

Linear programming notes

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Contents

| | | |
|----------|---------------------------|----------|
| 1 | Supply chain model | 1 |
| 1.1 | Production | 2 |
| 1.2 | Constraints | 3 |
| 1.3 | Value rate | 3 |
| 1.4 | Tensor form | 4 |
| 1.4.1 | Verifiche | 5 |
| 1.5 | Altro | 6 |

Abstract

This document contains notes about Linear programming.

1 Supply chain model

We define a simplified supply chain model as a system that produces products with a chain of product transformations performed by suppliers.

Let be

$A = a_1 \dots a_n$ the set of supplier types

$B = b_1 \dots b_m$ the set of product types.

The supplier can perform only a single transformation for the its duration.

Let us define

n_i the number of suppliers of type $i \in A$.

v_i the value of product $i \in B$

q_{ij} the quantity of produced product $j \in B$ for supplier $i \in A, j$ in production interval

$t_{ij}, i \in A, j \in B$ the production interval for product $j \in B$ in supplier $i \in A, j$

c_{ijk} the quantity of product $k \in B$ consumed to produce the product $j \in B$ in production interval of supplier $i \in A$.

1.1 Production

By now we do not consider the constraint on the availability of consuming products. It will be considered later.

When a supplier of type i is ready to produce we assign a production slot for the product j such that

y_{ij} is the weighted distribution of slots

The value y_{ij} is constrained to

$$0 \leq y_{ij} \leq 1 \quad (1)$$

Moreover we may put the supplier in idle state for a while

z_{i0} is the idle time during a whole production cycle

with

$$z_{i0} \geq 0 \quad (2)$$

The total time of production cycle for supplier i is

$$z_i = \sum_{j \in B} y_{ij} t_{ij} + z_{i0}$$

Let normalize the total time production to 1

$$z_i = \sum_{j \in B} y_{ij} t_{ij} + z_{i0} = 1 \quad (3)$$

During this interval all the suppliers i produce the product j at a slot rate

$$s_{ij} = n_i \frac{y_{ij}}{z_i} = n_i y_{ij} \Big|_{i \in A, j \in B} \quad (4)$$

The production rate of product j for supplier i is

$$p_{ij} = q_{ij} s_{ij} = n_i q_{ij} y_{ij} \Big|_{i \in A, j \in B}$$

While the production rate of product j is

$$p_{j \in B} = \sum_{i \in A,} p_{ij} = \sum_{i \in A,} n_i q_{ij} y_{ij} \quad (5)$$

1.2 Constraints

The (4) expresses the slot rate of a product for specific suppliers.

Therefore we can calculate the consumption rate of a product k to produce product j for suppliers i

$$e_{ijk} = c_{ijk}s_{ij}|_{i \in A, j, k \in B}$$

Then the total consumption rate for product k is

$$e_k = \sum_{i \in A, j \in B} e_{ijk} = \sum_{i \in A, j \in B} n_i c_{ijk} y_{ij} \quad (6)$$

If

$$e_i < p_i$$

the product s produced at a higher rate then it is consumed creating a surplus that can be sold.

On the other hane we cannot have

$$e_i > p_i$$

because it cannot consume more product than the produced.

So the system must satisfy the constraint

$$p_i - e_i \geq 0 \quad (7)$$

1.3 Value rate

We can compute the value rate for the whole supply chain

$$v = \sum_{i \in B} (p_i - e_i) v_i \quad (8)$$

The problem is to find the optimal production configuration that maximize the value rate. This is defined by the linear system composed by (1), (2), (3), (7), (8)

$$\begin{cases} \max_{(y_{ij}, z_{i0})}(v) & , i \in A, j \in B \\ \sum_{j \in B} y_{ij} t_{ij} + z_{i0} = 1 & , i \in A \\ z_{i0} \geq 0 & , i \in A \\ y_{ij} \geq 0 & , i \in A, j \in B \\ y_{ij} \leq 1 & , i \in A, j \in B \\ p_i - e_i \geq 0 & , i \in B \end{cases}$$

1.4 Tensor form

Trasformiamo ora la (5) in forma tensoriale:

$$P_i = A_{ijk} Y_{kj} = AY$$

che deve essere equivalente a

$$p_{j'} = \sum_{i'} n_{i'} q_{i'j'} y_{i'j'} \quad (9)$$

da cui

$$\begin{aligned} A_{iik} &= n_k q_{ki} \\ A_{ijk} &= 0, i \neq j \end{aligned} \quad (10)$$

Se poniamo

$$\begin{aligned} Q_{ij} &= q_{ij} \\ N_{ii} &= n_i \\ N_{ij} &= 0, i \neq j \end{aligned}$$

possiamo scrivere

$$B_{ij} = N_{ik} Q_{kj} = NQ = n_i q_{ij}$$

quindi

$$\begin{aligned} A_{iik} &= B_{ki} \\ A_{ijk} &= 0, i \neq j \end{aligned}$$

Poniamo

$$A_{ijk} = \delta_{ijl} B_{kl} = \delta B^T = B\delta = NQ\delta = N_{il} Q_{lm} \delta_{mjk} \quad (11)$$

dove

$$\begin{aligned} \delta_{ijl} &= 1, i = j = l \\ \delta_{ijl} &= 0 \end{aligned}$$

quindi

$$P_i = NQ\delta Y$$

Allo stesso modo trasformiamo la (6)

$$E_i = M_{ijk} Y_{kj} = MY$$

che deve essere equivalente a

$$e_{k'} = \sum_{i',j'} n_{i'} c_{i'j'k'} y_{i'j'} \quad (12)$$

da cui

$$M_{ijk} = n_k c_{kji} = N_{kl} C_{lji} = (NC)^T = C^T N^T = C^T N \quad (13)$$

dove

$$C^T = (C^T)_{ijk} = C_{kji}$$

quindi

$$E_i = C^T N Y$$

Abbiamo poi la (8)

$$v = V_i(P_i - E_i) = V \left(N Q \delta - C^T N \right) Y$$

Infine la (3) diventa

$$\sum_j t_{ij} * y_{ij} + z_{i0} = T_{ij} Y_{ij} + Z_i = T Y^T + Z = \vec{1} \quad (14)$$

1.4.1 Verifiche

Verifica (10)

$$\begin{aligned} P_1 &= A_{111} Y_{11} + A_{112} Y_{21} + A_{121} Y_{12} + A_{122} Y_{22} \\ P_2 &= A_{211} Y_{11} + A_{212} Y_{21} + A_{221} Y_{12} + A_{222} Y_{22} \\ A_{111} &= n_1 q_{11} \\ A_{112} &= n_2 q_{21} \\ A_{121} &= 0 \\ A_{122} &= 0 \\ A_{211} &= 0 \\ A_{212} &= 0 \\ A_{221} &= n_1 q_{12} \\ A_{222} &= n_2 q_{22} \\ P_1 &= n_1 q_{11} Y_{11} + n_2 q_{21} Y_{21} \\ P_2 &= n_1 q_{12} Y_{12} + n_2 q_{22} Y_{22} \end{aligned}$$

equivalente alla (9)

$$\begin{aligned} p_1 &= n_1 q_{11} y_{11} + n_2 q_{21} y_{21} \\ p_2 &= n_1 q_{12} y_{12} + n_2 q_{22} y_{22} \end{aligned}$$

cvd.

Verifica (11)

$$\begin{aligned} A_{ijk} &= \delta_{ijl} B_{lk} = \delta_{ijl} N_{km} Q_{ml} \\ A_{111} &= \delta_{111} N_{11} Q_{11} + \delta_{111} N_{12} Q_{21} \delta_{112} N_{11} Q_{12} + \delta_{112} N_{12} Q_{22} = N_{11} Q_{11} \\ A_{112} &= \delta_{111} N_{21} Q_{11} + \delta_{111} N_{22} Q_{21} \delta_{112} N_{21} Q_{12} + \delta_{112} N_{22} Q_{22} = N_{22} Q_{21} \\ A_{121} &= \delta_{121} N_{11} Q_{11} + \delta_{121} N_{12} Q_{21} \delta_{122} N_{11} Q_{12} + \delta_{122} N_{12} Q_{22} = 0 \\ A_{122} &= \delta_{121} N_{21} Q_{11} + \delta_{121} N_{22} Q_{21} \delta_{122} N_{21} Q_{12} + \delta_{122} N_{22} Q_{22} = 0 \\ A_{211} &= \delta_{211} N_{11} Q_{11} + \delta_{211} N_{12} Q_{21} \delta_{212} N_{11} Q_{12} + \delta_{212} N_{12} Q_{22} = 0 \\ A_{212} &= \delta_{211} N_{21} Q_{11} + \delta_{211} N_{22} Q_{21} \delta_{212} N_{21} Q_{12} + \delta_{212} N_{22} Q_{22} = 0 \\ A_{221} &= \delta_{221} N_{11} Q_{11} + \delta_{221} N_{12} Q_{21} \delta_{222} N_{11} Q_{12} + \delta_{222} N_{12} Q_{22} = N_{11} Q_{12} \\ A_{222} &= \delta_{221} N_{21} Q_{11} + \delta_{221} N_{22} Q_{21} \delta_{222} N_{21} Q_{12} + \delta_{222} N_{22} Q_{22} = N_{22} Q_{22} \end{aligned}$$

equivalente alla (10), cvd.

1.5 Altro

Let us normalize the problem such that

$$\begin{cases} \min C^T X \\ A_1 X = B_1 \\ A_2 X \leq B_2 \\ X \geq 0 \end{cases}$$

so

$$\begin{cases} \min(-v) = -V \left(NQ\delta - C^T N \right) Y \\ TY^T + Z = \vec{1} \\ -V \left(NQ\delta - C^T N \right) Y \leq \vec{0} \\ Y \geq 0 \\ Z \geq 0 \end{cases}$$