

## *Representation of All Possible Parenthesizations of Matrix-Chain Multiplication in a Directed Acyclic Graph*

Firstly, let's revise some notation. We define  $G$  as a directed acyclic graph (DAG) that will represent all possible parenthesizations for  $P(i, j)$  for each subproblem  $S(i, j)$  where  $1 \leq i \leq j \leq n$ . The set of vertices for this graph is:

$$V(G) = \{S(i, j) : 1 \leq i \leq j \leq n\}.$$

For  $i < j$ ,  $S(i, j)$  has exactly  $2(j - i)$  outgoing edges. For  $k = i, \dots, j - 1$  exactly two edges start from  $S(i, j)$  and end in  $S(i, k)$  and  $S(k + 1, j)$ . This is a *rigid pair* and we label these with index  $k$ .

For each vertex  $S(i, j)$ , we define the set  $P_G(i, j)$  of parenthesizations corresponding to  $S(i, j)$  in  $G$ . For  $i < j$ , let  $K_G(i, j)$  be the set of indexes of rigid pairs outgoing from  $S(i, j)$  in  $G$ , then

$$P_G(i, j) = \bigcup_{k \in K_G(i, j)} \{(p_1 \times p_2) : p_1 \in P_G(i, k), P_G(k + 1, j)\} \quad (1)$$

Let us consider the matrix chain  $A = A_1 \times A_2 \times A_3 \times A_4$ . We denote the problem of finding the parenthesizations of  $A_1 \times \dots \times A_4$  as  $S(1, 4)$ . The corresponding DAG for this problem can be seen in Figure 1. For  $i < j$  we can find the set of indexes  $K_G(i, j)$ ,

$$\begin{aligned} K_G(1, 4) &= \{1, 2, 3\} \\ K_G(1, 3) &= \{1, 2\} \\ K_G(2, 4) &= \{2, 3\} \\ K_G(1, 2) &= \{1\} \\ K_G(2, 3) &= \{2\} \\ K_G(3, 4) &= \{3\} \end{aligned}$$

These sets can then be used to find all possible parenthesizations  $P_G(i, j)$  for all subproblems  $S(i, j)$  using 1. When  $i = j$ , we know that  $P_G(i, j) = \{A_i\}$  trivially. Therefore,

$$\begin{aligned} P_G(1, 1) &= \{A_1\} \\ P_G(2, 2) &= \{A_2\} \\ P_G(3, 3) &= \{A_3\} \\ P_G(4, 4) &= \{A_4\} \end{aligned}$$

We now use 1 to find the parenthesizations for the other subproblems and then use those to find the parenthesizations for  $S(1, 4)$ , the parenthesizations  $P_G(1, 4)$  can be found using:

$$P_G(1, 4) = \{(P_G(1, 1) \times P_G(2, 4))\} \cup \{(P_G(1, 2) \times P_G(3, 4))\} \cup \{(P_G(1, 3) \times P_G(4, 4))\}$$

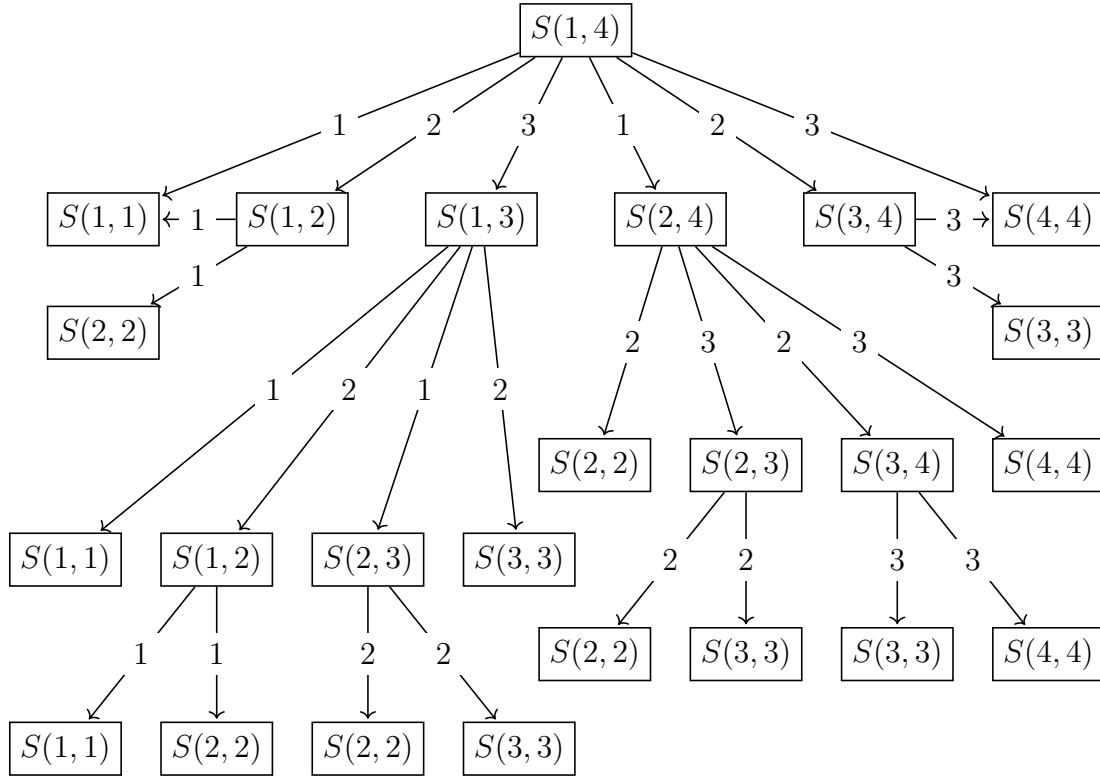


Figure 1: DAG for the matrix chain  $A_1 \times \cdots \times A_4$ .

Similarly, we can find parenthesizations for the subproblems. We first look at the simpler problems where  $i = j - 1$ , in this case  $|K_G(i, j)| = 1$ ,

$$\begin{aligned} P_G(1, 2) &= \{P_G(1, 1) \times P_G(2, 2)\} = \{(A_1 \times A_2)\} \\ P_G(2, 3) &= \{P_G(2, 2) \times P_G(3, 3)\} = \{(A_2 \times A_3)\} \\ P_G(3, 4) &= \{P_G(3, 3) \times P_G(4, 4)\} = \{(A_3 \times A_4)\} \end{aligned}$$

Next, we take a look at the more complicated cases where  $|K_G(i, j)| = 2$ ,

$$\begin{aligned} P_G(1, 3) &= \{(P_G(1, 1) \times P_G(2, 3))\} \cup \{(P_G(1, 2) \times P_G(3, 3))\} \\ &= \{(A_1 \times (A_2 \times A_3)), ((A_1 \times A_2) \times A_3)\} \\ P_G(2, 4) &= \{(P_G(2, 2) \times P_G(3, 4))\} \cup \{(P_G(2, 3) \times P_G(4, 4))\} \\ &= \{(A_2 \times (A_3 \times A_4)), ((A_2 \times A_3) \times A_4)\} \end{aligned}$$

We can now use these results to find  $P_G(1, 4)$ ,

$$\begin{aligned} P_G(1, 4) &= \{(A_1 \times (A_2 \times (A_3 \times A_4))), (A_1 \times ((A_2 \times A_3) \times A_4)), \\ &\quad ((A_1 \times A_2) \times (A_3 \times A_4)), ((A_1 \times (A_2 \times A_3)) \times A_4), (((A_1 \times A_2) \times A_3) \times A_4)\} \end{aligned}$$