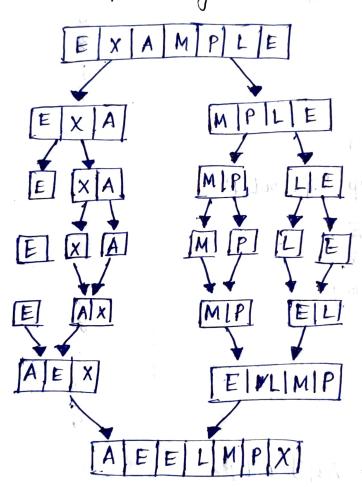
20 10 20 Q No.

HOMEWOR _ (Divide and conquer, Greedy approach)

1. Apply merge sort to sort the list E,X,A,M,P,L,E in alphabetical order. Check if the algorithm is stable.

Ans



In place-doesn't use extra space.

The algorithm is stable because the 1st E in EXAMPLE is before the 2nd E (EXAMPLE)
A spetting algorithm is said to be stable if two objects with equal keys appear in the same order in sorted output as they appear in the input array to be spetch.

2. Apply quick sort to sort the list A, B, C, D, E in alphabetical order. Check the number of comparisons by choosing (1) the first element as the pirot (11) last element as the pirot.

```
(9) ABEDE
                         No of compaisions
   p= A i= 0 j=4
     B>A go to j
  1=1
  E4 EYA
  j=3 D>A
  jed C>A
  j=1 8>A
  j=0 A > A (false)
  Since i 75 swap AGJ and P
  B C D E

P=8 i=0 j=3
  i=1 C>B go to s
  1°=3 €78
  j=2 D>B
  j=81 C>B
  j=o B=B
  Since o'75 Swap A GiJ and P
HIMA & DE
  p=c i=0 j=2
  i=1 D>C go to j
j=2 E>C
                          13
  J=1 DYC POM DE OF PROBE
                          14
  j=0 C5 e
  since it j swap Atj J and P.
```

```
(=1 E>D go to j
                                         16
   J=1 E>D
                                          17
                                          18
   j=0 D \leq D
   swap ALJJ and P
   E - no companision (1 element)
No. of companisions = 18.
         ABCDE
                                      No. of compairsions
     P=E i=0 1=4
PEO ASE
is B LE
1=2 CXE
1=3 DXE
1=4 E < E
    i out of bounds
   Swap ATJ and P.
     A B C D
    P=D 1=0 j=3.
i=o A & D
 1=1
     0 \leqslant 0
1=2
    CSD
 C=3
     D \leq D
     i out of bounds
     swap AGI and P
     A B C
P=C l=0 j=2
     AS C
      B & C
     cse
```

i out of bounds gwap A[f] and P

A B

$$f = B = 0 = 0 = 1$$

i=0 A & BB

i=0 wt of bounds
Swap A[f] and P.

B - A - single element (no compasission)

No of compasissions = 14.

3. Apply strassen's algorithm to compute

$$\begin{bmatrix}
1 & 0 & 1 \\
A & 1 & 1 & 0 \\
0 & 1 & 3 & 0 \\
5 & 0 & 2 & 1
\end{bmatrix}

= \begin{bmatrix}
0 & 1 & 0 & 1 \\
2 & 1 & 0 & 4 \\
2 & 0 & 1 & 1 \\
1 & 3 & 5 & 0
\end{bmatrix}

And

$$a_{00} = \begin{bmatrix}
1 & 0 \\
A & 1
\end{bmatrix}

= \begin{bmatrix}
0 & 1 & 0 & 1 \\
2 & 1 & 0 & 4 \\
2 & 0 & 1 & 1 \\
1 & 3 & 5 & 0
\end{bmatrix}

= \begin{bmatrix}
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0 &$$$$

$$= 4 * 1 = 4$$

$$m_{\ell} = (a_{10} - a_{00}) * (b_{00} + b_{01})$$

$$= 4 * 1 = 4$$

$$m_{1} = (a_{01} - a_{11}) * (b_{10} + b_{11})$$

$$= 0 * 3 = 0$$

$$m_{1} + m_{4} - m_{5} + m_{7} = 4 + 2 - 4 + 0 = 2$$

$$m_{3} + m_{5} = 0 + 4 = 4$$

$$m_{1} + m_{3} - m_{1} + m_{6} = 4 + 0 - 0 + 4 = 8$$

$$m_{1} = a_{00} * (b_{01} - b_{11})$$

$$= \begin{bmatrix} 1 & 0 \\ 4 & 1 \end{bmatrix} * \begin{bmatrix} 0 & 1 \\ 2 & 8 \end{bmatrix}$$

$$m_{1} = (a_{00} + a_{11}) * (b_{00} + b_{01})$$

$$= 2 * 3 = 6$$

$$m_{2} = (a_{10} + a_{11}) * (b_{00} + b_{01})$$

$$= 1 * (-4) = -4$$

$$m_{4} = a_{11} * (b_{10} - b_{00})$$

$$= 1 * (-4) = -4$$

$$m_{5} = (a_{00} + a_{01}) * (b_{10} + b_{01})$$

$$= 1 * 4 = 4$$

$$m_{6} = (a_{10} - a_{00}) * (b_{00} + b_{01})$$

$$= 3 * -1 = -3$$

$$m_{7} = (a_{01} - a_{11}) * (b_{10} + b_{11})$$

$$= -1 * -1 = 1$$

```
M_1 + M_4 - M_5 + M_7 = 6 + -4 - 4 + 1 = -1
    M_3 + M_5 = -4 + 4 = 0
   m2+ m4 = -5+-4=-9
   m_1 + m_3 - m_2 + m_6 = 6 + 14 - (-5) + (-3) = 4
  m4 = Qu & ( b10 - b00)
            = \begin{bmatrix} 3 & 0 \\ 2 & 1 \end{bmatrix} * \begin{bmatrix} 2 & 0 \\ 1 & 3 \end{bmatrix} - \begin{bmatrix} 0 & 1 \\ 2 & 1 \end{bmatrix}
               = \[ \frac{30}{21} \\ \frac{1}{-1} \\ \frac{2}{000} \\ \frac{1}{000} \\ \f
       M1 = ( a00 + a11) * (600 + b11)
                    = 4 \times 4 = 16
       m2 = (a10 + a11) * 600
                        z 3 x 2 = 6
      m_3 = 900 \times (601 - 64)
= 3 \times -3 = -9
       my = 911 * (b10 - 600) 100 1
                   = 1 \times (-3) = -3
      m5 = Lano + 201) x 611
                    = 3x2 = 61
    m6 = (a10 - a00) * (b00 + b01)
                   z - 1 + 1 = -1
  m= (a01 - a11) + (b10 + b11)
                 = -1 \times 1 = -1
m1+m4-m5+m7=16-3-6-1=6
 m3+ m52 -3
 m_1 + m_3 = 3
m_1 + m_3 = -m_2 + m_6 = 0
m_1 + m_3 = -m_2 + m_6 = 0
```

 \geq 0 x 4 = 0

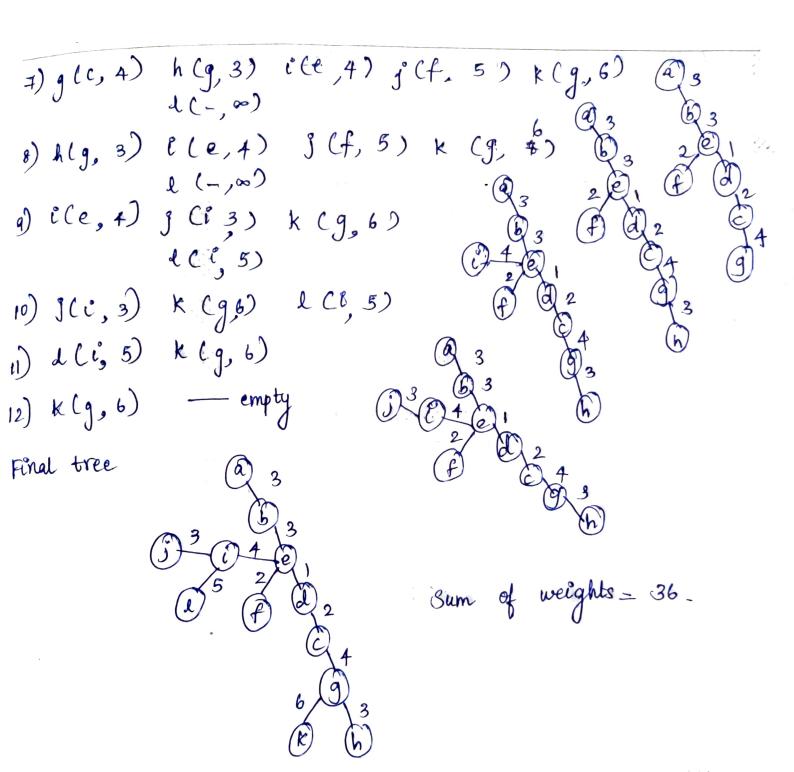
4. Find the minimum spanning tree of the graph given below using
(i) Prims algorithm

b (a,3)
e (a,5)
d (a,4)
e (:,
$$\infty$$
)
f (-, ∞)
h (-, ∞)
i (-, ∞)
j (-, ∞)
 λ (-, ∞)

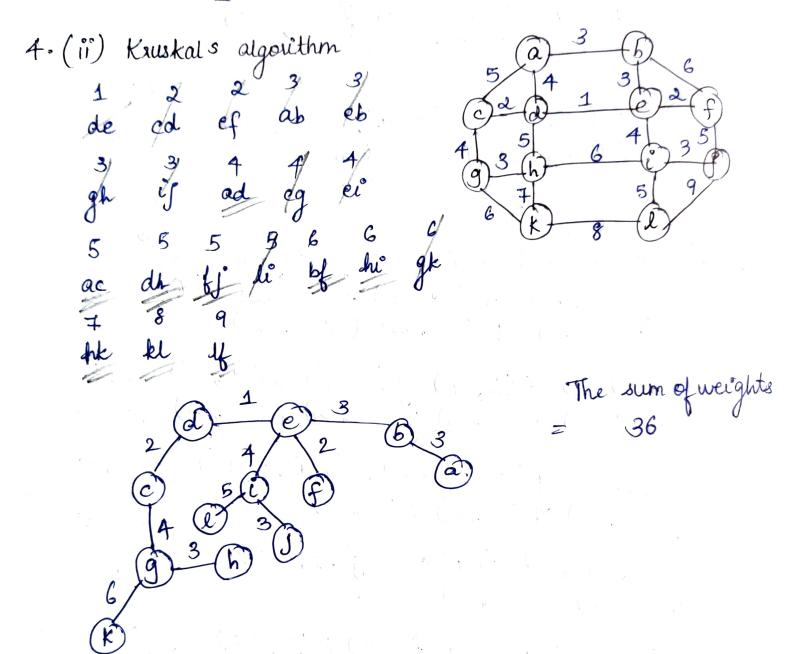
- 2) bla, 3) c (a,5) d (a1,4) e (b,3) (2) 3

 f (b,6) g (-,0) h (-,0) k (-,0) b

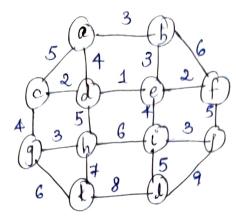
 i (-,0) J (-,0) k (-,0) b
- 3) ecb, 3) c(a,5) d(e,1) f(e,2) g(_,0) h(-,0) i(e,4) j(_,0) k(_,0) l(-,0)
- 4) a(e, 1) c(d, 2) f(p, 2) g(-, \omega) h(d, 5) 8 ce, 4) j(-, \omega) k(-, \omega) d(-, \omega)
- 5) c(d,2) f(e,2) g(c,4) h(d,5)i(e,4) g(-,0)k(-,0) d(-,0)
- 6) fle, 2) g(c,4) h(d,5) l(e,4) j(f,5) k,(-,0) l(-,0)



WS-4 Continuation



5. Find the shortest path in the graph given below considering the source vertex as a using dighter is algorithm.



(a) ac-, 0) b(a,3) c(a,5) d(a,4) (c)
$$e(-,\infty)$$
 f(-,\infty) g(-,\infty) $f(-,\infty)$ $f(-,\infty)$ $f(-,\infty)$ $f(-,\infty)$ $f(-,\infty)$

b(a, 3)
$$e(a, 5) d(a, 4) e(b, 6)$$

 $f(b, 9) g(-, \infty) h(c-, \infty)$
 $i(-, \infty) f(-, \infty) k(-, \infty)$
 $l(-, \infty)$

$$d(a,4) c(a,5) e(d,5) f(b,9)$$

$$g(-,\infty) h(d,9) c(-,\infty)$$

$$g(-,\infty) k(-,\infty) l(-,\infty)$$

c(a, 5)
$$e(d, 5)$$
 $f(b, 9)$ $g(c, 9)$ $h(d, 9)$ $f(c, 9$

$$f(e, 7) = \begin{cases} (c, 9) & (c, 9)$$

$$l(-9\%)$$
 $l(-9\%)$
 $l(-9\%)$

SCf, 12) K(g, 15) l(i, 14) k(g, 15)

K(g, 15)

Homework - 5 Dynamic Programming.

1. Apply the bottom-up dynamic programming algorithm to the following instance of the knapsack problem.

item	weight	value	r c	
1	3	\$25		
2	2	\$ 20	canacitu	W= 6.
3	1	\$15	capacity	102 0,
4	4	\$40	,	
5	5	\$ 50		

			ì	1			i .
i/j°	0	1	2	3	4	5	6
0	0	0	0	0	0	Ø	O
	0	0	0	25	25	25	25
2	O	Ô	20	25	25	45	45
3	0	15	20	85	40	45	60
4	0	15	20	35	40	55	60
5	O	15	20	35	40	55	65

$$F(i,j) = \begin{cases} \max \left(F(i-1,j), \forall i + F(i-1,j-w_i) \right) & \text{if } j-w_i \geqslant 0 \\ F(i-1,j) & \text{if } j-w_i < 0 \end{cases}$$

$$F(1,3)$$
 3-3>0 max(F(0,3), 25+F(0,0))=25

$$F(1,4)$$
 max(F(0,4), 25+F(0,1)) = 25

$$F(1,6)$$
 max($F(0,6)$, $25+F(0,8)$) = 25

$$F(2,3)$$
 max $(F(1,3), 20 + F(1,1)) = 25$

$$F(a,4)$$
 max($F(1,4)$, $a0+F(1,2)$) = 25

$$F(2,5)$$
 max($F(1,5)$, 20 + $F(1,3)$) = 45

```
max (F(2,1), 15+ F(2,0)) = 15
   F(3,1)
         max (F(2,2), 15+F(2, U) = 20
  f(3,2)
         \max(F(2,3), 15 + F(2,2)) = 35
  F(3,5)
         \max(F(2,4), 19+F(2,3)) = 40
 ((3,4)
         max ( F(25), 15+ F(2,4)) = 45
 F(3, 5)
         \max(F(2,6), 15+F(2,5)) = 60
 رو رو<sub>ا ۲</sub>
        F(3,1) = 15
 FL4, 1)
        F(3,2) = 20
 F(4,2)
        F(3,3) 235
 F(4, 3)
        \max(F(3,4), 40 + F(3,0)) = 40
 FC4, 4)
        \max(F(3,5), 40 + F(3,0)) = 55
 F(4,5)
        \max(F(3,6), 40 + F(3,2)) = 60
 F(A16)
F(5,1) F(4,1) = 15
· F(5,2) F(4,2) = 20
f(5,3) F(4,3) = 35
f(5,4) F(4,4) = 40
f(5,5) max (F(4,5), 50+F(4,0)) = 55
F(5,6) max(F(4,6), 50+ F(4,1)) = 65

⇒ Weight = 6. Value = 65.

  comparing F(5, 6) 4 F(4, 6) 65 > 60 =>, 5th element included
→ Weight = 1 Value = 15.
  Comparing F(4,1) & F(3,1) 15=15 > 4th element not included.
  Comparing F(3,1) 4 F(2,1) 15>0 > 3rd element included
⇒ Weight = 0 Value >0
Atu: Value=65 Items: 3,5.
```

Apply Warshall's algorithm to find the transitive closure of the digraph defined by the following adjacency matrix:

$$R^{\circ} = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$R^{1} = \begin{bmatrix} 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

$$R^{3} = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$R^{2} = \begin{bmatrix} 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix} \qquad R^{4} = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$R^{(k)}[i,j] \leftarrow R^{(k-1)}[i,j]$$
 or $(R^{(k-1)}[i,k])$ and $R^{(k-1)}[k,j]$

$$R^{3} = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

$$R^{4} = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

Solve the all-pains shortest-path problem for the digraph with the following weight matrix

$$D^{\circ} = \begin{bmatrix} 0 & 2 & \infty & 1 & 8 \\ 6 & 0 & 3 & 2 & \infty \\ \infty & \infty & 0 & 4 & \infty \\ \infty & \infty & 2 & 0 & 3 \\ 3 & \infty & \infty & \infty & 0 \end{bmatrix} D \begin{bmatrix} i, j \end{bmatrix} \leftarrow \min \left\{ D \begin{bmatrix} i, j \end{bmatrix} \right\}$$

$$D \begin{bmatrix} i, k \end{bmatrix} + D \begin{bmatrix} k, j \end{bmatrix} \right\}$$

$$D^{\downarrow} = \begin{bmatrix} 0 & 2 & 3 & 1 & 4 \\ 6 & 0 & 3 & 2 & 5 \end{bmatrix}$$

$$D^{2} = \begin{bmatrix} 0 & 2 & 5 & 1 & 8 \\ 6 & 0 & 3 & 2 & 14 \\ 8 & 8 & 0 & 4 & 8 \\ 8 & 8 & 8 & 4 & 0 \end{bmatrix}$$

$$D^{3} = \begin{bmatrix} 0 & 2 & 5 & 1 & 8 \\ 6 & 0 & 3 & 2 & 14 \\ 8 & 0 & 0 & 4 & 0 \\ 8 & 0 & 0 & 3 & 0 \\ 8 & 0 & 0 & 3 & 0 \\ 3 & 5 & 8 & 4 & 0 \end{bmatrix}$$

$$D^{4} = \begin{bmatrix} 0 & 2 & 3 & 1 & 4 \\ 6 & 0 & 3 & 2 & 5 \\ \infty & \infty & 0 & 4 & 7 \\ \infty & \infty & 2 & 0 & 3 \\ 0 & 5 & 6 & 4 & 0 \end{bmatrix}$$

$$D^{5} = \begin{bmatrix} 0 & 2 & 3 & 1 & 4 \\ 6 & 0 & 3 & 2 & 5 \\ 10 & 12 & 0 & 4 & 7 \\ 6 & 8 & 2 & 0 & 3 \\ 3 & 5 & 6 & 40 \end{bmatrix}$$