

Representation of All Possible Parenthesizations of Matrix-Chain Multiplication in a Directed Acyclic Graph

Firstly, let's revise some notation. We define G as a directed acyclic graph (DAG) that will represent all possible parenthesizations for $P(i, j)$ for each subproblem $S(i, j)$ where $1 \leq i \leq j \leq n$. The set of vertices for this graph is:

$$V(G) = \{S(i, j) : 1 \leq i \leq j \leq n\}.$$

For $i < j$, $S(i, j)$ has exactly $2(j - i)$ outgoing edges. For $k = i, \dots, j - 1$ exactly two edges start from $S(i, j)$ and end in $S(i, k)$ and $S(k + 1, j)$. This is a *rigid pair* and we label these with index k .

For each vertex $S(i, j)$, we define the set $P_G(i, j)$ of parenthesizations corresponding to $S(i, j)$ in G . For $i < j$, let $K_G(i, j)$ be the set of indexes of rigid pairs outgoing from $S(i, j)$ in G , then

$$P_G(i, j) = \bigcup_{k \in K_G(i, j)} \{(p_1 \times p_2) : p_1 \in P_G(i, k), P_G(k + 1, j)\} \quad (1)$$

Let us consider the matrix chain $A = A_1 \times A_2 \times A_3 \times A_4$. We denote the problem of finding the parenthesizations of $A_1 \times \dots \times A_4$ as $S(1, 4)$. The corresponding DAG for this problem can be seen in Figure 1. For $i < j$ we can find the set of indexes $K_G(i, j)$,

$$\begin{aligned} K_G(1, 4) &= \{1, 2, 3\} \\ K_G(1, 3) &= \{1, 2\} \\ K_G(2, 4) &= \{2, 3\} \\ K_G(1, 2) &= \{1\} \\ K_G(2, 3) &= \{2\} \\ K_G(3, 4) &= \{3\} \end{aligned}$$

These sets can then be used to find all possible parenthesizations $P_G(i, j)$ for all subproblems $S(i, j)$ using 1. When $i = j$, we know that $P_G(i, j) = \{A_i\}$ trivially. Therefore,

$$\begin{aligned} P_G(1, 1) &= \{A_1\} \\ P_G(2, 2) &= \{A_2\} \\ P_G(3, 3) &= \{A_3\} \\ P_G(4, 4) &= \{A_4\} \end{aligned}$$

We now use 1 to find the parenthesizations for the other subproblems and then use those to find the parenthesizations for $S(1, 4)$, the parenthesizations $P_G(1, 4)$ can be found using:

$$P_G(1, 4) = \{(P_G(1, 1) \times P_G(2, 4))\} \cup \{(P_G(1, 2) \times P_G(3, 4))\} \cup \{(P_G(1, 3) \times P_G(4, 4))\}$$

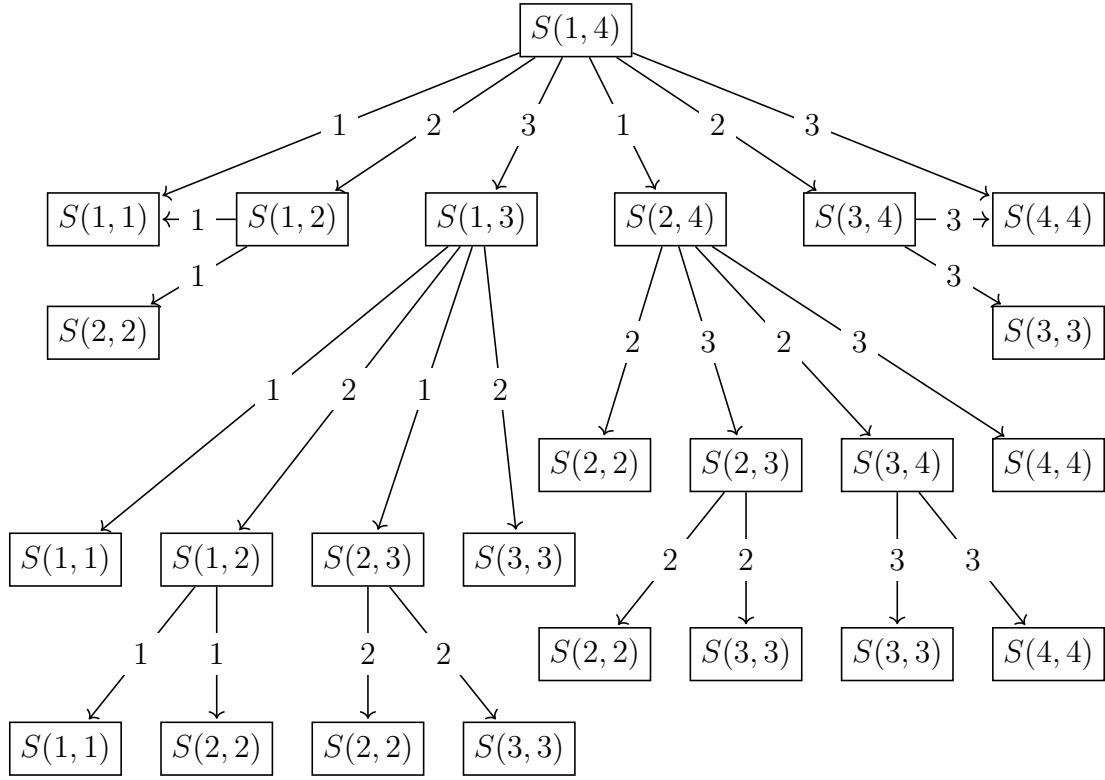


Figure 1: DAG for the matrix chain $A_1 \times \cdots \times A_4$.

Similarly, we can find parenthesizations for the subproblems. We first look at the simpler problems where $i = j - 1$, in this case $|K_G(i, j)| = 1$,

$$\begin{aligned} P_G(1, 2) &= \{P_G(1, 1) \times P_G(2, 2)\} = \{(A_1 \times A_2)\} \\ P_G(2, 3) &= \{P_G(2, 2) \times P_G(3, 3)\} = \{(A_2 \times A_3)\} \\ P_G(3, 4) &= \{P_G(3, 3) \times P_G(4, 4)\} = \{(A_3 \times A_4)\} \end{aligned}$$

Next, we take a look at the more complicated cases where $|K_G(i, j)| = 2$,

$$\begin{aligned} P_G(1, 3) &= \{(P_G(1, 1) \times P_G(2, 3))\} \cup \{(P_G(1, 2) \times P_G(3, 3))\} \\ &= \{(A_1 \times (A_2 \times A_3)), ((A_1 \times A_2) \times A_3)\} \\ P_G(2, 4) &= \{(P_G(2, 2) \times P_G(3, 4))\} \cup \{(P_G(2, 3) \times P_G(4, 4))\} \\ &= \{(A_2 \times (A_3 \times A_4)), ((A_2 \times A_3) \times A_4)\} \end{aligned}$$

We can now use these results to find $P_G(1, 4)$,

$$\begin{aligned} P_G(1, 4) &= \{(A_1 \times (A_2 \times (A_3 \times A_4))), (A_1 \times ((A_2 \times A_3) \times A_4)), \\ &\quad ((A_1 \times A_2) \times (A_3 \times A_4)), ((A_1 \times (A_2 \times A_3)) \times A_4), (((A_1 \times A_2) \times A_3) \times A_4)\} \end{aligned}$$