# Caffeine - The Habit-Forming Compiler

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#### What is Caffeine?

An interpreter and x86 compiler for the Habit language.

#### What is Habit?

#### Habit is ...

- Pure
- Functional

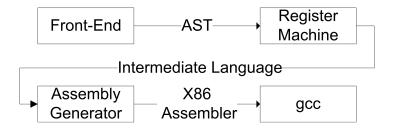
#### And Caffeine says it is . . .

- Syntactically identical to Haskell
- Strict

#### What Can Caffeine Do?

- Parameterized algebraic data types
- Higher-order functions, partial application
- Pattern-matching, conditionals, case analysis, guards, etc.
- Global and local definitions
- Mutually recursive definitions (mostly)

## Pipeline



# Register Machine

- Runs programs in "Intermediate Language" (IL); i.e., an interpreter.
- Simple easy to map to x86.
- Two special registers clo and arg.
- IL used as input to x86 compiler.

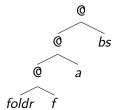
## x86 Compiler

- Produces x86 assembly from IL source.
- Relies on simple runtime for allocation support.
- Produces a result by dumping heap value produced by main function.

#### Consider

foldr f a bs

The front-end gives this AST:



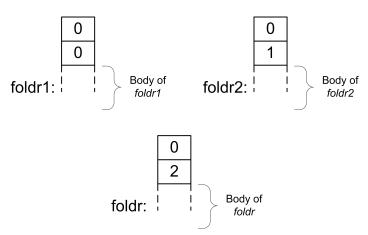
Three applications means three functions:

foldr f a bs 
$$\equiv$$
 ((foldr1 f) a) bs

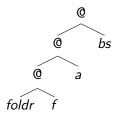
$$foldr1 \ f = foldr2 \ f$$
  
 $foldr2 \ f \ a = foldr \ f \ a$   
 $foldr \ f \ a \ bs = ...$ 

Register Machine representation:

#### Assembly representation:



foldr f a bs in IL. Recall:



Enter implements application. Top-level evaluation:

Enter foldr1 f result0 Enter result0 a result1 Enter result1 bs result2

- foldr1 Capture f.
- result0 Capture a.
- result1 Evalute foldr with bs.
- result2 Result of evaluation of foldr f a bs.

```
foldr1 captures f:
    AllocC result0 foldr2 1
    Store arg (result0, 0) -- arg special.
    Ret_result0
foldr2 copies f and captures a:
    Load (clo, 0) f -- clo special.
    AllocC result1 foldr 2
    Store f (result1,0)
    Store arg (result1, 1)
    Ret result1
```

```
foldr is:
     foldr\ f\ a\ Nil = a
     foldr f a (Cons b bs) = f b (foldr f a bs)
Check Nil first:
     Load (clo, 0) f
     Load (clo, 1) a
     FailT arg Nil lab28
     Ret a
```

Next, check Cons:

lab28 : FailT arg Cons lab32

Evaluate foldr f a bs and store in reg37.

Load (arg, 0) b Load (arg, 1) bs Enter foldr1 f reg35 Enter reg35 a reg36 Enter reg36 bs reg37

Finally, evaluate f a reg37.

Enter f b reg38 Enter reg38 reg37 result2

...

Ret result2

#### Foldr in Assembler

Three key instructions from IL to explore

- Enter Application.
- FailT Conditionals/Discrimination.
- AllocC Allocation.

### Enter in Assembler

```
# "Enter foldr1 f result0
pushl %esi
pushl %edi
movl foldr1, %edi
movl $f, %esi
call *(%edi)
```

#### FailT in Assembler

```
# "FailT arg Nil lab28
movl (%esi), %ecx
movl (%ecx), %ecx
cmpl $0x348, %ecx
jnz lab28
# "FailT arg Cons lab32
movl (%esi), %ecx
movl (%ecx), %ecx
cmpl $0x3b8, %ecx
jnz lab32
```

#### AllocC in Assembler

```
Recall foldr1 captures f and points to foldr2:
    AllocC result0 foldr2 1
    Store arg (result0,0) -- arg special.
    . . .
In assembler:
     # AllocC result0 foldr2 1
     pushl $foldr2
     call _alloc
     addl $0x4, %esp
     movl %eax, -4(%ebp)
     # Store arg (result0,0)
     movl %esi, %edx
     movl -4(\%ebp), \%ecx
```

movl %edx, 4(%ecx)

### AllocC in Assembler

Preceding \$foldr2 we find the *info table* for the closure:

```
.long 0x0
.long 0x1
foldr2:
```

# Composability

"Compiler as Service"

# Composable Signatures

• Register Compiler: Module/[Group]

 $compile :: Supply Int \rightarrow Module \rightarrow [Group]$ 

• Assembly Compiler: [Group]/Program<sup>1</sup>

assemble :: Supply Int  $\rightarrow$  [Group]  $\rightarrow$  Program



<sup>&</sup>lt;sup>1</sup>Dirty secret –  $Program \equiv [String]$ .

#### What it Does Not Do

- Primitive types (Int, String, etc.)
- Typeclasses
- Input/Output
- Multi-module programs
- Garbage Collection

## What it Could Do

All of the above ...

**REPL** 

Debugger

Staged compilation

#### Where to Find Out More

Code available in cg509. The README will tell you what to do.

# Samples

- tests\Even.hs
- tests\Fib.hs
- tests\Foldr2.hs