\notin\#$ *dom-m A* $\Longrightarrow$ *fmdrop x11 A = A*⟩
⟨*proof*⟩

lemma *PAC-checker-step-PAC-checker-specification2*:
fixes *a* :: ⟨*status*⟩
assumes *AB*: ⟨(($\mathcal{V}$, *A*),($\mathcal{V}_B$, *B*)) $\in$ *polys-rel-full*⟩ **and**
⟨$\neg$*is-failed a*⟩ **and**
[*simp*,*intro*]: ⟨*a = FOUND* $\Longrightarrow$ *spec* $\in$ *pac-ideal* (*set-mset* A_0)⟩ **and**
A_0B: ⟨*PAC-Format*** ($\mathcal{V}_0$, A_0) ($\mathcal{V}$, *B*)⟩ **and**
spec$_0$: ⟨*vars spec* $\subseteq$ $\mathcal{V}_0$⟩ **and**
vars-A$_0$: ⟨$\bigcup$ (*vars ‘ set-mset* A_0) $\subseteq$ $\mathcal{V}_0$⟩
shows ⟨*PAC-checker-step spec* (*a*, ($\mathcal{V}$, *A*)) *st* $\leq\Downarrow$ (*status-rel* $\times_r$ *polys-rel-full*) (*PAC-checker-specification-step2* ($\mathcal{V}_0$, A_0) *spec* ($\mathcal{V}$, *B*))⟩
⟨*proof*⟩

definition *PAC-checker*
:: ⟨*int-poly* $\Rightarrow$ *fpac-step* $\Rightarrow$ *status* $\Rightarrow$ (*int-poly*, *nat*, *nat*) *pac-step list* $\Rightarrow$
(*status* $\times$ *fpac-step*) *nres*⟩
where
⟨*PAC-checker spec A b st = do* {
(*S*, -) $\leftarrow$ *WHILE*$_T$
(λ((*b* :: *status*, *A* :: *fpac-step*), *st*). $\neg$*is-failed b* $\wedge$ *st* $\neq$ [])
(λ((*bA*), *st*). *do* {
ASSERT(*st* $\neq$ []);
S $\leftarrow$ *PAC-checker-step spec* (*bA*) (*hd st*);
RETURN (*S*, *tl st*)
})
((*b*, *A*), *st*);
RETURN S
}⟩

lemma *PAC-checker-specification-spec-trans*:
⟨*PAC-checker-specification-spec spec A* (*st*, *x2*) $\Longrightarrow$
PAC-checker-specification-step-spec A spec x2 (*st′*, *x1a*) $\Longrightarrow$
PAC-checker-specification-spec spec A (*st′*, *x1a*)⟩
⟨*proof*⟩

lemma *RES-SPEC-eq*:
⟨*RES* Φ = *SPEC*(λP. $P \in \Phi$)⟩
⟨*proof*⟩

lemma *is-failed-is-success-completeD*:
‹$\neg$ *is-failed* $x \Longrightarrow \neg$*is-success* $x \Longrightarrow$ *is-found* x›
⟨*proof*⟩

lemma *PAC-checker-PAC-checker-specification2*:
‹$(A, B) \in$ *polys-rel-full* $\Longrightarrow$
$\neg$*is-failed* $a \Longrightarrow$
$(a = FOUND \Longrightarrow spec \in$ *pac-ideal* (*set-mset* (*snd* B))) $\Longrightarrow$
$\bigcup$(*vars* ‘ *set-mset* (*ran-m* (*snd* A))) $\subseteq$ *fst* $B \Longrightarrow$
vars spec $\subseteq$ *fst* $B \Longrightarrow$
PAC-checker spec A a *st* $\leq \Downarrow$ (*status-rel* $\times_r$ *polys-rel-full*) (*PAC-checker-specification2 spec* B)›
⟨*proof*⟩

definition *remap-polys-polynom-bool* :: ‹*int mpoly* $\Rightarrow$ *nat set* $\Rightarrow$ (*nat*, *int-poly*) *fmap* $\Rightarrow$ (*status* $\times$ *fpac-step*) *nres*› **where**
‹*remap-polys-polynom-bool spec* = ($\lambda\mathcal{V}$ A.
SPEC(λ(*st*, $\mathcal{V}'$, A'). ($\neg$*is-failed st* $\longrightarrow$
dom-m A = *dom-m* A' $\wedge$
($\forall i \in\#$ *dom-m* A. *the* (*fmlookup* A i) $-$ *the* (*fmlookup* A' i) $\in$ *ideal polynom-bool*) $\wedge$
$\bigcup$(*vars* ‘ *set-mset* (*ran-m* A)) $\subseteq \mathcal{V}'$ $\wedge$
$\bigcup$(*vars* ‘ *set-mset* (*ran-m* A')) $\subseteq \mathcal{V}'$) $\wedge$
(*st* = *FOUND* $\longrightarrow$ *spec* $\in\#$ *ran-m* A')))›

definition *remap-polys-change-all* :: ‹*int mpoly* $\Rightarrow$ *nat set* $\Rightarrow$ (*nat*, *int-poly*) *fmap* $\Rightarrow$ (*status* $\times$ *fpac-step*) *nres*› **where**
‹*remap-polys-change-all spec* = ($\lambda\mathcal{V}$ A. *SPEC* (λ(*st*, $\mathcal{V}'$, A').
($\neg$*is-failed st* $\longrightarrow$
pac-ideal (*set-mset* (*ran-m* A)) = *pac-ideal* (*set-mset* (*ran-m* A')) $\wedge$
$\bigcup$(*vars* ‘ *set-mset* (*ran-m* A)) $\subseteq \mathcal{V}'$ $\wedge$
$\bigcup$(*vars* ‘ *set-mset* (*ran-m* A')) $\subseteq \mathcal{V}'$) $\wedge$
(*st* = *FOUND* $\longrightarrow$ *spec* $\in\#$ *ran-m* A')))›

lemma *fmap-eq-dom-iff*:
‹$A = A' \longleftrightarrow$ *dom-m* A = *dom-m* A' $\wedge$ ($\forall i \in\#$ *dom-m* A. *the* (*fmlookup* A i) = *the* (*fmlookup* A' i))›
⟨*proof*⟩

lemma *ideal-remap-incl*:
‹*finite* $A' \Longrightarrow (\forall a' \in A'. \exists a \in A. a - a' \in B) \Longrightarrow$ *ideal* $(A' \cup B) \subseteq$ *ideal* $(A \cup B)$›
⟨*proof*⟩

lemma *pac-ideal-remap-eq*:
‹*dom-m* b = *dom-m ba* $\Longrightarrow$
$\forall i \in\#$*dom-m ba*.
the (*fmlookup* b i) $-$ *the* (*fmlookup ba* i)
$\in$ *More-Modules.ideal polynom-bool* $\Longrightarrow$
pac-ideal ((λx. *the* (*fmlookup* b x)) ‘ *set-mset* (*dom-m ba*)) = *pac-ideal* ((λx. *the* (*fmlookup ba* x)) ‘ *set-mset* (*dom-m ba*))›
⟨*proof*⟩

lemma *remap-polys-polynom-bool-remap-polys-change-all*:
‹*remap-polys-polynom-bool spec* $\mathcal{V}$ $A \leq$ *remap-polys-change-all spec* $\mathcal{V}$ A›
⟨*proof*⟩

definition *remap-polys* :: ‹*int mpoly* ⇒ *nat set* ⇒ (*nat*, *int-poly*) *fmap* ⇒ (*status* × *fpac-step*) *nres*›
where
‹*remap-polys spec* = (λ𝒱 *A*. *do*{
dom ← *SPEC*(λ*dom*. *set-mset* (*dom-m A*) ⊆ *dom* ∧ *finite dom*);

failed ← *SPEC*(λ-::*bool*. *True*);
if failed
then do {
RETURN (*FAILED*, 𝒱, *fmempty*)
}
else do {
(*b*, *N*) ← *FOREACH dom*
(λ*i* (*b*, 𝒱, *A*′).
if i ∈# *dom-m A*
then do {
p ← *SPEC*(λ*p*. *the* (*fmlookup A i*) − *p* ∈ *ideal polynom-bool* ∧ *vars p* ⊆ *vars* (*the* (*fmlookup A i*)));
eq ← *SPEC*(λ*eq*. *eq* ⟶ *p* = *spec*);
𝒱 ← *SPEC*(λ𝒱′. 𝒱 ∪ *vars* (*the* (*fmlookup A i*)) ⊆ 𝒱′);
RETURN(*b* ∨ *eq*, 𝒱, *fmupd i p A*′)
} *else RETURN* (*b*, 𝒱, *A*′))
(*False*, 𝒱, *fmempty*);
RETURN (*if b then FOUND else SUCCESS*, *N*)
}
})›

lemma *remap-polys-spec*:
‹*remap-polys spec* 𝒱 *A* ≤ *remap-polys-polynom-bool spec* 𝒱 *A*›
⟨*proof*⟩

6.3 Full Checker

definition *full-checker*
:: ‹*int-poly* ⇒ (*nat*, *int-poly*) *fmap* ⇒ (*int-poly*, *nat*,*nat*) *pac-step list* ⇒ (*status* × -) *nres*›
where
‹*full-checker spec0 A pac* = *do* {
spec ← *normalize-poly-spec spec0*;
(*st*, 𝒱, *A*) ← *remap-polys-change-all spec* {} *A*;
if is-failed st then
RETURN (*st*, 𝒱, *A*)
else do {
𝒱 ← *SPEC*(λ𝒱′. 𝒱 ∪ *vars spec0* ⊆ 𝒱′);
PAC-checker spec (𝒱, *A*) *st pac*
}
}›

lemma *restricted-ideal-to-mono*:
‹*restricted-ideal-to*$_I$ 𝒱 *I* ⊆ *restricted-ideal-to*$_I$ 𝒱′ *J* ⟹
𝒰 ⊆ 𝒱 ⟹
restricted-ideal-to$_I$ 𝒰 *I* ⊆ *restricted-ideal-to*$_I$ 𝒰 *J*›
⟨*proof*⟩

lemma *full-checker-spec*:
assumes ‹(*A*, *A*′) ∈ *polys-rel*›
shows
‹*full-checker spec A pac* ≤ ⇓{((*st*, *G*), (*st*′, *G*′)). (*st*, *st*′) ∈ *status-rel* ∧

$(st \neq FAILED \longrightarrow (G,\ G') \in polys\text{-}rel\text{-}full)\}$
$(PAC\text{-}checker\text{-}specification\ spec\ (A'))$›
$\langle proof \rangle$

lemma *full-checker-spec′*:
shows
‹$(uncurry2\ full\text{-}checker,\ uncurry2\ (\lambda spec\ A\ \text{-}.\ PAC\text{-}checker\text{-}specification\ spec\ A)) \in$
$(Id \times_r polys\text{-}rel) \times_r Id \rightarrow_f \langle\{((st,\ G),\ (st',\ G')).\ (st,\ st') \in status\text{-}rel \wedge$
$(st \neq FAILED \longrightarrow (G,\ G') \in polys\text{-}rel\text{-}full)\}\rangle nres\text{-}rel$›
$\langle proof \rangle$

end
theory *PAC-Polynomials*
imports *PAC-Specification Finite-Map-Multiset*
begin

7 Polynomials of strings

Isabelle's definition of polynomials only work with variables of type *nat*. Therefore, we introduce a version that uses strings.

7.1 Polynomials and Variables

lemma *poly-embed-EX*:
‹$\exists \varphi.\ bij\ (\varphi :: string \Rightarrow nat)$›
$\langle proof \rangle$

Using a multiset instead of a list has some advantage from an abstract point of view. First, we can have monomials that appear several times and the coefficient can also be zero. Basically, we can represent un-normalised polynomials, which is very useful to talk about intermediate states in our program.

type-synonym *term-poly* = ‹*string multiset*›
type-synonym *mset-polynom* =
‹$(term\text{-}poly * int)\ multiset$›

definition *normalized-poly* :: ‹*mset-polynom* ⇒ *bool*› **where**
‹$normalized\text{-}poly\ p \longleftrightarrow$
$distinct\text{-}mset\ (fst\ \text{‘}\#\ p) \wedge$
$0 \notin\#\ snd\ \text{‘}\#\ p$›

lemma *normalized-poly-simps*[*simp*]:
‹$normalized\text{-}poly\ \{\#\}$›
‹$normalized\text{-}poly\ (add\text{-}mset\ t\ p) \longleftrightarrow snd\ t \neq 0 \wedge$
$fst\ t \notin\#\ fst\ \text{‘}\#\ p \wedge normalized\text{-}poly\ p$›
$\langle proof \rangle$

lemma *normalized-poly-mono*:
‹$normalized\text{-}poly\ B \Longrightarrow A \subseteq\#\ B \Longrightarrow normalized\text{-}poly\ A$›
$\langle proof \rangle$

definition *mult-poly-by-monom* :: ‹*term-poly* ∗ *int* ⇒ *mset-polynom* ⇒ *mset-polynom*› **where**
‹$mult\text{-}poly\text{-}by\text{-}monom\ = (\lambda ys\ q.\ image\text{-}mset\ (\lambda xs.\ (fst\ xs + fst\ ys,\ snd\ ys * snd\ xs))\ q)$›

definition *mult-poly-raw* :: ‹*mset-polynom* ⇒ *mset-polynom* ⇒ *mset-polynom*› **where**
‹*mult-poly-raw p q* =
(*sum-mset* ((λ*y*. *mult-poly-by-monom y q*) ‘# *p*))›

definition *remove-powers* :: ‹*mset-polynom* ⇒ *mset-polynom*› **where**
‹*remove-powers xs* = *image-mset* (*apfst remdups-mset*) *xs*›

definition *all-vars-mset* :: ‹*mset-polynom* ⇒ *string multiset*› **where**
‹*all-vars-mset p* = $\bigcup$# (*fst* ‘# *p*)›

abbreviation *all-vars* :: ‹*mset-polynom* ⇒ *string set*› **where**
‹*all-vars p* ≡ *set-mset* (*all-vars-mset p*)›

definition *add-to-coefficient* :: ‹- ⇒ *mset-polynom* ⇒ *mset-polynom*› **where**
‹*add-to-coefficient* = (λ(*a*, *n*) *b*. {#(*a*′, -) ∈# *b*. *a*′ ≠ *a*#} +
(*if n* + *sum-mset* (*snd* ‘# {#(*a*′, -) ∈# *b*. *a*′ = *a*#}) = *0 then* {#}
else {#(*a*, *n* + *sum-mset* (*snd* ‘# {#(*a*′, -) ∈# *b*. *a*′ = *a*#}))#}))›

definition *normalize-poly* :: ‹*mset-polynom* ⇒ *mset-polynom*› **where**
‹*normalize-poly p* = *fold-mset add-to-coefficient* {#} *p*›

lemma *add-to-coefficient-simps*:
‹*n* + *sum-mset* (*snd* ‘# {#(*a*′, -) ∈# *b*. *a*′ = *a*#}) ≠ *0* ⟹
add-to-coefficient (*a*, *n*) *b* = {#(*a*′, -) ∈# *b*. *a*′ ≠ *a*#} +
{#(*a*, *n* + *sum-mset* (*snd* ‘# {#(*a*′, -) ∈# *b*. *a*′ = *a*#}))#}›
‹*n* + *sum-mset* (*snd* ‘# {#(*a*′, -) ∈# *b*. *a*′ = *a*#}) = *0* ⟹
add-to-coefficient (*a*, *n*) *b* = {#(*a*′, -) ∈# *b*. *a*′ ≠ *a*#}› **and**
add-to-coefficient-simps-If:
‹*add-to-coefficient* (*a*, *n*) *b* = {#(*a*′, -) ∈# *b*. *a*′ ≠ *a*#} +
(*if n* + *sum-mset* (*snd* ‘# {#(*a*′, -) ∈# *b*. *a*′ = *a*#}) = *0 then* {#}
else {#(*a*, *n* + *sum-mset* (*snd* ‘# {#(*a*′, -) ∈# *b*. *a*′ = *a*#}))#})›
⟨*proof*⟩

interpretation *comp-fun-commute* ‹*add-to-coefficient*›
⟨*proof*⟩

lemma *normalized-poly-normalize-poly*[*simp*]:
‹*normalized-poly* (*normalize-poly p*)›
⟨*proof*⟩

7.2 Addition

inductive *add-poly-p* :: ‹*mset-polynom* × *mset-polynom* × *mset-polynom* ⇒ *mset-polynom* × *mset-polynom* × *mset-polynom* ⇒ *bool*› **where**
add-new-coeff-r:
‹*add-poly-p* (*p*, *add-mset x q*, *r*) (*p*, *q*, *add-mset x r*)› |
add-new-coeff-l:
‹*add-poly-p* (*add-mset x p*, *q*, *r*) (*p*, *q*, *add-mset x r*)› |
add-same-coeff-l:
‹*add-poly-p* (*add-mset* (*x*, *n*) *p*, *q*, *add-mset* (*x*, *m*) *r*) (*p*, *q*, *add-mset* (*x*, *n* + *m*) *r*)› |
add-same-coeff-r:
‹*add-poly-p* (*p*, *add-mset* (*x*, *n*) *q*, *add-mset* (*x*, *m*) *r*) (*p*, *q*, *add-mset* (*x*, *n* + *m*) *r*)› |
rem-0-coeff:

⟨*add-poly-p* (*p*, *q*, *add-mset* (*x*, *0*) *r*) (*p*, *q*, *r*)⟩

inductive-cases *add-poly-pE*: ⟨*add-poly-p* *S* *T*⟩

lemmas *add-poly-p-induct* =
add-poly-p.induct[*split-format*(*complete*)]

lemma *add-poly-p-empty-l*:
⟨*add-poly-p*** (*p*, *q*, *r*) ({#}, *q*, *p* + *r*)⟩
⟨*proof*⟩

lemma *add-poly-p-empty-r*:
⟨*add-poly-p*** (*p*, *q*, *r*) (*p*, {#}, *q* + *r*)⟩
⟨*proof*⟩

lemma *add-poly-p-sym*:
⟨*add-poly-p* (*p*, *q*, *r*) (*p'*, *q'*, *r'*) ⟷ *add-poly-p* (*q*, *p*, *r*) (*q'*, *p'*, *r'*)⟩
⟨*proof*⟩

lemma *wf-if-measure-in-wf*:
⟨*wf R* ⟹ (⋀*a b*. (*a*, *b*) ∈ *S* ⟹ (*ν a*, *ν b*)∈*R*) ⟹ *wf S*⟩
⟨*proof*⟩

lemma *lexn-n*:
⟨*n* > *0* ⟹ (*x* # *xs*, *y* # *ys*) ∈ *lexn r n* ⟷
(*length xs* = *n*−*1* ∧ *length ys* = *n*−*1*) ∧ ((*x*, *y*) ∈ *r* ∨ (*x* = *y* ∧ (*xs*, *ys*) ∈ *lexn r* (*n* − *1*)))⟩
⟨*proof*⟩

lemma *wf-add-poly-p*:
⟨*wf* {(*x*, *y*). *add-poly-p y x*}⟩
⟨*proof*⟩

lemma *mult-poly-by-monom-simps*[*simp*]:
⟨*mult-poly-by-monom t* {#} = {#}⟩
⟨*mult-poly-by-monom t* (*ps* + *qs*) = *mult-poly-by-monom t ps* + *mult-poly-by-monom t qs*⟩
⟨*mult-poly-by-monom a* (*add-mset p ps*) = *add-mset* (*fst a* + *fst p*, *snd a* ∗ *snd p*) (*mult-poly-by-monom a ps*)⟩
⟨*proof*⟩

inductive *mult-poly-p* :: ⟨*mset-polynom* ⇒ *mset-polynom* × *mset-polynom* ⇒ *mset-polynom* × *mset-polynom* ⇒ *bool*⟩
for *q* :: *mset-polynom* **where**
mult-step:
⟨*mult-poly-p q* (*add-mset* (*xs*, *n*) *p*, *r*) (*p*, (λ(*ys*, *m*). (*remdups-mset* (*xs* + *ys*), *n* ∗ *m*)) '# *q* + *r*)⟩

lemmas *mult-poly-p-induct* = *mult-poly-p.induct*[*split-format*(*complete*)]

7.3 Normalisation

inductive *normalize-poly-p* :: ⟨*mset-polynom* ⇒ *mset-polynom* ⇒ *bool*⟩**where**
rem-0-coeff[*simp*, *intro*]:
⟨*normalize-poly-p p q* ⟹ *normalize-poly-p* (*add-mset* (*xs*, *0*) *p*) *q*⟩ |
merge-dup-coeff[*simp*, *intro*]:
⟨*normalize-poly-p p q* ⟹ *normalize-poly-p* (*add-mset* (*xs*, *m*) (*add-mset* (*xs*, *n*) *p*)) (*add-mset* (*xs*, *m* + *n*) *q*)⟩ |

same[*simp*, *intro*]:
  ‹*normalize-poly-p p p*› |
keep-coeff[*simp*, *intro*]:
  ‹*normalize-poly-p p q* ⟹ *normalize-poly-p* (*add-mset x p*) (*add-mset x q*)›

7.4 Correctness

This locales maps string polynomials to real polynomials.

locale *poly-embed* =
 fixes *φ* :: ‹*string* ⇒ *nat*›
 assumes *φ-inj*: ‹*inj φ*›
begin

definition *poly-of-vars* :: *term-poly* ⇒ ('*a* :: {*comm-semiring-1*}) *mpoly* **where**
 ‹*poly-of-vars xs* = *fold-mset* (λ*a b*. *Var* (*φ a*) ∗ *b*) (*1* :: '*a mpoly*) *xs*›

lemma *poly-of-vars-simps*[*simp*]:
 shows
  ‹*poly-of-vars* (*add-mset x xs*) = *Var* (*φ x*) ∗ (*poly-of-vars xs* :: ('*a* :: {*comm-semiring-1*}) *mpoly*)› (**is** *?A*) **and**
  ‹*poly-of-vars* (*xs* + *ys*) = *poly-of-vars xs* ∗ (*poly-of-vars ys* :: ('*a* :: {*comm-semiring-1*}) *mpoly*)› (**is** *?B*)
⟨*proof*⟩

definition *mononom-of-vars* **where**
 ‹*mononom-of-vars* ≡ (λ(*xs*, *n*). (+) (*Const n* ∗ *poly-of-vars xs*))›

interpretation *comp-fun-commute* ‹*mononom-of-vars*›
 ⟨*proof*⟩

lemma [*simp*]:
 ‹*poly-of-vars* {#} = *1*›
 ⟨*proof*⟩

lemma *mononom-of-vars-add*[*simp*]:
 ‹*NO-MATCH 0 b* ⟹ *mononom-of-vars xs b* = *Const* (*snd xs*) ∗ *poly-of-vars* (*fst xs*) + *b*›
 ⟨*proof*⟩

definition *polynom-of-mset* :: ‹*mset-polynom* ⇒ -› **where**
 ‹*polynom-of-mset p* = *sum-mset* (*mononom-of-vars* '# *p*) *0*›

lemma *polynom-of-mset-append*[*simp*]:
 ‹*polynom-of-mset* (*xs* + *ys*) = *polynom-of-mset xs* + *polynom-of-mset ys*›
 ⟨*proof*⟩

lemma *polynom-of-mset-Cons*[*simp*]:
 ‹*polynom-of-mset* (*add-mset x ys*) = *Const* (*snd x*) ∗ *poly-of-vars* (*fst x*) + *polynom-of-mset ys*›
 ⟨*proof*⟩

lemma *polynom-of-mset-empty*[*simp*]:
 ‹*polynom-of-mset* {#} = *0*›
 ⟨*proof*⟩

lemma *polynom-of-mset-mult-poly-by-monom*[*simp*]:

‹polynom-of-mset (mult-poly-by-monom x ys) =
(Const (snd x) ∗ poly-of-vars (fst x) ∗ polynom-of-mset ys)›
⟨proof⟩

lemma *polynom-of-mset-mult-poly-raw*[*simp*]:
‹polynom-of-mset (mult-poly-raw xs ys) = polynom-of-mset xs ∗ polynom-of-mset ys›
⟨proof⟩

lemma *polynom-of-mset-uminus*:
‹polynom-of-mset {#case x of (a, b) ⇒ (a, − b). x ∈# za#} =
− polynom-of-mset za›
⟨proof⟩

lemma *X2-X-polynom-bool-mult-in*:
‹Var (x1) ∗ (Var (x1) ∗ p) − Var (x1) ∗ p ∈ More-Modules.ideal polynom-bool›
⟨proof⟩

lemma *polynom-of-list-remove-powers-polynom-bool*:
‹(polynom-of-mset xs) − polynom-of-mset (remove-powers xs) ∈ ideal polynom-bool›
⟨proof⟩

lemma *add-poly-p-polynom-of-mset*:
‹add-poly-p (p, q, r) (p′, q′, r′) ⟹
polynom-of-mset r + (polynom-of-mset p + polynom-of-mset q) =
polynom-of-mset r′ + (polynom-of-mset p′ + polynom-of-mset q′)›
⟨proof⟩

lemma *rtranclp-add-poly-p-polynom-of-mset*:
‹add-poly-p** (p, q, r) (p′, q′, r′) ⟹
polynom-of-mset r + (polynom-of-mset p + polynom-of-mset q) =
polynom-of-mset r′ + (polynom-of-mset p′ + polynom-of-mset q′)›
⟨proof⟩

lemma *rtranclp-add-poly-p-polynom-of-mset-full*:
‹add-poly-p** (p, q, {#}) ({#}, {#}, r′) ⟹
polynom-of-mset r′ = (polynom-of-mset p + polynom-of-mset q)›
⟨proof⟩

lemma *poly-of-vars-remdups-mset*:
‹poly-of-vars (remdups-mset (xs)) − (poly-of-vars xs)
∈ More-Modules.ideal polynom-bool›
⟨proof⟩

lemma *polynom-of-mset-mult-map*:
‹polynom-of-mset
{#case x of (ys, n) ⇒ (remdups-mset (ys + xs), n ∗ m). x ∈# q#} −
Const m ∗ (poly-of-vars xs ∗ polynom-of-mset q)
∈ More-Modules.ideal polynom-bool›
(**is** ‹?P q ∈ -›)
⟨proof⟩

lemma *mult-poly-p-mult-ideal*:

$\langle$*mult-poly-p q* (*p*, *r*) (*p*′, *r*′) $\Longrightarrow$
(*polynom-of-mset p*′ ∗ *polynom-of-mset q* + *polynom-of-mset r*′) − (*polynom-of-mset p* ∗ *polynom-of-mset q* + *polynom-of-mset r*)
∈ *ideal polynom-bool*$\rangle$
⟨*proof*⟩

lemma *rtranclp-mult-poly-p-mult-ideal*:
$\langle$(*mult-poly-p q*)** (*p*, *r*) (*p*′, *r*′) $\Longrightarrow$
(*polynom-of-mset p*′ ∗ *polynom-of-mset q* + *polynom-of-mset r*′) − (*polynom-of-mset p* ∗ *polynom-of-mset q* + *polynom-of-mset r*)
∈ *ideal polynom-bool*$\rangle$
⟨*proof*⟩

lemma *rtranclp-mult-poly-p-mult-ideal-final*:
$\langle$(*mult-poly-p q*)** (*p*, {#}) ({#}, *r*) $\Longrightarrow$
(*polynom-of-mset r*) − (*polynom-of-mset p* ∗ *polynom-of-mset q*)
∈ *ideal polynom-bool*$\rangle$
⟨*proof*⟩

lemma *normalize-poly-p-poly-of-mset*:
$\langle$*normalize-poly-p p q* $\Longrightarrow$ *polynom-of-mset p* = *polynom-of-mset q*$\rangle$
⟨*proof*⟩

lemma *rtranclp-normalize-poly-p-poly-of-mset*:
$\langle$*normalize-poly-p*** *p q* $\Longrightarrow$ *polynom-of-mset p* = *polynom-of-mset q*$\rangle$
⟨*proof*⟩

end

It would be nice to have the property in the other direction too, but this requires a deep dive into the definitions of polynomials.

locale *poly-embed-bij* = *poly-embed* +
fixes *V N*
assumes *φ-bij*: $\langle$*bij-betw φ V N*$\rangle$
begin

definition *φ*′ :: $\langle$*nat* ⇒ *string*$\rangle$ **where**
$\langle$*φ*′ = *the-inv-into V φ*$\rangle$

lemma *φ*′*-φ*[*simp*]:
$\langle$*x* ∈ *V* $\Longrightarrow$ *φ*′ (*φ x*) = *x*$\rangle$
⟨*proof*⟩

lemma *φ-φ*′[*simp*]:
$\langle$*x* ∈ *N* $\Longrightarrow$ *φ* (*φ*′ *x*) = *x*$\rangle$
⟨*proof*⟩

end

end
theory *PAC-Polynomials-Term*
imports *PAC-Polynomials*
Refine-Imperative-HOL.IICF
begin

8 Terms

We define some helper functions.

8.1 Ordering

lemma *fref-to-Down-curry-left*:
fixes f:: ‹$'a \Rightarrow 'b \Rightarrow 'c$ *nres*› **and**
A::‹$(('a \times 'b) \times 'd)$ *set*›
shows
‹(*uncurry* f, g) $\in [P]_f$ $A \to \langle B\rangle$*nres-rel* $\Longrightarrow$
$(\bigwedge a\ b\ x'.\ P\ x' \Longrightarrow ((a, b), x') \in A \Longrightarrow f\ a\ b \le \Downarrow B\ (g\ x'))$›
⟨*proof*⟩

lemma *fref-to-Down-curry-right*:
fixes g :: ‹$'a \Rightarrow 'b \Rightarrow 'c$ *nres*› **and** f :: ‹$'d \Rightarrow$ - *nres*› **and**
A::‹$('d \times ('a \times 'b))$ *set*›
shows
‹(f, *uncurry* g) $\in [P]_f$ $A \to \langle B\rangle$*nres-rel* $\Longrightarrow$
$(\bigwedge a\ b\ x'.\ P\ (a,b) \Longrightarrow (x', (a, b)) \in A \Longrightarrow f\ x' \le \Downarrow B\ (g\ a\ b))$›
⟨*proof*⟩

type-synonym *term-poly-list* = ‹*string list*›
type-synonym *llist-polynom* = ‹(*term-poly-list* × *int*) *list*›

We instantiate the characters with typeclass linorder to be able to talk abourt sorted and so on.

definition *less-eq-char* :: ‹*char* ⇒ *char* ⇒ *bool*› **where**
‹*less-eq-char* c d = (((*of-char* c) :: *nat*) ≤ *of-char* d)›

definition *less-char* :: ‹*char* ⇒ *char* ⇒ *bool*› **where**
‹*less-char* c d = (((*of-char* c) :: *nat*) < *of-char* d)›

global-interpretation *char*: *linorder less-eq-char less-char*
⟨*proof*⟩

abbreviation *less-than-char* :: ‹(*char* × *char*) *set*› **where**
‹*less-than-char* ≡ *p2rel less-char*›

lemma *less-than-char-def*:
‹(x,y) ∈ *less-than-char* ⟷ *less-char* x y›
⟨*proof*⟩

lemma *trans-less-than-char*[*simp*]:
‹*trans less-than-char*› **and**
irrefl-less-than-char:
‹*irrefl less-than-char*› **and**
antisym-less-than-char:
‹*antisym less-than-char*›
⟨*proof*⟩

8.2 Polynomials

definition *var-order-rel* :: ‹(*string* × *string*) *set*› **where**
‹*var-order-rel* ≡ *lexord less-than-char*›

abbreviation *var-order* :: ‹*string* ⇒ *string* ⇒ *bool*› **where**
‹*var-order* ≡ *rel2p var-order-rel*›

abbreviation *term-order-rel* :: ‹(*term-poly-list* × *term-poly-list*) *set*› **where**
‹*term-order-rel* ≡ *lexord var-order-rel*›

abbreviation *term-order* :: ‹*term-poly-list* ⇒ *term-poly-list* ⇒ *bool*› **where**
‹*term-order* ≡ *rel2p term-order-rel*›

definition *term-poly-list-rel* :: ‹(*term-poly-list* × *term-poly*) *set*› **where**
‹*term-poly-list-rel* = {(*xs*, *ys*).
ys = *mset xs* ∧
distinct xs ∧
sorted-wrt (*rel2p var-order-rel*) *xs*}›

definition *unsorted-term-poly-list-rel* :: ‹(*term-poly-list* × *term-poly*) *set*› **where**
‹*unsorted-term-poly-list-rel* = {(*xs*, *ys*).
ys = *mset xs* ∧ *distinct xs*}›

definition *poly-list-rel* :: ‹- ⇒ ((′*a* × *int*) *list* × *mset-polynom*) *set*› **where**
‹*poly-list-rel R* = {(*xs*, *ys*).
$(xs, ys) \in \langle R \times_r int\text{-}rel\rangle list\text{-}rel\ O\ list\text{-}mset\text{-}rel\ \wedge$
0 ∉# *snd* ‘# *ys*}›

definition *sorted-poly-list-rel-wrt* :: ‹(′*a* ⇒ ′*a* ⇒ *bool*)
⇒ (′*a* × *string multiset*) *set* ⇒ ((′*a* × *int*) *list* × *mset-polynom*) *set*› **where**
‹*sorted-poly-list-rel-wrt S R* = {(*xs*, *ys*).
$(xs, ys) \in \langle R \times_r int\text{-}rel\rangle list\text{-}rel\ O\ list\text{-}mset\text{-}rel\ \wedge$
sorted-wrt S (*map fst xs*) ∧
distinct (*map fst xs*) ∧
0 ∉# *snd* ‘# *ys*}›

abbreviation *sorted-poly-list-rel* **where**
‹*sorted-poly-list-rel R* ≡ *sorted-poly-list-rel-wrt R term-poly-list-rel*›

abbreviation *sorted-poly-rel* **where**
‹*sorted-poly-rel* ≡ *sorted-poly-list-rel term-order*›

definition *sorted-repeat-poly-list-rel-wrt* :: ‹(′*a* ⇒ ′*a* ⇒ *bool*)
⇒ (′*a* × *string multiset*) *set* ⇒ ((′*a* × *int*) *list* × *mset-polynom*) *set*› **where**
‹*sorted-repeat-poly-list-rel-wrt S R* = {(*xs*, *ys*).
$(xs, ys) \in \langle R \times_r int\text{-}rel\rangle list\text{-}rel\ O\ list\text{-}mset\text{-}rel\ \wedge$
sorted-wrt S (*map fst xs*) ∧
0 ∉# *snd* ‘# *ys*}›

abbreviation *sorted-repeat-poly-list-rel* **where**
‹*sorted-repeat-poly-list-rel R* ≡ *sorted-repeat-poly-list-rel-wrt R term-poly-list-rel*›

abbreviation *sorted-repeat-poly-rel* **where**
‹*sorted-repeat-poly-rel* ≡ *sorted-repeat-poly-list-rel* (*rel2p* (*Id* ∪ *lexord var-order-rel*))›

abbreviation *unsorted-poly-rel* **where**

‹*unsorted-poly-rel* ≡ *poly-list-rel term-poly-list-rel*›

lemma *sorted-poly-list-rel-empty-l*[*simp*]:
‹$([], s') \in$ *sorted-poly-list-rel-wrt* S $T \longleftrightarrow s' = \{\#\}$›
⟨*proof*⟩

definition *fully-unsorted-poly-list-rel* :: ‹- $\Rightarrow$ ((′*a* × *int*) *list* × *mset-polynom*) *set*› **where**
‹*fully-unsorted-poly-list-rel* $R = \{(xs, ys).$
$(xs, ys) \in \langle R \times_r$ *int-rel*⟩*list-rel O list-mset-rel*}›

abbreviation *fully-unsorted-poly-rel* **where**
‹*fully-unsorted-poly-rel* ≡ *fully-unsorted-poly-list-rel unsorted-term-poly-list-rel*›

lemma *fully-unsorted-poly-list-rel-empty-iff* [*simp*]:
‹$(p, \{\#\}) \in$ *fully-unsorted-poly-list-rel* $R \longleftrightarrow p = []$›
‹$([], p') \in$ *fully-unsorted-poly-list-rel* $R \longleftrightarrow p' = \{\#\}$›
⟨*proof*⟩

definition *poly-list-rel-with0* :: ‹- $\Rightarrow$ ((′*a* × *int*) *list* × *mset-polynom*) *set*› **where**
‹*poly-list-rel-with0* $R = \{(xs, ys).$
$(xs, ys) \in \langle R \times_r$ *int-rel*⟩*list-rel O list-mset-rel*}›

abbreviation *unsorted-poly-rel-with0* **where**
‹*unsorted-poly-rel-with0* ≡ *fully-unsorted-poly-list-rel term-poly-list-rel*›

lemma *poly-list-rel-with0-empty-iff* [*simp*]:
‹$(p, \{\#\}) \in$ *poly-list-rel-with0* $R \longleftrightarrow p = []$›
‹$([], p') \in$ *poly-list-rel-with0* $R \longleftrightarrow p' = \{\#\}$›
⟨*proof*⟩

definition *sorted-repeat-poly-list-rel-with0-wrt* :: ‹(′*a* ⇒ ′*a* ⇒ *bool*)
⇒ (′*a* × *string multiset*) *set* ⇒ ((′*a* × *int*) *list* × *mset-polynom*) *set*› **where**
‹*sorted-repeat-poly-list-rel-with0-wrt* S $R = \{(xs, ys).$
$(xs, ys) \in \langle R \times_r$ *int-rel*⟩*list-rel O list-mset-rel* ∧
sorted-wrt S (*map fst xs*)}›

abbreviation *sorted-repeat-poly-list-rel-with0* **where**
‹*sorted-repeat-poly-list-rel-with0* R ≡ *sorted-repeat-poly-list-rel-with0-wrt* R *term-poly-list-rel*›

abbreviation *sorted-repeat-poly-rel-with0* **where**
‹*sorted-repeat-poly-rel-with0* ≡ *sorted-repeat-poly-list-rel-with0* (*rel2p* (*Id* ∪ *lexord var-order-rel*))›

lemma *term-poly-list-relD*:
‹$(xs, ys) \in$ *term-poly-list-rel* ⟹ *distinct xs*›
‹$(xs, ys) \in$ *term-poly-list-rel* ⟹ *ys* = *mset xs*›
‹$(xs, ys) \in$ *term-poly-list-rel* ⟹ *sorted-wrt* (*rel2p var-order-rel*) *xs*›
‹$(xs, ys) \in$ *term-poly-list-rel* ⟹ *sorted-wrt* (*rel2p* (*Id* ∪ *var-order-rel*)) *xs*›
⟨*proof*⟩

end
theory *PAC-Polynomials-Operations*
imports *PAC-Polynomials-Term PAC-Checker-Specification*

begin

9 Polynomialss as Lists

9.1 Addition

In this section, we refine the polynomials to list. These lists will be used in our checker to represent the polynomials and execute operations.

There is one *key* difference between the list representation and the usual representation: in the former, coefficients can be zero and monomials can appear several times. This makes it easier to reason on intermediate representation where this has not yet been sanitized.

```
fun add-poly-l' :: ‹llist-polynom × llist-polynom ⇒ llist-polynom› where
  ‹add-poly-l' (p, []) = p› |
  ‹add-poly-l' ([], q) = q› |
  ‹add-poly-l' ((xs, n) # p, (ys, m) # q) =
          (if xs = ys then if n + m = 0 then add-poly-l' (p, q) else
               let pq = add-poly-l' (p, q) in
               ((xs, n + m) # pq)
          else if (xs, ys) ∈ term-order-rel
            then
               let pq = add-poly-l' (p, (ys, m) # q) in
               ((xs, n) # pq)
          else
               let pq = add-poly-l' ((xs, n) # p, q) in
               ((ys, m) # pq)
          )›

definition add-poly-l :: ‹llist-polynom × llist-polynom ⇒ llist-polynom nres› where
  ‹add-poly-l = REC_T
      (λadd-poly-l (p, q).
        case (p,q) of
          (p, []) ⇒ RETURN p
        | ([], q) ⇒ RETURN q
        | ((xs, n) # p, (ys, m) # q) ⇒
          (if xs = ys then if n + m = 0 then add-poly-l (p, q) else
             do {
               pq ← add-poly-l (p, q);
               RETURN ((xs, n + m) # pq)
            }
          else if (xs, ys) ∈ term-order-rel
            then do {
               pq ← add-poly-l (p, (ys, m) # q);
               RETURN ((xs, n) # pq)
          }
          else do {
               pq ← add-poly-l ((xs, n) # p, q);
               RETURN ((ys, m) # pq)
          }))›

definition nonzero-coeffs where
  ‹nonzero-coeffs a ⟷ 0 ∉# snd ‘# a›

lemma nonzero-coeffs-simps[simp]:
  ‹nonzero-coeffs {#}›
```

‹*nonzero-coeffs* (*add-mset* (*xs*, *n*) *a*) $\longleftrightarrow$ *nonzero-coeffs a* $\wedge$ *n* $\neq$ *0*›
⟨*proof*⟩

lemma *nonzero-coeffsD*:
‹*nonzero-coeffs a* $\Longrightarrow$ (*x*, *n*) $\in\#$ *a* $\Longrightarrow$ *n* $\neq$ *0*›
⟨*proof*⟩

lemma *sorted-poly-list-rel-ConsD*:
‹((*ys*, *n*) # *p*, *a*) $\in$ *sorted-poly-list-rel S* $\Longrightarrow$ (*p*, *remove1-mset* (*mset ys*, *n*) *a*) $\in$ *sorted-poly-list-rel S* $\wedge$
(*mset ys*, *n*) $\in\#$ *a* $\wedge$ ($\forall$ *x* $\in$ *set p*. *S ys* (*fst x*)) $\wedge$ *sorted-wrt* (*rel2p var-order-rel*) *ys* $\wedge$
distinct ys $\wedge$ *ys* $\notin$ *set* (*map fst p*) $\wedge$ *n* $\neq$ *0* $\wedge$ *nonzero-coeffs a*›
⟨*proof*⟩

lemma *sorted-poly-list-rel-Cons-iff*:
‹((*ys*, *n*) # *p*, *a*) $\in$ *sorted-poly-list-rel S* $\longleftrightarrow$ (*p*, *remove1-mset* (*mset ys*, *n*) *a*) $\in$ *sorted-poly-list-rel S* $\wedge$
(*mset ys*, *n*) $\in\#$ *a* $\wedge$ ($\forall$ *x* $\in$ *set p*. *S ys* (*fst x*)) $\wedge$ *sorted-wrt* (*rel2p var-order-rel*) *ys* $\wedge$
distinct ys $\wedge$ *ys* $\notin$ *set* (*map fst p*) $\wedge$ *n* $\neq$ *0* $\wedge$ *nonzero-coeffs a*›
⟨*proof*⟩

lemma *sorted-repeat-poly-list-rel-ConsD*:
‹((*ys*, *n*) # *p*, *a*) $\in$ *sorted-repeat-poly-list-rel S* $\Longrightarrow$ (*p*, *remove1-mset* (*mset ys*, *n*) *a*) $\in$ *sorted-repeat-poly-list-rel S* $\wedge$
(*mset ys*, *n*) $\in\#$ *a* $\wedge$ ($\forall$ *x* $\in$ *set p*. *S ys* (*fst x*)) $\wedge$ *sorted-wrt* (*rel2p var-order-rel*) *ys* $\wedge$
distinct ys $\wedge$ *n* $\neq$ *0* $\wedge$ *nonzero-coeffs a*›
⟨*proof*⟩

lemma *sorted-repeat-poly-list-rel-Cons-iff*:
‹((*ys*, *n*) # *p*, *a*) $\in$ *sorted-repeat-poly-list-rel S* $\longleftrightarrow$ (*p*, *remove1-mset* (*mset ys*, *n*) *a*) $\in$ *sorted-repeat-poly-list-rel S* $\wedge$
(*mset ys*, *n*) $\in\#$ *a* $\wedge$ ($\forall$ *x* $\in$ *set p*. *S ys* (*fst x*)) $\wedge$ *sorted-wrt* (*rel2p var-order-rel*) *ys* $\wedge$
distinct ys $\wedge$ *n* $\neq$ *0* $\wedge$ *nonzero-coeffs a*›
⟨*proof*⟩

lemma *add-poly-p-add-mset-sum-0*:
‹*n* + *m* = *0* $\Longrightarrow$*add-poly-p*** (*A*, *Aa*, {#}) ({#}, {#}, *r*) $\Longrightarrow$
add-poly-p**
(*add-mset* (*mset ys*, *n*) *A*, *add-mset* (*mset ys*, *m*) *Aa*, {#})
({#}, {#}, *r*)›
⟨*proof*⟩

lemma *monoms-add-poly-l′D*:
‹(*aa*, *ba*) $\in$ *set* (*add-poly-l′ x*) $\Longrightarrow$ *aa* $\in$ *fst* ‘ *set* (*fst x*) $\vee$ *aa* $\in$ *fst* ‘ *set* (*snd x*)›
⟨*proof*⟩

lemma *add-poly-p-add-to-result*:
‹*add-poly-p*** (*A*, *B*, *r*) (*A′*, *B′*, *r′*) $\Longrightarrow$
add-poly-p**
(*A*, *B*, *p* + *r*) (*A′*, *B′*, *p* + *r′*)›
⟨*proof*⟩

lemma *add-poly-p-add-mset-comb*:
‹$add\text{-}poly\text{-}p^{**}$ $(A,\ Aa,\ \{\#\})$ $(\{\#\},\ \{\#\},\ r)$ $\Longrightarrow$
$add\text{-}poly\text{-}p^{**}$
$(add\text{-}mset\ (xs,\ n)\ A,\ Aa,\ \{\#\})$
$(\{\#\},\ \{\#\},\ add\text{-}mset\ (xs,\ n)\ r)$›
⟨*proof*⟩

lemma *add-poly-p-add-mset-comb2*:
‹$add\text{-}poly\text{-}p^{**}$ $(A,\ Aa,\ \{\#\})$ $(\{\#\},\ \{\#\},\ r)$ $\Longrightarrow$
$add\text{-}poly\text{-}p^{**}$
$(add\text{-}mset\ (ys,\ n)\ A,\ add\text{-}mset\ (ys,\ m)\ Aa,\ \{\#\})$
$(\{\#\},\ \{\#\},\ add\text{-}mset\ (ys,\ n\ +\ m)\ r)$›
⟨*proof*⟩

lemma *add-poly-p-add-mset-comb3*:
‹$add\text{-}poly\text{-}p^{**}$ $(A,\ Aa,\ \{\#\})$ $(\{\#\},\ \{\#\},\ r)$ $\Longrightarrow$
$add\text{-}poly\text{-}p^{**}$
$(A,\ add\text{-}mset\ (ys,\ m)\ Aa,\ \{\#\})$
$(\{\#\},\ \{\#\},\ add\text{-}mset\ (ys,\ m)\ r)$›
⟨*proof*⟩

lemma *total-on-lexord*:
‹*Relation.total-on UNIV R* $\Longrightarrow$ *Relation.total-on UNIV* (*lexord R*)›
⟨*proof*⟩

lemma *antisym-lexord*:
‹*antisym R* $\Longrightarrow$ *irrefl R* $\Longrightarrow$ *antisym* (*lexord R*)›
⟨*proof*⟩

lemma *less-than-char-linear*:
‹$(a,\ b) \in less\text{-}than\text{-}char\ \vee$
$a\ =\ b\ \vee\ (b,\ a) \in less\text{-}than\text{-}char$›
⟨*proof*⟩

lemma *total-on-lexord-less-than-char-linear*:
‹$xs \neq ys \Longrightarrow (xs,\ ys) \notin lexord\ (lexord\ less\text{-}than\text{-}char) \longleftrightarrow$
$(ys,\ xs) \in lexord\ (lexord\ less\text{-}than\text{-}char)$›
⟨*proof*⟩

lemma *sorted-poly-list-rel-nonzeroD*:
‹$(p,\ r) \in sorted\text{-}poly\text{-}list\text{-}rel\ term\text{-}order \Longrightarrow$
$nonzero\text{-}coeffs\ (r)$›
‹$(p,\ r) \in sorted\text{-}poly\text{-}list\text{-}rel\ (rel2p\ (lexord\ (lexord\ less\text{-}than\text{-}char))) \Longrightarrow$
$nonzero\text{-}coeffs\ (r)$›
⟨*proof*⟩

lemma *add-poly-l′-add-poly-p*:
assumes ‹$(pq,\ pq') \in sorted\text{-}poly\text{-}rel \times_r sorted\text{-}poly\text{-}rel$›
shows ‹$\exists\, r.\ (add\text{-}poly\text{-}l'\ pq,\ r) \in sorted\text{-}poly\text{-}rel\ \wedge$
$add\text{-}poly\text{-}p^{**}\ (fst\ pq',\ snd\ pq',\ \{\#\})\ (\{\#\},\ \{\#\},\ r)$›
⟨*proof*⟩

lemma *add-poly-l-add-poly*:
  ‹*add-poly-l* x = *RETURN* (*add-poly-l′* x)›
  ⟨*proof*⟩

lemma *add-poly-l-spec*:
  ‹(*add-poly-l*, *uncurry* ($\lambda p\ q$. *SPEC*(λr. *add-poly-p*** (p, q, {#}) ({#}, {#}, r)))) ∈
    sorted-poly-rel $\times_r$ *sorted-poly-rel* $\rightarrow_f$ ⟨*sorted-poly-rel*⟩*nres-rel*›
  ⟨*proof*⟩

definition *sort-poly-spec* :: ‹*llist-polynom* ⇒ *llist-polynom nres*› **where**
‹*sort-poly-spec* p =
  SPEC($\lambda p'$. *mset* p = *mset* p' ∧ *sorted-wrt* (*rel2p* (*Id* ∪ *term-order-rel*)) (*map fst* p'))›

lemma *sort-poly-spec-id*:
  assumes ‹(p, p') ∈ *unsorted-poly-rel*›
  shows ‹*sort-poly-spec* p ≤ ⇓ (*sorted-repeat-poly-rel*) (*RETURN* p')›
⟨*proof*⟩

9.2 Multiplication

fun *mult-monoms* :: ‹*term-poly-list* ⇒ *term-poly-list* ⇒ *term-poly-list*› **where**
  ‹*mult-monoms* p [] = p› |
  ‹*mult-monoms* [] p = p› |
  ‹*mult-monoms* (x # p) (y # q) =
    (*if* $x = y$ *then* x # *mult-monoms* p q
     else if (x, y) ∈ *var-order-rel then* x # *mult-monoms* p (y # q)
      else y # *mult-monoms* (x # p) q)›

lemma *term-poly-list-rel-empty-iff* [*simp*]:
  ‹([], q') ∈ *term-poly-list-rel* ⟷ q' = {#}›
  ⟨*proof*⟩

lemma *term-poly-list-rel-Cons-iff*:
  ‹(y # p, p') ∈ *term-poly-list-rel* ⟷
    (p, *remove1-mset* y p') ∈ *term-poly-list-rel* ∧
    y ∈# p' ∧ y ∉ *set* p ∧ y ∉# *remove1-mset* y p' ∧
    (∀ x∈#*mset* p. (y, x) ∈ *var-order-rel*)›
  ⟨*proof*⟩

lemma *var-order-rel-antisym*[*simp*]:
  ‹(y, y) ∉ *var-order-rel*›
  ⟨*proof*⟩

lemma *term-poly-list-rel-remdups-mset*:
  ‹(p, p') ∈ *term-poly-list-rel* ⟹
    (p, *remdups-mset* p') ∈ *term-poly-list-rel*›
  ⟨*proof*⟩

lemma *var-notin-notin-mult-monomsD*:
  ‹y ∈ *set* (*mult-monoms* p q) ⟹ y ∈ *set* p ∨ y ∈ *set* q›
  ⟨*proof*⟩

lemma *term-poly-list-rel-set-mset*:
  ‹(p, q) ∈ *term-poly-list-rel* ⟹ *set* p = *set-mset* q›
  ⟨*proof*⟩

lemma *mult-monoms-spec*:
‹$(mult\text{-}monoms, (\lambda a\ b.\ remdups\text{-}mset\ (a + b))) \in term\text{-}poly\text{-}list\text{-}rel \to term\text{-}poly\text{-}list\text{-}rel \to term\text{-}poly\text{-}list\text{-}rel$›
⟨*proof*⟩

definition *mult-monomials* :: ‹$term\text{-}poly\text{-}list \times int \Rightarrow term\text{-}poly\text{-}list \times int \Rightarrow term\text{-}poly\text{-}list \times int$› **where**
‹$mult\text{-}monomials = (\lambda(x,\ a)\ (y,\ b).\ (mult\text{-}monoms\ x\ y,\ a * b))$›

definition *mult-poly-raw* :: ‹$llist\text{-}polynom \Rightarrow llist\text{-}polynom \Rightarrow llist\text{-}polynom$› **where**
‹$mult\text{-}poly\text{-}raw\ p\ q = foldl\ (\lambda b\ x.\ map\ (mult\text{-}monomials\ x)\ q\ @\ b)\ []\ p$›

fun *map-append* **where**
‹$map\text{-}append\ f\ b\ [] = b$› |
‹$map\text{-}append\ f\ b\ (x\ \#\ xs) = f\ x\ \#\ map\text{-}append\ f\ b\ xs$›

lemma *map-append-alt-def*:
‹$map\text{-}append\ f\ b\ xs = map\ f\ xs\ @\ b$›
⟨*proof*⟩

lemma *foldl-append-empty*:
‹$NO\text{-}MATCH\ []\ xs \Longrightarrow foldl\ (\lambda b\ x.\ f\ x\ @\ b)\ xs\ p = foldl\ (\lambda b\ x.\ f\ x\ @\ b)\ []\ p\ @\ xs$›
⟨*proof*⟩

lemma *poly-list-rel-empty-iff*[*simp*]:
‹$([],\ r) \in poly\text{-}list\text{-}rel\ R \longleftrightarrow r = \{\#\}$›
⟨*proof*⟩

lemma *mult-poly-raw-simp*[*simp*]:
‹$mult\text{-}poly\text{-}raw\ []\ q = []$›
‹$mult\text{-}poly\text{-}raw\ (x\ \#\ p)\ q = mult\text{-}poly\text{-}raw\ p\ q\ @\ map\ (mult\text{-}monomials\ x)\ q$›
⟨*proof*⟩

lemma *sorted-poly-list-relD*:
‹$(q,\ q') \in sorted\text{-}poly\text{-}list\text{-}rel\ R \Longrightarrow q' = (\lambda(a,\ b).\ (mset\ a,\ b))\ `\#\ mset\ q$›
⟨*proof*⟩

lemma *list-all2-in-set-ExD*:
‹$list\text{-}all2\ R\ p\ q \Longrightarrow x \in set\ p \Longrightarrow \exists\, y \in set\ q.\ R\ x\ y$›
⟨*proof*⟩

inductive-cases *mult-poly-p-elim*: ‹$mult\text{-}poly\text{-}p\ q\ (A,\ r)\ (B,\ r')$›

lemma *mult-poly-p-add-mset-same*:
‹$(mult\text{-}poly\text{-}p\ q')^{**}\ (A,\ r)\ (B,\ r') \Longrightarrow (mult\text{-}poly\text{-}p\ q')^{**}\ (add\text{-}mset\ x\ A,\ r)\ (add\text{-}mset\ x\ B,\ r')$›
⟨*proof*⟩

lemma *mult-poly-raw-mult-poly-p*:
assumes ‹$(p,\ p') \in sorted\text{-}poly\text{-}rel$› **and** ‹$(q,\ q') \in sorted\text{-}poly\text{-}rel$›
shows ‹$\exists\, r.\ (mult\text{-}poly\text{-}raw\ p\ q,\ r) \in unsorted\text{-}poly\text{-}rel \wedge (mult\text{-}poly\text{-}p\ q')^{**}\ (p',\ \{\#\})\ (\{\#\},\ r)$›
⟨*proof*⟩

fun *merge-coeffs* :: ‹$llist\text{-}polynom \Rightarrow llist\text{-}polynom$› **where**
‹$merge\text{-}coeffs\ [] = []$› |
‹$merge\text{-}coeffs\ [(xs,\ n)] = [(xs,\ n)]$› |

‹merge-coeffs ((xs, n) # (ys, m) # p) =
  (if xs = ys
  then if n + m ≠ 0 then merge-coeffs ((xs, n + m) # p) else merge-coeffs p
  else (xs, n) # merge-coeffs ((ys, m) # p))›

abbreviation (**in** −)*mononoms* :: ‹llist-polynom ⇒ term-poly-list set› **where**
‹mononoms p ≡ fst ‘set p›

lemma *fst-normalize-polynom-subset*:
‹mononoms (merge-coeffs p) ⊆ mononoms p›
⟨proof⟩

lemma *fst-normalize-polynom-subsetD*:
‹(a, b) ∈ set (merge-coeffs p) ⟹ a ∈ mononoms p›
⟨proof⟩

lemma *distinct-merge-coeffs*:
assumes ‹sorted-wrt R (map fst xs)› **and** ‹transp R› ‹antisymp R›
shows ‹distinct (map fst (merge-coeffs xs))›
⟨proof⟩

lemma *in-set-merge-coeffsD*:
‹(a, b) ∈ set (merge-coeffs p) ⟹∃ b. (a, b) ∈ set p›
⟨proof⟩

lemma *rtranclp-normalize-poly-add-mset*:
‹normalize-poly-p** A r ⟹ normalize-poly-p** (add-mset x A) (add-mset x r)›
⟨proof⟩

lemma *nonzero-coeffs-diff*:
‹nonzero-coeffs A ⟹ nonzero-coeffs (A − B)›
⟨proof⟩

lemma *merge-coeffs-is-normalize-poly-p*:
‹(xs, ys) ∈ sorted-repeat-poly-rel ⟹ ∃ r. (merge-coeffs xs, r) ∈ sorted-poly-rel ∧ normalize-poly-p** ys r›
⟨proof⟩

9.3 Normalisation

definition *normalize-poly* **where**
‹normalize-poly p = do {
  p ← sort-poly-spec p;
  RETURN (merge-coeffs p)
}›

definition *sort-coeff* :: ‹string list ⇒ string list nres› **where**
‹sort-coeff ys = SPEC(λxs. mset xs = mset ys ∧ sorted-wrt (rel2p (Id ∪ var-order-rel)) xs)›

lemma *distinct-var-order-Id-var-order*:
‹distinct a ⟹ sorted-wrt (rel2p (Id ∪ var-order-rel)) a ⟹
    sorted-wrt var-order a›
⟨proof⟩

definition *sort-all-coeffs* :: ‹*llist-polynom* ⇒ *llist-polynom nres*› **where**
‹*sort-all-coeffs xs* = *monadic-nfoldli xs* (λ-. *RETURN True*) (λ(*a*, *n*) *b*. *do* {*a* ← *sort-coeff a*; *RETURN* ((*a*, *n*) # *b*)}) []›

lemma *sort-all-coeffs-gen*:
  assumes ‹(∀ *xs* ∈ *mononoms xs′*. *sorted-wrt* (*rel2p* (*var-order-rel*)) *xs*)› **and**
    ‹∀ *x* ∈ *mononoms* (*xs* @ *xs′*). *distinct x*›
  shows ‹*monadic-nfoldli xs* (λ-. *RETURN True*) (λ(*a*, *n*) *b*. *do* {*a* ← *sort-coeff a*; *RETURN* ((*a*, *n*) # *b*)}) *xs′* ≤
    ⇓*Id* (*SPEC*(λ*ys*. *map* (λ(*a*,*b*). (*mset a*, *b*)) (*rev xs* @ *xs′*) = *map* (λ(*a*,*b*). (*mset a*, *b*)) (*ys*) ∧
    (∀ *xs* ∈ *mononoms ys*. *sorted-wrt* (*rel2p* (*var-order-rel*)) *xs*)))›
  ⟨*proof*⟩

definition *shuffle-coefficients* **where**
  ‹*shuffle-coefficients xs* = (*SPEC*(λ*ys*. *map* (λ(*a*,*b*). (*mset a*, *b*)) (*rev xs*) = *map* (λ(*a*,*b*). (*mset a*, *b*)) *ys* ∧
    (∀ *xs* ∈ *mononoms ys*. *sorted-wrt* (*rel2p* (*var-order-rel*)) *xs*)))›

lemma *sort-all-coeffs*:
  ‹∀ *x* ∈ *mononoms xs*. *distinct x* ⟹
  sort-all-coeffs xs ≤ ⇓ *Id* (*shuffle-coefficients xs*)›
  ⟨*proof*⟩

lemma *unsorted-term-poly-list-rel-mset*:
  ‹(*ys*, *aa*) ∈ *unsorted-term-poly-list-rel* ⟹ *mset ys* = *aa*›
  ⟨*proof*⟩

lemma *RETURN-map-alt-def*:
  ‹*RETURN o* (*map f*) =
    REC_T (λ*g xs*.
      case xs of
        [] ⇒ *RETURN* []
      | *x* # *xs* ⇒ *do* {*xs* ← *g xs*; *RETURN* (*f x* # *xs*)})›
  ⟨*proof*⟩

lemma *fully-unsorted-poly-rel-Cons-iff*:
  ‹((*ys*, *n*) # *p*, *a*) ∈ *fully-unsorted-poly-rel* ⟷
    (*p*, *remove1-mset* (*mset ys*, *n*) *a*) ∈ *fully-unsorted-poly-rel* ∧
    (*mset ys*, *n*) ∈# *a* ∧ *distinct ys*›
  ⟨*proof*⟩

lemma *map-mset-unsorted-term-poly-list-rel*:
  ‹(⋀*a*. *a* ∈ *mononoms s* ⟹ *distinct a*) ⟹ ∀ *x* ∈ *mononoms s*. *distinct x* ⟹
    (∀ *xs* ∈ *mononoms s*. *sorted-wrt* (*rel2p* (*Id* ∪ *var-order-rel*)) *xs*) ⟹
    (*s*, *map* (λ(*a*, *y*). (*mset a*, *y*)) *s*)
        ∈ ⟨*term-poly-list-rel* $\times_r$ *int-rel*⟩*list-rel*›
  ⟨*proof*⟩

lemma *list-rel-unsorted-term-poly-list-relD*:
  ‹(*p*, *y*) ∈ ⟨*unsorted-term-poly-list-rel* $\times_r$ *int-rel*⟩*list-rel* ⟹
   mset y = (λ(*a*, *y*). (*mset a*, *y*)) ‘# *mset p* ∧ (∀ *x* ∈ *mononoms p*. *distinct x*)›
  ⟨*proof*⟩

lemma *shuffle-terms-distinct-iff*:

assumes ‹$map\ (\lambda(a,\ y).\ (mset\ a,\ y))\ p = map\ (\lambda(a,\ y).\ (mset\ a,\ y))\ s$›
shows ‹$(\forall x{\in}set\ p.\ distinct\ (fst\ x)) \longleftrightarrow (\forall x{\in}set\ s.\ distinct\ (fst\ x))$›
⟨*proof*⟩

lemma
‹$(p,\ y) \in \langle unsorted\text{-}term\text{-}poly\text{-}list\text{-}rel \times_r int\text{-}rel\rangle list\text{-}rel \Longrightarrow$
$(a,\ b) \in set\ p \Longrightarrow\ distinct\ a$›
⟨*proof*⟩

lemma *sort-all-coeffs-unsorted-poly-rel-with0*:
assumes ‹$(p,\ p') \in fully\text{-}unsorted\text{-}poly\text{-}rel$›
shows ‹$sort\text{-}all\text{-}coeffs\ p \le \Downarrow (unsorted\text{-}poly\text{-}rel\text{-}with0)\ (RETURN\ p')$›
⟨*proof*⟩

lemma *sort-poly-spec-id′*:
assumes ‹$(p,\ p') \in unsorted\text{-}poly\text{-}rel\text{-}with0$›
shows ‹$sort\text{-}poly\text{-}spec\ p \le \Downarrow (sorted\text{-}repeat\text{-}poly\text{-}rel\text{-}with0)\ (RETURN\ p')$›
⟨*proof*⟩

fun *merge-coeffs0* :: ‹*llist-polynom* ⇒ *llist-polynom*› **where**
‹$merge\text{-}coeffs0\ [] = []$› |
‹$merge\text{-}coeffs0\ [(xs,\ n)] = (if\ n = 0\ then\ []\ else\ [(xs,\ n)])$› |
‹$merge\text{-}coeffs0\ ((xs,\ n)\ \#\ (ys,\ m)\ \#\ p) =$
$(if\ xs = ys$
$then\ if\ n + m \ne 0\ then\ merge\text{-}coeffs0\ ((xs,\ n + m)\ \#\ p)\ else\ merge\text{-}coeffs0\ p$
$else\ if\ n = 0\ then\ merge\text{-}coeffs0\ ((ys,\ m)\ \#\ p)$
$else(xs,\ n)\ \#\ merge\text{-}coeffs0\ ((ys,\ m)\ \#\ p))$›

lemma *sorted-repeat-poly-list-rel-with0-wrt-ConsD*:
‹$((ys,\ n)\ \#\ p,\ a) \in sorted\text{-}repeat\text{-}poly\text{-}list\text{-}rel\text{-}with0\text{-}wrt\ S\ term\text{-}poly\text{-}list\text{-}rel \Longrightarrow$
$(p,\ remove1\text{-}mset\ (mset\ ys,\ n)\ a) \in sorted\text{-}repeat\text{-}poly\text{-}list\text{-}rel\text{-}with0\text{-}wrt\ S\ term\text{-}poly\text{-}list\text{-}rel \wedge$
$(mset\ ys,\ n) \in\#\ a \wedge (\forall x \in set\ p.\ S\ ys\ (fst\ x)) \wedge sorted\text{-}wrt\ (rel2p\ var\text{-}order\text{-}rel)\ ys \wedge$
$distinct\ ys$›
⟨*proof*⟩

lemma *sorted-repeat-poly-list-rel-with0-wrtl-Cons-iff*:
‹$((ys,\ n)\ \#\ p,\ a) \in sorted\text{-}repeat\text{-}poly\text{-}list\text{-}rel\text{-}with0\text{-}wrt\ S\ term\text{-}poly\text{-}list\text{-}rel \longleftrightarrow$
$(p,\ remove1\text{-}mset\ (mset\ ys,\ n)\ a) \in sorted\text{-}repeat\text{-}poly\text{-}list\text{-}rel\text{-}with0\text{-}wrt\ S\ term\text{-}poly\text{-}list\text{-}rel \wedge$
$(mset\ ys,\ n) \in\#\ a \wedge (\forall x \in set\ p.\ S\ ys\ (fst\ x)) \wedge sorted\text{-}wrt\ (rel2p\ var\text{-}order\text{-}rel)\ ys \wedge$
$distinct\ ys$›
⟨*proof*⟩

lemma *fst-normalize0-polynom-subsetD*:
‹$(a,\ b) \in set\ (merge\text{-}coeffs0\ p) \Longrightarrow a \in mononoms\ p$›
⟨*proof*⟩

lemma *in-set-merge-coeffs0D*:
‹$(a,\ b) \in set\ (merge\text{-}coeffs0\ p) \Longrightarrow \exists b.\ (a,\ b) \in set\ p$›
⟨*proof*⟩

lemma *merge-coeffs0-is-normalize-poly-p*:
‹$(xs,\ ys) \in sorted\text{-}repeat\text{-}poly\text{-}rel\text{-}with0 \Longrightarrow \exists r.\ (merge\text{-}coeffs0\ xs,\ r) \in sorted\text{-}poly\text{-}rel \wedge normalize\text{-}poly\text{-}p^{**}$

ys r›
⟨*proof*⟩

definition *full-normalize-poly* **where**
‹*full-normalize-poly p* = *do* {
p ← *sort-all-coeffs p*;
p ← *sort-poly-spec p*;
RETURN (*merge-coeffs0 p*)
}›

fun *sorted-remdups* **where**
‹*sorted-remdups* (*x* # *y* # *zs*) =
(*if x* = *y then sorted-remdups* (*y* # *zs*) *else x* # *sorted-remdups* (*y* # *zs*))› |
‹*sorted-remdups zs* = *zs*›

lemma *set-sorted-remdups*[*simp*]:
‹*set* (*sorted-remdups xs*) = *set xs*›
⟨*proof*⟩

lemma *distinct-sorted-remdups*:
‹*sorted-wrt R xs* $\Longrightarrow$ *transp R* $\Longrightarrow$ *Restricted-Predicates.total-on R UNIV* $\Longrightarrow$
antisymp R $\Longrightarrow$ *distinct* (*sorted-remdups xs*)›
⟨*proof*⟩

lemma *full-normalize-poly-normalize-poly-p*:
assumes ‹(p, p') ∈ *fully-unsorted-poly-rel*›
shows ‹*full-normalize-poly p* ≤ ⇓ (*sorted-poly-rel*) (*SPEC* (λ*r*. *normalize-poly-p*** p' *r*))›
(**is** ‹*?A* ≤ ⇓ *?R ?B*›)
⟨*proof*⟩

definition *mult-poly-full* :: ‹-› **where**
‹*mult-poly-full p q* = *do* {
let pq = *mult-poly-raw p q*;
normalize-poly pq
}›

lemma *normalize-poly-normalize-poly-p*:
assumes ‹(p, p') ∈ *unsorted-poly-rel*›
shows ‹*normalize-poly p* ≤ ⇓ (*sorted-poly-rel*) (*SPEC* (λ*r*. *normalize-poly-p*** p' *r*))›
⟨*proof*⟩

9.4 Multiplication and normalisation

definition *mult-poly-p'* :: ‹-› **where**
‹*mult-poly-p'* p' q' = *do* {
pq ← *SPEC*(λ*r*. (*mult-poly-p* q')** (p', {#}) ({#}, *r*));
SPEC (λ*r*. *normalize-poly-p*** *pq r*)
}›

lemma *unsorted-poly-rel-fully-unsorted-poly-rel*:
‹*unsorted-poly-rel* ⊆ *fully-unsorted-poly-rel*›
⟨*proof*⟩

lemma *mult-poly-full-mult-poly-p'*:
assumes ‹(p, p') ∈ *sorted-poly-rel*› ‹(q, q') ∈ *sorted-poly-rel*›
shows ‹*mult-poly-full p q* ≤ ⇓ (*sorted-poly-rel*) (*mult-poly-p'* p' q')›

$\langle proof \rangle$

definition *add-poly-spec* :: ‹-› **where**
‹$add\text{-}poly\text{-}spec\ p\ q = SPEC\ (\lambda r.\ p + q - r \in ideal\ polynom\text{-}bool)$›

definition *add-poly-p′* :: ‹-› **where**
‹$add\text{-}poly\text{-}p'\ p\ q = SPEC(\lambda r.\ add\text{-}poly\text{-}p^{**}\ (p,\ q,\ \{\#\})\ (\{\#\},\ \{\#\},\ r))$›

lemma *add-poly-l-add-poly-p′*:
assumes ‹$(p,\ p') \in sorted\text{-}poly\text{-}rel$› ‹$(q,\ q') \in sorted\text{-}poly\text{-}rel$›
shows ‹$add\text{-}poly\text{-}l\ (p,\ q) \leq \Downarrow\ (sorted\text{-}poly\text{-}rel)\ (add\text{-}poly\text{-}p'\ p'\ q')$›
$\langle proof \rangle$

9.5 Correctness

context *poly-embed*
begin

definition *mset-poly-rel* **where**
‹$mset\text{-}poly\text{-}rel = \{(a,\ b).\ b = polynom\text{-}of\text{-}mset\ a\}$›

definition *var-rel* **where**
‹$var\text{-}rel = br\ \varphi\ (\lambda\text{-}.\ True)$›

lemma *normalize-poly-p-normalize-poly-spec*:
‹$(p,\ p') \in mset\text{-}poly\text{-}rel \Longrightarrow$
$SPEC\ (\lambda r.\ normalize\text{-}poly\text{-}p^{**}\ p\ r) \leq \Downarrow mset\text{-}poly\text{-}rel\ (normalize\text{-}poly\text{-}spec\ p')$›
$\langle proof \rangle$

lemma *mult-poly-p′-mult-poly-spec*:
‹$(p,\ p') \in mset\text{-}poly\text{-}rel \Longrightarrow (q,\ q') \in mset\text{-}poly\text{-}rel \Longrightarrow$
$mult\text{-}poly\text{-}p'\ p\ q \leq \Downarrow mset\text{-}poly\text{-}rel\ (mult\text{-}poly\text{-}spec\ p'\ q')$›
$\langle proof \rangle$

lemma *add-poly-p′-add-poly-spec*:
‹$(p,\ p') \in mset\text{-}poly\text{-}rel \Longrightarrow (q,\ q') \in mset\text{-}poly\text{-}rel \Longrightarrow$
$add\text{-}poly\text{-}p'\ p\ q \leq \Downarrow mset\text{-}poly\text{-}rel\ (add\text{-}poly\text{-}spec\ p'\ q')$›
$\langle proof \rangle$

end

definition *weak-equality-l* :: ‹$llist\text{-}polynom \Rightarrow llist\text{-}polynom \Rightarrow bool\ nres$› **where**
‹$weak\text{-}equality\text{-}l\ p\ q = RETURN\ (p = q)$›

definition *weak-equality* :: ‹$int\ mpoly \Rightarrow int\ mpoly \Rightarrow bool\ nres$› **where**
‹$weak\text{-}equality\ p\ q = SPEC\ (\lambda r.\ r \longrightarrow p = q)$›

definition *weak-equality-spec* :: ‹$mset\text{-}polynom \Rightarrow mset\text{-}polynom \Rightarrow bool\ nres$› **where**
‹$weak\text{-}equality\text{-}spec\ p\ q = SPEC\ (\lambda r.\ r \longrightarrow p = q)$›

lemma *term-poly-list-rel-same-rightD*:
‹$(a,\ aa) \in term\text{-}poly\text{-}list\text{-}rel \Longrightarrow (a,\ ab) \in term\text{-}poly\text{-}list\text{-}rel \Longrightarrow aa = ab$›
$\langle proof \rangle$

lemma *list-rel-term-poly-list-rel-same-rightD*:
‹$(xa, y) \in \langle$*term-poly-list-rel* $\times_r$ *int-rel*$\rangle$*list-rel* $\Longrightarrow$
$(xa, ya) \in \langle$*term-poly-list-rel* $\times_r$ *int-rel*$\rangle$*list-rel* $\Longrightarrow$
$y = ya$›
⟨*proof*⟩

lemma *weak-equality-l-weak-equality-spec*:
‹(*uncurry weak-equality-l*, *uncurry weak-equality-spec*) $\in$
sorted-poly-rel $\times_r$ *sorted-poly-rel* $\rightarrow_f$ ⟨*bool-rel*⟩*nres-rel*›
⟨*proof*⟩

end
theory *PAC-Checker*
imports *PAC-Polynomials-Operations*
PAC-Checker-Specification
PAC-Map-Rel
Show.Show
Show.Show-Instances
begin

10 Executable Checker

In this layer we finally refine the checker to executable code.

10.1 Definitions

Compared to the previous layer, we add an error message when an error is discovered. We do not attempt to prove anything on the error message (neither that there really is an error, nor that the error message is correct).

Extended error message **datatype** *'a code-status* =
is-cfailed: *CFAILED* (*the-error*: *'a*) |
CSUCCESS |
is-cfound: *CFOUND*

In the following function, we merge errors. We will never merge an error message with an another error message; hence we do not attempt to concatenate error messages.

fun *merge-cstatus* **where**
‹*merge-cstatus* (*CFAILED a*) - = *CFAILED a*› |
‹*merge-cstatus* - (*CFAILED a*) = *CFAILED a*› |
‹*merge-cstatus CFOUND* - = *CFOUND*› |
‹*merge-cstatus* - *CFOUND* = *CFOUND*› |
‹*merge-cstatus* - - = *CSUCCESS*›

definition *code-status-status-rel* :: ‹(*'a code-status* $\times$ *status*) *set*› **where**
‹*code-status-status-rel* =
{(*CFOUND*, *FOUND*), (*CSUCCESS*, *SUCCESS*)} $\cup$
{(*CFAILED a*, *FAILED*)| *a*. *True*}›

lemma *in-code-status-status-rel-iff* [*simp*]:
‹(*CFOUND*, *b*) $\in$ *code-status-status-rel* $\longleftrightarrow$ *b* = *FOUND*›
‹(*a*, *FOUND*) $\in$ *code-status-status-rel* $\longleftrightarrow$ *a* = *CFOUND*›

‹(CSUCCESS, b) ∈ code-status-status-rel ⟷ b = SUCCESS›
‹(a, SUCCESS) ∈ code-status-status-rel ⟷ a = CSUCCESS›
‹(a, FAILED) ∈ code-status-status-rel ⟷ is-cfailed a›
‹(CFAILED C, b) ∈ code-status-status-rel ⟷ b = FAILED›
⟨proof⟩

Refinement relation **fun** *pac-step-rel-raw* :: ‹('olbl × 'lbl) set ⇒ ('a × 'b) set ⇒ ('c × 'd) set ⇒ ('a, 'c, 'olbl) pac-step ⇒ ('b, 'd, 'lbl) pac-step ⇒ bool› **where**
‹pac-step-rel-raw R1 R2 R3 (Add p1 p2 i r) (Add p1' p2' i' r') ⟷
  (p1, p1') ∈ R1 ∧ (p2, p2') ∈ R1 ∧ (i, i') ∈ R1 ∧
  (r, r') ∈ R2› |
‹pac-step-rel-raw R1 R2 R3 (Mult p1 p2 i r) (Mult p1' p2' i' r') ⟷
  (p1, p1') ∈ R1 ∧ (p2, p2') ∈ R2 ∧ (i, i') ∈ R1 ∧
  (r, r') ∈ R2› |
‹pac-step-rel-raw R1 R2 R3 (Del p1) (Del p1') ⟷
  (p1, p1') ∈ R1› |
‹pac-step-rel-raw R1 R2 R3 (Extension i x p1) (Extension j x' p1') ⟷
  (i, j) ∈ R1 ∧ (x, x') ∈ R3 ∧ (p1, p1') ∈ R2› |
‹pac-step-rel-raw R1 R2 R3 - - ⟷ False›

fun *pac-step-rel-assn* :: ‹('olbl ⇒ 'lbl ⇒ assn) ⇒ ('a ⇒ 'b ⇒ assn) ⇒ ('c ⇒ 'd ⇒ assn) ⇒ ('a, 'c, 'olbl) pac-step ⇒ ('b, 'd, 'lbl) pac-step ⇒ assn› **where**
‹pac-step-rel-assn R1 R2 R3 (Add p1 p2 i r) (Add p1' p2' i' r') =
  R1 p1 p1' * R1 p2 p2' * R1 i i' *
  R2 r r'› |
‹pac-step-rel-assn R1 R2 R3 (Mult p1 p2 i r) (Mult p1' p2' i' r') =
  R1 p1 p1' * R2 p2 p2' * R1 i i' *
  R2 r r'› |
‹pac-step-rel-assn R1 R2 R3 (Del p1) (Del p1') =
  R1 p1 p1'› |
‹pac-step-rel-assn R1 R2 R3 (Extension i x p1) (Extension i' x' p1') =
  R1 i i' * R3 x x' * R2 p1 p1'› |
‹pac-step-rel-assn R1 R2 - - - = false›

lemma *pac-step-rel-assn-alt-def*:
‹pac-step-rel-assn R1 R2 R3 x y = (
case (x, y) of
    (Add p1 p2 i r, Add p1' p2' i' r') ⇒
      R1 p1 p1' * R1 p2 p2' * R1 i i' * R2 r r'
  | (Mult p1 p2 i r, Mult p1' p2' i' r') ⇒
      R1 p1 p1' * R2 p2 p2' * R1 i i' * R2 r r'
  | (Del p1, Del p1') ⇒ R1 p1 p1'
  | (Extension i x p1, Extension i' x' p1') ⇒ R1 i i' * R3 x x' * R2 p1 p1'
  | - ⇒ false
)›
  ⟨proof⟩

Addition checking **definition** *error-msg* **where**
‹error-msg i msg = CFAILED (''s CHECKING failed at line '' @ show i @ '' with error '' @ msg)›

definition *error-msg-notin-dom-err* **where**
‹error-msg-notin-dom-err = '' notin domain''›

definition *error-msg-notin-dom* :: ‹nat ⇒ string› **where**
‹error-msg-notin-dom i = show i @ error-msg-notin-dom-err›

definition *error-msg-reused-dom* **where**
⟨*error-msg-reused-dom i = show i @ '' already in domain''*⟩

definition *error-msg-not-equal-dom* **where**
⟨*error-msg-not-equal-dom p q pq r = show p @ '' + '' @ show q @ '' = '' @ show pq @ '' not equal'' @ show r*⟩

definition *check-not-equal-dom-err* :: ⟨*llist-polynom ⇒ llist-polynom ⇒ llist-polynom ⇒ llist-polynom ⇒ string nres*⟩ **where**
⟨*check-not-equal-dom-err p q pq r = SPEC (λ-. True)*⟩

definition *vars-llist* :: ⟨*llist-polynom ⇒ string set*⟩ **where**
⟨*vars-llist xs* = $\bigcup$(*set ' fst ' set xs*)⟩

definition *check-addition-l* :: ⟨*- ⇒ - ⇒ string set ⇒ nat ⇒ nat ⇒ nat ⇒ llist-polynom ⇒ string code-status nres*⟩ **where**
⟨*check-addition-l spec A 𝒱 p q i r = do {*
let b = p ∈# dom-m A ∧ q ∈# dom-m A ∧ i ∉# dom-m A ∧ vars-llist r ⊆ 𝒱;
if ¬b
then RETURN (error-msg i ((if p ∉# dom-m A then error-msg-notin-dom p else []) @ (if q ∉# dom-m A then error-msg-notin-dom p else []) @
(if i ∈# dom-m A then error-msg-reused-dom p else [])))
else do {
ASSERT (p ∈# dom-m A);
let p = the (fmlookup A p);
ASSERT (q ∈# dom-m A);
let q = the (fmlookup A q);
pq ← add-poly-l (p, q);
b ← weak-equality-l pq r;
b′ ← weak-equality-l r spec;
if b then (if b′ then RETURN CFOUND else RETURN CSUCCESS)
else do {
c ← check-not-equal-dom-err p q pq r;
RETURN (error-msg i c)}
}
}⟩

Multiplication checking **definition** *check-mult-l-dom-err* :: ⟨*bool ⇒ nat ⇒ bool ⇒ nat ⇒ string nres*⟩ **where**
⟨*check-mult-l-dom-err p-notin p i-already i = SPEC (λ-. True)*⟩

definition *check-mult-l-mult-err* :: ⟨*llist-polynom ⇒ llist-polynom ⇒ llist-polynom ⇒ llist-polynom ⇒ string nres*⟩ **where**
⟨*check-mult-l-mult-err p q pq r = SPEC (λ-. True)*⟩

definition *check-mult-l* :: ⟨*- ⇒ - ⇒ - ⇒ nat ⇒llist-polynom ⇒ nat ⇒ llist-polynom ⇒ string code-status nres*⟩ **where**
⟨*check-mult-l spec A 𝒱 p q i r = do {*

let $b = p \in\#$ *dom-m* $A \wedge i \notin\#$ *dom-m* $A \wedge$ *vars-llist* $q \subseteq \mathcal{V} \wedge$ *vars-llist* $r \subseteq \mathcal{V}$;
if $\neg b$
then do {
c ← *check-mult-l-dom-err* ($p \notin\#$ *dom-m* A) p ($i \in\#$ *dom-m* A) i;
RETURN (*error-msg i c*)}
else do {
ASSERT ($p \in\#$ *dom-m* A);
let p = *the* (*fmlookup* A p);
pq ← *mult-poly-full p q*;
b ← *weak-equality-l pq r*;
b′ ← *weak-equality-l r spec*;
if b then (*if b′ then RETURN CFOUND else RETURN CSUCCESS*) *else do* {
c ← *check-mult-l-mult-err p q pq r*;
RETURN (*error-msg i c*)
}
}
}›

Deletion checking **definition** *check-del-l* :: ‹- ⇒ - ⇒ *nat* ⇒ *string code-status nres*› **where**
‹*check-del-l spec A p* = *RETURN CSUCCESS*›

Extension checking **definition** *check-extension-l-dom-err* :: ‹*nat* ⇒ *string nres*› **where**
‹*check-extension-l-dom-err p* = *SPEC* (λ-. *True*)›

definition *check-extension-l-no-new-var-err* :: ‹*llist-polynom* ⇒ *string nres*› **where**
‹*check-extension-l-no-new-var-err p* = *SPEC* (λ-. *True*)›

definition *check-extension-l-new-var-multiple-err* :: ‹*string* ⇒ *llist-polynom* ⇒ *string nres*› **where**
‹*check-extension-l-new-var-multiple-err v p* = *SPEC* (λ-. *True*)›

definition *check-extension-l-side-cond-err*
:: ‹*string* ⇒ *llist-polynom* ⇒ *llist-polynom* ⇒ *llist-polynom* ⇒ *string nres*›
where
‹*check-extension-l-side-cond-err v p p′ q* = *SPEC* (λ-. *True*)›

definition *check-extension-l*
:: ‹- ⇒ - ⇒ *string set* ⇒ *nat* ⇒ *string* ⇒ *llist-polynom* ⇒ (*string code-status*) *nres*›
where
‹*check-extension-l spec A* $\mathcal{V}$ *i v p* = *do* {
let $b = i \notin\#$ *dom-m* $A \wedge v \notin \mathcal{V} \wedge ([v], -1) \in$ *set* p;
if $\neg b$
then do {
c ← *check-extension-l-dom-err i*;
RETURN (*error-msg i c*)
} *else do* {
let p' = *remove1* ($[v], -1$) p;
let b = *vars-llist* $p' \subseteq \mathcal{V}$;
if $\neg b$
then do {
c ← *check-extension-l-new-var-multiple-err v p′*;
RETURN (*error-msg i c*)
}
else do {
p2 ← *mult-poly-full p′ p′*;

```
        let p' = map (λ(a,b). (a, −b)) p';
        q ← add-poly-l (p2, p');
        eq ← weak-equality-l q [];
        if eq then do {
          RETURN (CSUCCESS)
        } else do {
         c ← check-extension-l-side-cond-err v p p' q;
         RETURN (error-msg i c)
       }
     }
   }
 }›
```

lemma *check-extension-alt-def*:

```
 ‹check-extension A 𝒱 i v p ≥ do {
   b ← SPEC(λb. b ⟶ i ∉# dom-m A ∧ v ∉ 𝒱);
   if ¬b
   then RETURN (False)
   else do {
        p' ← RETURN (p + Var v);
        b ← SPEC(λb. b ⟶ vars p' ⊆ 𝒱);
        if ¬b
        then RETURN (False)
        else do {
          pq ← mult-poly-spec p' p';
          let p' = − p';
          p ← add-poly-spec pq p';
          eq ← weak-equality p 0;
          if eq then RETURN(True)
          else RETURN (False)
      }
    }
  }›
⟨proof⟩
```

lemma *RES-RES-RETURN-RES*: ‹*RES A* ⤜ (*λT. RES* (*f T*)) = *RES* (⋃(*f* ‘ *A*))›
⟨*proof*⟩

lemma *check-add-alt-def*:

```
 ‹check-add A 𝒱 p q i r ≥
   do {
    b ← SPEC(λb. b ⟶ p ∈# dom-m A ∧ q ∈# dom-m A ∧ i ∉# dom-m A ∧ vars r ⊆ 𝒱);
    if ¬b
    then RETURN False
    else do {
      ASSERT (p ∈# dom-m A);
      let p = the (fmlookup A p);
      ASSERT (q ∈# dom-m A);
      let q = the (fmlookup A q);
      pq ← add-poly-spec p q;
      eq ← weak-equality pq r;
      RETURN eq
```

}
}› (**is** ‹- ≥ *?A*›)
⟨*proof*⟩

lemma *check-mult-alt-def*:
‹*check-mult A* $\mathcal{V}$ *p q i r* ≥
do {
b ← *SPEC*(λ*b*. *b* ⟶ *p* ∈# *dom-m A* ∧ *i* ∉# *dom-m A* ∧ *vars q* ⊆ $\mathcal{V}$ ∧ *vars r* ⊆ $\mathcal{V}$);
if ¬*b*
then RETURN False
else do {
ASSERT (*p* ∈# *dom-m A*);
let p = *the* (*fmlookup A p*);
pq ← *mult-poly-spec p q*;
p ← *weak-equality pq r*;
RETURN p
}
}›
⟨*proof*⟩

primrec *insort-key-rel* :: ('*b* ⇒ '*b* ⇒ *bool*) ⇒ '*b* ⇒ '*b list* ⇒ '*b list* **where**
insort-key-rel f x [] = [*x*] |
insort-key-rel f x (*y*#*ys*) =
(*if f x y then* (*x*#*y*#*ys*) *else y*#(*insort-key-rel f x ys*))

lemma *set-insort-key-rel*[*simp*]: ‹*set* (*insort-key-rel R x xs*) = *insert x* (*set xs*)›
⟨*proof*⟩

lemma *sorted-wrt-insort-key-rel*:
‹*total-on R* (*insert x* (*set xs*)) ⟹ *transp R* ⟹ *reflp R* ⟹
sorted-wrt R xs ⟹ *sorted-wrt R* (*insort-key-rel R x xs*)›
⟨*proof*⟩

lemma *sorted-wrt-insort-key-rel2*:
‹*total-on R* (*insert x* (*set xs*)) ⟹ *transp R* ⟹ *x* ∉ *set xs* ⟹
sorted-wrt R xs ⟹ *sorted-wrt R* (*insort-key-rel R x xs*)›
⟨*proof*⟩

Step checking **definition** *PAC-checker-l-step* :: ‹- ⇒ *string code-status* × *string set* × - ⇒ (*llist-polynom*, *string*, *nat*) *pac-step* ⇒ -› **where**
‹*PAC-checker-l-step* = (λ*spec* (*st'*, $\mathcal{V}$, *A*) *st*. *case st of*
Add - - - - ⇒
do {
r ← *full-normalize-poly* (*pac-res st*);
eq ← *check-addition-l spec A* $\mathcal{V}$ (*pac-src1 st*) (*pac-src2 st*) (*new-id st*) *r*;
let - = *eq*;
if ¬*is-cfailed eq*
then RETURN (*merge-cstatus st' eq*,
$\mathcal{V}$, *fmupd* (*new-id st*) *r A*)
else RETURN (*eq*, $\mathcal{V}$, *A*)
}
| *Del* - ⇒
do {
eq ← *check-del-l spec A* (*pac-src1 st*);
let - = *eq*;

if ¬is-cfailed eq
then RETURN (merge-cstatus st′ eq, $\mathcal{V}$, fmdrop (pac-src1 st) A)
else RETURN (eq, $\mathcal{V}$, A)
}
| *Mult - - - - ⇒*
do {
r ← full-normalize-poly (pac-res st);
q ← full-normalize-poly (pac-mult st);
eq ← check-mult-l spec A $\mathcal{V}$ (pac-src1 st) q (new-id st) r;
let - = eq;
if ¬is-cfailed eq
then RETURN (merge-cstatus st′ eq,
$\mathcal{V}$, fmupd (new-id st) r A)
else RETURN (eq, $\mathcal{V}$, A)
}
| *Extension - - - ⇒*
do {
r ← full-normalize-poly (([new-var st], −1) # (pac-res st));
(eq) ← check-extension-l spec A $\mathcal{V}$ (new-id st) (new-var st) r;
if ¬is-cfailed eq
then do {
RETURN (st′,
insert (new-var st) $\mathcal{V}$, fmupd (new-id st) r A)}
else RETURN (eq, $\mathcal{V}$, A)
}
)›

lemma *pac-step-rel-raw-def*:
‹⟨K, V, R⟩ *pac-step-rel-raw = pac-step-rel-raw K V R*›
⟨*proof*⟩

definition *mononoms-equal-up-to-reorder* **where**
‹*mononoms-equal-up-to-reorder xs ys ⟷*
map (λ(a, b). (mset a, b)) xs = map (λ(a, b). (mset a, b)) ys›

definition *normalize-poly-l* **where**
‹*normalize-poly-l p = SPEC (λp′.*
*normalize-poly-p** ((λ(a, b). (mset a, b)) ‘# mset p) ((λ(a, b). (mset a, b)) ‘# mset p′) ∧*
0 ∉# snd ‘# mset p′ ∧
sorted-wrt (rel2p (term-order-rel $\times_r$ int-rel)) p′ ∧
(∀ x ∈ mononoms p′. sorted-wrt (rel2p var-order-rel) x))›

definition *remap-polys-l-dom-err* :: ‹*string nres*› **where**
‹*remap-polys-l-dom-err = SPEC (λ-. True)*›

definition *remap-polys-l* :: ‹*llist-polynom ⇒ string set ⇒ (nat, llist-polynom) fmap ⇒*
(- code-status × string set × (nat, llist-polynom) fmap) nres› **where**
‹*remap-polys-l spec = (λ$\mathcal{V}$ A. do*{
dom ← SPEC(λdom. set-mset (dom-m A) ⊆ dom ∧ finite dom);
failed ← SPEC(λ-::bool. True);
if failed
then do {

$c \leftarrow$ *remap-polys-l-dom-err*;
$RETURN$ (*error-msg* (*0* :: *nat*) *c*, $\mathcal{V}$, *fmempty*)
}
else do {
(b, $\mathcal{V}$, A) $\leftarrow$ *FOREACH dom*
(λi (b, $\mathcal{V}$, A').
if $i \in\#$ *dom-m* A
then do {
$p \leftarrow$ *full-normalize-poly* (*the* (*fmlookup* A i));
$eq \leftarrow$ *weak-equality-l p spec*;
$\mathcal{V} \leftarrow RETURN(\mathcal{V} \cup$ *vars-llist* (*the* (*fmlookup* A i)));
$RETURN(b \vee eq$, $\mathcal{V}$, *fmupd i p* A')
} *else RETURN* (b, $\mathcal{V}$, A'))
(*False*, $\mathcal{V}$, *fmempty*);
RETURN (*if b then CFOUND else CSUCCESS*, $\mathcal{V}$, A)
}})›

definition *PAC-checker-l* **where**
‹*PAC-checker-l spec A b st* = *do* {
(S, -) $\leftarrow$ $WHILE_T$
(λ((b, A), n). $\neg$*is-cfailed* $b \wedge n \neq []$)
(λ((bA), n). *do* {
$ASSERT(n \neq [])$;
$S \leftarrow$ *PAC-checker-l-step spec bA* (*hd* n);
RETURN (S, *tl* n)
})
((b, A), *st*);
RETURN S
}›

10.2 Correctness

We now enter the locale to reason about polynomials directly.

context *poly-embed*
begin

abbreviation *pac-step-rel* **where**
‹*pac-step-rel* $\equiv$ *p2rel* (⟨*Id*, *fully-unsorted-poly-rel O mset-poly-rel*, *var-rel*⟩ *pac-step-rel-raw*)›

abbreviation *fmap-polys-rel* **where**
‹*fmap-polys-rel* $\equiv$ ⟨*nat-rel*, *sorted-poly-rel O mset-poly-rel*⟩*fmap-rel*›

lemma
‹*normalize-poly-p s0 s* $\Longrightarrow$
(*s0*, *p*) $\in$ *mset-poly-rel* $\Longrightarrow$
(*s*, *p*) $\in$ *mset-poly-rel*›
⟨*proof*⟩

lemma *vars-poly-of-vars*:
‹*vars* (*poly-of-vars a* :: *int mpoly*) $\subseteq$ (φ ‘ *set-mset a*)›
⟨*proof*⟩

lemma *vars-polynom-of-mset*:
‹*vars* (*polynom-of-mset za*) $\subseteq \bigcup$(*image* φ ‘ (*set-mset o fst*) ‘ *set-mset za*)›
⟨*proof*⟩

lemma *fully-unsorted-poly-rel-vars-subset-vars-llist*:
‹$(A, B) \in$ *fully-unsorted-poly-rel O mset-poly-rel* $\Longrightarrow$ *vars* $B \subseteq \varphi$ ‘ *vars-llist* A›
⟨*proof*⟩

lemma *fully-unsorted-poly-rel-extend-vars*:
‹$(A, B) \in$ *fully-unsorted-poly-rel O mset-poly-rel* $\Longrightarrow$
$(x1c, x1a) \in \langle var\text{-}rel\rangle set\text{-}rel \Longrightarrow$
RETURN $(x1c \cup$ *vars-llist* $A)$
$\leq \Downarrow (\langle var\text{-}rel\rangle set\text{-}rel)$
$(SPEC\ ((\subseteq)\ (x1a \cup vars\ (B))))$›
⟨*proof*⟩

lemma *remap-polys-l-remap-polys*:
assumes
AB: ‹$(A, B) \in \langle nat\text{-}rel,$ *fully-unsorted-poly-rel O mset-poly-rel*$\rangle fmap\text{-}rel$› **and**
spec: ‹$(spec, spec') \in$ *sorted-poly-rel O mset-poly-rel*› **and**
V: ‹$(\mathcal{V}, \mathcal{V}') \in \langle var\text{-}rel\rangle set\text{-}rel$›
shows ‹*remap-polys-l spec* $\mathcal{V}$ $A \leq$
$\Downarrow$(*code-status-status-rel* $\times_r$ $\langle var\text{-}rel\rangle set\text{-}rel$ $\times_r$ *fmap-polys-rel*) (*remap-polys spec'* $\mathcal{V}'$ B)›
(**is** ‹- $\leq \Downarrow$ *?R* -›)
⟨*proof*⟩

lemma *fref-to-Down-curry*:
‹$(uncurry\ f, uncurry\ g) \in [P]_f\ A \rightarrow \langle B\rangle nres\text{-}rel \Longrightarrow$
$(\bigwedge x\ x'\ y\ y'.\ P\ (x', y') \Longrightarrow ((x, y), (x', y')) \in A \Longrightarrow f\ x\ y \leq \Downarrow B\ (g\ x'\ y'))$›
⟨*proof*⟩

lemma *weak-equality-spec-weak-equality*:
‹$(p, p') \in$ *mset-poly-rel* $\Longrightarrow$
$(r, r') \in$ *mset-poly-rel* $\Longrightarrow$
weak-equality-spec $p\ r \leq$ *weak-equality* $p'\ r'$›
⟨*proof*⟩

lemma *weak-equality-l-weak-equality-l'*[*refine*]:
‹*weak-equality-l* $p\ q \leq \Downarrow$ *bool-rel* (*weak-equality* $p'\ q'$)›
if ‹$(p, p') \in$ *sorted-poly-rel O mset-poly-rel*›
‹$(q, q') \in$ *sorted-poly-rel O mset-poly-rel*›
for $p\ p'\ q\ q'$
⟨*proof*⟩

lemma *error-msg-ne-SUCCES*[*iff*]:
‹*error-msg* $i\ m \neq$ *CSUCCESS*›
‹*error-msg* $i\ m \neq$ *CFOUND*›
‹*is-cfailed* (*error-msg* $i\ m$)›
‹$\neg$*is-cfound* (*error-msg* $i\ m$)›
⟨*proof*⟩

lemma *sorted-poly-rel-vars-llist*:
‹$(r, r') \in$ *sorted-poly-rel O mset-poly-rel* $\Longrightarrow$
vars $r' \subseteq \varphi$ ‘ *vars-llist* r›
⟨*proof*⟩

lemma *check-addition-l-check-add*:
assumes ‹$(A, B) \in$ *fmap-polys-rel*› **and** ‹$(r, r') \in$ *sorted-poly-rel O mset-poly-rel*›
‹$(p, p') \in$ *Id*› ‹$(q, q') \in$ *Id*› ‹$(i, i') \in$ *nat-rel*›
‹$(\mathcal{V}', \mathcal{V}) \in \langle$*var-rel*$\rangle$*set-rel*›
shows
‹*check-addition-l spec A* $\mathcal{V}'$ *p q i r* $\le \Downarrow \{(st, b).\ (\neg$*is-cfailed st* $\longleftrightarrow b) \wedge$
(*is-cfound st* $\longrightarrow$ *spec* $= r$)} (*check-add B* $\mathcal{V}$ p' q' i' r')›
⟨*proof*⟩

lemma *check-del-l-check-del*:
‹$(A, B) \in$ *fmap-polys-rel* $\Longrightarrow$ (*x3*, *x3a*) $\in$ *Id* $\Longrightarrow$ *check-del-l spec A* (*pac-src1* (*Del x3*))
$\le \Downarrow \{(st, b).\ (\neg$*is-cfailed st* $\longleftrightarrow b) \wedge (b \longrightarrow st =$ *CSUCCESS*)} (*check-del B* (*pac-src1* (*Del x3a*)))›
⟨*proof*⟩

lemma *check-mult-l-check-mult*:
assumes ‹$(A, B) \in$ *fmap-polys-rel*› **and** ‹$(r, r') \in$ *sorted-poly-rel O mset-poly-rel*› **and**
‹$(q, q') \in$ *sorted-poly-rel O mset-poly-rel*›
‹$(p, p') \in$ *Id*› ‹$(i, i') \in$ *nat-rel*› ‹$(\mathcal{V}, \mathcal{V}') \in \langle$*var-rel*$\rangle$*set-rel*›
shows
‹*check-mult-l spec A* $\mathcal{V}$ *p q i r* $\le \Downarrow$ $\{(st, b).\ (\neg$*is-cfailed st* $\longleftrightarrow b) \wedge$
(*is-cfound st* $\longrightarrow$ *spec* $= r$)} (*check-mult B* $\mathcal{V}'$ p' q' i' r')›
⟨*proof*⟩

lemma *normalize-poly-normalize-poly-spec*:
assumes ‹$(r, t) \in$ *unsorted-poly-rel O mset-poly-rel*›
shows
‹*normalize-poly r* $\le \Downarrow$(*sorted-poly-rel O mset-poly-rel*) (*normalize-poly-spec t*)›
⟨*proof*⟩

lemma *remove1-list-rel*:
‹$(xs, ys) \in \langle R\rangle$ *list-rel* $\Longrightarrow$
$(a, b) \in R \Longrightarrow$
IS-RIGHT-UNIQUE R $\Longrightarrow$
IS-LEFT-UNIQUE R $\Longrightarrow$
(*remove1 a xs*, *remove1 b ys*) $\in \langle R\rangle$*list-rel*›
⟨*proof*⟩

lemma *remove1-list-rel2*:
‹$(xs, ys) \in \langle R\rangle$ *list-rel* $\Longrightarrow$
$(a, b) \in R \Longrightarrow$
$(\bigwedge c.\ (a, c) \in R \Longrightarrow c = b) \Longrightarrow$
$(\bigwedge c.\ (c, b) \in R \Longrightarrow c = a) \Longrightarrow$
(*remove1 a xs*, *remove1 b ys*) $\in \langle R\rangle$*list-rel*›
⟨*proof*⟩

lemma *remove1-sorted-poly-rel-mset-poly-rel*:
assumes
‹$(r, r') \in$ *sorted-poly-rel O mset-poly-rel*› **and**
‹$([a], 1) \in$ *set r*›
shows
‹(*remove1* $([a], 1)$ *r*, $r' -$ *Var* $(\varphi\ a)$)
$\in$ *sorted-poly-rel O mset-poly-rel*›
⟨*proof*⟩

lemma *remove1-sorted-poly-rel-mset-poly-rel-minus*:
  assumes
    ‹$(r,\ r') \in$ *sorted-poly-rel O mset-poly-rel*› **and**
    ‹$([a],\ -1) \in$ *set r*›
  shows
    ‹$(remove1\ ([a],\ -1)\ r,\ r' +\ Var\ (\varphi\ a))$
        $\in$ *sorted-poly-rel O mset-poly-rel*›
⟨*proof*⟩

lemma *insert-var-rel-set-rel*:
  ‹$(\mathcal{V},\ \mathcal{V}') \in \langle var\text{-}rel\rangle set\text{-}rel \Longrightarrow$
  $(yb,\ x2) \in var\text{-}rel \Longrightarrow$
  $(insert\ yb\ \mathcal{V},\ insert\ x2\ \mathcal{V}') \in \langle var\text{-}rel\rangle set\text{-}rel$›
  ⟨*proof*⟩

lemma *var-rel-set-rel-iff*:
  ‹$(\mathcal{V},\ \mathcal{V}') \in \langle var\text{-}rel\rangle set\text{-}rel \Longrightarrow$
  $(yb,\ x2) \in var\text{-}rel \Longrightarrow$
  $yb \in \mathcal{V} \longleftrightarrow x2 \in\ \mathcal{V}'$›
  ⟨*proof*⟩

lemma *check-extension-l-check-extension*:
  assumes ‹$(A,\ B) \in$ *fmap-polys-rel*› **and** ‹$(r,\ r') \in$ *sorted-poly-rel O mset-poly-rel*› **and**
    ‹$(i,\ i') \in$ *nat-rel*› ‹$(\mathcal{V},\ \mathcal{V}') \in \langle var\text{-}rel\rangle set\text{-}rel$› ‹$(x,\ x') \in$ *var-rel*›
  shows
    ‹$check\text{-}extension\text{-}l\ spec\ A\ \mathcal{V}\ i\ x\ r \le$
      $\Downarrow\{((st),\ (b)).$
        $(\neg is\text{-}cfailed\ st \longleftrightarrow b)\ \wedge$
      $(is\text{-}cfound\ st \longrightarrow spec = r)\}\ (check\text{-}extension\ B\ \mathcal{V}'\ i'\ x'\ r')$›
⟨*proof*⟩

lemma *full-normalize-poly-diff-ideal*:
  fixes *dom*
  assumes ‹$(p,\ p') \in$ *fully-unsorted-poly-rel O mset-poly-rel*›
  shows
    ‹*full-normalize-poly p*
    $\le \Downarrow$ (*sorted-poly-rel O mset-poly-rel*)
      $(normalize\text{-}poly\text{-}spec\ p')$›
⟨*proof*⟩

lemma *insort-key-rel-decomp*:
  ‹$\exists\ ys\ zs.\ xs = ys\ @\ zs \wedge insort\text{-}key\text{-}rel\ R\ x\ xs = ys\ @\ x\ \#\ zs$›
  ⟨*proof*⟩

lemma *list-rel-append-same-length*:
  ‹$length\ xs = length\ xs' \Longrightarrow (xs\ @\ ys,\ xs'\ @\ ys') \in \langle R\rangle list\text{-}rel \longleftrightarrow (xs,\ xs') \in \langle R\rangle list\text{-}rel \wedge (ys,\ ys') \in$
$\langle R\rangle list\text{-}rel$›
  ⟨*proof*⟩

lemma *term-poly-list-rel-list-relD*: ‹$(ys,\ cs) \in \langle term\text{-}poly\text{-}list\text{-}rel \times_r int\text{-}rel\rangle list\text{-}rel \Longrightarrow$
    $cs = map\ (\lambda(a,\ y).\ (mset\ a,\ y))\ ys$›

⟨*proof*⟩

lemma *term-poly-list-rel-single*: ‹$([x32], \{\#x32\#\}) \in$ *term-poly-list-rel*›
⟨*proof*⟩

lemma *unsorted-poly-rel-list-rel-list-rel-uminus*:
‹$(map\ (\lambda(a, b).\ (a, -\ b))\ r,\ yc)$
$\in \langle$*unsorted-term-poly-list-rel* $\times_r$ *int-rel*$\rangle$*list-rel* $\Longrightarrow$
$(r,\ map\ (\lambda(a, b).\ (a, -\ b))\ yc)$
$\in \langle$*unsorted-term-poly-list-rel* $\times_r$ *int-rel*$\rangle$*list-rel*›
⟨*proof*⟩

lemma *mset-poly-rel-minus*: ‹$(\{\#(a, b)\#\}, v') \in$ *mset-poly-rel* $\Longrightarrow$
$(mset\ yc,\ r') \in$ *mset-poly-rel* $\Longrightarrow$
$(r,\ yc)$
$\in \langle$*unsorted-term-poly-list-rel* $\times_r$ *int-rel*$\rangle$*list-rel* $\Longrightarrow$
(*add-mset* (a, b) $(mset\ yc)$,
$v' + r'$)
$\in$ *mset-poly-rel*›
⟨*proof*⟩

lemma *fully-unsorted-poly-rel-diff*:
‹$([v], v') \in$ *fully-unsorted-poly-rel O mset-poly-rel* $\Longrightarrow$
$(r, r') \in$ *fully-unsorted-poly-rel O mset-poly-rel* $\Longrightarrow$
$(v\ \#\ r,$
$v' + r')$
$\in$ *fully-unsorted-poly-rel O mset-poly-rel*›
⟨*proof*⟩

lemma *PAC-checker-l-step-PAC-checker-step*:
assumes
‹$(Ast, Bst) \in$ *code-status-status-rel* $\times_r$ ⟨*var-rel*⟩*set-rel* $\times_r$ *fmap-polys-rel*› **and**
‹$(st, st') \in$ *pac-step-rel*› **and**
spec: ‹$(spec, spec') \in$ *sorted-poly-rel O mset-poly-rel*›
shows
‹*PAC-checker-l-step spec Ast st* $\leq \Downarrow$ (*code-status-status-rel* $\times_r$ ⟨*var-rel*⟩*set-rel* $\times_r$ *fmap-polys-rel*) (*PAC-checker-step spec′ Bst st′*)›
⟨*proof*⟩

lemma *code-status-status-rel-discrim-iff*:
‹$(x1a, x1c) \in$ *code-status-status-rel* $\Longrightarrow$ *is-cfailed x1a* $\longleftrightarrow$ *is-failed x1c*›
‹$(x1a, x1c) \in$ *code-status-status-rel* $\Longrightarrow$ *is-cfound x1a* $\longleftrightarrow$ *is-found x1c*›
⟨*proof*⟩

lemma *PAC-checker-l-PAC-checker*:
assumes
‹$(A, B) \in$ ⟨*var-rel*⟩*set-rel* $\times_r$ *fmap-polys-rel*› **and**
‹$(st, st') \in$ ⟨*pac-step-rel*⟩*list-rel*› **and**
‹$(spec, spec') \in$ *sorted-poly-rel O mset-poly-rel*› **and**
‹$(b, b') \in$ *code-status-status-rel*›
shows
‹*PAC-checker-l spec A b st* $\leq \Downarrow$ (*code-status-status-rel* $\times_r$ ⟨*var-rel*⟩*set-rel* $\times_r$ *fmap-polys-rel*) (*PAC-checker spec′ B b′ st′*)›
⟨*proof*⟩

end

lemma *less-than-char-of-char*[*code-unfold*]:
‹$(x,\ y) \in$ *less-than-char* $\longleftrightarrow$ (*of-char* x :: *nat*) $<$ *of-char* y›
⟨*proof*⟩

lemmas [*code*] =
add-poly-l′.simps[*unfolded var-order-rel-def*]

export-code *add-poly-l′* **in** *SML* **module-name** *test*

definition *full-checker-l*
:: ‹*llist-polynom* ⇒ (*nat*, *llist-polynom*) *fmap* ⇒ (-, *string*, *nat*) *pac-step list* ⇒
(*string code-status* × -) *nres*›
where
‹*full-checker-l spec A st* = *do* {
spec′ ← *full-normalize-poly spec*;
(b, $\mathcal{V}$, A) ← *remap-polys-l spec′* {} A;
if is-cfailed b
then RETURN (b, $\mathcal{V}$, A)
else do {
let $\mathcal{V} = \mathcal{V} \cup$ *vars-llist spec*;
PAC-checker-l spec′ ($\mathcal{V}$, A) *b st*
}
}›

context *poly-embed*
begin

term *normalize-poly-spec*
thm *full-normalize-poly-diff-ideal*[*unfolded normalize-poly-spec-def*[*symmetric*]]
abbreviation *unsorted-fmap-polys-rel* **where**
‹*unsorted-fmap-polys-rel* ≡ ⟨*nat-rel*, *fully-unsorted-poly-rel O mset-poly-rel*⟩*fmap-rel*›

lemma *full-checker-l-full-checker*:
assumes
‹(A, B) ∈ *unsorted-fmap-polys-rel*› **and**
‹(*st*, *st′*) ∈ ⟨*pac-step-rel*⟩*list-rel*› **and**
‹(*spec*, *spec′*) ∈ *fully-unsorted-poly-rel O mset-poly-rel*›
shows
‹*full-checker-l spec A st* ≤ ⇓ (*code-status-status-rel* $\times_r$ ⟨*var-rel*⟩*set-rel* $\times_r$ *fmap-polys-rel*) (*full-checker spec′ B st′*)›
⟨*proof*⟩

lemma *full-checker-l-full-checker′*:
‹(*uncurry2 full-checker-l*, *uncurry2 full-checker*) ∈
((*fully-unsorted-poly-rel O mset-poly-rel*) $\times_r$ *unsorted-fmap-polys-rel*) $\times_r$ ⟨*pac-step-rel*⟩*list-rel* $\rightarrow_f$
⟨(*code-status-status-rel* $\times_r$ ⟨*var-rel*⟩*set-rel* $\times_r$ *fmap-polys-rel*)⟩*nres-rel*›
⟨*proof*⟩

end

definition *remap-polys-l2* :: ‹*llist-polynom* ⇒ *string set* ⇒ (*nat*, *llist-polynom*) *fmap* ⇒ - *nres*› **where**
‹*remap-polys-l2 spec* = (λ𝒱 *A*. *do*{
  n ← *upper-bound-on-dom A*;
  b ← *RETURN* (*n* ≥ *2^64*);
  if b
  then do {
    c ← *remap-polys-l-dom-err*;
    RETURN (*error-msg* (*0* ::*nat*) *c*, 𝒱, *fmempty*)
  }
  else do {
      (*b*, 𝒱, *A*) ← *nfoldli* ([*0*..<*n*]) (λ-. *True*)
      (λ*i* (*b*, 𝒱, *A'*).
        if i ∈# *dom-m A*
        then do {
          ASSERT(*fmlookup A i* ≠ *None*);
          p ← *full-normalize-poly* (*the* (*fmlookup A i*));
          eq ← *weak-equality-l p spec*;
          𝒱 ← *RETURN* (𝒱 ∪ *vars-llist* (*the* (*fmlookup A i*)));
          RETURN(*b* ∨ *eq*, 𝒱, *fmupd i p A'*)
        } *else RETURN* (*b*, 𝒱, *A'*)
)
      (*False*, 𝒱, *fmempty*);
    RETURN (*if b then CFOUND else CSUCCESS*, 𝒱, *A*)
  }
})›

lemma *remap-polys-l2-remap-polys-l*:
‹*remap-polys-l2 spec* 𝒱 *A* ≤ ⇓ *Id* (*remap-polys-l spec* 𝒱 *A*)›
⟨*proof*⟩

end
theory *PAC-Checker-Relation*
  imports *PAC-Checker WB-Sort Native-Word.Uint64*
begin

11 Various Refinement Relations

When writing this, it was not possible to share the definition with the IsaSAT version.

definition *uint64-nat-rel* :: (*uint64* × *nat*) *set* **where**
‹*uint64-nat-rel* = *br nat-of-uint64* (λ-. *True*)›

abbreviation *uint64-nat-assn* **where**
‹*uint64-nat-assn* ≡ *pure uint64-nat-rel*›

instantiation *uint32* :: *hashable*
begin
definition *hashcode-uint32* :: ‹*uint32* ⇒ *uint32*› **where**
‹*hashcode-uint32 n* = *n*›

definition *def-hashmap-size-uint32* :: ‹*uint32 itself* ⇒ *nat*› **where**
‹*def-hashmap-size-uint32* = (λ-. *16*)›
— same as *nat*
instance

⟨*proof*⟩
end

instantiation *uint64* :: *hashable*
begin
definition *hashcode-uint64* :: ‹*uint64* ⇒ *uint32*› **where**
‹$hashcode\text{-}uint64\ n = (uint32\text{-}of\text{-}nat\ (nat\text{-}of\text{-}uint64\ ((n)\ AND\ ((2 :: uint64)\hat{\ }32\ -1))))$›

definition *def-hashmap-size-uint64* :: ‹*uint64 itself* ⇒ *nat*› **where**
‹$def\text{-}hashmap\text{-}size\text{-}uint64 = (\lambda\text{-}.\ 16)$›
— same as *nat*
instance
⟨*proof*⟩
end

lemma *word-nat-of-uint64-Rep-inject*[*simp*]: ‹$nat\text{-}of\text{-}uint64\ ai = nat\text{-}of\text{-}uint64\ bi \longleftrightarrow ai = bi$›
⟨*proof*⟩

instance *uint64* :: *heap*
⟨*proof*⟩

instance *uint64* :: *semiring-numeral*
⟨*proof*⟩

lemma *nat-of-uint64-012*[*simp*]: ‹$nat\text{-}of\text{-}uint64\ 0 = 0$› ‹$nat\text{-}of\text{-}uint64\ 2 = 2$› ‹$nat\text{-}of\text{-}uint64\ 1 = 1$›
⟨*proof*⟩

definition *uint64-of-nat-conv* **where**
[*simp*]: ‹$uint64\text{-}of\text{-}nat\text{-}conv\ (x :: nat) = x$›
lemma *less-upper-bintrunc-id*: ‹$n < 2\ \hat{\ }b \Longrightarrow n \geq 0 \Longrightarrow bintrunc\ b\ n = n$›
⟨*proof*⟩

lemma *nat-of-uint64-uint64-of-nat-id*: ‹$n < 2\hat{\ }64 \Longrightarrow nat\text{-}of\text{-}uint64\ (uint64\text{-}of\text{-}nat\ n) = n$›
⟨*proof*⟩

lemma [*sepref-fr-rules*]:
‹$(return\ o\ uint64\text{-}of\text{-}nat,\ RETURN\ o\ uint64\text{-}of\text{-}nat\text{-}conv) \in [\lambda a.\ a < 2\ \hat{\ }64]_a\ nat\text{-}assn^k \rightarrow uint64\text{-}nat\text{-}assn$›
⟨*proof*⟩

definition *string-rel* :: ‹$(String.literal \times string)\ set$› **where**
‹$string\text{-}rel = \{(x,\ y).\ y = String.explode\ x\}$›

abbreviation *string-assn* :: ‹$string \Rightarrow String.literal \Rightarrow assn$› **where**
‹$string\text{-}assn \equiv pure\ string\text{-}rel$›

lemma *eq-string-eq*:
‹$((=),\ (=)) \in string\text{-}rel \rightarrow string\text{-}rel \rightarrow bool\text{-}rel$›
⟨*proof*⟩

lemmas *eq-string-eq-hnr* =
eq-string-eq[*sepref-import-param*]

definition *string2-rel* :: ‹$(string \times string)\ set$› **where**
‹$string2\text{-}rel \equiv \langle Id\rangle list\text{-}rel$›

abbreviation *string2-assn* :: ‹*string* ⇒ *string* ⇒ *assn*› **where**
‹*string2-assn* ≡ *pure string2-rel*›

abbreviation *monom-rel* **where**
‹*monom-rel* ≡ ⟨*string-rel*⟩*list-rel*›

abbreviation *monom-assn* **where**
‹*monom-assn* ≡ *list-assn string-assn*›

abbreviation *monomial-rel* **where**
‹*monomial-rel* ≡ *monom-rel* $\times_r$ *int-rel*›

abbreviation *monomial-assn* **where**
‹*monomial-assn* ≡ *monom-assn* $\times_a$ *int-assn*›

abbreviation *poly-rel* **where**
‹*poly-rel* ≡ ⟨*monomial-rel*⟩*list-rel*›

abbreviation *poly-assn* **where**
‹*poly-assn* ≡ *list-assn monomial-assn*›

lemma *poly-assn-alt-def*:
‹*poly-assn* = *pure poly-rel*›
⟨*proof*⟩

abbreviation *polys-assn* **where**
‹*polys-assn* ≡ *hm-fmap-assn uint64-nat-assn poly-assn*›

lemma *string-rel-string-assn*:
‹(↑ ((*c*, *a*) ∈ *string-rel*)) = *string-assn a c*›
⟨*proof*⟩

lemma *single-valued-string-rel*:
‹*single-valued string-rel*›
⟨*proof*⟩

lemma *IS-LEFT-UNIQUE-string-rel*:
‹*IS-LEFT-UNIQUE string-rel*›
⟨*proof*⟩

lemma *IS-RIGHT-UNIQUE-string-rel*:
‹*IS-RIGHT-UNIQUE string-rel*›
⟨*proof*⟩

lemma *single-valued-monom-rel*: ‹*single-valued monom-rel*›
⟨*proof*⟩

lemma *single-valued-monomial-rel*:
‹*single-valued monomial-rel*›
⟨*proof*⟩

lemma *single-valued-monom-rel'*: ‹*IS-LEFT-UNIQUE monom-rel*›
⟨*proof*⟩

lemma *single-valued-monomial-rel'*:
 ‹*IS-LEFT-UNIQUE monomial-rel*›
 ⟨*proof*⟩

lemma [*safe-constraint-rules*]:
 ‹*Sepref-Constraints.CONSTRAINT single-valued string-rel*›
 ‹*Sepref-Constraints.CONSTRAINT IS-LEFT-UNIQUE string-rel*›
 ⟨*proof*⟩

lemma *eq-string-monom-hnr*[*sepref-fr-rules*]:
 ‹$(uncurry\ (return\ oo\ (=)),\ uncurry\ (RETURN\ oo\ (=))) \in monom\text{-}assn^k *_a monom\text{-}assn^k \rightarrow_a bool\text{-}assn$›
 ⟨*proof*⟩

definition *term-order-rel'* **where**
 [*simp*]: ‹$term\text{-}order\text{-}rel'\ x\ y = ((x,\ y) \in term\text{-}order\text{-}rel)$›

lemma *term-order-rel*[*def-pat-rules*]:
 ‹(∈)\$(*x*,*y*)\$*term-order-rel* ≡ *term-order-rel'*\$*x*\$*y*›
 ⟨*proof*⟩

lemma *term-order-rel-alt-def*:
 ‹*term-order-rel* = *lexord* (*p2rel char.lexordp*)›
 ⟨*proof*⟩

instantiation *char* :: *linorder*
begin
 definition *less-char* **where** [*symmetric*, *simp*]: *less-char* = *PAC-Polynomials-Term.less-char*
 definition *less-eq-char* **where** [*symmetric*, *simp*]: *less-eq-char* = *PAC-Polynomials-Term.less-eq-char*
instance
 ⟨*proof*⟩
end

instantiation *list* :: (*linorder*) *linorder*
begin
 definition *less-list* **where** *less-list* = *lexordp* (<)
 definition *less-eq-list* **where** *less-eq-list* = *lexordp-eq*

instance
 ⟨*proof*⟩

end

lemma *term-order-rel'-alt-def-lexord*:
  ‹$term\text{-}order\text{-}rel'\ x\ y = ord\text{-}class.lexordp\ x\ y$› **and**
 term-order-rel'-alt-def:
  ‹$term\text{-}order\text{-}rel'\ x\ y \longleftrightarrow x < y$›
⟨*proof*⟩

lemma *list-rel-list-rel-order-iff*:
 assumes ‹$(a,\ b) \in \langle string\text{-}rel\rangle list\text{-}rel$› ‹$(a',\ b') \in \langle string\text{-}rel\rangle list\text{-}rel$›

shows ‹$a < a' \longleftrightarrow b < b'$›
⟨*proof*⟩

lemma *string-rel-le*[*sepref-import-param*]:
shows ‹$((<), (<)) \in \langle$*string-rel*$\rangle$*list-rel* $\rightarrow$ $\langle$*string-rel*$\rangle$*list-rel* $\rightarrow$ *bool-rel*›
⟨*proof*⟩

lemma [*sepref-import-param*]:
assumes ‹*CONSTRAINT IS-LEFT-UNIQUE R*› ‹*CONSTRAINT IS-RIGHT-UNIQUE R*›
shows ‹(*remove1*, *remove1*) $\in R \rightarrow \langle R\rangle$*list-rel* $\rightarrow \langle R\rangle$*list-rel*›
⟨*proof*⟩

instantiation *pac-step* :: (*heap*, *heap*, *heap*) *heap*
begin

instance
⟨*proof*⟩

end

end
theory *PAC-Checker-Init*
imports *PAC-Checker WB-Sort PAC-Checker-Relation*
begin

12 Initial Normalisation of Polynoms

12.1 Sorting

Adapted from the theory *HOL−ex.MergeSort* by Tobias. We did not change much, but we refine it to executable code and try to improve efficiency.

fun *merge* :: - ⇒ *'a list* ⇒ *'a list* ⇒ *'a list*
where
merge f (*x#xs*) (*y#ys*) =
(*if f x y then x # merge f xs* (*y#ys*) *else y # merge f* (*x#xs*) *ys*)
| *merge f xs* [] = *xs*
| *merge f* [] *ys* = *ys*

lemma *mset-merge* [*simp*]:
mset (*merge f xs ys*) = *mset xs* + *mset ys*
⟨*proof*⟩

lemma *set-merge* [*simp*]:
set (*merge f xs ys*) = *set xs* ∪ *set ys*
⟨*proof*⟩

lemma *sorted-merge*:
transp f $\Longrightarrow$ ($\bigwedge x\ y.\ f\ x\ y \vee f\ y\ x$) $\Longrightarrow$
sorted-wrt f (*merge f xs ys*) $\longleftrightarrow$ *sorted-wrt f xs* ∧ *sorted-wrt f ys*
⟨*proof*⟩

fun *msort* :: - ⇒ *'a list* ⇒ *'a list*

where
  msort f [] = []
| *msort f* [*x*] = [*x*]
| *msort f xs* = *merge f*
      (*msort f* (*take* (*size xs div 2*) *xs*))
      (*msort f* (*drop* (*size xs div 2*) *xs*))

fun *swap-ternary* :: ‹-⇒*nat*⇒*nat*⇒ ('*a* × '*a* × '*a*) ⇒ ('*a* × '*a* × '*a*)› **where**
  ‹*swap-ternary f m n* =
    (*if* (*m* = *0* ∧ *n* = *1*)
    then (λ(*a*, *b*, *c*). *if f a b then* (*a*, *b*, *c*)
      else (*b*,*a*,*c*))
    else if (*m* = *0* ∧ *n* = *2*)
    then (λ(*a*, *b*, *c*). *if f a c then* (*a*, *b*, *c*)
      else (*c*,*b*,*a*))
    else if (*m* = *1* ∧ *n* = *2*)
    then (λ(*a*, *b*, *c*). *if f b c then* (*a*, *b*, *c*)
      else (*a*,*c*,*b*))
    else (λ(*a*, *b*, *c*). (*a*,*b*,*c*)))›

fun *msort2* :: - ⇒ '*a list* ⇒ '*a list*
where
  msort2 f [] = []
| *msort2 f* [*x*] = [*x*]
| *msort2 f* [*x*,*y*] = (*if f x y then* [*x*,*y*] *else* [*y*,*x*])
| *msort2 f xs* = *merge f*
      (*msort f* (*take* (*size xs div 2*) *xs*))
      (*msort f* (*drop* (*size xs div 2*) *xs*))

lemmas [*code del*] =
  msort2.simps

declare *msort2.simps*[*simp del*]
lemmas [*code*] =
  msort2.simps[*unfolded swap-ternary.simps*, *simplified*]

declare *msort2.simps*[*simp*]

lemma *msort-msort2*:
  fixes *xs* :: ‹'*a* :: *linorder list*›
  shows ‹*msort* (≤) *xs* = *msort2* (≤) *xs*›
  ⟨*proof*⟩

lemma *sorted-msort*:
  transp f ⟹ (⋀*x y*. *f x y* ∨ *f y x*) ⟹
   sorted-wrt f (*msort f xs*)
  ⟨*proof*⟩

lemma *mset-msort*[*simp*]:
  mset (*msort f xs*) = *mset xs*
  ⟨*proof*⟩

12.2 Sorting applied to monomials

lemma *merge-coeffs-alt-def*:
  ‹(*RETURN o merge-coeffs*) *p* =

$REC_T(\lambda f\ p.$
  $(case\ p\ of$
    $[] \Rightarrow RETURN\ []$
  $|\ [\text{-}] => RETURN\ p$
  $|\ ((xs,\ n)\ \#\ (ys,\ m)\ \#\ p) \Rightarrow$
   $(if\ xs = ys$
    $then\ if\ n + m \neq 0\ then\ f\ ((xs,\ n + m)\ \#\ p)\ else\ f\ p$
    $else\ do\ \{p \leftarrow f\ ((ys,\ m)\ \#\ p);\ RETURN\ ((xs,\ n)\ \#\ p)\})))$
  p›
⟨*proof*⟩

lemma *hn-invalid-recover*:
‹$is\text{-}pure\ R \Longrightarrow hn\text{-}invalid\ R = (\lambda x\ y.\ R\ x\ y * true)$›
‹$is\text{-}pure\ R \Longrightarrow invalid\text{-}assn\ R = (\lambda x\ y.\ R\ x\ y * true)$›
⟨*proof*⟩

lemma *safe-poly-vars*:
shows
  [*safe-constraint-rules*]:
    is-pure (*poly-assn*) **and**
  [*safe-constraint-rules*]:
    is-pure (*monom-assn*) **and**
  [*safe-constraint-rules*]:
    is-pure (*monomial-assn*) **and**
  [*safe-constraint-rules*]:
    is-pure string-assn
⟨*proof*⟩

lemma *invalid-assn-distrib*:
‹$invalid\text{-}assn\ monom\text{-}assn \times_a invalid\text{-}assn\ int\text{-}assn = invalid\text{-}assn\ (monom\text{-}assn \times_a int\text{-}assn)$›
  ⟨*proof*⟩

lemma *WTF-RF-recover*:
‹$hn\text{-}ctxt\ (invalid\text{-}assn\ monom\text{-}assn \times_a invalid\text{-}assn\ int\text{-}assn)\ xb$
    $x'a \vee_A$
   $hn\text{-}ctxt\ monomial\text{-}assn\ xb\ x'a \Longrightarrow_t$
   $hn\text{-}ctxt\ (monomial\text{-}assn)\ xb\ x'a$›
⟨*proof*⟩

lemma *WTF-RF*:
‹$hn\text{-}ctxt\ (invalid\text{-}assn\ monom\text{-}assn \times_a invalid\text{-}assn\ int\text{-}assn)\ xb\ x'a *$
   $(hn\text{-}invalid\ poly\text{-}assn\ la\ l'a * hn\text{-}invalid\ int\text{-}assn\ a2'\ a2 *$
    $hn\text{-}invalid\ monom\text{-}assn\ a1'\ a1 *$
    $hn\text{-}invalid\ poly\text{-}assn\ l\ l' *$
    $hn\text{-}invalid\ monomial\text{-}assn\ xa\ x' *$
    $hn\text{-}invalid\ poly\text{-}assn\ ax\ px) \Longrightarrow_t$
   $hn\text{-}ctxt\ (monomial\text{-}assn)\ xb\ x'a *$
   $hn\text{-}ctxt\ poly\text{-}assn$
    $la\ l'a *$
   $hn\text{-}ctxt\ poly\text{-}assn\ l\ l' *$
   $(hn\text{-}invalid\ int\text{-}assn\ a2'\ a2 *$
    $hn\text{-}invalid\ monom\text{-}assn\ a1'\ a1 *$
    $hn\text{-}invalid\ monomial\text{-}assn\ xa\ x' *$
    $hn\text{-}invalid\ poly\text{-}assn\ ax\ px)$›
‹$hn\text{-}ctxt\ (invalid\text{-}assn\ monom\text{-}assn \times_a invalid\text{-}assn\ int\text{-}assn)\ xa\ x' *$

$(hn\text{-}ctxt\ poly\text{-}assn\ l\ l' * hn\text{-}invalid\ poly\text{-}assn\ ax\ px) \Longrightarrow_t$
$hn\text{-}ctxt\ (monomial\text{-}assn)\ xa\ x' *$
$hn\text{-}ctxt\ poly\text{-}assn\ l\ l' *$
$hn\text{-}ctxt\ poly\text{-}assn\ ax\ px *$
emp›
⟨*proof*⟩

The refinement frameword is completely lost here when synthesizing the constants – it does not understant what is pure (actually everything) and what must be destroyed.

sepref-definition *merge-coeffs-impl*
is ‹$RETURN\ o\ merge\text{-}coeffs$›
:: ‹$poly\text{-}assn^d \rightarrow_a poly\text{-}assn$›
⟨*proof*⟩

definition *full-quicksort-poly* **where**
‹$full\text{-}quicksort\text{-}poly = full\text{-}quicksort\text{-}ref\ (\lambda x\ y.\ x = y \lor (x, y) \in term\text{-}order\text{-}rel)\ fst$›

lemma *down-eq-id-list-rel*: ‹$\Downarrow(\langle Id\rangle list\text{-}rel)\ x = x$›
⟨*proof*⟩

definition *quicksort-poly*:: ‹$nat \Rightarrow nat \Rightarrow llist\text{-}polynom \Rightarrow (llist\text{-}polynom)\ nres$› **where**
‹$quicksort\text{-}poly\ x\ y\ \ z = quicksort\text{-}ref\ (\le)\ fst\ (x, y, z)$›

term *partition-between-ref*

definition *partition-between-poly* :: ‹$nat \Rightarrow nat \Rightarrow llist\text{-}polynom \Rightarrow (llist\text{-}polynom \times nat)\ nres$› **where**
‹$partition\text{-}between\text{-}poly = partition\text{-}between\text{-}ref\ (\le)\ fst$›

definition *partition-main-poly* :: ‹$nat \Rightarrow nat \Rightarrow llist\text{-}polynom \Rightarrow (llist\text{-}polynom \times nat)\ nres$› **where**
‹$partition\text{-}main\text{-}poly = partition\text{-}main\ (\le)\ \ fst$›

lemma *string-list-trans*:
‹$(xa :: char\ list\ list, ya) \in lexord\ (lexord\ \{(x, y).\ x < y\}) \Longrightarrow$
$(ya, z) \in lexord\ (lexord\ \{(x, y).\ x < y\}) \Longrightarrow$
$(xa, z) \in lexord\ (lexord\ \{(x, y).\ x < y\})$›
⟨*proof*⟩

lemma *full-quicksort-sort-poly-spec*:
‹$(full\text{-}quicksort\text{-}poly, sort\text{-}poly\text{-}spec) \in \langle Id\rangle list\text{-}rel \rightarrow_f \langle\langle Id\rangle list\text{-}rel\rangle nres\text{-}rel$›
⟨*proof*⟩

12.3 Lifting to polynomials

definition *merge-sort-poly* :: ‹-› **where**
‹$merge\text{-}sort\text{-}poly = msort\ (\lambda a\ b.\ fst\ a \le fst\ b)$›

definition *merge-monoms-poly* :: ‹-› **where**
‹$merge\text{-}monoms\text{-}poly = msort\ (\le)$›

definition *merge-poly* :: ‹-› **where**
‹$merge\text{-}poly = merge\ (\lambda a\ b.\ fst\ a \le fst\ b)$›

definition *merge-monoms* :: ‹-› **where**
‹$merge\text{-}monoms = merge\ (\le)$›

definition *msort-poly-impl* :: ‹(*String.literal list* × *int*) *list* ⇒ -› **where**
‹*msort-poly-impl* = *msort* (λ*a b*. *fst a* ≤ *fst b*)›

definition *msort-monoms-impl* :: ‹(*String.literal list*) ⇒ -› **where**
‹*msort-monoms-impl* = *msort* (≤)›

lemma *msort-poly-impl-alt-def*:
  ‹*msort-poly-impl xs* =
    (*case xs of*
      [] ⇒ []
    | [*a*] ⇒ [*a*]
    | [*a*,*b*] ⇒ *if fst a* ≤ *fst b then* [*a*,*b*]*else* [*b*,*a*]
    | *xs* ⇒ *merge-poly*
                (*msort-poly-impl* (*take* ((*length xs*) *div 2*) *xs*))
                (*msort-poly-impl* (*drop* ((*length xs*) *div 2*) *xs*)))›
  ⟨*proof*⟩

lemma *le-term-order-rel′*:
  ‹(≤) = (λ*x y*. *x* = *y* ∨ *term-order-rel′ x y*)›
  ⟨*proof*⟩

fun *lexord-eq* **where**
  ‹*lexord-eq* [] - = *True*› |
  ‹*lexord-eq* (*x* # *xs*) (*y* # *ys*) = (*x* < *y* ∨ (*x* = *y* ∧ *lexord-eq xs ys*))› |
  ‹*lexord-eq* - - = *False*›

lemma [*simp*]:
  ‹*lexord-eq* [] [] = *True*›
  ‹*lexord-eq* (*a* # *b*)[] = *False*›
  ‹*lexord-eq* [] (*a* # *b*) = *True*›
  ⟨*proof*⟩

lemma *var-order-rel′*:
  ‹(≤) = (λ*x y*. *x* = *y* ∨ (*x*,*y*) ∈ *var-order-rel*)›
  ⟨*proof*⟩

lemma *var-order-rel″*:
  ‹(*x*,*y*) ∈ *var-order-rel* ⟷ *x* < *y*›
  ⟨*proof*⟩

lemma *lexord-eq-alt-def1*:
  ‹*a* ≤ *b* = *lexord-eq a b*› **for** *a b* :: ‹*String.literal list*›
  ⟨*proof*⟩

lemma *lexord-eq-alt-def2*:
  ‹(*RETURN oo lexord-eq*) *xs ys* =
    REC_T (λ*f* (*xs*, *ys*).
      case (*xs*, *ys*) *of*
        ([], -) ⇒ *RETURN True*
      | (*x* # *xs*, *y* # *ys*) ⇒
          if x < *y then RETURN True*
          else if x = *y then f* (*xs*, *ys*) *else RETURN False*
      | - ⇒ *RETURN False*)
      (*xs*, *ys*)›

⟨*proof*⟩

definition *var-order'* **where**
[*simp*]: ‹*var-order'* = *var-order*›

lemma *var-order-rel*[*def-pat-rules*]:
‹(∈)$(*x*,*y*)$*var-order-rel* ≡ *var-order'*$*x*$*y*›
⟨*proof*⟩

lemma *var-order-rel-alt-def*:
‹*var-order-rel* = *p2rel char.lexordp*›
⟨*proof*⟩

lemma *var-order-rel-var-order*:
‹(*x*, *y*) ∈ *var-order-rel* ⟷ *var-order x y*›
⟨*proof*⟩

lemma *var-order-string-le*[*sepref-import-param*]:
‹((<), *var-order'*) ∈ *string-rel* → *string-rel* → *bool-rel*›
⟨*proof*⟩

lemma [*sepref-import-param*]:
‹((≤), (≤)) ∈ *monom-rel* → *monom-rel* →*bool-rel*›
⟨*proof*⟩

lemma [*sepref-import-param*]:
‹((<), (<)) ∈ *string-rel* → *string-rel* →*bool-rel*›
⟨*proof*⟩

lemma [*sepref-import-param*]:
‹((≤), (≤)) ∈ *string-rel* → *string-rel* →*bool-rel*›
⟨*proof*⟩

sepref-register *lexord-eq*
sepref-definition *lexord-eq-term*
is ‹*uncurry* (*RETURN oo lexord-eq*)›
:: ‹$monom\text{-}assn^k *_a monom\text{-}assn^k \to_a bool\text{-}assn$›
⟨*proof*⟩

declare *lexord-eq-term.refine*[*sepref-fr-rules*]

lemmas [*code del*] = *msort-poly-impl-def msort-monoms-impl-def*
lemmas [*code*] =
msort-poly-impl-def[*unfolded lexord-eq-alt-def1*[*abs-def*]]
msort-monoms-impl-def[*unfolded msort-msort2*]

lemma *term-order-rel-trans*:
‹ (*a*, *aa*) ∈ *term-order-rel* ⟹
(*aa*, *ab*) ∈ *term-order-rel* ⟹ (*a*, *ab*) ∈ *term-order-rel*›
⟨*proof*⟩

lemma *merge-sort-poly-sort-poly-spec*:
‹(*RETURN o merge-sort-poly*, *sort-poly-spec*) ∈ ⟨*Id*⟩*list-rel* $\to_f$ ⟨⟨*Id*⟩*list-rel*⟩*nres-rel*›

⟨proof⟩

lemma *msort-alt-def*:
‹$RETURN\ o\ (msort\ f) =$
$REC_T\ (\lambda g\ xs.$
$case\ xs\ of$
$[] \Rightarrow RETURN\ []$
$|\ [x] \Rightarrow RETURN\ [x]$
$|\ \text{-} \Rightarrow do\ \{$
$a \leftarrow g\ (take\ (size\ xs\ div\ 2)\ xs);$
$b \leftarrow g\ (drop\ (size\ xs\ div\ 2)\ xs);$
$RETURN\ (merge\ f\ a\ b)\})$›
⟨proof⟩

lemma *monomial-rel-order-map*:
‹$(x,\ a,\ b) \in monomial\text{-}rel \Longrightarrow$
$(y,\ aa,\ bb) \in monomial\text{-}rel \Longrightarrow$
$fst\ x \leq fst\ y \longleftrightarrow a \leq aa$›
⟨proof⟩

lemma *step-rewrite-pure*:
fixes K :: ‹$('olbl \times 'lbl)\ set$›
shows
‹$pure\ (p2rel\ (\langle K,\ V,\ R\rangle pac\text{-}step\text{-}rel\text{-}raw)) = pac\text{-}step\text{-}rel\text{-}assn\ (pure\ K)\ (pure\ V)\ (pure\ R)$›
‹$monomial\text{-}assn = pure\ (monom\text{-}rel \times_r int\text{-}rel)$› **and**
poly-assn-list:
‹$poly\text{-}assn = pure\ (\langle monom\text{-}rel \times_r int\text{-}rel\rangle list\text{-}rel)$›
⟨proof⟩

lemma *safe-pac-step-rel-assn*[*safe-constraint-rules*]:
$is\text{-}pure\ K \Longrightarrow is\text{-}pure\ V \Longrightarrow is\text{-}pure\ R \Longrightarrow is\text{-}pure\ (pac\text{-}step\text{-}rel\text{-}assn\ K\ V\ R)$
⟨proof⟩

lemma *merge-poly-merge-poly*:
‹$(merge\text{-}poly,\ merge\text{-}poly)$
$\in poly\text{-}rel \rightarrow poly\text{-}rel \rightarrow poly\text{-}rel$›
⟨proof⟩

lemmas [*fcomp-norm-unfold*] =
poly-assn-list[*symmetric*]
step-rewrite-pure(*1*)

lemma *merge-poly-merge-poly2*:
‹$(a,\ b) \in poly\text{-}rel \Longrightarrow (a',\ b') \in poly\text{-}rel \Longrightarrow$
$(merge\text{-}poly\ a\ a',\ merge\text{-}poly\ b\ b') \in poly\text{-}rel$›
⟨proof⟩

lemma *list-rel-takeD*:
‹$(a,\ b) \in \langle R\rangle list\text{-}rel \Longrightarrow (n,\ n') \in Id \Longrightarrow (take\ n\ a,\ take\ n'\ b) \in \langle R\rangle list\text{-}rel$›
⟨proof⟩

lemma *list-rel-dropD*:
‹$(a,\ b) \in \langle R\rangle list\text{-}rel \Longrightarrow (n,\ n') \in Id \Longrightarrow (drop\ n\ a,\ drop\ n'\ b) \in \langle R\rangle list\text{-}rel$›

⟨*proof*⟩

lemma *merge-sort-poly*[*sepref-import-param*]:
‹(*msort-poly-impl*, *merge-sort-poly*)
∈ *poly-rel* → *poly-rel*›
⟨*proof*⟩

lemmas [*sepref-fr-rules*] = *merge-sort-poly*[*FCOMP merge-sort-poly-sort-poly-spec*]

sepref-definition *partition-main-poly-impl*
is ‹*uncurry2 partition-main-poly*›
:: ‹$nat\text{-}assn^k *_a nat\text{-}assn^k *_a poly\text{-}assn^k \rightarrow_a prod\text{-}assn\ poly\text{-}assn\ nat\text{-}assn$ ›
⟨*proof*⟩

declare *partition-main-poly-impl.refine*[*sepref-fr-rules*]

sepref-definition *partition-between-poly-impl*
is ‹*uncurry2 partition-between-poly*›
:: ‹$nat\text{-}assn^k *_a nat\text{-}assn^k *_a poly\text{-}assn^k \rightarrow_a prod\text{-}assn\ poly\text{-}assn\ nat\text{-}assn$ ›
⟨*proof*⟩

declare *partition-between-poly-impl.refine*[*sepref-fr-rules*]

sepref-definition *quicksort-poly-impl*
is ‹*uncurry2 quicksort-poly*›
:: ‹$nat\text{-}assn^k *_a nat\text{-}assn^k *_a poly\text{-}assn^k \rightarrow_a poly\text{-}assn$›
⟨*proof*⟩

lemmas [*sepref-fr-rules*] = *quicksort-poly-impl.refine*

sepref-register *quicksort-poly*
sepref-definition *full-quicksort-poly-impl*
is ‹*full-quicksort-poly*›
:: ‹$poly\text{-}assn^k \rightarrow_a poly\text{-}assn$›
⟨*proof*⟩

lemmas *sort-poly-spec-hnr* =
full-quicksort-poly-impl.refine[*FCOMP full-quicksort-sort-poly-spec*]

declare *merge-coeffs-impl.refine*[*sepref-fr-rules*]

sepref-definition *normalize-poly-impl*
is ‹*normalize-poly*›
:: ‹$poly\text{-}assn^k \rightarrow_a poly\text{-}assn$›
⟨*proof*⟩

declare *normalize-poly-impl.refine*[*sepref-fr-rules*]

definition *full-quicksort-vars* **where**
‹*full-quicksort-vars* = *full-quicksort-ref* ($\lambda x\ y.\ x = y \lor (x, y) \in$ *var-order-rel*) *id*›

definition *quicksort-vars*:: ‹$nat \Rightarrow nat \Rightarrow string\ list \Rightarrow (string\ list)\ nres$› **where**
‹$quicksort\text{-}vars\ x\ y\ \ z = quicksort\text{-}ref\ (\leq)\ id\ (x,\ y,\ z)$›

definition *partition-between-vars* :: ‹$nat \Rightarrow nat \Rightarrow string\ list \Rightarrow (string\ list \times nat)\ nres$› **where**
‹$partition\text{-}between\text{-}vars = partition\text{-}between\text{-}ref\ (\leq)\ id$›

definition *partition-main-vars* :: ‹$nat \Rightarrow nat \Rightarrow string\ list \Rightarrow (string\ list \times nat)\ nres$› **where**
‹$partition\text{-}main\text{-}vars = partition\text{-}main\ (\leq)\ id$›

lemma *total-on-lexord-less-than-char-linear2*:
‹$xs \neq ys \Longrightarrow (xs,\ ys) \notin lexord\ (less\text{-}than\text{-}char) \longleftrightarrow$
$(ys,\ xs) \in lexord\ less\text{-}than\text{-}char$›
$\langle proof \rangle$

lemma *string-trans*:
‹$(xa,\ ya) \in lexord\ \{(x::char,\ y::char).\ x < y\} \Longrightarrow$
$(ya,\ z) \in lexord\ \{(x::char,\ y::char).\ x < y\} \Longrightarrow$
$(xa,\ z) \in lexord\ \{(x::char,\ y::char).\ x < y\}$›
$\langle proof \rangle$

lemma *full-quicksort-sort-vars-spec*:
‹$(full\text{-}quicksort\text{-}vars,\ sort\text{-}coeff) \in \langle Id\rangle list\text{-}rel \rightarrow_f \langle\langle Id\rangle list\text{-}rel\rangle nres\text{-}rel$›
$\langle proof \rangle$

sepref-definition *partition-main-vars-impl*
is ‹$uncurry2\ partition\text{-}main\text{-}vars$›
:: ‹$nat\text{-}assn^k *_a nat\text{-}assn^k *_a (monom\text{-}assn)^k \rightarrow_a prod\text{-}assn\ (monom\text{-}assn)\ nat\text{-}assn$›
$\langle proof \rangle$

declare *partition-main-vars-impl.refine*[*sepref-fr-rules*]

sepref-definition *partition-between-vars-impl*
is ‹$uncurry2\ partition\text{-}between\text{-}vars$›
:: ‹$nat\text{-}assn^k *_a nat\text{-}assn^k *_a monom\text{-}assn^k \rightarrow_a prod\text{-}assn\ monom\text{-}assn\ nat\text{-}assn$ ›
$\langle proof \rangle$

declare *partition-between-vars-impl.refine*[*sepref-fr-rules*]

sepref-definition *quicksort-vars-impl*
is ‹$uncurry2\ quicksort\text{-}vars$›
:: ‹$nat\text{-}assn^k *_a nat\text{-}assn^k *_a monom\text{-}assn^k \rightarrow_a monom\text{-}assn$›
$\langle proof \rangle$

lemmas [*sepref-fr-rules*] = *quicksort-vars-impl.refine*

sepref-register *quicksort-vars*

lemma *le-var-order-rel*:
‹$(\leq) = (\lambda x\ y.\ x = y \vee (x,\ y) \in var\text{-}order\text{-}rel)$›
$\langle proof \rangle$

sepref-definition *full-quicksort-vars-impl*
  is ‹*full-quicksort-vars*›
  :: ‹$monom\text{-}assn^k \rightarrow_a monom\text{-}assn$›
  ⟨*proof*⟩

lemmas *sort-vars-spec-hnr* =
  full-quicksort-vars-impl.refine[*FCOMP full-quicksort-sort-vars-spec*]

lemma *string-rel-order-map*:
  ‹$(x, a) \in string\text{-}rel \Longrightarrow$
    $(y, aa) \in string\text{-}rel \Longrightarrow$
    $x \leq y \longleftrightarrow a \leq aa$›
  ⟨*proof*⟩

lemma *merge-monoms-merge-monoms*:
  ‹$(merge\text{-}monoms, merge\text{-}monoms) \in monom\text{-}rel \rightarrow monom\text{-}rel \rightarrow monom\text{-}rel$›
  ⟨*proof*⟩

lemma *merge-monoms-merge-monoms2*:
  ‹$(a, b) \in monom\text{-}rel \Longrightarrow (a', b') \in monom\text{-}rel \Longrightarrow$
    $(merge\text{-}monoms\ a\ a', merge\text{-}monoms\ b\ b') \in monom\text{-}rel$›
  ⟨*proof*⟩

lemma *msort-monoms-impl*:
  ‹$(msort\text{-}monoms\text{-}impl, merge\text{-}monoms\text{-}poly)$
   $\in monom\text{-}rel \rightarrow monom\text{-}rel$›
  ⟨*proof*⟩

lemma *merge-sort-monoms-sort-monoms-spec*:
  ‹$(RETURN\ o\ merge\text{-}monoms\text{-}poly, sort\text{-}coeff) \in \langle Id\rangle list\text{-}rel \rightarrow_f \langle\langle Id\rangle list\text{-}rel\rangle nres\text{-}rel$›
  ⟨*proof*⟩

sepref-register *sort-coeff*
lemma [*sepref-fr-rules*]:
  ‹$(return\ o\ msort\text{-}monoms\text{-}impl, sort\text{-}coeff) \in monom\text{-}assn^k \rightarrow_a monom\text{-}assn$›
  ⟨*proof*⟩

sepref-definition *sort-all-coeffs-impl*
  is ‹*sort-all-coeffs*›
  :: ‹$poly\text{-}assn^k \rightarrow_a poly\text{-}assn$›
  ⟨*proof*⟩

declare *sort-all-coeffs-impl.refine*[*sepref-fr-rules*]

lemma *merge-coeffs0-alt-def*:
  ‹$(RETURN\ o\ merge\text{-}coeffs0)\ p =$
   $REC_T(\lambda f\ p.$
    $(case\ p\ of$
      $[] \Rightarrow RETURN\ []$
    $|\ [p] => if\ snd\ p = 0\ then\ RETURN\ []\ else\ RETURN\ [p]$
    $|\ ((xs, n)\ \#\ (ys, m)\ \#\ p) \Rightarrow$
     $(if\ xs = ys$
      $then\ if\ n + m \neq 0\ then\ f\ ((xs, n + m)\ \#\ p)\ else\ f\ p$

else if $n = 0$ *then*
  do $\{p \leftarrow f\ ((ys,\ m)\ \#\ p)$;
    RETURN $p\}$
else do $\{p \leftarrow f\ ((ys,\ m)\ \#\ p)$;
    RETURN $((xs,\ n)\ \#\ p)\})))$
p›
⟨*proof*⟩

Again, Sepref does not understand what is going here.

sepref-definition *merge-coeffs0-impl*
  is ‹*RETURN o merge-coeffs0*›
  :: ‹$poly\text{-}assn^k \rightarrow_a poly\text{-}assn$›
  ⟨*proof*⟩

declare *merge-coeffs0-impl.refine*[*sepref-fr-rules*]

sepref-definition *fully-normalize-poly-impl*
  is ‹*full-normalize-poly*›
  :: ‹$poly\text{-}assn^k \rightarrow_a poly\text{-}assn$›
  ⟨*proof*⟩

declare *fully-normalize-poly-impl.refine*[*sepref-fr-rules*]

end
theory *PAC-Version*
  imports *Main*
begin

This code was taken from IsaFoR and adapted to git.

local-setup ‹
  let
    val version = *2020−AFP*
(∗   *trim-line* (*#1* (*Isabelle-System.bash-output* (*cd* $*ISAFOL/* && *git rev−parse −−short HEAD* ||
echo unknown))) ∗)
  in
    Local-Theory.define
      ((***binding***‹*version*›, *NoSyn*),
        ((***binding***‹*version-def*›, []), *HOLogic.mk-literal version*)) *#> #2*
  end
›

declare *version-def* [*code*]

end
theory *PAC-Checker-Synthesis*
  imports *PAC-Checker WB-Sort PAC-Checker-Relation*
    PAC-Checker-Init More-Loops PAC-Version
begin

13 Code Synthesis of the Complete Checker

We here combine refine the full checker, using the initialisation provided in another file.

abbreviation *vars-assn* **where**
‹*vars-assn* ≡ *hs.assn string-assn*›

fun *vars-of-monom-in* **where**
‹*vars-of-monom-in* [] - = *True*› |
‹*vars-of-monom-in* (*x* # *xs*) $\mathcal{V}$ ⟷ $x \in \mathcal{V}$ ∧ *vars-of-monom-in xs* $\mathcal{V}$›

fun *vars-of-poly-in* **where**
‹*vars-of-poly-in* [] - = *True*› |
‹*vars-of-poly-in* ((*x*, -) # *xs*) $\mathcal{V}$ ⟷ *vars-of-monom-in x* $\mathcal{V}$ ∧ *vars-of-poly-in xs* $\mathcal{V}$›

lemma *vars-of-monom-in-alt-def*:
‹*vars-of-monom-in xs* $\mathcal{V}$ ⟷ *set xs* ⊆ $\mathcal{V}$›
⟨*proof*⟩

lemma *vars-llist-alt-def*:
‹*vars-llist xs* ⊆ $\mathcal{V}$ ⟷ *vars-of-poly-in xs* $\mathcal{V}$›
⟨*proof*⟩

lemma *vars-of-monom-in-alt-def2*:
‹*vars-of-monom-in xs* $\mathcal{V}$ ⟷ *fold* (λ*x b*. *b* ∧ $x \in \mathcal{V}$) *xs True*›
⟨*proof*⟩

sepref-definition *vars-of-monom-in-impl*
is ‹*uncurry* (*RETURN oo vars-of-monom-in*)›
:: ‹(*list-assn string-assn*)^k $*_a$ *vars-assn*^{k} $\rightarrow_a$ *bool-assn*›
⟨*proof*⟩

declare *vars-of-monom-in-impl.refine*[*sepref-fr-rules*]

lemma *vars-of-poly-in-alt-def2*:
‹*vars-of-poly-in xs* $\mathcal{V}$ ⟷ *fold* (λ(*x*, -) *b*. *b* ∧ *vars-of-monom-in x* $\mathcal{V}$) *xs True*›
⟨*proof*⟩

sepref-definition *vars-of-poly-in-impl*
is ‹*uncurry* (*RETURN oo vars-of-poly-in*)›
:: ‹(*poly-assn*)^k $*_a$ *vars-assn*^{k} $\rightarrow_a$ *bool-assn*›
⟨*proof*⟩

declare *vars-of-poly-in-impl.refine*[*sepref-fr-rules*]

definition *union-vars-monom* :: ‹*string list* ⇒ *string set* ⇒ *string set*› **where**
‹*union-vars-monom xs* $\mathcal{V}$ = *fold insert xs* $\mathcal{V}$›

definition *union-vars-poly* :: ‹*llist-polynom* ⇒ *string set* ⇒ *string set*› **where**
‹*union-vars-poly xs* $\mathcal{V}$ = *fold* (λ(*xs*, -) $\mathcal{V}$. *union-vars-monom xs* $\mathcal{V}$) *xs* $\mathcal{V}$›

lemma *union-vars-monom-alt-def*:
‹*union-vars-monom xs* $\mathcal{V}$ = $\mathcal{V}$ ∪ *set xs*›
⟨*proof*⟩

lemma *union-vars-poly-alt-def*:
‹*union-vars-poly xs* $\mathcal{V}$ = $\mathcal{V}$ ∪ *vars-llist xs*›

⟨*proof*⟩

sepref-definition *union-vars-monom-impl*
is ‹*uncurry* (*RETURN oo union-vars-monom*)›
:: ‹$monom\text{-}assn^k *_a vars\text{-}assn^d \rightarrow_a vars\text{-}assn$›
⟨*proof*⟩

declare *union-vars-monom-impl.refine*[*sepref-fr-rules*]

sepref-definition *union-vars-poly-impl*
is ‹*uncurry* (*RETURN oo union-vars-poly*)›
:: ‹$poly\text{-}assn^k *_a vars\text{-}assn^d \rightarrow_a vars\text{-}assn$›
⟨*proof*⟩

declare *union-vars-poly-impl.refine*[*sepref-fr-rules*]

hide-const (**open**) *Autoref-Fix-Rel.CONSTRAINT*

fun *status-assn* **where**
‹*status-assn* - *CSUCCESS CSUCCESS* = *emp*› |
‹*status-assn* - *CFOUND CFOUND* = *emp*› |
‹*status-assn R* (*CFAILED a*) (*CFAILED b*) = *R a b*› |
‹*status-assn* - - - = *false*›

lemma *SUCCESS-hnr*[*sepref-fr-rules*]:
‹(*uncurry0* (*return CSUCCESS*), *uncurry0* (*RETURN CSUCCESS*)) ∈ $unit\text{-}assn^k \rightarrow_a status\text{-}assn\ R$›
⟨*proof*⟩

lemma *FOUND-hnr*[*sepref-fr-rules*]:
‹(*uncurry0* (*return CFOUND*), *uncurry0* (*RETURN CFOUND*)) ∈ $unit\text{-}assn^k \rightarrow_a status\text{-}assn\ R$›
⟨*proof*⟩

lemma *is-success-hnr*[*sepref-fr-rules*]:
‹*CONSTRAINT is-pure R* ⟹
((*return o is-cfound*), (*RETURN o is-cfound*)) ∈ $(status\text{-}assn\ R)^k \rightarrow_a bool\text{-}assn$›
⟨*proof*⟩

lemma *is-cfailed-hnr*[*sepref-fr-rules*]:
‹*CONSTRAINT is-pure R* ⟹
((*return o is-cfailed*), (*RETURN o is-cfailed*)) ∈ $(status\text{-}assn\ R)^k \rightarrow_a bool\text{-}assn$›
⟨*proof*⟩

lemma *merge-cstatus-hnr*[*sepref-fr-rules*]:
‹*CONSTRAINT is-pure R* ⟹
(*uncurry* (*return oo merge-cstatus*), *uncurry* (*RETURN oo merge-cstatus*)) ∈
$(status\text{-}assn\ R)^k *_a (status\text{-}assn\ R)^k \rightarrow_a status\text{-}assn\ R$›
⟨*proof*⟩

sepref-definition *add-poly-impl*
is ‹*add-poly-l*›
:: ‹$(poly\text{-}assn \times_a poly\text{-}assn)^k \rightarrow_a poly\text{-}assn$›
⟨*proof*⟩

declare *add-poly-impl.refine*[*sepref-fr-rules*]

sepref-register *mult-monomials*
lemma *mult-monoms-alt-def*:
‹(*RETURN oo mult-monoms*) *x y* = REC_T
  (λf (*p*, *q*).
    case (*p*, *q*) *of*
      ([], -) ⇒ *RETURN q*
    | (-, []) ⇒ *RETURN p*
    | (*x* # *p*, *y* # *q*) ⇒
      (*if x* = *y then do* {
        pq ← *f* (*p*, *q*);
         RETURN (*x* # *pq*)}
      else if (*x*, *y*) ∈ *var-order-rel*
      then do {
        pq ← *f* (*p*, *y* # *q*);
        RETURN (*x* # *pq*)}
      else do {
        pq ← *f* (*x* # *p*, *q*);
        RETURN (*y* # *pq*)}))
   (*x*, *y*)›
⟨*proof*⟩

sepref-definition *mult-monoms-impl*
is ‹*uncurry* (*RETURN oo mult-monoms*)›
:: ‹$(monom\text{-}assn)^k *_a (monom\text{-}assn)^k \rightarrow_a (monom\text{-}assn)$›
⟨*proof*⟩

declare *mult-monoms-impl.refine*[*sepref-fr-rules*]

sepref-definition *mult-monomials-impl*
is ‹*uncurry* (*RETURN oo mult-monomials*)›
:: ‹$(monomial\text{-}assn)^k *_a (monomial\text{-}assn)^k \rightarrow_a (monomial\text{-}assn)$›
⟨*proof*⟩

lemma *map-append-alt-def2*:
‹(*RETURN o* (*map-append f b*)) *xs* = REC_T
  (λg *xs*. *case xs of* [] ⇒ *RETURN b*
   | *x* # *xs* ⇒ *do* {
       y ← *g xs*;
       RETURN (*f x* # *y*)
   }) *xs*›
 ⟨*proof*⟩

definition *map-append-poly-mult* **where**
‹*map-append-poly-mult x* = *map-append* (*mult-monomials x*)›

declare *mult-monomials-impl.refine*[*sepref-fr-rules*]

sepref-definition *map-append-poly-mult-impl*
is ‹*uncurry2* (*RETURN ooo map-append-poly-mult*)›

:: ‹$monomial\text{-}assn^k *_a poly\text{-}assn^k *_a poly\text{-}assn^k \rightarrow_a poly\text{-}assn$›
⟨*proof*⟩

declare *map-append-poly-mult-impl.refine*[*sepref-fr-rules*]

TODO *foldl* ($\lambda l\ x.\ l$ @ [*?f x*]) [] *?l* = *map ?f ?l* is the worst possible implementation of map!

sepref-definition *mult-poly-raw-impl*
is ‹*uncurry* (*RETURN oo mult-poly-raw*)›
:: ‹$poly\text{-}assn^k *_a poly\text{-}assn^k \rightarrow_a poly\text{-}assn$›
⟨*proof*⟩

declare *mult-poly-raw-impl.refine*[*sepref-fr-rules*]

sepref-definition *mult-poly-impl*
is ‹*uncurry mult-poly-full*›
:: ‹$poly\text{-}assn^k *_a poly\text{-}assn^k \rightarrow_a poly\text{-}assn$›
⟨*proof*⟩

declare *mult-poly-impl.refine*[*sepref-fr-rules*]

lemma *inverse-monomial*:
‹$monom\text{-}rel^{-1} \times_r int\text{-}rel = (monom\text{-}rel \times_r int\text{-}rel)^{-1}$›
⟨*proof*⟩

lemma *eq-poly-rel-eq*[*sepref-import-param*]:
‹((=), (=)) ∈ *poly-rel* → *poly-rel* → *bool-rel*›
⟨*proof*⟩

sepref-definition *weak-equality-l-impl*
is ‹*uncurry weak-equality-l*›
:: ‹$poly\text{-}assn^k *_a poly\text{-}assn^k \rightarrow_a bool\text{-}assn$›
⟨*proof*⟩

declare *weak-equality-l-impl.refine*[*sepref-fr-rules*]
sepref-register *add-poly-l mult-poly-full*

abbreviation *raw-string-assn* :: ‹*string* ⇒ *string* ⇒ *assn*› **where**
‹*raw-string-assn* ≡ *list-assn id-assn*›

definition *show-nat* :: ‹*nat* ⇒ *string*› **where**
‹*show-nat i* = *show i*›

lemma [*sepref-import-param*]:
‹(*show-nat*, *show-nat*) ∈ *nat-rel* → ⟨*Id*⟩*list-rel*›
⟨*proof*⟩

lemma *status-assn-pure-conv*:
‹*status-assn* (*id-assn*) *a b* = *id-assn a b*›
⟨*proof*⟩

lemma [*sepref-fr-rules*]:
‹(*uncurry3* ($\lambda x\ y$. *return oo* (*error-msg-not-equal-dom x y*)), *uncurry3 check-not-equal-dom-err*) ∈ $poly\text{-}assn^k *_a poly\text{-}assn^k *_a poly\text{-}assn^k *_a poly\text{-}assn^k \rightarrow_a raw\text{-}string\text{-}assn$›

⟨proof⟩

lemma [*sepref-fr-rules*]:
‹(*return o* (*error-msg-notin-dom o nat-of-uint64*), *RETURN o error-msg-notin-dom*)
$\in$ *uint64-nat-assn*^{k} $\rightarrow_a$ *raw-string-assn*›
‹(*return o* (*error-msg-reused-dom o nat-of-uint64*), *RETURN o error-msg-reused-dom*)
$\in$ *uint64-nat-assn*^{k} $\rightarrow_a$ *raw-string-assn*›
‹(*uncurry* (*return oo* ($\lambda i.$ *error-msg* (*nat-of-uint64 i*))), *uncurry* (*RETURN oo error-msg*))
$\in$ *uint64-nat-assn*^{k} $*_a$ *raw-string-assn*^{k} $\rightarrow_a$ *status-assn raw-string-assn*›
‹(*uncurry* (*return oo error-msg*), *uncurry* (*RETURN oo error-msg*))
$\in$ *nat-assn*^{k} $*_a$ *raw-string-assn*^{k} $\rightarrow_a$ *status-assn raw-string-assn*›
⟨proof⟩

sepref-definition *check-addition-l-impl*
is ‹*uncurry6 check-addition-l*›
:: ‹*poly-assn*^{k} $*_a$ *polys-assn*^{k} $*_a$ *vars-assn*^{k} $*_a$ *uint64-nat-assn*^{k} $*_a$ *uint64-nat-assn*^{k} $*_a$
uint64-nat-assn^k $*_a$ *poly-assn*^{k} $\rightarrow_a$ *status-assn raw-string-assn*›
⟨proof⟩

declare *check-addition-l-impl.refine*[*sepref-fr-rules*]

sepref-register *check-mult-l-dom-err*

definition *check-mult-l-dom-err-impl* **where**
‹*check-mult-l-dom-err-impl pd p ia i* =
(*if pd then ''The polynomial with id '' @ show* (*nat-of-uint64 p*) @ *'' was not found'' else ''''*) @
(*if ia then ''The id of the resulting id '' @ show* (*nat-of-uint64 i*) @ *'' was already given'' else ''''*)›

definition *check-mult-l-mult-err-impl* **where**
‹*check-mult-l-mult-err-impl p q pq r* =
''Multiplying '' @ show p @ '' by '' @ show q @ '' gives '' @ show pq @ '' and not '' @ show r›

lemma [*sepref-fr-rules*]:
‹(*uncurry3* (($\lambda x\ y.$ *return oo* (*check-mult-l-dom-err-impl x y*))),
uncurry3 (*check-mult-l-dom-err*)) $\in$ *bool-assn*^{k} $*_a$ *uint64-nat-assn*^{k} $*_a$ *bool-assn*^{k} $*_a$ *uint64-nat-assn*^{k}
$\rightarrow_a$ *raw-string-assn*›
⟨proof⟩

lemma [*sepref-fr-rules*]:
‹(*uncurry3* (($\lambda x\ y.$ *return oo* (*check-mult-l-mult-err-impl x y*))),
uncurry3 (*check-mult-l-mult-err*)) $\in$ *poly-assn*^{k} $*_a$ *poly-assn*^{k} $*_a$ *poly-assn*^{k} $*_a$ *poly-assn*^{k} $\rightarrow_a$ *raw-string-assn*›
⟨proof⟩

sepref-definition *check-mult-l-impl*
is ‹*uncurry6 check-mult-l*›
:: ‹*poly-assn*^{k} $*_a$ *polys-assn*^{k} $*_a$ *vars-assn*^{k} $*_a$ *uint64-nat-assn*^{k} $*_a$ *poly-assn*^{k} $*_a$ *uint64-nat-assn*^{k} $*_a$
poly-assn^k $\rightarrow_a$ *status-assn raw-string-assn*›
⟨proof⟩

declare *check-mult-l-impl.refine*[*sepref-fr-rules*]

definition *check-ext-l-dom-err-impl* :: ‹*uint64* ⇒ -› **where**
‹*check-ext-l-dom-err-impl p* =

$''There\ is\ already\ a\ polynomial\ with\ index\ ''$ @ *show* (*nat-of-uint64 p*)›

lemma [*sepref-fr-rules*]:
‹(((*return o* (*check-ext-l-dom-err-impl*))),
(*check-extension-l-dom-err*)) ∈ *uint64-nat-assn*^{k} $\rightarrow_a$ *raw-string-assn*›
⟨*proof*⟩

definition *check-extension-l-no-new-var-err-impl* :: ‹- ⇒ -› **where**
‹*check-extension-l-no-new-var-err-impl p* =
$''No\ new\ variable\ could\ be\ found\ in\ polynomial\ ''$ @ *show p*›

lemma [*sepref-fr-rules*]:
‹(((*return o* (*check-extension-l-no-new-var-err-impl*))),
(*check-extension-l-no-new-var-err*)) ∈ *poly-assn*^{k} $\rightarrow_a$ *raw-string-assn*›
⟨*proof*⟩

definition *check-extension-l-side-cond-err-impl* :: ‹- ⇒ -› **where**
‹*check-extension-l-side-cond-err-impl v p r s* =
$''Error\ while\ checking\ side\ conditions\ of\ extensions\ polynow,\ var\ is\ ''$ @ *show v* @
$''\ polynomial\ is\ ''$ @ *show p* @ $''side\ condition\ p{*}p - p = ''$ @ *show s* @ $''\ and\ should\ be\ 0''$›

lemma [*sepref-fr-rules*]:
‹((*uncurry3* (λ*x y*. *return oo* (*check-extension-l-side-cond-err-impl x y*))),
uncurry3 (*check-extension-l-side-cond-err*)) ∈ *string-assn*^{k} $*_a$ *poly-assn*^{k} $*_a$ *poly-assn*^{k} $*_a$ *poly-assn*^{k}
$\rightarrow_a$ *raw-string-assn*›
⟨*proof*⟩

definition *check-extension-l-new-var-multiple-err-impl* :: ‹- ⇒ -› **where**
‹*check-extension-l-new-var-multiple-err-impl v p* =
$''Error\ while\ checking\ side\ conditions\ of\ extensions\ polynow,\ var\ is\ ''$ @ *show v* @
$''\ but\ it\ either\ appears\ at\ least\ once\ in\ the\ polynomial\ or\ another\ new\ variable\ is\ created\ ''$ @
show p @ $''\ but\ should\ not.''$›

lemma [*sepref-fr-rules*]:
‹((*uncurry* (*return oo* (*check-extension-l-new-var-multiple-err-impl*))),
uncurry (*check-extension-l-new-var-multiple-err*)) ∈ *string-assn*^{k} $*_a$ *poly-assn*^{k} $\rightarrow_a$ *raw-string-assn*›
⟨*proof*⟩

sepref-register *check-extension-l-dom-err fmlookup′*
check-extension-l-side-cond-err check-extension-l-no-new-var-err
check-extension-l-new-var-multiple-err

definition *uminus-poly* :: ‹*llist-polynom* ⇒ *llist-polynom*› **where**
‹*uminus-poly p′* = *map* (λ(*a*, *b*). (*a*, − *b*)) *p′*›

sepref-register *uminus-poly*
lemma [*sepref-import-param*]:
‹(*map* (λ(*a*, *b*). (*a*, − *b*)), *uminus-poly*) ∈ *poly-rel* → *poly-rel*›
⟨*proof*⟩

sepref-register *vars-of-poly-in*
weak-equality-l

lemma [*safe-constraint-rules*]:
‹*Sepref-Constraints.CONSTRAINT single-valued* (*the-pure monomial-assn*)› **and**
single-valued-the-monomial-assn:
‹*single-valued* (*the-pure monomial-assn*)›
‹*single-valued* ((*the-pure monomial-assn*)$^{-1}$)›
⟨*proof*⟩

sepref-definition *check-extension-l-impl*
is ‹*uncurry5 check-extension-l*›
:: ‹*poly-assn*^{k} $*_a$ *polys-assn*^{k} $*_a$ *vars-assn*^{k} $*_a$ *uint64-nat-assn*^{k} $*_a$ *string-assn*^{k} $*_a$ *poly-assn*^{k} $\rightarrow_a$ *status-assn raw-string-assn*›
⟨*proof*⟩

declare *check-extension-l-impl.refine*[*sepref-fr-rules*]

sepref-definition *check-del-l-impl*
is ‹*uncurry2 check-del-l*›
:: ‹*poly-assn*^{k} $*_a$ *polys-assn*^{k} $*_a$ *uint64-nat-assn*^{k} $\rightarrow_a$ *status-assn raw-string-assn*›
⟨*proof*⟩

lemmas [*sepref-fr-rules*] = *check-del-l-impl.refine*

abbreviation *pac-step-rel* **where**
‹*pac-step-rel* ≡ *p2rel* (⟨*Id*, ⟨*monomial-rel*⟩*list-rel*, *Id*⟩ *pac-step-rel-raw*)›

sepref-register *PAC-Polynomials-Operations.normalize-poly*
pac-src1 pac-src2 new-id pac-mult case-pac-step check-mult-l
check-addition-l check-del-l check-extension-l

lemma *pac-step-rel-assn-alt-def2*:
‹*hn-ctxt* (*pac-step-rel-assn nat-assn poly-assn id-assn*) *b bi* =
hn-val
(*p2rel*
((⟨*nat-rel*, *poly-rel*, *Id* :: (*string* × -) *set*⟩*pac-step-rel-raw*)) *b bi*›
⟨*proof*⟩

lemma *is-AddD-import*[*sepref-fr-rules*]:
assumes ‹*CONSTRAINT is-pure K*› ‹*CONSTRAINT is-pure V*›
shows
‹(*return o pac-res*, *RETURN o pac-res*) ∈ [λ*x*. *is-Add x* ∨ *is-Mult x* ∨ *is-Extension x*]$_a$ (*pac-step-rel-assn K V R*)^k → *V*›
‹(*return o pac-src1*, *RETURN o pac-src1*) ∈ [λ*x*. *is-Add x* ∨ *is-Mult x* ∨ *is-Del x*]$_a$ (*pac-step-rel-assn K V R*)^k → *K*›
‹(*return o new-id*, *RETURN o new-id*) ∈ [λ*x*. *is-Add x* ∨ *is-Mult x* ∨ *is-Extension x*]$_a$ (*pac-step-rel-assn K V R*)^k → *K*›
‹(*return o is-Add*, *RETURN o is-Add*) ∈ (*pac-step-rel-assn K V R*)^k $\rightarrow_a$ *bool-assn*›
‹(*return o is-Mult*, *RETURN o is-Mult*) ∈ (*pac-step-rel-assn K V R*)^k $\rightarrow_a$ *bool-assn*›
‹(*return o is-Del*, *RETURN o is-Del*) ∈ (*pac-step-rel-assn K V R*)^k $\rightarrow_a$ *bool-assn*›
‹(*return o is-Extension*, *RETURN o is-Extension*) ∈ (*pac-step-rel-assn K V R*)^k $\rightarrow_a$ *bool-assn*›
⟨*proof*⟩

lemma [*sepref-fr-rules*]:
‹*CONSTRAINT is-pure K* ⟹

$(return\ o\ pac\text{-}src2,\ RETURN\ o\ pac\text{-}src2) \in [\lambda x.\ is\text{-}Add\ x]_a\ (pac\text{-}step\text{-}rel\text{-}assn\ K\ V\ R)^k \to K$›
‹$CONSTRAINT\ is\text{-}pure\ V \Longrightarrow$
$(return\ o\ pac\text{-}mult,\ RETURN\ o\ pac\text{-}mult) \in [\lambda x.\ is\text{-}Mult\ x]_a\ (pac\text{-}step\text{-}rel\text{-}assn\ K\ V\ R)^k \to V$›
‹$CONSTRAINT\ is\text{-}pure\ R \Longrightarrow$
$(return\ o\ new\text{-}var,\ RETURN\ o\ new\text{-}var) \in [\lambda x.\ is\text{-}Extension\ x]_a\ (pac\text{-}step\text{-}rel\text{-}assn\ K\ V\ R)^k \to R$›
⟨*proof*⟩

lemma *is-Mult-lastI*:
‹$\neg\ is\text{-}Add\ b \Longrightarrow \neg is\text{-}Mult\ b \Longrightarrow \neg is\text{-}Extension\ b \Longrightarrow is\text{-}Del\ b$›
⟨*proof*⟩

sepref-register *is-cfailed is-Del*

definition *PAC-checker-l-step′* :: - **where**
‹*PAC-checker-l-step′ a b c d = PAC-checker-l-step a (b, c, d)*›

lemma *PAC-checker-l-step-alt-def*:
‹*PAC-checker-l-step a bcd e = (let (b,c,d) = bcd in PAC-checker-l-step′ a b c d e)*›
⟨*proof*⟩

sepref-decl-intf (*′k*) *acode-status* **is** (*′k*) *code-status*
sepref-decl-intf (*′k*, *′b*, *′lbl*) *apac-step* **is** (*′k*, *′b*, *′lbl*) *pac-step*

sepref-register *merge-cstatus full-normalize-poly new-var is-Add*

lemma *poly-rel-the-pure*:
‹*poly-rel = the-pure poly-assn*› **and**
nat-rel-the-pure:
‹*nat-rel = the-pure nat-assn*› **and**
WTF-RF: ‹*pure (the-pure nat-assn) = nat-assn*›
⟨*proof*⟩

lemma [*safe-constraint-rules*]:
‹*CONSTRAINT IS-LEFT-UNIQUE uint64-nat-rel*› **and**
single-valued-uint64-nat-rel[*safe-constraint-rules*]:
‹*CONSTRAINT single-valued uint64-nat-rel*›
⟨*proof*⟩

sepref-definition *check-step-impl*
is ‹*uncurry4 PAC-checker-l-step′*›
:: ‹$poly\text{-}assn^k *_a (status\text{-}assn\ raw\text{-}string\text{-}assn)^d *_a vars\text{-}assn^d *_a polys\text{-}assn^d *_a (pac\text{-}step\text{-}rel\text{-}assn$
$(uint64\text{-}nat\text{-}assn)\ poly\text{-}assn\ (string\text{-}assn :: string \Rightarrow -))^d \to_a$
$status\text{-}assn\ raw\text{-}string\text{-}assn \times_a vars\text{-}assn \times_a polys\text{-}assn$›
⟨*proof*⟩

declare *check-step-impl.refine*[*sepref-fr-rules*]

sepref-register *PAC-checker-l-step PAC-checker-l-step′ fully-normalize-poly-impl*

definition *PAC-checker-l′* **where**
‹*PAC-checker-l′ p 𝒱 A status steps = PAC-checker-l p (𝒱, A) status steps*›

lemma *PAC-checker-l-alt-def*:
‹*PAC-checker-l p 𝒱A status steps =*

$(let\ (\mathcal{V},\ A) = \mathcal{V}A\ in\ PAC\text{-}checker\text{-}l'\ p\ \mathcal{V}\ A\ status\ steps)$›
⟨*proof*⟩

sepref-definition *PAC-checker-l-impl*
is ‹*uncurry4 PAC-checker-l'*›
:: ‹$poly\text{-}assn^k *_a vars\text{-}assn^d *_a polys\text{-}assn^d *_a (status\text{-}assn\ raw\text{-}string\text{-}assn)^d *_a$
$(list\text{-}assn\ (pac\text{-}step\text{-}rel\text{-}assn\ (uint64\text{-}nat\text{-}assn)\ poly\text{-}assn\ string\text{-}assn))^k \rightarrow_a$
$status\text{-}assn\ raw\text{-}string\text{-}assn \times_a vars\text{-}assn \times_a polys\text{-}assn$›
⟨*proof*⟩

declare *PAC-checker-l-impl.refine*[*sepref-fr-rules*]

abbreviation *polys-assn-input* **where**
‹$polys\text{-}assn\text{-}input \equiv iam\text{-}fmap\text{-}assn\ nat\text{-}assn\ poly\text{-}assn$›

definition *remap-polys-l-dom-err-impl* :: ‹-› **where**
‹*remap-polys-l-dom-err-impl* =
″*Error during initialisation. Too many polynomials where provided. If this happens,*″ @
″*please report the example to the authors, because something went wrong during* ″ @
″*code generation* (*code generation to arrays is likely to be broken*).″›

lemma [*sepref-fr-rules*]:
‹((*uncurry0* (*return* (*remap-polys-l-dom-err-impl*))),
uncurry0 (*remap-polys-l-dom-err*)) $\in unit\text{-}assn^k \rightarrow_a raw\text{-}string\text{-}assn$›
⟨*proof*⟩

MLton is not able to optimise the calls to pow.

lemma *pow-2-64*: ‹$(2::nat)\ \hat{}\ 64 = 18446744073709551616$›
⟨*proof*⟩

sepref-register *upper-bound-on-dom op-fmap-empty*

sepref-definition *remap-polys-l-impl*
is ‹*uncurry2 remap-polys-l2*›
:: ‹$poly\text{-}assn^k *_a vars\text{-}assn^d *_a polys\text{-}assn\text{-}input^d \rightarrow_a$
$status\text{-}assn\ raw\text{-}string\text{-}assn \times_a vars\text{-}assn \times_a polys\text{-}assn$›
⟨*proof*⟩

lemma *remap-polys-l2-remap-polys-l*:
‹(*uncurry2 remap-polys-l2*, *uncurry2 remap-polys-l*) $\in (Id \times_r \langle Id\rangle set\text{-}rel) \times_r Id \rightarrow_f \langle Id\rangle nres\text{-}rel$›
⟨*proof*⟩

lemma [*sepref-fr-rules*]:
‹(*uncurry2 remap-polys-l-impl*,
uncurry2 remap-polys-l) $\in poly\text{-}assn^k *_a vars\text{-}assn^d *_a polys\text{-}assn\text{-}input^d \rightarrow_a$
$status\text{-}assn\ raw\text{-}string\text{-}assn \times_a vars\text{-}assn \times_a polys\text{-}assn$›
⟨*proof*⟩

sepref-register *remap-polys-l*

sepref-definition *full-checker-l-impl*
is ‹*uncurry2 full-checker-l*›
:: ‹$poly\text{-}assn^k *_a polys\text{-}assn\text{-}input^d *_a (list\text{-}assn\ (pac\text{-}step\text{-}rel\text{-}assn\ (uint64\text{-}nat\text{-}assn)\ poly\text{-}assn\ string\text{-}assn))^k$
$\rightarrow_a$
$status\text{-}assn\ raw\text{-}string\text{-}assn \times_a vars\text{-}assn \times_a polys\text{-}assn$›

⟨*proof*⟩

sepref-definition *PAC-update-impl*
is ‹*uncurry2* (*RETURN ooo fmupd*)›
:: ‹$nat\text{-}assn^k *_a poly\text{-}assn^k *_a (polys\text{-}assn\text{-}input)^d \rightarrow_a polys\text{-}assn\text{-}input$›
⟨*proof*⟩

sepref-definition *PAC-empty-impl*
is ‹*uncurry0* (*RETURN fmempty*)›
:: ‹$unit\text{-}assn^k \rightarrow_a polys\text{-}assn\text{-}input$›
⟨*proof*⟩

sepref-definition *empty-vars-impl*
is ‹*uncurry0* (*RETURN* {})›
:: ‹$unit\text{-}assn^k \rightarrow_a vars\text{-}assn$›
⟨*proof*⟩

This is a hack for performance. There is no need to recheck that that a char is valid when working on chars coming from strings... It is not that important in most cases, but in our case the preformance difference is really large.

definition *unsafe-asciis-of-literal* :: ‹-› **where**
‹*unsafe-asciis-of-literal xs* = *String.asciis-of-literal xs*›

definition *unsafe-asciis-of-literal′* :: ‹-› **where**
[*simp*, *symmetric*, *code*]: ‹*unsafe-asciis-of-literal′* = *unsafe-asciis-of-literal*›

code-printing
constant *unsafe-asciis-of-literal′* ⇀
(*SML*) !(*List.map* (*fn c* => *let val k* = *Char.ord c in IntInf.fromInt k end*) /*o String.explode*)

Now comes the big and ugly and unsafe hack.

Basically, we try to avoid the conversion to IntInf when calculating the hash. The performance gain is roughly 40%, which is a LOT and definitively something we need to do. We are aware that the SML semantic encourages compilers to optimise conversions, but this does not happen here, corroborating our early observation on the verified SAT solver IsaSAT.x

definition *raw-explode* **where**
[*simp*]: ‹*raw-explode* = *String.explode*›
code-printing
constant *raw-explode* ⇀
(*SML*) *String.explode*

definition ‹*hashcode-literal′ s* ≡
foldl (λ*h x*. *h* ∗ *33* + *uint32-of-int* (*of-char x*)) *5381*
(*raw-explode s*)›

lemmas [*code*] =
hashcode-literal-def[*unfolded String.explode-code*
unsafe-asciis-of-literal-def[*symmetric*]]

definition *uint32-of-char* **where**
[*symmetric*, *code-unfold*]: ‹*uint32-of-char x* = *uint32-of-int* (*int-of-char x*)›

code-printing

constant *uint32-of-char* ⇀
(*SML*) !(*Word32.fromInt* /*o* (*Char.ord*))

lemma [*code*]: ‹*hashcode s* = *hashcode-literal′ s*›
⟨*proof*⟩

We do not include

export-code *PAC-checker-l-impl PAC-update-impl PAC-empty-impl the-error is-cfailed is-cfound int-of-integer Del Add Mult nat-of-integer String.implode remap-polys-l-impl fully-normalize-poly-impl union-vars-poly-impl empty-vars-impl full-checker-l-impl check-step-impl CSUCCESS Extension hashcode-literal′ version*
in *SML-imp* **module-name** *PAC-Checker*
file-prefix *checker*

compile-generated-files -
external-files
‹*code/parser.sml*›
‹*code/pasteque.sml*›
‹*code/pasteque.mlb*›
where ‹*fn dir* =>
let
val exec = *Generated-Files.execute* (*Path.append dir* (*Path.basic code*));
val - = *exec* ‹*rename file*› *mv checker.ML checker.sml*
val - =
exec ‹*Compilation*›
(*File.bash-path* ***path*** ‹$*ISABELLE-MLTON*› ^ ^
−*const* '*MLton.safe false*' −*verbose 1* −*default*−*type int64* −*output pasteque* ^
−*codegen native* −*inline 700* −*cc*−*opt* −*O3 pasteque.mlb*);
in () *end*›

14 Correctness theorem

context *poly-embed*
begin

definition *full-poly-assn* **where**
‹*full-poly-assn* = *hr-comp poly-assn* (*fully-unsorted-poly-rel O mset-poly-rel*)›

definition *full-poly-input-assn* **where**
‹*full-poly-input-assn* = *hr-comp*
(*hr-comp polys-assn-input*
(⟨*nat-rel*, *fully-unsorted-poly-rel O mset-poly-rel*⟩*fmap-rel*))
polys-rel›

definition *fully-pac-assn* **where**
‹*fully-pac-assn* = (*list-assn*
(*hr-comp* (*pac-step-rel-assn uint64-nat-assn poly-assn string-assn*)
(*p2rel*
(⟨*nat-rel*,
fully-unsorted-poly-rel O
mset-poly-rel, *var-rel*⟩*pac-step-rel-raw*))))›

definition *code-status-assn* **where**

‹$code\text{-}status\text{-}assn = hr\text{-}comp\ (status\text{-}assn\ raw\text{-}string\text{-}assn)$
$code\text{-}status\text{-}status\text{-}rel$›

definition *full-vars-assn* **where**
‹$full\text{-}vars\text{-}assn = hr\text{-}comp\ (hs.assn\ string\text{-}assn)$
$(\langle var\text{-}rel\rangle set\text{-}rel)$›

lemma *polys-rel-full-polys-rel*:
‹$polys\text{-}rel\text{-}full = Id \times_r polys\text{-}rel$›
⟨*proof*⟩

definition *full-polys-assn* :: ‹-› **where**
‹$full\text{-}polys\text{-}assn = hr\text{-}comp\ (hr\text{-}comp\ polys\text{-}assn$
$(\langle nat\text{-}rel,$
$sorted\text{-}poly\text{-}rel\ O\ mset\text{-}poly\text{-}rel\rangle fmap\text{-}rel))$
$polys\text{-}rel$›

Below is the full correctness theorems. It basically states that:

1. assuming that the input polynomials have no duplicate variables

Then:

1. if the checker returns *CFOUND*, the spec is in the ideal and the PAC file is correct
2. if the checker returns *CSUCCESS*, the PAC file is correct (but there is no information on the spec, aka checking failed)
3. if the checker return *CFAILED err*, then checking failed (and *err might* give you an indication of the error, but the correctness theorem does not say anything about that).

   The input parameters are:
4. the specification polynomial represented as a list
5. the input polynomials as hash map (as an array of option polynom)
6. a represention of the PAC proofs.

lemma *PAC-full-correctness*:
‹$(uncurry2\ full\text{-}checker\text{-}l\text{-}impl,$
$uncurry2\ (\lambda spec\ A\ \text{-}.\ PAC\text{-}checker\text{-}specification\ spec\ A))$
$\in (full\text{-}poly\text{-}assn)^k *_a (full\text{-}poly\text{-}input\text{-}assn)^d *_a (fully\text{-}pac\text{-}assn)^k \rightarrow_a hr\text{-}comp$
$(code\text{-}status\text{-}assn \times_a full\text{-}vars\text{-}assn \times_a hr\text{-}comp\ polys\text{-}assn$
$(\langle nat\text{-}rel,\ sorted\text{-}poly\text{-}rel\ O\ mset\text{-}poly\text{-}rel\rangle fmap\text{-}rel))$
$\{((st,\ G),\ st',\ G').$
$st = st' \wedge (st \neq FAILED \longrightarrow (G,\ G') \in Id \times_r polys\text{-}rel)\}$›
⟨*proof*⟩

It would be more efficient to move the parsing to Isabelle, as this would be more memory efficient (and also reduce the TCB). But now comes the fun part: It cannot work. A stream (of a file) is consumed by side effects. Assume that this would work. The code could look like:
Let (*read-file file*) *f*
This code is equal to (in the HOL sense of equality): *let - = read-file file in Let* (*read-file file*) *f*

However, as an hypothetical *read-file* changes the underlying stream, we would get the next token. Remark that this is already a weird point of ML compilers. Anyway, I see currently two solutions to this problem:

1. The meta-argument: use it only in the Refinement Framework in a setup where copies are disallowed. Basically, this works because we can express the non-duplication constraints on the type level. However, we cannot forbid people from expressing things directly at the HOL level.

2. On the target language side, model the stream as the stream and the position. Reading takes two arguments. First, the position to read. Second, the stream (and the current position) to read. If the position to read does not match the current position, return an error. This would fit the correctness theorem of the code generation (roughly "if it terminates without exception, the answer is the same"), but it is still unsatisfactory.

end

definition φ :: ‹*string* ⇒ *nat*› **where**
‹$\varphi = (SOME\ \varphi.\ bij\ \varphi)$›

lemma *bij-*φ: ‹*bij* φ›
⟨*proof*⟩

global-interpretation *PAC*: *poly-embed* **where**
$\varphi = \varphi$
⟨*proof*⟩

The full correctness theorem is (*uncurry2 full-checker-l-impl*, *uncurry2* (λ*spec A -. PAC-checker-specification spec A*)) $\in$ *PAC.full-poly-assn*^{k} $*_a$ *PAC.full-poly-input-assn*^{d} $*_a$ *PAC.fully-pac-assn*^{k} $\rightarrow_a$ *hr-comp* (*PAC.code-status-assn* $\times_a$ *PAC.full-vars-assn* $\times_a$ *hr-comp polys-assn* (⟨*nat-rel*, *sorted-poly-rel O PAC.mset-poly-rel*⟩*fmap-rel*)) $\{((st,\ G),\ st',\ G').\ st = st' \wedge (st \neq FAILED \longrightarrow (G,\ G') \in Id \times_r polys\text{-}rel)\}$.

end

References

[1] D. Kaufmann, M. Fleury, and A. Biere. The proof checkers pacheck and pasteque for the practical algebraic calculus. In O. Strichman and A. Ivrii, editors, *Formal Methods in Computer-Aided Design, FMCAD 2020, September 21-24, 2020.* IEEE, 2020.